$\frac{\frac{D_{j}\vdash D_{j}}{D_{j} < < D_{j}\vdash \bot} W <_{R}}{\frac{D_{j}\vdash \bot; D_{j}}{I\vdash D_{j} >> \bot; D_{j}}} W >_{R}}{\frac{I\vdash D_{j} >> \bot; D_{j}}{I\vdash D_{j} >> \bot} C_{R}} \frac{W >_{R}}{Grishin_{R}} \frac{H\vdash (D_{j} >> \bot; D_{j})}{H\vdash D_{j} >> \bot} (>,;)}{\frac{D_{j} >> I\vdash D_{j} >> \bot}{D_{j} >> I\vdash D_{j} >> \bot} (>,;)}{\frac{(D_{j} \rightarrow \bot) \rightarrow \bot\vdash (D_{j} >> I) >> D_{j}}{(D_{j} >> I); (D_{j} \rightarrow \bot) \rightarrow \bot\vdash D_{j}} E_{L}} \frac{(D_{j} >> I); (D_{j} \rightarrow \bot) \rightarrow \bot\vdash D_{j}}{(D_{j} >> I\vdash (D_{j} \rightarrow \bot) \rightarrow \bot >> D_{j}} (:,>)} \frac{(I_{j} >> I) \rightarrow \bot\vdash (D_{j} >> I) \rightarrow \bot\vdash D_{j}}{\frac{I\vdash D_{j}; ((D_{j} \rightarrow \bot) \rightarrow \bot >> D_{j})}{I\vdash ((D_{j} \rightarrow \bot) \rightarrow \bot >> D_{j}); D_{j}}} Grishin_{R}} \frac{I\vdash (D_{j} \rightarrow \bot) \rightarrow \bot\vdash D_{j}; D_{j}}{I\vdash (D_{j} \rightarrow \bot) \rightarrow \bot\vdash D_{j}; D_{j}} Grishin_{R}} \frac{(D_{j} \rightarrow \bot) \rightarrow \bot\vdash D_{j}; D_{j}}{(D_{j} \rightarrow \bot) \rightarrow \bot\vdash D_{j}} C_{R}} D_{j} \vdash \{father_{n}\}\{n\} \frac{(D_{j} \rightarrow \bot) \rightarrow \bot\vdash \{father_{n}\}\{n\}}{(D_{j} \rightarrow \bot) \rightarrow \bot\vdash \{father_{n}\}\{n\}} Cut}$ $\begin{array}{c} D_j \vdash D_i & \stackrel{1}{\longmapsto} \stackrel{1}{\vdash} \stackrel{1}{\vdash} \stackrel{1}{\vdash} \\ D_j \rightarrow 1 \vdash D_j > \bot \rightarrow \kappa \\ \hline D_j \rightarrow 1 \vdash D_j > \bot \rightarrow \kappa \\ \hline D_j \rightarrow 1 \vdash D_j > \bot \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_{\mathcal{F}} \vdash D_j \rightarrow \bot \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_{\mathcal{F}} \vdash D_j \rightarrow \bot \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_{\mathcal{F}} \vdash D_j \rightarrow \bot \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_{\mathcal{F}} \vdash D_j \rightarrow \bot \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_{\mathcal{F}} \vdash D_j \rightarrow \bot \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_{\mathcal{F}} \vdash D_j \rightarrow \bot \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_{\mathcal{F}} \vdash (D_j \rightarrow 1) \rightarrow [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_{\mathcal{F}} \vdash (\operatorname{father}_{\kappa})[\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_{\mathcal{F}} \vdash (\operatorname{father}_{\kappa})[\operatorname{po}_{\kappa}]^{k-1}D_j \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_{\mathcal{F}} \vdash (\operatorname{father}_{\kappa})[\operatorname{po}_{\kappa}]^{k-1}D_j \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_{\mathcal{F}} \vdash (\operatorname{father}_{\kappa})[\operatorname{po}_{\kappa}]^{k-1}D_j \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_j \vdash (\operatorname{father}_{\kappa})[\operatorname{po}_{\kappa}]^{k-1}D_j \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_j \vdash (\operatorname{father}_{\kappa})[\operatorname{po}_{\kappa}]^{k-1}D_j \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_j \vdash (\operatorname{father}_{\kappa})[\operatorname{po}_{\kappa}]^{k-1}D_j \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_j \vdash (\operatorname{father}_{\kappa})[\operatorname{po}_{\kappa}]^{k-1}D_j \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_j \vdash (\operatorname{father}_{\kappa})[\operatorname{po}_{\kappa}]^{k-1}D_j \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_j \vdash (\operatorname{father}_{\kappa})[\operatorname{po}_{\kappa}]^{k-1}D_j \vdash (\operatorname{father}_{\kappa})[\operatorname{po}_{\kappa}]^{k-1}D_j \rightarrow \kappa \\ \hline (D_j \rightarrow 1) = [\operatorname{father}_{\kappa}][\operatorname{po}_{\kappa}]^{k-1}\Box_{\mathcal{F}}D_j \vdash (\operatorname{father}_{\kappa})[\operatorname{po}_{\kappa}]^{k-1}D_j \vdash (\operatorname{father}_{\kappa})[\operatorname{po}_{\kappa}]^{k-1}$ $D_j \vdash \{\mathsf{father}_n\}\{\mathsf{no}_n\}^{k-1}D_j \quad Cut$ $\frac{((D_j \to \bot) \to [\mathsf{father}_n][\mathsf{no}_n]^{k-1} \square_{j'} D_{j'}) \wedge \langle \mathsf{father}_n \rangle \langle \mathsf{no}_n \rangle^{k-1} \mathsf{no}_n \vdash (D_j \to \bot) \to \bot}{((D_j \to \bot) \to [\mathsf{father}_n][\mathsf{no}_n]^{k-1} \square_{j'} D_{j'}) \wedge \langle \mathsf{father}_n \rangle \langle \mathsf{no}_n \rangle^{k-1} \mathsf{no}_n \vdash \{\mathsf{father}_n\} \{\mathsf{no}_n\}^{k-1} D_j} Cut}$ $\frac{\frac{\Phi_{\mathsf{no}_n} \vdash 1_{\mathsf{no}_n} \qquad D_j \vdash \{\mathsf{no}_n\}D_j}{1_{\mathsf{no}_n} \to D_j \vdash \Phi_{\mathsf{no}_n} >> \{\mathsf{no}_n\}D_j}}{1_{\mathsf{no}_n} \to D_j \vdash \{\mathsf{no}_n\}D_j} \xrightarrow{Reduce'R}$ $\underbrace{\widehat{\log}_n \stackrel{k-1}{\text{father}}_n (D_j \to \bot) \to [\mathsf{father}_n][\mathsf{no}_n]^{k-1} \square_{j'} D_{j'} \vdash \{\mathsf{no}_n\} D_j}_{F}$ $\widehat{\operatorname{ho}}_n^{k-1}\operatorname{father}_n(D_j o ot) o [\operatorname{father}_n][\operatorname{no}_n]^{k-1} \Box_{j'} D_{j'} dash [\operatorname{no}_n] D_j$ $\underbrace{\operatorname{father}_n\left(D_j\to\bot\right)\to\left[\operatorname{father}_n\right][\operatorname{no}_n]^{k-1}\Box_{j'}D_{j'}\vdash\left\{\operatorname{no}_n\right\}^{k-1}[\operatorname{no}_n]D_j}_{FboxA_R}$ $\widehat{\mathsf{father}}_n\left(D_j\to\bot\right)\to[\mathsf{father}_n][\mathsf{no}_n]^{k-1}\Box_{j'}D_{j'}\vdash[\mathsf{no}_n]^{k-1}[\mathsf{no}_n]D_j$ $\frac{\overbrace{(D_j \to \bot) \to [\mathsf{father}_n][\mathsf{no}_n]^{k-1}\Box_{j'}D_{j'} \vdash [\mathsf{no}_n]^kD_j}}{\underbrace{(D_j \to \bot) \to [\mathsf{father}_n][\mathsf{no}_n]^{k-1}\Box_{j'}D_{j'} \vdash \{\mathsf{father}_n\}[\mathsf{no}_n]^kD_j}}_{(D_j \to \bot) \to [\mathsf{father}_n][\mathsf{no}_n]^{k-1}\Box_{j'}D_{j'} \vdash [\mathsf{father}_n][\mathsf{no}_n]^kD_j}} \underbrace{Fbox A_R}_{Fbox A_R}$ $\frac{\mathsf{father}_n \to (D_j \to \bot) \to [\mathsf{father}_n][\mathsf{no}_n]^{k-1} \square_{j'} D_{j'} \vdash \mathsf{father}_n >> [\mathsf{father}_n][\mathsf{no}_n]^k D_j}{Cut}$ $\frac{(D_j \to \bot) \land \mathsf{father}_n \to [\mathsf{father}_n][\mathsf{no}_n]^{k-1} \square_{j'} D_{j'} \vdash \mathsf{father}_n >> [\mathsf{father}_n][\mathsf{no}_n]^k D_j}{(D_j \to \bot) \land \mathsf{father}_n \to [\mathsf{father}_n][\mathsf{no}_n]^{k-1} \square_{j'} D_{j'} \vdash \mathsf{father}_n \to [\mathsf{father}_n][\mathsf{no}_n]^k D_j} \xrightarrow{Fbox K_L} \frac{\square_j ((D_j \to \bot) \land \mathsf{father}_n \to [\mathsf{father}_n][\mathsf{no}_n]^{k-1} \square_{j'} D_{j'}) \vdash \{j\} \mathsf{father}_n \to [\mathsf{father}_n][\mathsf{no}_n]^k D_j}{\square_j ((D_j \to \bot) \land \mathsf{father}_n \to [\mathsf{father}_n][\mathsf{no}_n]^{k-1} \square_{j'} D_{j'}) \vdash \square_j (\mathsf{father}_n \to [\mathsf{father}_n][\mathsf{no}_n]^k D_j)} \xrightarrow{Fbox K_R} \frac{\mathsf{father}_n \to [\mathsf{father}_n][\mathsf{no}_n]^k D_j}{\mathsf{father}_n \to [\mathsf{father}_n][\mathsf{no}_n]^k D_j} \xrightarrow{Fbox K_R} \frac{\mathsf{father}_n \to [\mathsf{father}_n][\mathsf{no}_n]^k D_j}{\mathsf{father}_n \to [\mathsf{father}_n][\mathsf{no}_n]^k D_j} \xrightarrow{Fbox K_R} \frac{\mathsf{father}_n \to [\mathsf{father}_n][\mathsf{no}_n]^k D_j}{\mathsf{father}_n \to [\mathsf{father}_n][\mathsf{no}_n]^k D_j} \xrightarrow{Fbox K_R} \frac{\mathsf{father}_n}{\mathsf{father}_n} = \frac{\mathsf{father}_n}{\mathsf{father}_n}$

 $\frac{(D_j \to \bot) \land \mathsf{father}_n \to [\mathsf{father}_n][\mathsf{no}_n]^{k-1} \square_{j'} D_{j'} \vdash \mathsf{father}_n >> (D_j \to \bot) \to [\mathsf{father}_n][\mathsf{no}_n]^{k-1} \square_{j'} D_{j'}}{(D_j \to \bot) \land \mathsf{father}_n \to [\mathsf{father}_n][\mathsf{no}_n]^{k-1} \square_{j'} D_{j'} \vdash \mathsf{father}_n \to (D_j \to \bot) \to [\mathsf{father}_n][\mathsf{no}_n]^{k-1} \square_{j'} D_{j'}} \xrightarrow{P_R}$