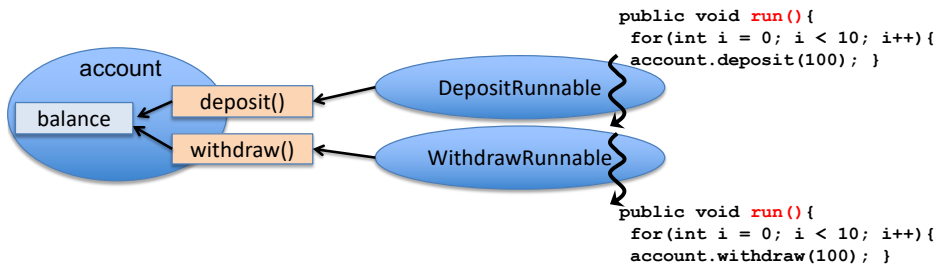


# Thread Safety Issues

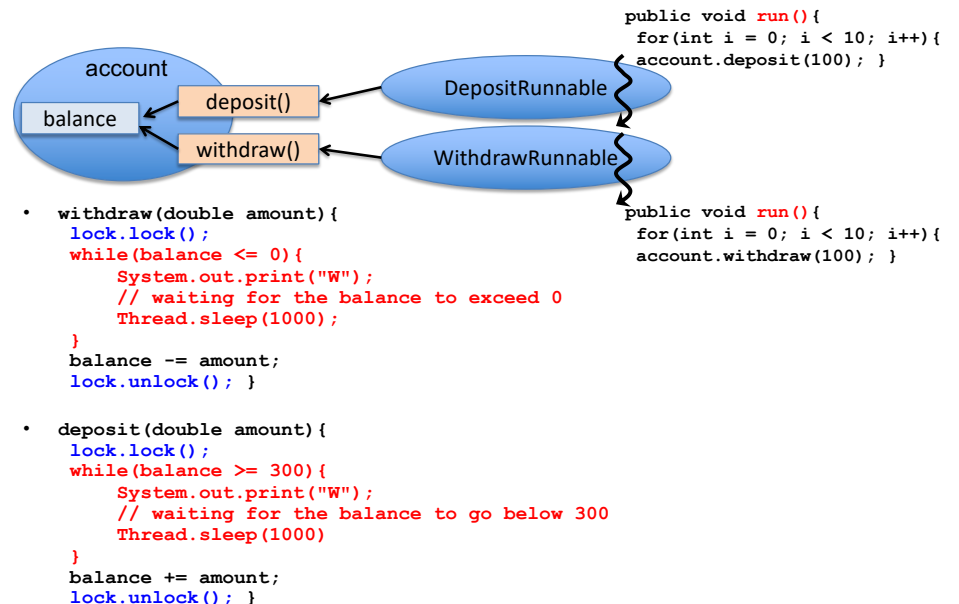
- Race conditions
- Deadlocks
- Thread-safe code is free from both race conditions and deadlocks.

