

Sabanci University, CS 302 AUTOMATA THEORY 17.1.2023

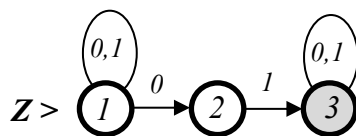
Answers to Final Examination

Answer 1 (25 pts)

(a) (5 pts)

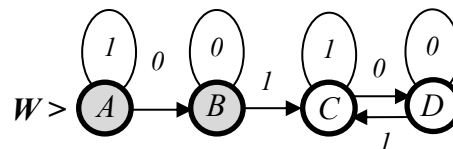
See the relevant slides.

(b) (10 pts) We first construct the NFA Z that accepts the complement language L^c as shown below.



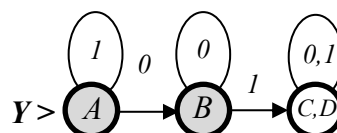
Using the table below we compute the DFA (hence NFA) W that accepts L

state	input	state'
$A^* := I$	0	$B := I, 2$
A^*	1	A
B^*	0	B
B^*	1	$C := I, 3$
C	0	$D := I, 2, 3$
C	1	C
D	0	D
D	1	C



(c) (10 pts) Minimal state Y is computed using the table filling algorithm as below

	C	D	A	B
C			0	0
D			0	0
A				1
B				



Answer 2 (25 pts)

(a) (12 pts)

$$E = (0+1)^*. (0.0.0+0.0.1 + 0.1.0 + 1.0.0) . (0+1)^* + (1+0).(1+0) + 1+0 + e$$

(b) (13 pts)

The language L in question can be expressed as below :

$$L = \{ w \in \{0,1\}^* \mid \text{if } k := \#0's \text{ in } w ; m := \#1's \text{ in } w ; \text{ then } |k-m| \leq 1 \}$$

L is NOT a regular language as demonstrated by the pumping lemma (PL) as shown below :

Choose $w = 0^{n+1}1^n \in L$ where n is the PL constant then by the PL

$$w = xyz \text{ and } |xy| \leq n ; |y| = q > 0 ; \text{ and } xy^jz \in L \text{ for } j=0, 1, 2, 3 \dots$$

Hence $xy = 0^p$, $p \leq n$ and so for $j=2$, $xy^2z \in L$

$$\text{But } x = 0^{p-q} , y = 0^q \text{ and } z = 0^{n+1-p}1^n \text{ so } xy^2z = 0^{p-q} 0^{2q} 0^{n+1-p} 1^n = 0^{q+n+1} 1^n \notin L$$

since $|q+n+1-n| = q+1 > 1$ and so PL is contradicted and L is not regular. and there is no regular expression E corresponding to L .

Answer 3 (25 pts)

(a) Given CFG : $S \rightarrow A \mid B ; A \rightarrow aAb \mid C \mid e ; C \rightarrow cCd \mid cC \mid e ; B \rightarrow Bb$
convert into CNF

(i) Eliminate null productions

$$S \rightarrow A \mid B ; A \rightarrow aAb \mid C \mid ab ; C \rightarrow cCd \mid cC \mid cd \mid c ; B \rightarrow Bb$$

(ii) Eliminate unit pairs

$$S \rightarrow aAb \mid cCd \mid cC \mid cd \mid c \mid ab \mid Bb ; A \rightarrow aAb \mid cCd \mid cC \mid cd \mid c \mid ab ;$$

$$C \rightarrow cCd \mid cC \mid cd \mid c ; B \rightarrow Bb$$

(iii) Eliminate the nongenerative variable B

$$S \rightarrow aAb \mid cCd \mid cC \mid cd \mid c \mid ab ; A \rightarrow aAb \mid cCd \mid cC \mid cd \mid c \mid ab ;$$

$$C \rightarrow cCd \mid cC \mid cd \mid c$$

(iv) Replace terminals by nonterminals

$$S \rightarrow XAY \mid UCV \mid UC \mid UV \mid c \mid XY ; A \rightarrow XAY \mid UCV \mid UC \mid UV \mid c \mid XY ;$$

$$C \rightarrow UCV \mid UC \mid UV \mid c ; X \rightarrow a ; Y \rightarrow b ; U \rightarrow c ; V \rightarrow d ;$$

