Computer Organization

Sequential Logic & RTL

Pyeongsu Park (pyeongsu@snu.ac.kr)

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Objective

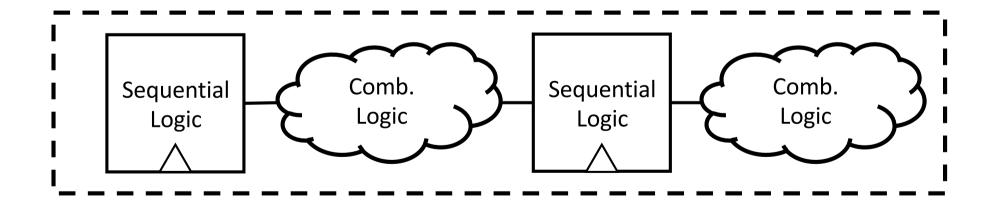
- Understand the Register Transfer Level (RTL) concept of HDL
- Understand how to make a sequential logic in Verilog

• If you are familiar with FSM and how to implement it in Verilog, then you can skip to page #29.

RTL: a way to represent a circuit

Digital Circuit

• Combinational logic + Sequential logic



Combinational Logic (CL)

 The last assignment, ALU, was a very common example of a combinational circuit.

- Outputs are a function of inputs ONLY.
 - Outputs = f(inputs)

Sequential Logic (SL)

- Outputs of a logic are a function of inputs & states.
 - Outputs = f(inputs, state)
- SLs have memory elements to save their states.
 - Remember? SR-latches, JK-FF, D-FF, RAM, ROM
- SLs can contain CLs, but CLs cannot not contain SLs.

Synchronous Circuit

- A circuit that is synchronized to a clock signal.
- Consist of memory elements and combinational logics.
- The changes in the value of memory elements are synchronized to a clock signal (positive or negative edge).
- Easy to debug, and we focus this for the project.
- Opposite: Asynchronous circuit
 - It consists of combinational logics and memory elements, but both are not synchronized to a clock.
 - Instead, memory elements change their values according to external events (e.g., previous calculation is done).

Verilog Representation – Gate Level

 Represent a logic as a <u>set of logic gates</u> ONLY (AND, NOT ...)

- (+) More realistic code than RTL
- (-) Hard to program

Verilog Representation – Register Transfer Level (RTL)

- Represent a logic as "registers" and "data manipulations" between the registers.
- Programmers can concentrate on the data flow between registers, not real gates.
- Compilers convert/optimize RTL into gates.
- (+) Can focus functional aspects, not gates
- (+) Easier and faster development than gate level
- (-) Not all the codes can be made into a real HW

Gate Level vs. RTL

```
module FullAdder_GL (A,B,Cin,Cout,S);
input A, B, Cin;
output Cout, S;
wire w1, w2, w3, w4;

xor(w1, A, B);
and(w2, A, B);
and (w3, w1, Cin);
xor (S, w1, Cin)
or (Cout, w3, w2);
endmodule
```

VS

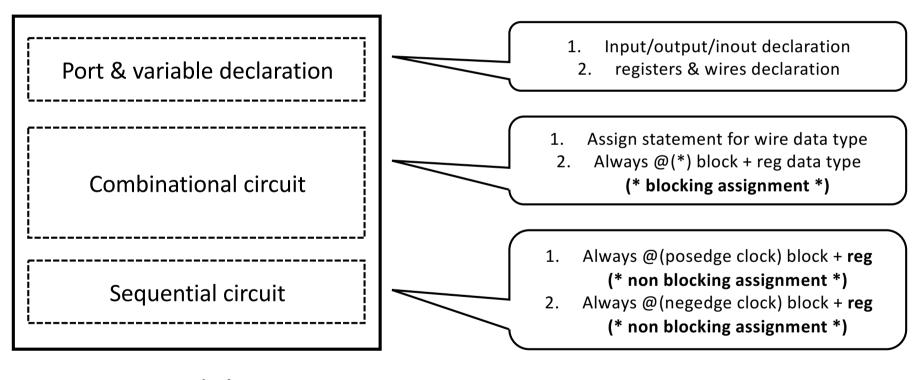
```
module FullAdder_RTL (A,B,S,C);
input A, B;
output S, C;

assign S = A ^ B & Cin;
assign Cout = (A & B) | (Cin & (A ^ B));
endmodule
```

