Milestone 1

Group Members

Anthony Mezzatesta

Amezzatesta@ufl.edu
Github - AnthonyMezzatesta
Role - Developer

Jonathan Palma

palma.jonathan@ufl.edu
Github - gitttjonzzz
Role - Developer

Luis Gonzalez

Igonzaherman@gmail.com Github - gonzaherman99 Role - Developer

Mert Dayi

imert.dayi@ufl.edu
Github - mertdayi01
Role - Project Manager / Developer

Communication Method

Slack

Github Repository Link

https://github.com/gonzaherman99/C4533-GroupProject

Gantt Chart

	TACK TITLE	TASK	START	DUE	DUDATION	PCT OF TASK	T OF TASK WEEK 1 WEEK 2 WEEK 3 WEEK 4 WEEK 5 WEEK 6			WEEK 2 WEEK 3						5 WEEK 6 WEEK 7 WE			EEK 8																			
	TASK TITLE	OWNER	DATE	DUE DATE	DURATION	COMPLETE		T W	/ R	F	м	r w	R	F	м т	w	R	F	м	T W	R	F	м	T W	R	F	м	тΝ	N R	F	м	тΝ	N R	F	м	Т	W R	F
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Step 1)

Input Matrix:

ck 2
ck 3
ck 4

Step 2) STOCK 1:		
BUY DAY	SELL DAY	PROFIT
1	2	-11
1	3	-7
1	4	-9
1	5	4
2	3	4
2	4	2
2	5	15
3	4	-2
3	5	11
4	5	13

STOCK 2:

BUY DAY		SELL DAY		PROFIT
1		2	İ	0
1		3	ĺ	9
1	İ	4	j	0
1		5	İ	5
2	İ	3	į	9
2	İ	4	į	0
2	İ	5	į	5
3	İ	4	ĺ	-9
3		5	ĺ	-4
4	İ	5	İ	5

STOCK 3:		
BUY DAY	SELL DAY	PROFIT
1	2	2
1	3	0
1	4	-5
1	5	-4
2	3	-2
2	4	-7
2	5	-6
3	4	-5
3	5	-4
4	5	1

STOCK 4:

BUY DAY	SELL DAY	PROFIT
1	2	-11
1	3	10
1	4	-6
1	5	-4
2	3	1
2	4	5
2	5	7
3	4	4
3	5	6
4	5	2

Step 3)

Stock 1 most profitable transaction:

Buy at day 2 for 1 and sell at day 5 for 16. Profit = 15

Stock 2 most profitable transaction:

Buy at day 1 for 4 and sell at day 3 for 13. Profit = 9

Stock 3 most profitable transaction:

Buy at day 1 for 6 and sell at day 2 for 8. Profit = 2

Stock 4 most profitable transaction:

Buy at day 2 for 3 and sell at day 5 for 10. Profit = 7

Step 4)

Maximum profit is 15 from Stock 1 – Buy day 2 and sell day 5 OUTPUT: (1, 2, 5, 15)

Problem Statement 2

You are given a matrix A of dimensions $m \times n$, where each element represents the predicted prices of m different stocks for n consecutive days. Additionally, you are given an integer k $(1 \le k \le n)$. Your task is to manually find a sequence of at most k transactions, each involving the purchase and sale of a single stock, that yields the maximum profit.

Step 1)

Input Matrix:

```
A = [[25, 30, 15, 40, 50], // Stock 1

[10, 20, 30, 25, 5], // Stock 2

[30, 45, 35, 10, 15], // Stock 3

[5, 50, 35, 25, 45]] // Stock 4
```

Step 2)

STOCK 1:

BUY DAY	SELL DAY	PROFIT
1	2	5
1	3	-10
1	4	15
1	5	25
2	3	-15

2	4	10
2	5	20
3	4	25
3	5	35
4	5	10

STOCK 2:

BUY DAY	SELL DAY	PROFIT
1	2	10
1	3	20
1	4	15
1	5	-5
2	3	10
2	4	5
2	5	-15
3	4	-5
3	5	-25
4	5	-20

STOCK 3:

BUY DAY	SELL DAY	PROFIT
1	2	15
1	3	5
1	4	-20
1	5	-15
2	3	-10
2	4	-35
2	5	-30
3	4	-25
3	5	-20
4	5	5

STOCK 4:

BUY DAY	SELL DAY	PROFIT
1	2	45
1	3	30
1	4	20
1	5	40
2	3	-15

2 4 -25
2 5 -5
3 4 -10
3 5 10
4 5 20

Step 3)

Finding optimal non-overlapping transactions (k=3):

Best profitable transactions:

- 1. Stock 4: $(1,2) \rightarrow \text{Profit} = 45$
- 2. Stock 1: $(3,5) \rightarrow \text{Profit} = 35$
- 3. Stock 4: $(1,5) \rightarrow \text{Profit} = 40$
- 4. Stock 1: $(1,5) \rightarrow \text{Profit} = 25$
- 5. Stock 1: $(3,4) \rightarrow \text{Profit} = 25$