## Deadlock

### DeadlockedBankAccount.java

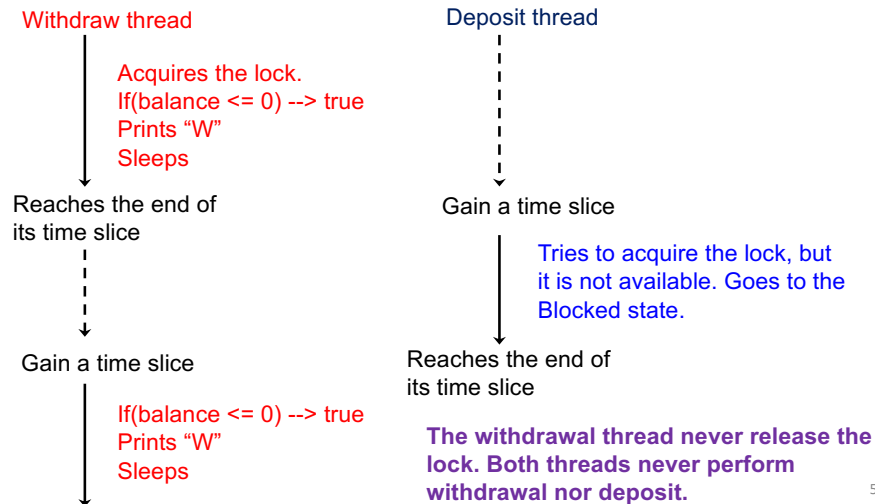


### DeadlockedBankAccount.java



# How Can a Deadlock Occur?

- Suppose the withdrawal thread goes ahead.



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## Note

- A JVM can perform context switches even when a thread runs atomic code.
  - A lock guarantees that only one thread exclusively runs atomic code at a time.
    - It does NOT control when to (or when not to) perform context switches.
  - Some resources explicitly/implicitly say that context switches never occur when a thread runs atomic code.
    - It is WRONG!

## DeadlockedBankAccount2.java

- Previous version

```
- withdraw(double amount){
    lock.lock();
    while( balance <= 0 ){
        System.out.print("W");
        Thread.sleep(1000);
    }
    balance -= amount;
    lock.unlock();
}

- deposit(double amount){
    lock.lock();
    while( balance >= 300 ){
        System.out.print("W");
        Thread.sleep(1000);
    }
    balance += amount;
    lock.unlock();
}
```

- New version

```
- withdraw(double amount){
    while( balance <= 0 ){
        System.out.print("W");
        Thread.sleep(1000);
    }
    lock.lock();
    balance -= amount;
    lock.unlock();
}

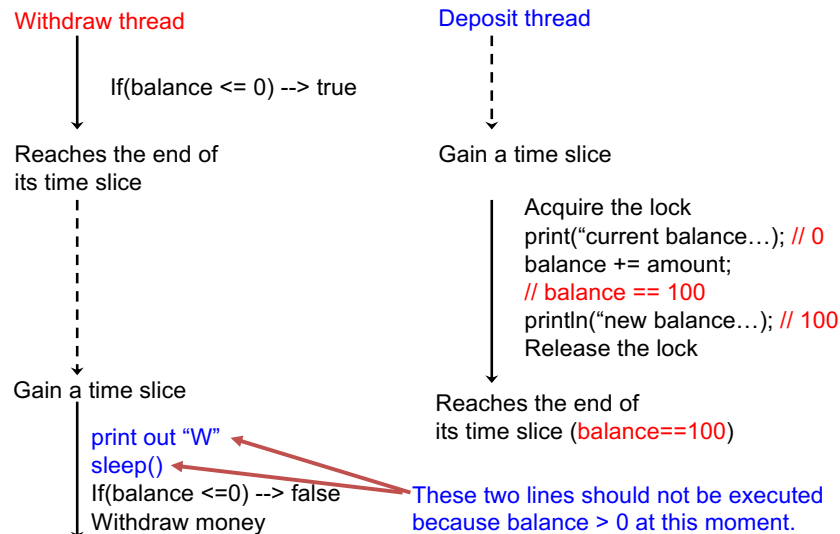
- deposit(double amount){
    while( balance >= 300 ){
        System.out.print("W");
        Thread.sleep(1000);
    }
    lock.lock();
    balance += amount;
    lock.unlock();
}
```

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- Has no deadlock problems.
- Can generate race conditions.