Gate Level

RTL

General RTL Structure



Module

RTL Programming Guide

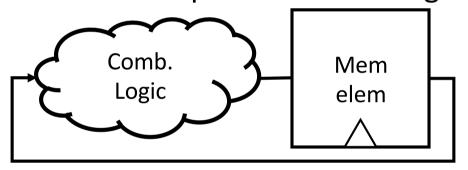
- How to make an RTL code from the given block diagram?
- (1) Understand inputs & outputs, and the functionality of it.
- (2) Determine what to be saved in the registers.
- (3) Think which values should be synchronized to the clock.
- (4) Think about how the data move between the registers.
- (5) Consider what operations should be done between the data movement.

(* step-by-step example will follow after the next few slides. *)

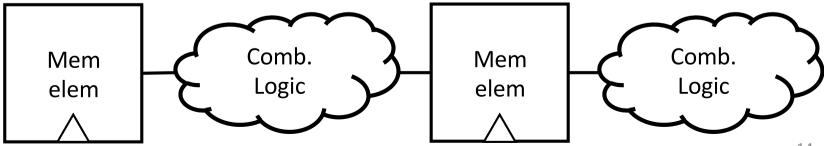
How to write a sequential logic in Verilog

Two Possible SL in Verilog

- Finite State Machine (FSM)
 - Should know how to implement this using Verilog

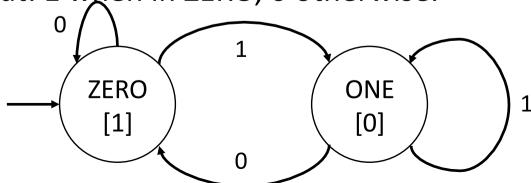


- Pipeline
 - Straight forward: connect each logic using mem elem.



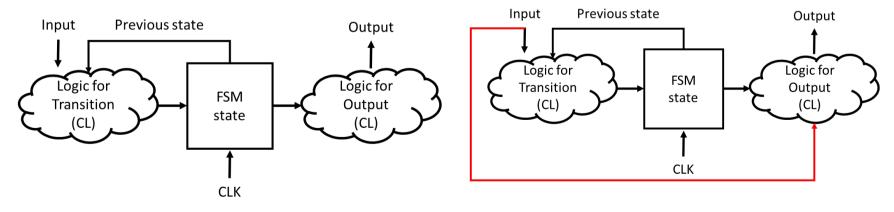
Finite State Machine (FSM)

- Represent a logic as a finite set of states (or nodes), outputs, and transition functions.
 - Transition: movement between nodes for given inputs
- In below example,
 - State = ZERO, ONE
 - Transition function: ZERO=t(ZERO, 0), ZERO=t(ONE, 0), ONE=t(ZERO, 1), ONE=t(ONE, 1)
 - Output: 1 when in ZERO, 0 otherwise.



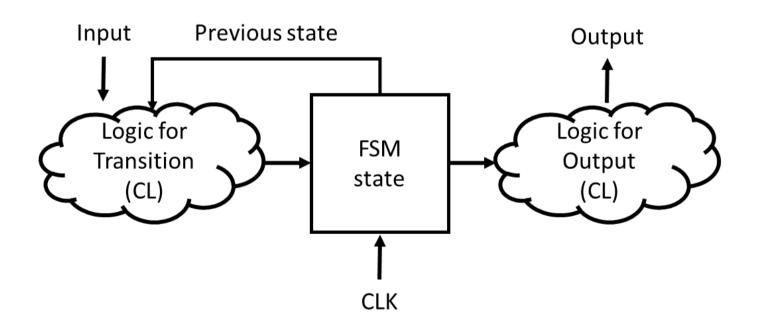
Two Ways to Represent FSM

- We can classify an FSM based on the relationship between outputs and inputs.
- The State is synchronized to CLK.
 - The state is changed when clock is positive edge or negative edge.
 - Assume the state is saved in D-flip-flops.
- Two types of representation of an FSM: Moore & Mealy
 - You can implement the same logic in both ways.



