Interprocess Communication

Interprocess Communciation

- What is IPC?
 - Mechanisms to transfer data between processes
- Why is it needed?
 - Not all important procedures can be easily built in a single process



Interprocess Communication

- Cooperating processes
 - Can affect or be affected by other processes, including sharing data
 - Just like cooperating threads!
 - Benefits
 - Information sharing
 - Computation speedup
 - Modularity
 - Convenience



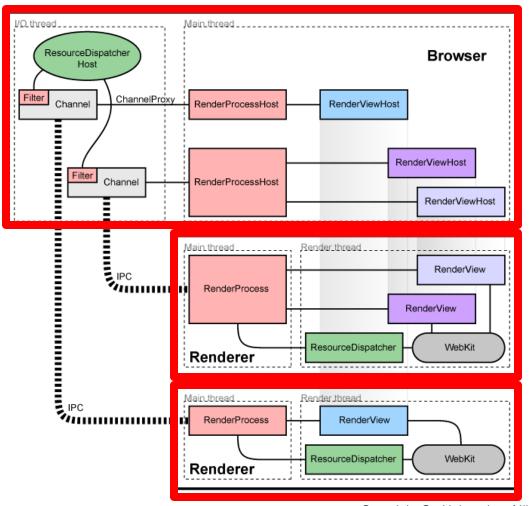
Interprocess Communication

Can you think of a common use of IPC?

Can you think of any large applications that use IPC?



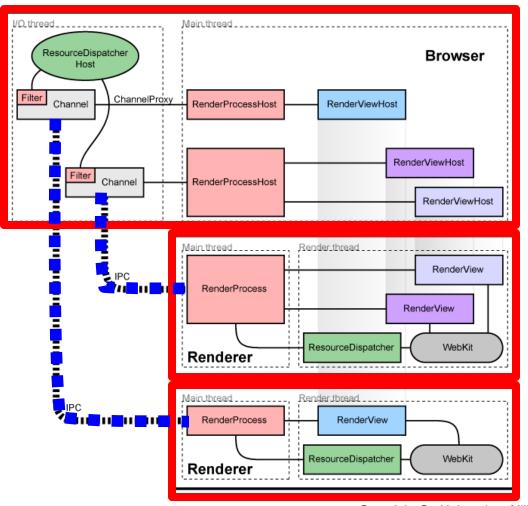
Google Chrome architecture (figure borrowed from Google)



- Separate processes for browser tabs to protect the overall application from bugs and glitches in the rendering engine
- Restricted access from each rendering engine process to others and to the rest of the system

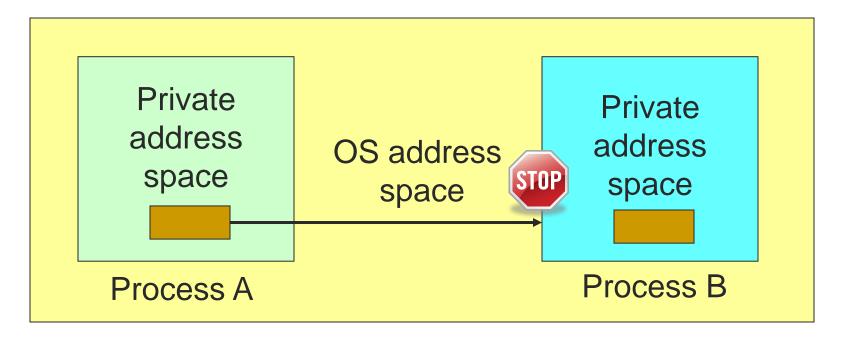


Google Chrome architecture (figure borrowed from Google)



- A named pipe is allocated for each renderer process for communication with the browser process
- Pipes are used in asynchronous mode to ensure that neither end is blocked waiting for the other

IPC Communications Model



- Each process has a private address space
- No process can write to another process's space
- How can we get data from process A to process B?



