Give or Take: Non-Determinism, Linear Logic and Session Types

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December 17, 2018

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Example: dramatis personæ



Mary



John



Pâtisserie

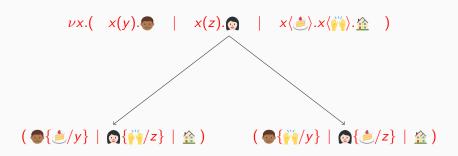


Cake



No Cake

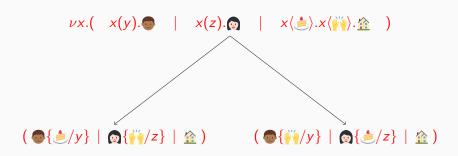
Example: buying cake at the pâtisserie



The π -calculus: syntax

```
Term P, Q, R
::= \nu x.P \qquad \text{create a new channel } x, \text{ then run } P
\mid x(y).P \qquad \text{receive } y \text{ over } x, \text{ then run } P
\mid x\langle y\rangle.P \qquad \text{send } y \text{ over } x, \text{ then run } P
\mid (P\mid Q) \qquad \text{run } P \text{ and } Q \text{ in parallel}
\mid 0 \qquad \text{halt}
```

Example: buying cake at the pâtisserie

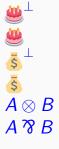


Pitfall: illegal resale of cake



$$\nu x.(x(s).x(s).x(s).0 | x(s).x(s).0)$$

Type
$$A, B ::= \alpha \mid \alpha^{\perp} \mid A \otimes B \mid A \aleph B \mid \dots$$



channel from which you get cake
channel in which you put cake
channel from which you get money
channel in which you put money
pair of independent channels
pair of possibly dependent channels

Type
$$A, B ::= \alpha \mid \alpha^{\perp} \mid A \otimes B \mid A \aleph B \mid \dots$$

$$\frac{P \vdash \Gamma, y : A \quad Q \vdash \Delta, x : B}{\nu y. x \langle y \rangle. (P \mid Q) \vdash \Gamma, \Delta, x : A \otimes B} \otimes$$

$$\frac{P \vdash \Gamma, y : A, x : B}{x(y).P \vdash \Gamma, x : A \aleph B} \aleph$$

Type
$$A, B ::= \alpha \mid \alpha^{\perp} \mid A \otimes B \mid A \otimes B \mid \ldots$$

$$\frac{P \vdash \Gamma, x : A \qquad Q \vdash \Delta, x : A^{\perp}}{\nu x \cdot (P \mid Q) \vdash \Gamma, \Delta} CUT$$

Type of x used by \bigcirc Type of x used by \bigcirc















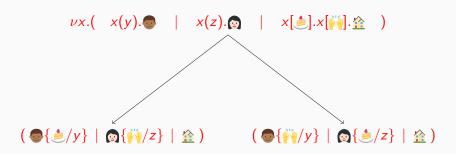


$$\nu x.(x(y). \bigcirc | \nu w.x \langle w \rangle.(w \leftrightarrow \overline{s} | x(z). \bigcirc))$$

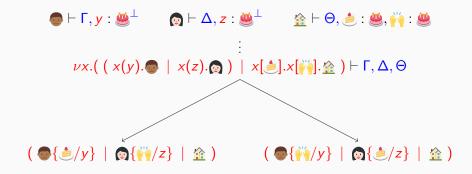
$$x[y].P ::= \nu y.x\langle y \rangle.P$$

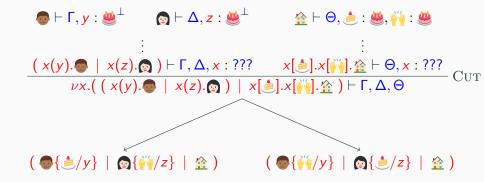
$$\nu x.(x(y). \bigcirc | x[w].(w \leftrightarrow \bigcirc | x(z). \bigcirc))$$