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## A Potential Race Condition in DeadlockedBankAccount2



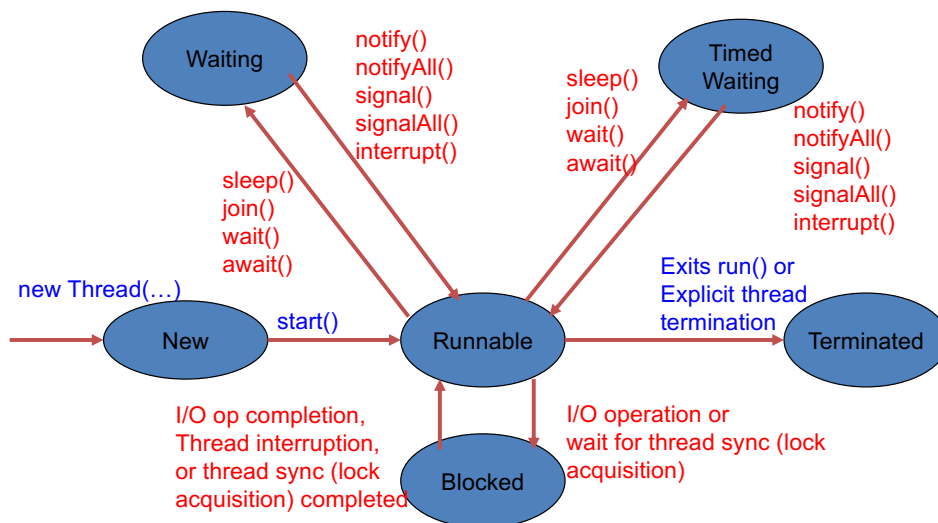
9

## Avoiding Deadlocks and Race Conditions

- Use a **Condition** object(s).
  - Allows a thread to
    - Temporarily release a lock so that another thread can acquire it and proceed.
    - Re-acquire the lock later.
- `java.util.concurrent.locks.Condition`
  - Use a lock to obtain a Condition instance
    - ```
ReentrantLock lock = new ReentrantLock();
Condition condition = lock.newCondition(); // Factory method
condition.await(); // Temporarily releases the lock
// Goes to the Waiting state until getting
// signaled.
```

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## States of a Thread



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## ThreadSafeBankAccount2.java

- ```
Condition sufficientFundsCondition = lock.newCondition();
Condition belowUpperLimitFundsCondition = lock.newCondition();
```
- ```
withdraw(double amount){
    lock.lock();
    while(balance <= 0){
        // Wait for the balance to exceed 0
        sufficientFundsCondition.await();
    }
    balance -= amount;
    belowUpperLimitFundsCondition.signalAll();
    lock.unlock();
}
```
- ```
deposit(double amount){
    lock.lock();
    while(balance >= 300){
        // Wait for the balance to go below 300.
        belowUpperLimitFundsCondition.await();
    }
    balance += amount;
    sufficientFundsCondition.signalAll();
    lock.unlock();
}
```

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# ThreadSafeBankAccount2.java

```
• Condition sufficientFundsCondition = lock.newCondition();
  Condition belowUpperLimitFundsCondition = lock.newCondition();

• withdraw(double amount){
  lock.lock();
  while(balance <= 0){
    // Wait for the balance to exceed 0
    sufficientFundsCondition.await();
  }
  balance -= amount;
  belowUpperLimitFundsCondition.signalAll();
  lock.unlock(); }

• deposit(double amount){
  lock.lock();
  while(balance >= 300){
    // Wait for the balance to go below 300.
    belowUpperLimitFundsCondition.await();
  }
  balance += amount;
  sufficientFundsCondition.signalAll();
  lock.unlock(); }
```

A "deposit" thread calls signalAll() to wake up a thread(s) that is/are waiting until balance > 0.

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# ThreadSafeBankAccount2.java

```
• Condition sufficientFundsCondition = lock.newCondition();
  Condition belowUpperLimitFundsCondition = lock.newCondition();

• withdraw(double amount){
  lock.lock();
  while(balance <= 0){
    // Wait for the balance to exceed 0
    sufficientFundsCondition.await();
  }
  balance -= amount;
  belowUpperLimitFundsCondition.signalAll();
  lock.unlock(); }