Selected non-overlapping transactions:

- Transaction 1: Stock 4, Buy Day 1, Sell Day 2 → Profit = 45
- Transaction 2: Stock 1, Buy Day 3, Sell Day 5 → Profit = 35

Total Maximum Profit = 80

Output: [(4,1,2), (1,3,5)]

Problem 3

Problem Statement

You are given a matrix A of dimensions $m \times n$, where each element represents the predicted prices of m different stocks for n consecutive days. Additionally, you are given an integer c ($1 \le c \le n-2$). Your task is to determine the maximum profit achievable under the given trading restrictions, where you cannot buy any stock for c days after selling any stock. If you sell a stock on day i, you are not allowed to buy any stock until day i+c+1.

Input

	Day 1	Day 2	Day 3	Day 4	Day 5	Day 6	Day 7
Stock 1	7	1	5	3	6	8	9
Stock 2	2	4	3	7	9	1	8
Stock 3	5	8	9	1	2	3	10
Stock 4	9	3	4	8	7	4	1
Stock 5	3	1	5	8	9	6	4

Cooldown: c = 2

First let's define best(t) as the maximum profit achievable with one buy, sell between days t and 7.

t	Window	Best trade	best ₁ (t)
1–4	Days 1 to 7	buy Stock 3 on day 4 at 1 then sell on day 7 at 10	9
5	Days 5 to 7	buy Stock 3 on day 5 at 2 then sell on day 7 at 10	8
6	Days 6 to 7	buy Stock 2 on day 6 at 1 then sell on day 7 at 8	7
7	day 7	no transaction	0

Now for each possible first trade (i, j, l), let's compute $Profit_1 = P[i, l] - P[i, j]$, earliest next buy B = l + 3, $Profit_2 = best_1(B)$, and $Total = Profit_1 + Profit_2$. Here are the top contenders:

Buy j	First trade (i,1)	Profit ₁	Next B=l +3	Profit ₂	Total
j = 1	Stock 2: buy	7	8 to 0	0	7

	1, sell 5				
j = 2	Stock 1: buy 2, sell 3	4	6 to 7	7	11
j = 3	Stock 2: buy 3, sell 5	6	8 to 0	0	6
j = 4	Stock 3: buy 4, sell 7	9	10 to 0	0	9
j = 5	Stock 3: buy 5, sell 7	8	10 to 0	0	8
j = 6	Stock 2: buy 6, sell 7	7	10 to 0	0	7

The best total (11) come from:

- 1. Stock 1: buy on day 2 at 1 then sell on day 3 at 5 (Profit₁ = 4) with the cooldown until day 6(3+2+1)
- 2. Stock 2: buy on day 6 at 1 then sell on day 7 at 8 ($Profit_2 = 7$)

Total profit = 4 + 7 = 11.

Final Answer

Maximum profit: 11

Trades: (i=1,j=2,1=3) and (i=2,j=6,1=7)

Milestone 2

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Anthony Mezzatesta

Amezzatesta@ufl.edu
Github - AnthonyMezzatesta
Role - Developer

Jonathan Palma

palma.jonathan@ufl.eduGithub - gitttjonzzzRole - Developer

Luis Gonzalez

Igonzaherman@gmail.com Github - gonzaherman99 Role - Developer

Mert Dayi

imert.dayi@ufl.edu
Github - mertdayi01
Role - Project Manager / Developer

Communication Method

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Programming Language

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```
Brute Force Problem 1)
function brute force prob1(A):
       m = rows in A
       n = columns in A
       best profit = 0
       best stock = 0
       best_buy = 0
       best sell = 0
       # Iterate over each stock
       for i from 0 to m-1:
              # Buy day index
               for j from 0 to n-2:
                      # Sell day index (must be after buy)
                      for k from j+1 to n-1:
                             profit = A[i][k] - A[i][j]
                              if profit > best_profit:
                                     best profit = profit
                                     best stock = i
                                     best buy = j
                                     best sell = k
       # No profitable transaction
       if best profit <= 0:
               return (0, 0, 0, 0)
       else:
               # Convert to 1-based indexing
               return (best stock+1, best buy+1, best sell+1, best profit)
```

```
Greedy Problem 1)
function greedy prob1(A):
  m = rows in A
  n = columns in A
  best profit = 0
  best stock = -1
  best buy = -1
  best sell = -1
       # Iterate over each stock
       for i from 0 to m-1:
               min price = A[i][0]
               min_index = 0
               current profit = 0
               current buy = 0
               current sell = 0
               # Start from day 1
               for j from 1 to n-1:
                      # Found new minimum price
                      if A[i][j] < min_price:
                             min_price = A[i][j]
                              min index = i
                      else:
                              profit = A[i][j] - min price
                              # Update if better profit
                              if profit > current profit:
                                     current profit = profit
                                     current buy = min index
                                     current sell = j
               # Update global best
               if current profit > best profit:
                      best profit = current profit
                      best stock = i
                      best buy = current buy
                      best sell = current sell
       if best_profit <= 0:
               return (0, 0, 0, 0)
       else:
               return (best stock+1, best buy+1, best sell+1, best profit)
```