IPC Solutions

- Two options
 - Support some form of shared address space
 - Shared memory, memory mapped files
 - Use OS mechanisms to transport data from one address space to another
 - Pipes, FIFOs
 - Messages, signals

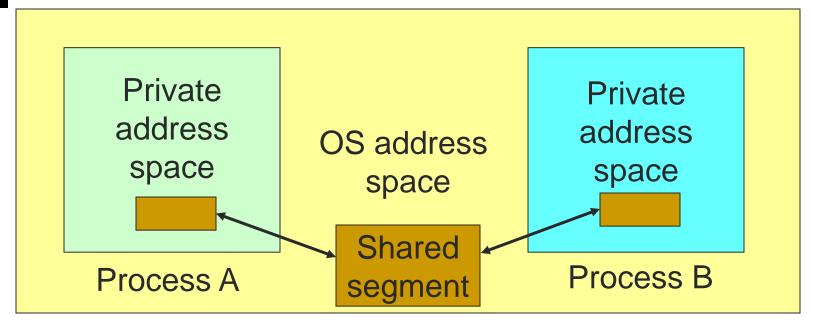


Shared Memory

- Processes share the same segment of memory directly
 - Memory is mapped into the address space of each sharing process
 - Memory is persistent beyond the lifetime of the creating or modifying processes (until deleted)
- Mutual exclusion must be provided by processes using the shared memory



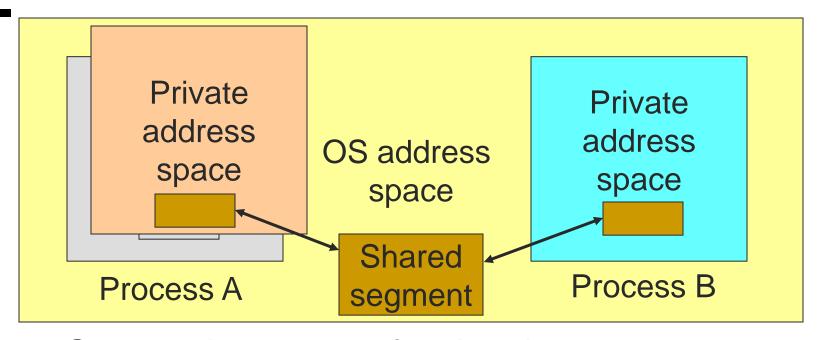
Shared Memory



- Processes request the segment
- OS maintains the segment
- Processes can attach/detach the segment



Shared Memory



Can mark segment for deletion on last detach



POSIX Shared Memory

```
#include <sys/types.h>
#include <sys/shm.h>
```

Create identifier ("key") for a shared memory segment

```
key_t ftok(const char *pathname, int proj_id);
k = ftok("/my/file", 0xaa);
```

Create shared memory segment

```
int shmget(key_t key, size_t size, int shmflg);
id = shmget(key, size, 0644 | IPC_CREAT);
```

Access to shared memory requires an attach

```
void *shmat(int shmid, const void *shmaddr, int shmflg);
shared_memory = (char *) shmat(id, (void *) 0, 0);
```



POSIX Shared Memory

Write to the shared memory using normal system calls

```
sprintf(shared_memory, "Writing to shared
    memory");
```

Detach the shared memory from its address space

```
int shmdt(const void *shmaddr);
shmdt(shared memory);
```



Shared Memory example

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <sys/types.h>
#include <sys/ipc.h>
#include <sys/shm.h>
#define SHM SIZE 1024 /* a 1K shared memory segment */
int main(int argc, char *argv[]) {
   key t key;
   int shmid;
   char *data;
   int mode;
```



Shared Memory example

```
/* make the key: */
if ((key = ftok("shmdemo.c", 'R')) == -1) {
    perror("ftok");
    exit(1);
/* connect to (and possibly create) the segment: */
if ((shmid = shmget(key, SHM SIZE, 0644 | IPC CREAT))
    perror("shmget");
    exit(1);
/* attach to the segment to get a pointer to it: */
data = shmat(shmid, (void *)0, 0);
if (data == (char *)(-1)) {
    perror("shmat");
    exit(1);
```

Shared Memory example