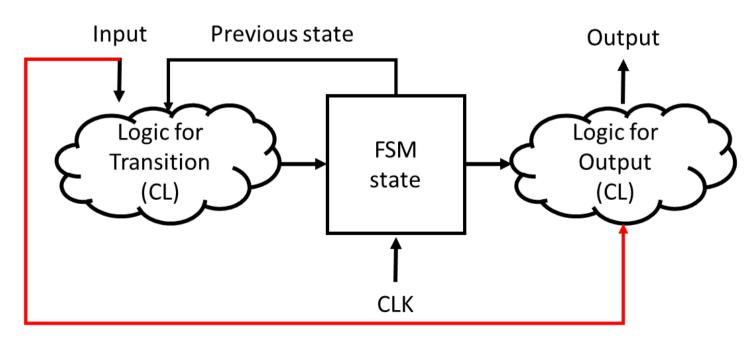
Moore Machine

- Output: function of only the current state
- **Synchronous** machine: the outputs are changed only when the state is changed.



Mealy Machine

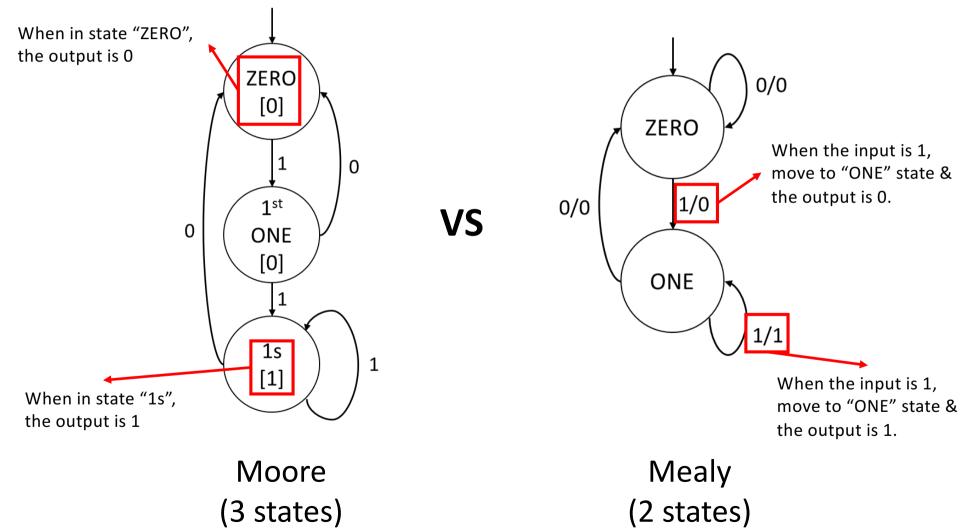
- Output: function of the current state and inputs
- **Asynchronous** machine: the outputs are changed as soon as the inputs are changed.



Example: Reduce 1s

- Input: binary stream
- Output: Convert first 1 to 0 when a stream of 1s is detected.
- Example:
 - 0000**111**00 -> 0000**011**00
 - 010101010 -> 000000000
 - 011001100 -> 001000100

Moore vs. Mealy



Moore vs. Mealy State declaration & update

```
Module Moore (input CLK, input in,
                                               Module Mealy (input CLK, input in, output
output reg out)
                                               reg out)
  parameter zero=0, firstOne=1, ones=2;
                                                 parameter zero=0, one=1;
  reg [1:0] state;
                                                 reg state;
                         2 bits vs 1 bit
  reg [1:0] nextState;
                                                 reg nextState;
                                         VS
                                                 initial begin state
 initial begin state
                                                    state <= zero;
    state <= zero;
 end
                                                end
 always @(posedge CLK)
                                                always @(posedge CLK)
                                                   begin
    begin
                            2 D-FFs
                                                                            1 D-FF
                                                     state <= nextState:
      state <= nextState:
    end
                                                   end
```

Moore Mealy

Moore vs. Mealy Output & state logic

```
always @(in or state)
   begin
      case (state)
        zero:
        begin
           if (n == 1) nextState <= firstOne;</pre>
           else nextState <= zero:
         end
        firstOne:
        begin
           if (n == 1) nextState <= ones;</pre>
           else nextState <= zero:
         end
        ones:
        begin
           if (n == 1) nextState <= ones;
           else nextState <= zero;</pre>
         end
        default: nextState <= zero;</pre>
      endcase
   end
```

```
begin
                                                     case (state)
                                                       zero:
                                                       begin
                                                         out <= 0:
                                                         if (n == 1) nextState <= one:
                                Output logic
                                                         else nextState <= zero;
                                                        end
always @(state)
                                                       one:
    begin
                                                       begin
      case (state)
                                                              if (n == 1) begin
              zero: out <= 0;
                                                                out <= 1:
              firstOne: out <= 0;
                                                                nextState <= one;</pre>
              ones: out <= 1;
                                                              end else begin
              default: out <= 0;
                                                                out <= 0;
      endcase
                                                                nextState <= zero;</pre>
    end
                                                              end
                                                        end
      Output logic
                                                     endcase
                                                   end
```

