Gorken YAR 27970 Gorlens

Question 1:

They are some since + operator means union we can simplify by ensing one of them

$$=) \left(\left(0^{*} \cdot 1^{*} \right)^{*} + 1^{*} \right)^{*}$$

$$=) \left(\left(0+1 \right)^{*} + 1^{*} \right)^{*}$$

$$\Rightarrow (a+1) \Rightarrow$$

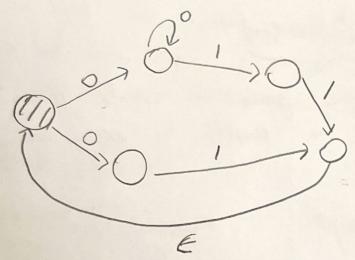
$$(L + M^*)^* \Rightarrow (L^* + (M^*)^*)^* \Rightarrow (L^* + M^*)^* \Rightarrow (L + M)^*$$

0)
$$(L, M^*)^*$$
 if $e \in L \Rightarrow (L, M^*)^* = (L^*, M^*)^* = (L+M)^*$

if e & L I here is no simplification.

Gorlan YAR 27970 Corbert

Question _ 2 ;



 Question 3.1.1

C) We can divide this regular expression into two parts:

1- the set of thrings does not have any consecutive 1's

2- the set of strings that consist exactly one pair of consecutive 1's.

1- (1+e). (0+01)* Part Charles

2- (0+10)*.1.1. (0+01)*

Any string that does not include include consecutive consecutive consecutive is and of one with o

RE = (1+e), (0+01)* + (0+10)*.1.1. (0+01)*

CS-302 HW2

Görken YAR 27970 60The

Question 3.1.4

6) (0+1+)* 000 (0+1)*

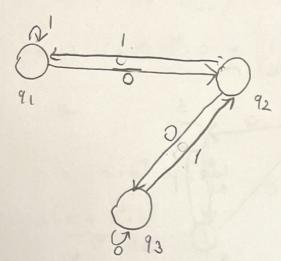
(0+1)*.000.(0+1)*

A language that accepts strings that contains three consecution zeros at least ones. Allow I Take Charles

c) (0 +10)*.1*

A language that does not include consecutive I's before the last o element.

Question 3.2.1



 $R_{11}^{\circ} = 1 + e$ $R_{12}^{\circ} = e$ $R_{12}^{\circ} = 0$ $R_{21}^{\circ} = 1$ $R_{33}^{\circ} = 0 + e$ $R_{31}^{\circ} = \emptyset$ $R_{13}^{\circ} = \emptyset$ $R_{23}^{\circ} = 0$ $R_{23}^{\circ} = 0$

a)

 $R_{11}^{1} = R_{11}^{\circ} + R_{11}^{\circ} \cdot (R_{11}^{\circ})^{*} \cdot R_{11} = (+e) + (+e) \cdot (+e)^{*} \cdot (+e) = 1^{*}$ $R_{22}^{1} = R_{22}^{\circ} + R_{21}^{\circ} \cdot (R_{11}^{\circ})^{*} \cdot R_{12}^{\circ} = e + 1 \cdot (1+e)^{*} \cdot 0 = 1^{*} \cdot 0^{\circ}$ $R_{12}^{1} = R_{12}^{\circ} + R_{11}^{\circ} \cdot (R_{11}^{\circ})^{*} \cdot R_{12}^{\circ} = 0 + (1+e) \cdot (1+e)^{*} \cdot 0 = 1^{*} \cdot 0^{\circ}$ $R_{21}^{1} = R_{21}^{\circ} + R_{21}^{\circ} \cdot (R_{11}^{\circ})^{*} \cdot R_{12}^{\circ} = 1 + 1 \cdot (1+e)^{*} \cdot (+e) = 1 \cdot 1^{*}$ $R_{33}^{1} = R_{33}^{\circ} + R_{31}^{\circ} \cdot (R_{11}^{\circ})^{*} \cdot R_{13}^{\circ} = 0 + e$ $R_{31}^{1} = R_{31}^{\circ} + R_{31}^{\circ} \cdot (R_{11}^{\circ})^{*} \cdot R_{13}^{\circ} = 0 + e$ $R_{13}^{1} = R_{13}^{\circ} + R_{11}^{\circ} \cdot (R_{11}^{\circ})^{*} \cdot R_{13}^{\circ} = 0 + e$ $R_{23}^{1} = R_{23}^{\circ} + R_{21}^{\circ} \cdot (R_{11}^{\circ})^{*} \cdot R_{13}^{\circ} = 0 + e$ $R_{32}^{1} = R_{23}^{\circ} + R_{21}^{\circ} \cdot (R_{11}^{\circ})^{*} \cdot R_{13}^{\circ} = 0 + e$ $R_{32}^{1} = R_{32}^{\circ} + R_{31}^{\circ} \cdot (R_{11}^{\circ})^{*} \cdot R_{13}^{\circ} = 0 + e$ $R_{32}^{1} = R_{32}^{\circ} + R_{31}^{\circ} \cdot (R_{11}^{\circ})^{*} \cdot R_{13}^{\circ} = 0 + e$ $R_{32}^{1} = R_{32}^{\circ} + R_{31}^{\circ} \cdot (R_{11}^{\circ})^{*} \cdot R_{13}^{\circ} = 0 + e$

 $R_{11}^{2} = R_{11}^{1} + R_{12}^{1} \cdot (R_{22}^{1})^{2} \cdot R_{21} = 1 + (1.0) \cdot (1.1.0)^{2} \cdot (1.1.$

d) The language for the automation is $e_{13}^{3} = R_{13}^{2} + R_{13}^{2} \cdot (\epsilon_{33}^{2})^{*} \cdot R_{23}^{2}$ $\Rightarrow (l^{*}, b) \cdot (l \cdot l^{*}, o)^{*} \cdot o + (l^{*}, o) \cdot (l \cdot l^{*}, o)^{*} \cdot o \cdot (l \cdot l^{*}, o)^{*} \cdot o)$ $\Rightarrow a + a \cdot b^{*} \cdot b + b$ $\Rightarrow a \cdot b^{*}$ $\Rightarrow (l^{*}, o) \cdot (l \cdot l^{*}, o)^{*} \cdot o \cdot ((o + e) + 1 \cdot (l \cdot l^{*}, o)^{*}, o)^{*}$ The legytor expression that accepts the language.

Question 3.2.3