```
/* read or modify the segment, based on the command line: */
if (argc == 2) {
    printf("writing to segment: \"%s\"\n", argv[1]);
    strncpy(data, argv[1], SHM_SIZE);
} else
    printf("segment contains: \"%s\"\n", data);
```

```
/* detach from the segment: */
if (shmdt(data) == -1) {
   perror("shmdt");
   exit(1);
}
```

```
return 0;
```

I

Memory Mapped Files

- Memory-mapped file I/O
 - Map a disk block to a page in memory
 - Allows file I/O to be treated as routine memory access
- Use
 - File is initially read using demand paging
 - When needed, a page-sized portion of the file is read from the file system into a physical page of memory
 - Subsequent reads/writes to/from that page are treated as ordinary memory accesses

OK, we haven't really talked about memory yet, so bear with us ...

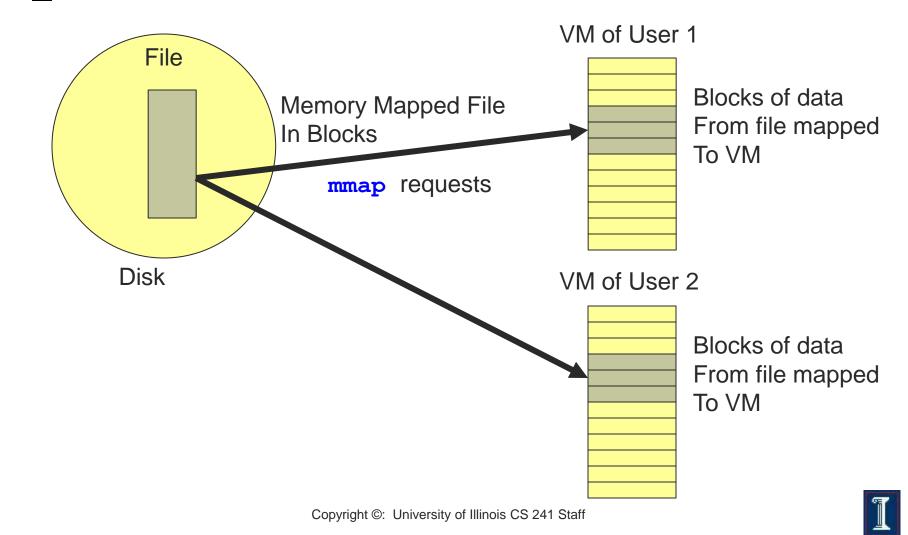


Memory Mapped Files

- Traditional File I/O
 - Calls to file I/O functions (e.g., read() and write())
 - First copy data to a kernel's intermediary buffer
 - Then transfer data to the physical file or the process
 - Intermediary buffering is slow and expensive
- Memory Mapping
 - Eliminate intermediary buffering
 - Significantly improve performance



Memory Mapped Files



-Memory Mapped Files: Benefits

- Treats file I/O like memory access rather than read(), write() system calls
 - Simplifies file access; e.g., no need to fseek()
- Streamlining file access
 - Access a file mapped into a memory region via pointers
 - Same as accessing ordinary variables and objects
- Dynamic loading
 - Map executable files and shared libraries into address space
 - Programs can load and unload executable code sections dynamically



-Memory Mapped Files: Benefits

- Several processes can map the same file
 - Allows pages in memory to be shared -- saves memory space
- Memory persistence
 - Enables processes to share memory sections that persist independently of the lifetime of a certain process



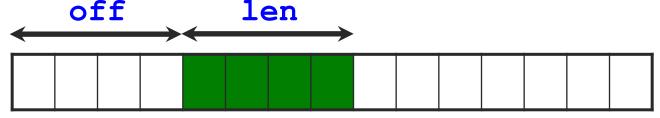


- Memory map a file
 - Establish mapping from the address space of the process to the object represented by the file descriptor
- Parameters:
 - addr: the starting memory address into which to map the file
 - len: the length of the data to map into memory
 - o prot: the kind of access to the memory mapped region
 - flags: flags that can be set for the system call
 - fd: file descriptor
 - off: the offset in the file to start mapping from



- Memory map a file
 - Establish mapping from the address space of the process to the object represented by the file descriptor

File fd

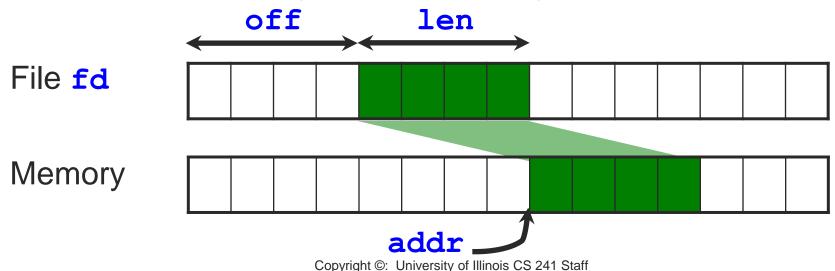




- Memory map a file
 - Establish a mapping between the address space of the process to the memory object represented by the file descriptor
- Return value: pointer to mapped region
 - On success, implementation-defined function of addr and flags.
 - On failure, sets errno and returns **MAP FAILED**



- Memory map a file
 - Establish a mapping between the address space of the process to the memory object represented by the file descriptor





mmap options

Protection Flags

O PROT READ

O PROT WRITE

O PROT_EXEC

O PROT NONE

Data can be read

Data can be written

Data can be executed

Data cannot be accessed

Flags

MAP_SHARED

O MAP PRIVATE

O MAP FIXED

Changes are shared.

Changes are private.

Interpret addr exactly



mmap Example

Map first 4kb of file and read an integer