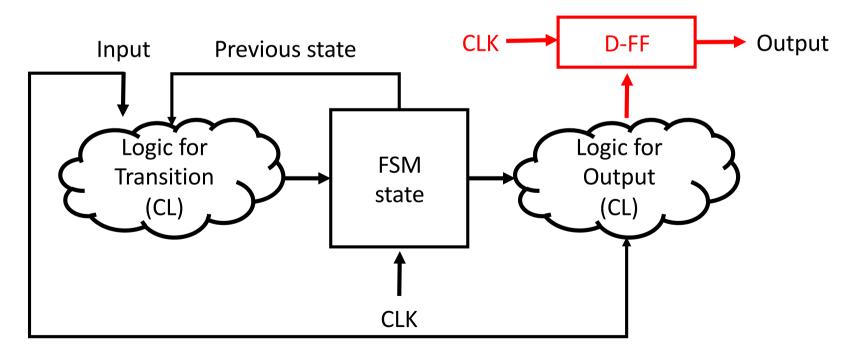
always @(in or state)

State logic

State & output logic

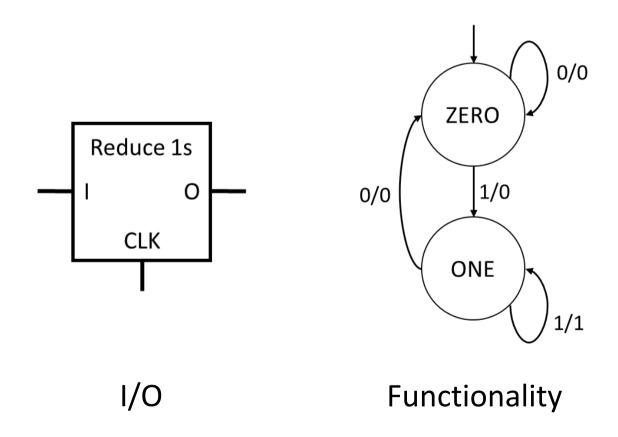
Synchronous Mealy Machine

 If you want to get the synchronized output in a Mealy Machine, buffer the output logic's value using D-FF.



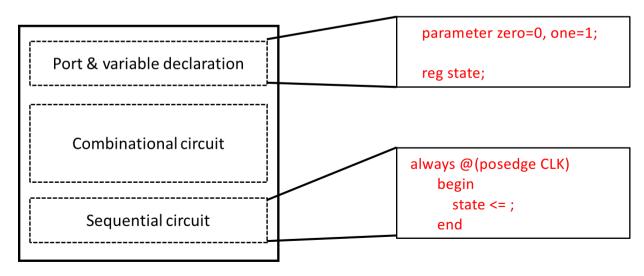
Step by Step Programming (1)

• Understand I/Os and the logic's functionality.



Step by Step Programming (2-3)

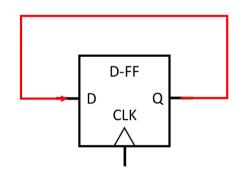
- (2) Determine what to be saved in the logic.
 - FSM: 2 states => 1 bit registers
- (3) What to be synchronized to the clock?
 - FSM: The state update should be synchronized to the positive edge of the clock.



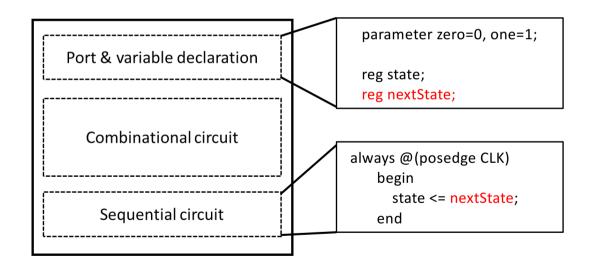
Result after (2-3)

Step by Step Programming (4)

• Think about how the data move between registers; that is think about data flow.