• deposit(double amount){
  lock.lock();
  while(balance >= 300){
    // Wait for the balance to go below 300.
    belowUpperLimitFundsCondition.await();
  }
  balance += amount;
  sufficientFundsCondition.signalAll();
  lock.unlock(); }
```

A "withdraw" thread calls signalAll() to wake up a thread(s) that is/are waiting until balance < 300.

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# ThreadSafeBankAccount2.java

```
• Condition sufficientFundsCondition = lock.newCondition();
  Condition belowUpperLimitFundsCondition = lock.newCondition();

• withdraw(double amount){
  lock.lock();
  while(balance <= 0){
    // Wait for the balance to exceed 0
    sufficientFundsCondition.await();
  }
  balance -= amount;
  belowUpperLimitFundsCondition.signalAll();
  lock.unlock(); }

• deposit(double amount){
  lock.lock();
  while(balance >= 300){
    // Wait for the balance to go below 300.
    belowUpperLimitFundsCondition.await();
  }
  balance += amount;
  sufficientFundsCondition.signalAll();
  lock.unlock(); }
```

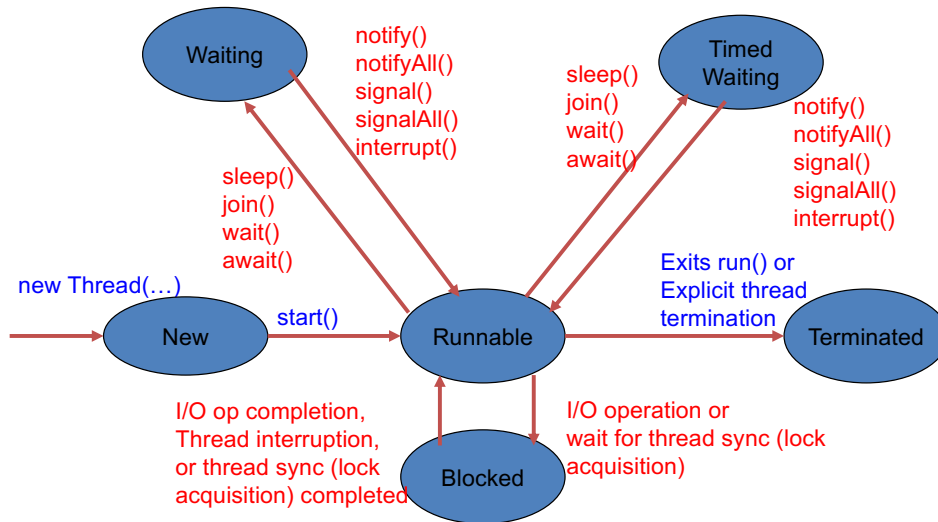
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# Condition

- **await()**
  - Will wait until it is signaled or interrupted
  - Will wait until it is signaled or interrupted, or until a specified waiting time (relative time) elapsed.
  - Will wait until it is signaled or interrupted, or until a specified deadline (absolute time).
  - goes to the Runnable state and re-acquires a lock, if signaled.
    - Will be "blocked" if the thread re-acquisition fails.
  - Throws an **InterruptedException**, if interrupted.
    - c.f. A previous lecture note on thread interruption
- **signalAll()**
  - Wakes up all waiting threads on a condition object.
    - All of them go to the "runnable" state.
    - One of them will re-acquire a lock.

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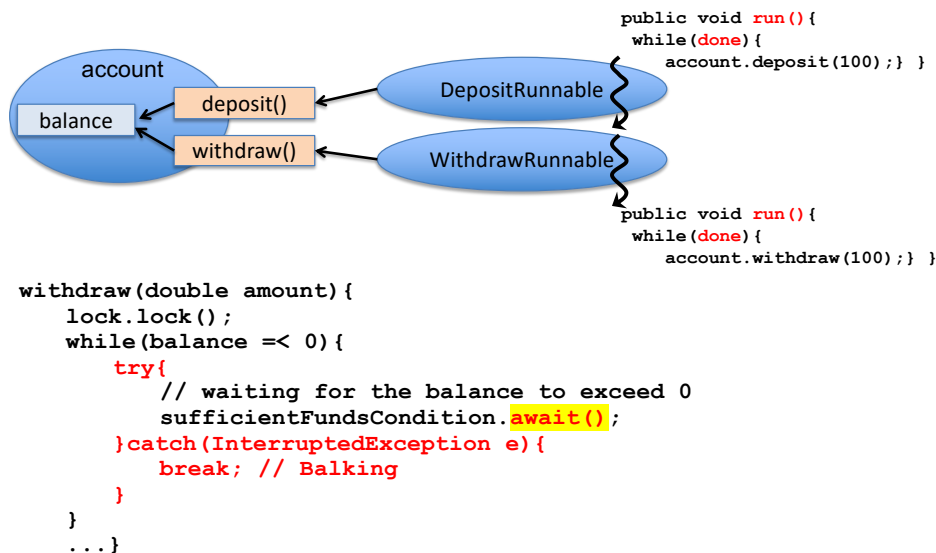
# States of a Thread



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- When a thread calls `await()`, `signalAll()` or `signal()` on a Condition object,
  - the thread is assumed to hold a lock associated with the Condition object.
  - If the thread does not, an `IllegalMonitorStateException` is thrown.

## 2-Step Thread Termination



## HW 14

- Implement 2-step termination for “deposit” and “withdraw” threads.
  - Have the main thread call `interrupt()` on “deposit” and “withdraw” threads
    - To let those threads to wake up in case they are in the Waiting state due to `await()` or `sleep()`.
- Due: April 18 (Thu) midnight

# signalAll() Before or After a State Change?