```
#include <errno.h>
#include <fcntl.h>
#include <sys/mman.h>
#include <sys/types.h>
int main(int argc, char *argv[]) {
   int fd:
   void *pregion;
   if (fd = open(argv[1], O RDONLY) <0) {</pre>
       perror("failed on open");
       return -1;
   write(fd,"\0",1); // make sure at least 1 page is mapped
```

mmap Example

```
pregion = mmap(NULL, 4096, PROT READ,
               MAP SHARED, fd, 0);
if (pregion == MAP FAILED) {
  perror("mmap failed")
  return -1;
                 /* close the physical file */
close(fd);
/* access mapped memory; read the first int in
 * the mapped file */
int val = *((int*) pregion);
```

munmap



msync

```
#include <sys/mman.h>
int msync(void *addr, size t len, int flags);
    Write all modified data to permanent storage locations
    Return value
        0 on success
        -1 on error, sets errno
   Parameters:
        addr: returned from mmap ()
        len: same as the len passed to mmap ()
        flags:
           MS ASYNC = Perform asynchronous writes
           MS SYNC = Perform synchronous writes
           MS INVALIDATE = Invalidate cached data
```



Example 2: Shared memory using mmap

```
#include <stdio.h>
#include <stdlib.h>
#include <errno.h>
#include <fcntl.h>
#include <string.h>
#include <sys/mman.h>
#include <sys/types.h>
int main(int argc, char** argv) {
  int
          fd:
  char * shared mem;
  fd = open(argv[1], O RDWR | O CREAT);
  write(fd,"\0",1); // make sure at least 1 page is mapped
  shared mem = mmap(NULL, 10, PROT READ | PROT WRITE,
                      MAP SHARED, fd, 0);
  close(fd);
```



Example 2: Shared memory using mmap

```
Reader
           if (!strcmp(argv[2], "read")) {
             while (1) {
               printf("%s\n", shared mem);
               sleep(1);
Writer
                     while (1)
            scanf("%s\n", shared mem);
```



Recall POSIX Shared Mem...