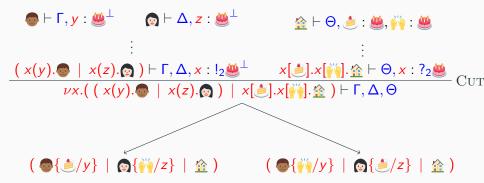
Example: buying cake at the pâtisserie



$$\frac{P \vdash \Gamma, x : A \qquad Q \vdash \Delta, x : A^{\perp}}{\nu x \cdot (P \mid Q) \vdash \Gamma, \Delta} CUT$$







Type A, B ::= $\cdots | ?_n A | !_n A$

Type
$$A, B ::= \cdots \mid ?_n A \mid !_n A$$

$$\frac{P \vdash \Gamma, y : A}{x[y].P \vdash \Gamma, x : ?_1 A} ?_1 \qquad \frac{P \vdash \Gamma, y : A}{x(y).P \vdash \Gamma, x : !_1 A} !_1$$

$$\frac{P \vdash \Gamma, x : ?_m A, y : ?_n A}{P\{x/y\} \vdash \Gamma, x : ?_{m+n} A} \text{ Contract}$$

$$\frac{P \vdash \Gamma, x : !_m A \qquad Q \vdash \Delta, x : !_n A}{(P \mid Q) \vdash \Gamma, \Delta, x : !_{m+n} A} \text{ Pool}$$

Cut-elimination/Communication

$$\frac{P \vdash \Gamma, x : ?_n A \qquad Q \vdash \Delta, x : !_n A^{\perp}}{\nu x. (P \mid Q) \vdash \Gamma, \Delta}$$
 CUT

Cut-elimination/Communication

$$\frac{P \vdash \Gamma, x : ?_{n}A}{x \uparrow y_{1} \cdots y_{n}.P \vdash \Gamma, y_{1} : A \cdots y_{n} : A} \text{EXP}$$

$$\frac{P \vdash \Gamma, y_{1} : A \cdots y_{n} : A}{x \downarrow y_{1} \cdots y_{n}.Q \vdash \Delta} \text{INT}$$

Type A, B ::= $\cdots | ?_n A | !_n A$

Cut-elimination/Communication

Cut-elimination/Communication

$$\frac{P \vdash \Gamma, x : ?_{n}A \qquad Q \vdash \Delta, x : !_{n}A^{\perp}}{\nu x . (P \mid Q) \vdash \Gamma, \Delta}$$
CUT

$$\frac{P \vdash \Gamma, x : ?_{n}A}{x \uparrow y_{1} \cdots y_{n}.P \vdash \Gamma, y_{1} : A \cdots y_{n} : A} \text{Exp} \qquad Q \vdash \Delta, x : !_{n}A^{\perp}}{x \downarrow y_{1} \cdots y_{n}.(x \uparrow y_{1} \cdots y_{n}.P \mid Q) \vdash \Delta} \text{Int}$$

Related and Future Work

I should mention...

- type system implemented in Agda;
- cut-elimination procedure implemented in Agda;

And I'm looking into...

- extension with resource quantifiers;
- embedding CP back into CP_{NP}.

Type
$$A, B ::= \cdots \mid ?_n A \mid !_n A$$

$$\frac{P \vdash \Gamma, y : A}{x[y].P \vdash \Gamma, x : ?_1 A} ?_1 \qquad \frac{P \vdash \Gamma, y : A}{x(y).P \vdash \Gamma, x : !_1 A} !_1$$

$$\frac{P \vdash \Gamma, x : ?_m A, y : ?_n A}{P\{x/y\} \vdash \Gamma, x : ?_{m+n} A} \text{ Contract}$$

$$\frac{P \vdash \Gamma, x : !_m A \qquad Q \vdash \Delta, x : !_n A}{(P \mid Q) \vdash \Gamma, \Delta, x : !_{m+n} A} \text{ Pool}$$

Resource Quantifiers

Type
$$A, B ::= \ldots \mid \forall n.A \mid \exists n.A$$

$$!A := \forall n.!_n A$$

 $?A := \exists n.?_n A$

$$Server(A) := \forall n.?_n A$$

$$Client(A) := \exists n.!_n A$$