

$C: \quad \rightarrow 011298247 \rightarrow 01898742 \rightarrow 01298784 \rightarrow 01248987$   
 $\rightarrow 01247988 \rightarrow 01347898 \rightarrow 01147889$

$2. \text{ have } T(n) \text{ as } n \log n + \frac{1}{2} [2n \log \frac{n}{2} + 2n \log \frac{n}{2}] + \frac{(n-1)}{2} \log(n-1) + bn$  have  $k = \frac{n}{2}, T_k = \frac{1}{2} 2k \log k + \frac{k}{2} \log k + \dots$   
 $= \log 2k \cdot k^{\frac{1}{2}} \cdot (k+1)^{\frac{1}{2}} + 4kk = f(k); f(1) = 1.4 + 4k; f(2) = 5.6 + 8k; f(3) = 10.7 + 12k$   
 $2k \log k = g(k); g(1) = 2.4; g(2) = 4.2; g(3) = 13.9$   
 $\Rightarrow g(k)$  might increase faster than  $f(k)$ , so  $Ch \log h \geq T(n)$ ,  $T(n)$  is  $O(n \log n)$   
 $3. \text{ have } v \text{ as the leaf of the tree: have to go through } h \text{ paths}$

$h$        $h-1$        $h-2$        $h-3$        $\dots$   
 $\text{operations} = h + (h-1) + (h-2) + \dots + 1 = \frac{(h+1)h}{2} = \frac{h^2 + h}{2}$   
 $\Rightarrow O(h^2)$