- ```

withdraw(double amount){
    lock.lock();
    while(balance <= 0){
        // waiting for the balance to exceed 0
        sufficientFundsCondition.await(); }
        balance -= amount;
        belowUpperLimitFundsCondition.signalAll();
    lock.unlock(); }

```
- ```

deposit(double amount){
    lock.lock();
    while(balance >= 300){
        // waiting for the balance to go below 300.
        belowUpperLimitFundsCondition.await(); }
        balance += amount;
        sufficientFundsCondition.signalAll();
    lock.unlock(); }

```
- What if you call signalAll() first and then update the balance? Will any thread safety issues come out?

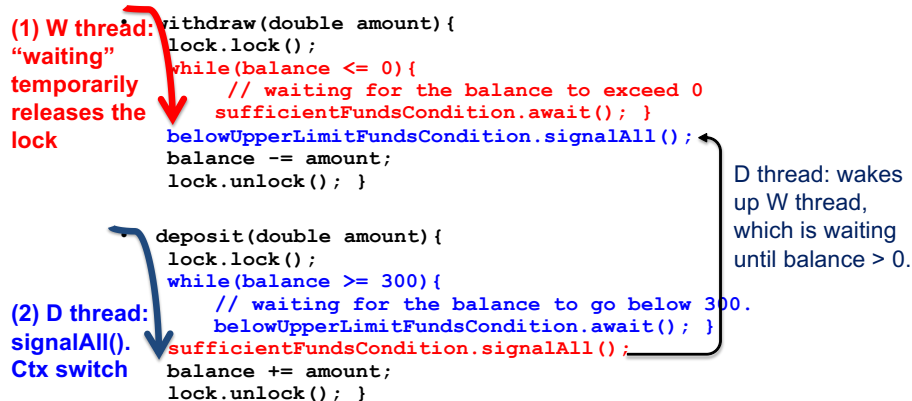
- ```

withdraw(double amount){
    lock.lock();
    while(balance <= 0){
        // waiting for the balance to exceed 0
        sufficientFundsCondition.await(); }
        belowUpperLimitFundsCondition.signalAll();
        balance -= amount;
    lock.unlock(); }

