Agenda: - Chandy- Lamport snapshot algorithm - Limitations I assumptions l properties - Centralised vs decentralized algorithms Given events A and B, $A \rightarrow B$, if B is in snapshot, A should be too! Recording a snapshot The initiator process; (one or more) - Records its own state - Sends a marker message out on all its obtgoing channels. - stant neconding the onsys it neceives on all its incoming channels. When process Pi receives a monker message - if it is the first marker Pi has seen -Pi gerords its state - Pi mark channel Cki as empty -Pi sends a marker out on all its outgoing channels

- Pi storts sucording on all incoming

channels except CKi - Otherwise: -Pi stop recording on Cxi. Consistent snatshot Stop recording on Cal count this NOT happen in changly-compate Why Channels are FIFOD Moorker missage from Pa to P, cannot happen before mag From c to p.

More interesting et algorithm Future X- Empty = [m]Total number of marker msgs. N(N-1)N- No. OF PROCESSES - Even with multiple initiators, the Final snapshot is a consistent snapshot, - It can run even when app messages are being sent /happening

CL snupshots: BIG PICTURE
- ctannels: - How msgs toware C - FIFO behaviour
- This FIFO assumption is a nequinement for CL algorithm!
- No pausing application messages in CL.
- C-L assumes messages one NOT lost, consupted on duplicated Also, processes don't crash ! C-L: Grood things
- snapshots are consistent (+ A, B
if $A \rightarrow B$ & B is in smpsh So is A)
- Grassanteed to teaminate (given assumptions about msgs not being lost, etc.) - klosies fine with more than one initiation!