```
#include <sys/shm.h>
int shmget(key t key, size t size, int
  shmflq);
Create shared memory segment
  id = shmget(key, size, 0644 | IPC CREAT);
void *shmat(int shmid, const void
  *shmaddr, int shmflg);

    Access to shared memory requires an attach

  shared memory = (char *) shmat(id, (void
     *) O, O);
```

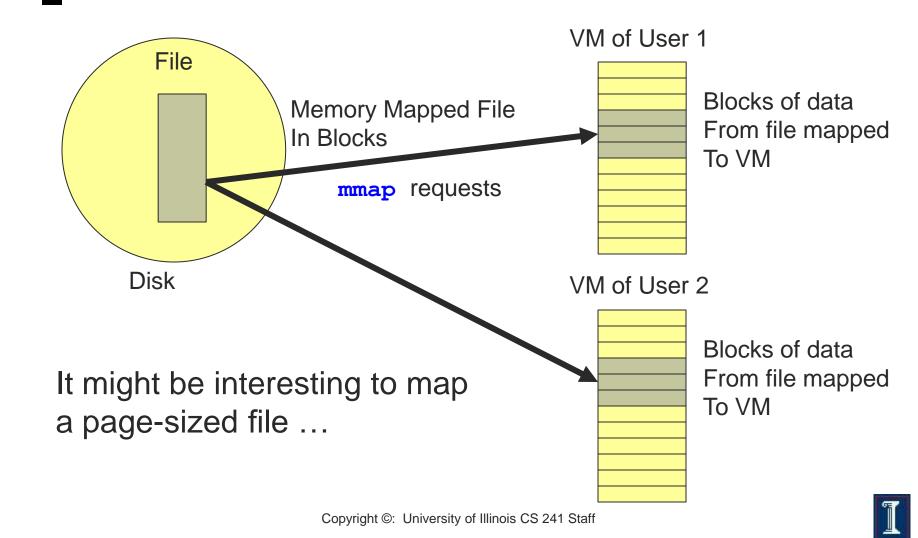


How do mmap and POSIX shared memory compare?

- Persistence
 - shm memory kept in memory
 - Remains available until system is shut down
 - mmap backed by a file
 - Persists even after programs quit or machine reboots



Memory mapped files and virtual memory



-Memory mapped files and virtual memory

```
#include <unistd.h>
long sysconf(int name);
```

- Determine the current value of a configurable system variable
- Return value
 - 0 on success
 - -1 on error, sets errno
- Parameters:
 - name: the system variable to be queried
 - SC_PAGESIZE



sysconf: Creating page-sized memory mapped segments

```
#include <errno.h>
#include <fcntl.h>
#include <unistd.h>
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <sys/types.h>
#include <sys/mman.h>
main(void) {
   size t bytesWritten = 0;
   int fd;
   int PageSize;
   const char text = "This is a test";
```



Example

```
if ((PageSize = sysconf( SC PAGE SIZE)) < 0) {</pre>
    perror("sysconf() Error=");
    return -1;
}
fd = open("/tmp/mmsyncTest", (O CREAT | O TRUNC |
          O RDWR), (S IRWXU | S IRWXG | S IRWXO));
if (fd < 0) {
    perror("open() error");
    return fd;
}
off t lastoffset = lseek(fd, PageSize, SEEK SET);
bytesWritten = write(fd, "x", 1);
if (bytesWritten != 1 ) {
    perror("write error. ");
    return -1;
}
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```



More Examples

```
/* mmap the file. */
void *address;
int len;
off t my offset = 0;
len = PageSize;
/* Map one page */
address = mmap(NULL, len, PROT WRITE, MAP SHARED, fd,
               my offset);
if (address == MAP FAILED) {
   perror("mmap error.");
    return -1;
}
```



More Examples

```
/* Move some data into the file using memory map. */
(void) strcpy((char*) address, text);
/* use msync to write changes to disk. */
if (msync(address, PageSize , MS SYNC) < 0 ) {</pre>
    perror("msync failed with error:");
    return -1;
} else
    (void) printf("%s","msync completed successfully.");
close(fd);
unlink("/tmp/msyncTest");
```

Run demo



Illegal Memory Access

- Use signals!
 - SIGSEGV signal allows you to catch references to memory that have the wrong protection mode
- Coming soon... signals!