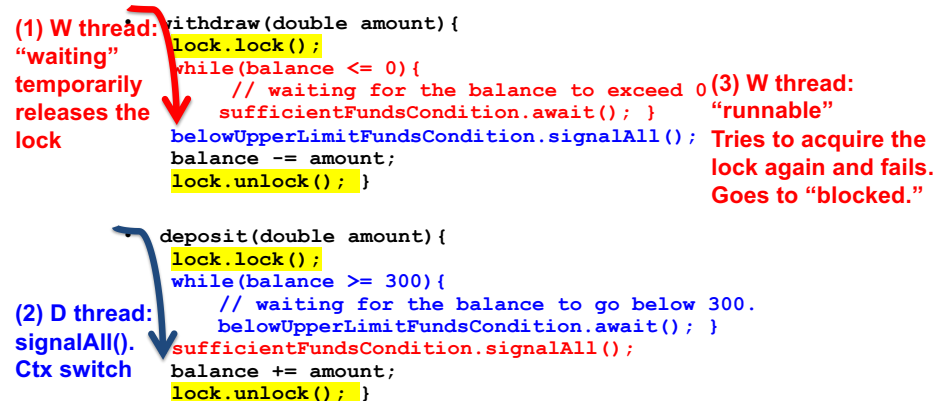
```
- ```

deposit(double amount){
    lock.lock();
    while(balance >= 300){
        // waiting for the balance to go below 300.
        belowUpperLimitFundsCondition.await(); }
        sufficientFundsCondition.signalAll();
        balance += amount;
    lock.unlock(); }

```
- For example, do you need to worry about race conditions in this case?

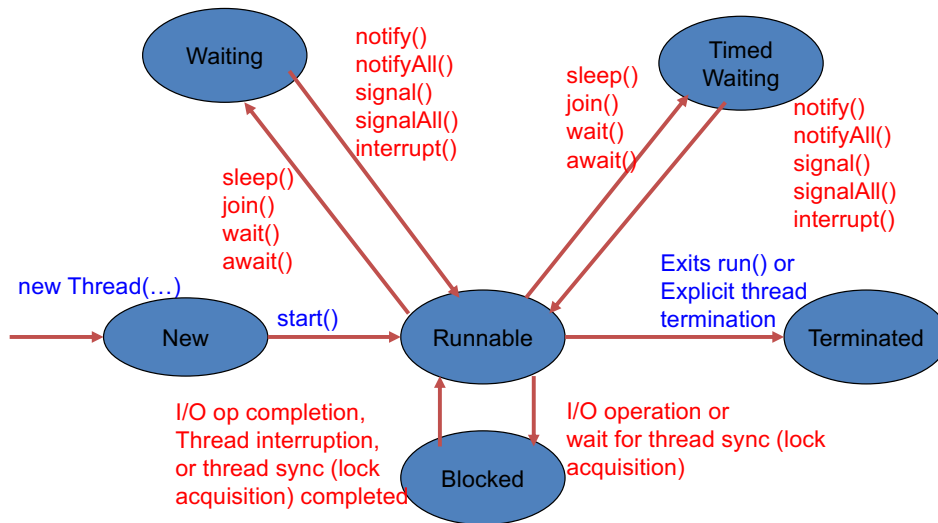


- Can the "W" thread withdraw money before the "D" thread deposits money?
  - Can the balance have a negative value?
    - The answer is NO.



- "W" thread CANNOT withdraw money before "D" thread deposits money.
- "D" thread CANNOT deposit money before "W" thread withdraws money.

## States of a Thread



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## Two Important Things (1)

- You can safely change the state/value of a shared variable after calling `signalAll()`.
  - AS FAR AS the state changes in atomic code
- That said, common programming convention/practice is:
  - A state change first, followed by `signalAll()`.

## Two Important Things (2)

- A JVM can perform context switches even when a thread runs atomic code.
  - A lock guarantees that only one thread exclusively runs atomic code at a time.
  - Some resources (books, online materials, etc.) explicitly/implicitly say that context switches never occur when a thread runs atomic code.
    - It is WRONG!

## `signal()` and `signalAll()`

- `signalAll()`
  - Wakes up all waiting threads on a condition object.
    - All of them go to the “runnable” state.
    - One of them will re-acquire a lock. The others will go to the “blocked” state.
- `signal()`
  - Wakes up one of waiting threads on a condition object.
    - The selected thread goes to the “runnable” state. The others stay at the “waiting” state.
    - JVM’s thread scheduler selects one of them. Assume random selection.
      - Not predictable which waiting thread to be selected.

## signal() and signalAll()?

- Either one works well.
- signalAll() is favored in many cases/projects.
  - I prefer signalAll() in my personal taste.

## ThreadSafeBankAccount2.java