,

(iv) Reduce triple nonterminals

$S \rightarrow WY \mid ZV \mid UC \mid UV \mid c \mid XY ; A \rightarrow WY \mid ZV \mid UC \mid UV \mid c \mid XY ;$

$C \rightarrow ZV \mid UC \mid UV \mid c ; X \rightarrow a ; Y \rightarrow b ; U \rightarrow c ; V \rightarrow d ;$

$W \rightarrow XA ; Z \rightarrow UC$

(b) $S \Rightarrow WY \Rightarrow XAY \Rightarrow aAY \Rightarrow aWYY \Rightarrow aXAYY \Rightarrow aaAYY \Rightarrow aaZVYY$
 $\Rightarrow aaUCVYY \Rightarrow aacCVYY \Rightarrow aaccVYY \Rightarrow aaccdYY \Rightarrow aacddbY \Rightarrow$
 $aacddb$

(c) $L = \{ a^n c^m d^k b^n \mid n, m, k \geq 0, n+m+k > 0, m \geq k \}$

Answer 4 (25 pts)

(a) (5 pts) L_1 is a CFL generated by $G = (\{S, A, B, C\}, \{a, b, c\}, R, S)$ where R is given as below.

$S \rightarrow AC ; A \rightarrow aAb \mid B ; B \rightarrow aB \mid a ; C \rightarrow cC \mid c$

(b) (10 pts) L_2 is not a CFL as proved by the PL below.

Given $n > 0$ choose $z = a^{n+1} b^{n+1} c^n \in L_2$ and $|z| = 3n+2 > n$

Then by PL $z = uvwxy$, $|vwx| \leq n$, $|vx| = q > 0$ and $uv^jwx^jy \in L_2$ for $j=0,1,2,3,\dots$

But (i) $vwx = a^k$ or $= b^k$ or (ii) $= c^k$ where $k \leq n$. If (i) prevails then for $j=0$

$uwy = a^{n+1-q} b^{n+1} c^n$ or $= a^{n+1} b^{n+1-q} c^n$ both not in L_2 since $n+1-q \neq n+1$ else if

(ii) prevails then for $j=3$ namely, $uv^3wx^3y = a^{n+1} b^{n+1} c^{n+2q}$ not in L_2 since $n+2q > n+1$.

Other two possibilities are (iii) $vwx = a^i b^j$ or $= b^i c^j$ with $i+j \leq n$ which again leads to

conclusion that either uwy is not in L_2 when $vwx = a^i b^j$; or when $vwx = b^i c^j$ with $i, j > 0$;

uwy is not in L_2 . Hence PL is contradicted and L_2 is not a CFL.

(c) (10 pts)

<i>Label</i>	<i>Condition</i>	<i>TM</i>
<i>M</i> >	-	$R_{\{a, b, c, \#\}} . A$
<i>A</i>	$\sigma = a$	$x . R_{\{b, c, \#\}} . B$
	<i>else</i>	h_{NO}
<i>B</i>	$\sigma = b$	$x . R_{\{a, c, \#\}} . C$
	<i>else</i>	h_{NO}
<i>C</i>	$\sigma = c$	$x . R_{\{a, b, \#\}} . D$
	<i>else</i>	h_{NO}
<i>D</i>	$\sigma = \#$	$L_{\#} . E$
	<i>else</i>	h_{NO}
<i>E</i>	$\sigma = a$	$x . R_{\{b, c, \#\}} . F$
	$\sigma = \#$	h_{YES}
	<i>else</i>	h_{NO}
<i>F</i>	$\sigma = b$	$x . R_{\{a, c, \#\}} . G$
	<i>else</i>	h_{NO}
<i>G</i>	$\sigma = c$	$x . R_{\{a, b, \#\}} . H$
	<i>else</i>	h_{NO}
<i>H</i>	$\sigma = \#$	$L_{\#} . E$
	<i>else</i>	h_{NO}