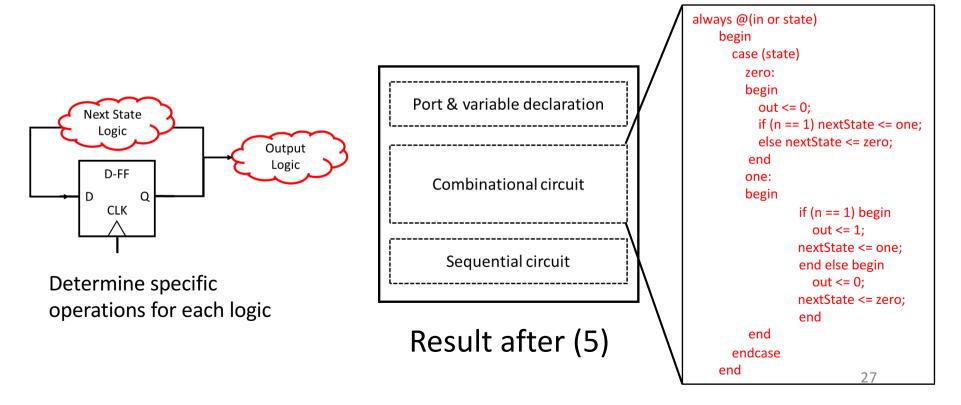
D-FF's data moves from D-FF to D-FF (i.e., same FF).



Result after (4)

Step by Step Programming (5)

- Which operations should be done during a data movement?
 - Next state and output should be computed.



Assignment #2: 101 Detector & Register File

Assign 2-1: 101 Detector

- Goal: Implement a synchronous 101 detector
 - Input: A sequence of binary numbers, 1 bit per 1 clock
 - Output: 1 if 101 is detected, 0 otherwise
 - Example sequence:
 - Input: 0001010100101000011111
 - Output: 000001010000100000000
- How to start?
 - First draw an FSM (either Moore or Mealy).
 - Then, implement the FSM following the steps in the previous slides.

Assign 2-1: Register File (RF)

- Registers are the fundamental components of CPUs, and they have data to be computed or the state of the CPU.
 - E.g., \$v0, \$a1, \$t4, \$s3, \$k0, and \$ra in MIPS
- A register file is a collection of registers.
- * You will use RF in the next assignment.

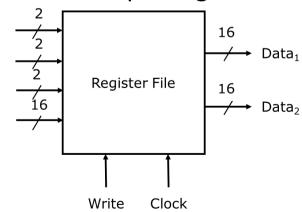
Assign 2-2: Register File

- Goal: Implement a 16-bit 2-read/1-write register file
 - Input: Three addresses (each 2 bits),
 one data to write (16 bits)
 one write signal (1 bit), one clock signal (1 bit)
 - Output: Two read data (each 16 bits)
 - Note1: Data_x is for Addr_x
 - Note2: Addr₁ and Addr₂ are for read while Addr₃ is for write.
 - Note3: Only when "write" signal is on, write data into the register.
 - Note4: Write should be synchronized. That is, the register write should be happened only if "Write" is on and the clock is posedge.

Addr₂

Addr₃

- Q: how many registers in this RF?
- How to start?
 - Understand the I/Os and operations of RF.Data3
 - Make a module for RF and specify its I/Os.



Note 1)

- Do NOT consider Verilog programming (or other HDL) as a variation of C programming.
 - The purpose of HDL is "describing a circuit".
 - The result of the HDL programming is the circuit itself, so think about how your codes would be when it becomes a real hardware (e.g., FF, latch, gates)
 - (c.f., SW programming is for procedures.)
- "reg" type
 - "reg" does not mean either "register" in CPU or flip-flops. It can be equivalent to even "wire" based on its usage.
 - When a "reg" is synchronized to a clock, it becomes FFs, or equivalent to FFs.

Note 2)

Module

- Make your module as small as possible.
- Connect each small module to make a bigger module (i.e., bottom-up fashion).
- Verify each module separately.

Inter-module

- Be careful on endians and overflows
- wire [15:0] a != wire [0:15] a

```
module A()
wire b_out;
B submodule(b_out); // you connected wrongly
endmodule

module B(output wire [2:0] out);
assign out = 3b'11;
endmodule
```

Inside module

 Separate "always block" as much as possible based on the functionality.

Thanks