# Exercise 1 — assigned Thursday 1 February — due Tuesday 13 February

The project will be carried out in teams. The first task is to form and announce teams.

The implementation language is Haskell. More specifically, the project can be implemented in the Haskell 2010 standard language. If your team wishes to use GHC language extensions, you are free to do so; when you present your project, you will teach the class about any language extensions you have used.

The project will be carried out in several parts, called exercises.

For most parts of the project, you will be provided with some skeleton code in the form of type declarations, function signatures, or working components.

#### 1.1 Core lambda language

We start the project with a small core lambda language, consisting of the lambda calculus with booleans and integers. Later we will add further syntactic forms.

The reference for this exercise is Pierce, *Types and Programming Languages*, Chapters 5 and 9.

Here is the concrete syntax in BNF:

```
Type -->
    arrow lpar Type comma Type rpar
  | Bool_keyword
  | Int_keyword
Term -->
    identifier
  | abs_keyword | lpar identifier colon Type fullstop Term rpar
  | app_keyword lpar Term comma Term rpar
  | true_keyword
  | false_keyword
  | if_keyword Term then_keyword Term else_keyword Term fi_keyword
  | intliteral
  | plus lpar Term comma Term rpar
  | minus lpar Term comma Term rpar
  | mul lpar Term comma Term rpar
  | div lpar Term comma Term rpar
  | nand lpar Term comma Term rpar
  | equal lpar Term comma Term rpar
```

```
| lt lpar Term comma Term rpar
| lpar Term rpar
```

Here is an informal description of the terminal symbols (tokens) used in the grammar above:

```
arrow
                ->
                (
lpar
comma
                )
rpar
Bool_keyword
                Bool
Int_keyword
                Int
identifier
                an identifier, e.g., as in Haskell
abs_keyword
                abs
colon
fullstop
app_keyword
                app
true_keyword
                true
false_keyword false
if_keyword
                if
then_keyword
                then
                else
else_keyword
                fi
fi_keyword
intliteral
                a non-negative decimal numeral
plus
minus
mul
div
nand
equal
lt
```

White space, such as space, tab, and newline characters, is permitted between tokens. White space is required between adjacent keyword tokens.

Comment syntax is as in Haskell.

Here are some example programs:

```
app (abs (x: Int . 1234), 10) -- would look better without the space

if true then true else false fi

-- evaluates to true
```

```
if =(0,0) then 8 else 9 fi -- evaluates to 8
```

```
{- a division example -} /(4294967295,76)
```

## 1.2 Parsing

Write a parser for this language. (You may split this task into a scanner and a parser proper.)

### 1.3 Binding and free variables

Here you will define a handful of functions to manipulate the abstract syntax. Place them in the module *AbstractSyntax*.

Implement a function to enumerate the free variables of a term:

```
fv :: Term \rightarrow [Var]
```

Substitution: *subst* x s t, or in writing [ $x \mapsto s$ ]t, is the result of substituting s for the free occurrences of x in t. Implement it.

```
subst:: Var \rightarrow Term \rightarrow Term \rightarrow Term
```

#### 1.4 Arithmetic

The provided module *IntegerArithmetic* formalizes the primitive operators for integer arithmetic, which is 64-bit 2's complement, with silent overflow.

# 1.5 Type checker

It is always good to be sure a program is well-typed before we try to evaluate it. You can use the following type checker or write your own.

```
module Typing where
import Data.Maybe
import Data.List
import qualified AbstractSyntax as S
data Context = Empty
```

# | Bind Context S.Var S.Type deriving Eq

```
instance Show Context where
   show Empty
                               = "<>"
   show (Bind \Gamma x \tau) = show \Gamma + "," + x + ":" + show \tau
contextLookup :: S.Var \rightarrow Context \rightarrow Maybe S.Type
contextLookup \ x \ Empty = Nothing
contextLookup x (Bind \Gamma \gamma \tau)
      x \equiv y
                     = Just \tau
     | otherwise = contextLookup x \Gamma
typing :: Context \rightarrow S.Term \rightarrow Maybe S.Type
typing \Gamma t = \mathbf{case} \ t \ \mathbf{of}
   S.Var x \rightarrow contextLookup x \Gamma
   S.Abs x \tau_1 t_2 \to \mathbf{do} \tau_2 \leftarrow typing (Bind \Gamma x \tau_1) t_2; Just (S.TypeArrow \tau_1 \tau_2)
   S.App t_1 \ t_2 \rightarrow \mathbf{do} \ S.TypeArrow \ \tau_{11} \ \tau_{12} \leftarrow typing \ \Gamma \ t_1
       \tau \leftarrow typing \Gamma t_2
       if \tau \equiv \tau_{11} then Just \tau_{12} else Nothing
   S.Tru \rightarrow Just S.TypeBool
   S.Fls \rightarrow Just S.TypeBool
   S.If t_1 t_2 t_3 \rightarrow \mathbf{do} S.TypeBool \leftarrow typing \Gamma t_1
       \tau \leftarrow typing \Gamma t_2
       tau' \leftarrow typing \Gamma t_3
       if tau' \equiv \tau then Just \tau else Nothing
   S.IntConst \_ \rightarrow Just S.TypeInt
   S.IntAdd t_1 t_2 \rightarrow arith t_1 t_2
   S.IntSub\ t_1\ t_2 \rightarrow arith\ t_1\ t_2
   S.IntMul t_1 t_2 \rightarrow arith t_1 t_2
   S.IntDiv t_1 t_2 \rightarrow arith t_1 t_2
   S.IntNand t_1 t_2 \rightarrow arith t_1 t_2
   S.IntEq t_1 t_2 \rightarrow rel t_1 t_2
   S.IntLt \ t_1 \ t_2 \rightarrow rel \ t_1 \ t_2
   where
       arith t_1 t_2 = \mathbf{do} S.TypeInt \leftarrow typing \Gamma t_1; S.TypeInt \leftarrow typing \Gamma t_2; Just S.TypeInt
       rel t_1 t_2 = \mathbf{do} S.TypeInt \leftarrow typing \Gamma t_1; S.TypeInt \leftarrow typing \Gamma t_2; Just S.TypeBool
typeCheck :: S.Term \rightarrow S.Type
typeCheck\ t =
   case typing Empty t of
       Just \tau \rightarrow \tau
       _- \rightarrow error "type error"
```

#### 1.6 Main program

Write a main program which will (1) read the program text from a file into a string, (2) invoke the parser to produce an abstract syntax tree for the program, (3) print the free variables of the program, (4) in case the program is closed (has no free variables), type-check the program, and (5) in case the program has a type, evaluate it with respect to its call-by-value structural operational semantics, as provided.

Can you produce informative error messages from the parser and the type checker? Can you display the proof of typing? Can you display the individual evaluation steps and their proofs?

#### 1.7 Provided code

exercise1.lhs
AbstractSyntax.lhs
IntegerArithmetic.lhs
Typing.lhs
StructuralOperationalSemantics\_CBV.lhs
Latex.lhs

#### 1.8 What to submit

Prepare a project report (PDF). If you want, you can prepare the code and the report together in literate Haskell (like this assignment). Use Canvas to submit a tar file with all the code and the report.