```
DP Problem 1)
function dp prob1(A):
       m = rows in A
       n = columns in A
       best profit = 0
       best stock = -1
       best_buy = -1
       best sell = -1
       # Iterate over each stock
       for i from 0 to m-1:
               min price = A[i][0]
               min index = 0
               max profit i = 0
               buy i = 0
               sell i = 0
               # Start from day 1
               for j from 1 to n-1:
                      # Update minimum price if lower found
                      if A[i][j] < min_price:
                             min price = A[i][j]
                              min index = i
                      # Calculate profit if selling today
                      profit = A[i][j] - min price
                      if profit > max_profit_i:
                                                  # Update if better profit
                              \max profit i = profit
                              buy i = min index
                             sell i = j
               # Update global best
               if max profit i > best profit:
                      best profit = max profit i
                      best stock = i
                      best buy = buy i
                      best sell = sell i
       if best profit \leq 0:
               return (0, 0, 0, 0)
       else:
               return (best stock+1, best buy+1, best sell+1, best profit)
```

```
Dynamic Programming Algorithm for Problem 2 (O(m · n<sup>2</sup>k) Time)
function dp prob2 n2k(A):
  m = number of rows in A
  n = number of columns in A
  k = [number of operations]
  DP = 3D array of size [m+1][n+1][k+1], initialized to -infinity
  DP[0][0][0] = 0 # or base case value
  for i from 1 to m:
     for j from 1 to n:
       for x from 0 to k:
          for p from 1 to n:
            for q from 0 to k:
               if valid transition(p, q, j, x):
                 DP[i][j][x] = max(
                    DP[i][j][x],
                    DP[i-1][p][q] + cost(B, i, j, x)
  answer = maximum DP[m][j][x] over all valid j, x
  return answer
Dynamic Programming Algorithm for Problem 2 (O(m \cdot n \cdot k) Time)
function dp prob2 nk(A):
  m = number of rows in A
  n = number of columns in A
  k = [number of operations]
  DP = 3D array of size [m+1][n+1][k+1], initialized to -infinity
  DP[0][0][0] = 0 # or base case value
  for i from 1 to m:
     for j from 1 to n:
       for x from 0 to k:
          DP[i][j][x] = max(
            DP[i-1][j][x],
            DP[i-1][j-1][x-1]
         ) + cost(B, i, j, x)
  answer = maximum DP[m][j][x] over all valid j, x
  return answer
```

```
function dp_prob3_slow(A, c):
  m = rows in A
  n = columns in A
  dp = array of size n initialized to 0
  choice = array of size n initialized to null
  for i from 1 to n-1:
     dp[i] = dp[i-1]
     choice[i] = null
     for j from 0 to i-1:
        prev profit = (j-c-1 >= 0)? dp[j-c-1]: 0
        for stock from 0 to m-1:
           profit = A[stock][i] - A[stock][j]
           total = prev profit + profit
           if total > dp[i]:
             dp[i] = total
             choice[i] = (stock, j, prev_profit)
  if dp[n-1] <= 0:
     return []
  result = []
  i = n-1
  while i > 0 and choice[i] != null:
     stock, buy_day, prev_profit = choice[i]
     result.add((stock+1, buy_day+1, i+1))
     j = buy day - c - 1
     while j >= 0 and dp[j] != prev profit:
       j = j - 1
     i = j
  reverse(result)
  return result
function dp_prob3_fast(A, c):
  m = rows in A
  n = columns in A
  hold = -infinity
```

```
sold = 0
rest = 0
transactions = []
hold_info = null
for i from 0 to n-1:
  prev_hold = hold
  prev sold = sold
  prev_rest = rest
  rest = max(prev_rest, prev_sold)
  sold = prev hold
  sell_stock = -1
  for stock from 0 to m-1:
     profit = prev_hold + A[stock][i]
     if profit > sold:
       sold = profit
       sell_stock = stock
  if sell_stock != -1 and sold > prev_sold:
     if hold_info != null:
       transactions.add((hold info.stock+1, hold info.day+1, i+1))
     hold_info = null
  hold = prev hold
  buy_stock = -1
  buy profit = hold
  for stock from 0 to m-1:
     profit = rest - A[stock][i]
     if profit > buy_profit:
       buy profit = profit
       buy stock = stock
  if buy_stock != -1:
     hold = buy_profit
     hold_info = (stock: buy_stock, day: i)
final profit = max(sold, rest)
if final_profit <= 0:
  return []
```

return transactions