```
• Condition sufficientFundsCondition = lock.newCondition();
  Condition belowUpperLimitFundsCondition = lock.newCondition();

• withdraw(double amount){
    lock.lock();
    while(balance <= 0){
        // waiting for the balance to exceed 0
        sufficientFundsCondition.await(); }
    balance -= amount;
    belowUpperLimitFundsCondition.signalAll();
    lock.unlock(); }

• deposit(double amount){
    lock.lock();
    while(balance >= 300){
        // waiting for the balance to go below 300.
        belowUpperLimitFundsCondition.await(); }
    balance += amount;
    sufficientFundsCondition.signalAll();
    lock.unlock(); }
```

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## “while” or “if” to Surround await()?

- ```
withdraw(double amount){
    lock.lock();
    while(balance <= 0){
        // waiting for the balance to exceed 0
        sufficientFundsCondition.await(); }
    balance -= amount;
    belowUpperLimitFundsCondition.signalAll();
    lock.unlock(); }
```
- ```
deposit(double amount){
    lock.lock();
    while(balance >= 300){
        // waiting for the balance to go below 300.
        belowUpperLimitFundsCondition.await(); }
    balance += amount;
    sufficientFundsCondition.signalAll();
    lock.unlock(); }
```
- “while” should be used rather than “if” when multiple threads call withdraw() concurrently. Why?

## A Potential Problem

(1)  $b==0$ . Two W threads: “waiting”

```
withdraw(double amount){
    lock.lock();
    if(balance <= 0){
        // waiting for the balance to exceed 0
        sufficientFundsCondition.await(); }
    balance -= amount;
    belowUpperLimitFundsCondition.signalAll();
    lock.unlock(); }
```

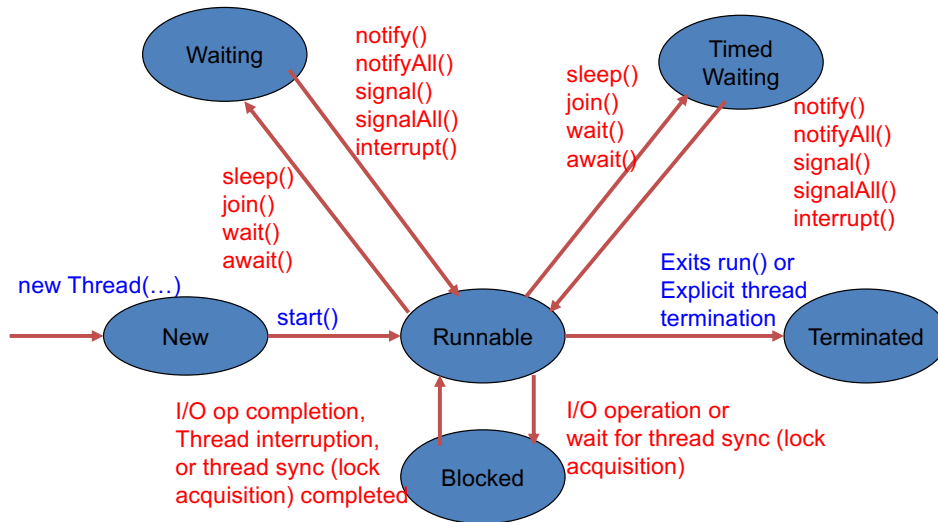
(3) Two W threads: Go to “runnable” One of them acquires the lock again. The other W thread: Go to “blocked” on acquiring the lock.

```
deposit(double amount){
    lock.lock();
    if(balance >= 300){
        // waiting for the balance to go below 300.
        belowUpperLimitFundsCondition.await(); }
    balance += amount;
    sufficientFundsCondition.signalAll();
    lock.unlock(); }
```

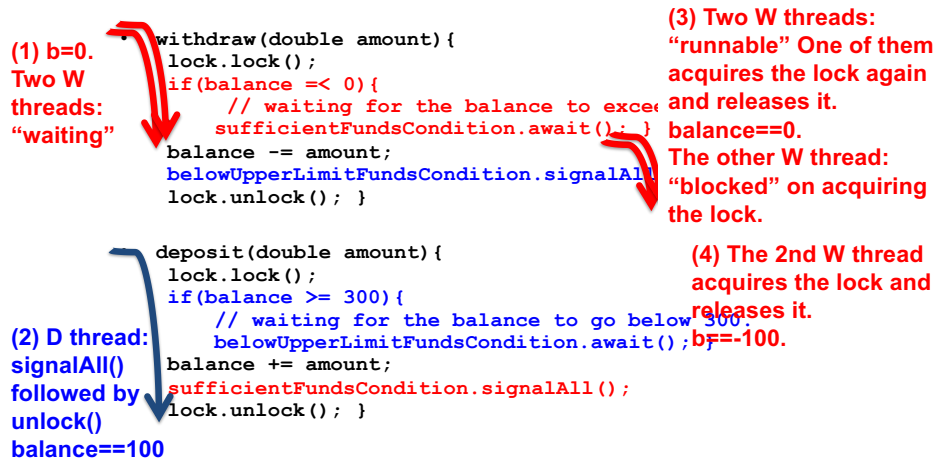
(2) D thread: signalAll() followed by unlock()  $b==100$



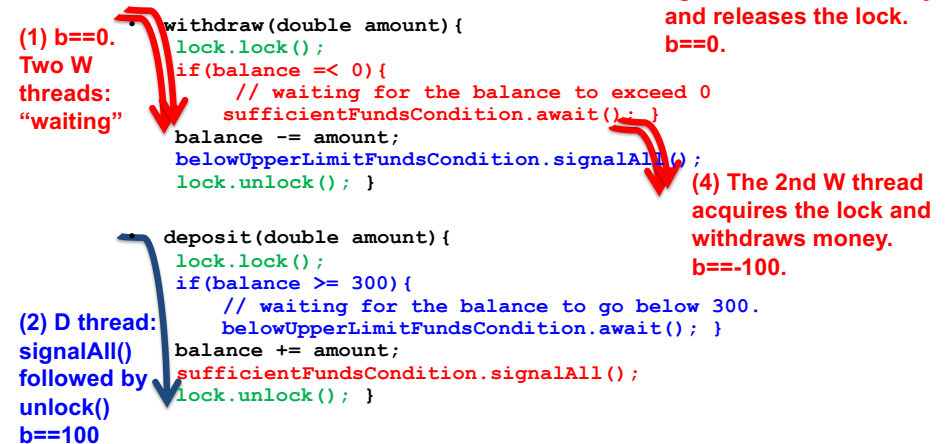
# States of a Thread



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- The 2nd "W" thread **should have made sure** if **balance>0**.
- If only one "W" thread runs, this problem does not occur.
- Just always use a **while loop** regardless of the number of threads you use.



## "if" or "while" in Atomic Code?

- You can use "if", rather than "while," for a conditional check
  - if you use `signal()`, not `signalAll()`.
- However, in practice, the **while-signalAll** pair is more common than the **if-signal** pair.