

The pitfalls of inconsistency in modern systems

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Jonas Boner *(inspired by Pat Helland)* <http://jonasboner.com/talks/> - Life beyond (the) Illusion

of Present / Pat Helland – Life Beyond Distributed Transactions: an Apostate's Opinion / Pat Helland, Dave Campbell – Building on Quicksand (Memories, Guesses, Apologies)

**IF WE CAN'T
COORDINATE AND
CANNOT BE CERTAIN -
IN REAL LIFE WE TAKE
AN EDUCATED GUESS -
AND WHEN WE ARE
WRONG WE APOLOGISE**



Are apologies always enough?

**Tesla driver killed in
crash with Autopilot
active, NHTSA
investigating**

<http://www.theverge.com/2016/6/30/12072408/tesla-autopilot-car-crash-death-autonomous-model-s>

Apologies...
Enough?

Are apologies always enough?

**IF WE CAN'T
COORDINATE AND
CANNOT BE CERTAIN -
IN REAL LIFE WE TAKE
AN EDUCATED GUESS -
AND WHEN WE ARE
WRONG WE GO TO JAIL**

Sometimes people go to jail for what they do... this is also “Real life”, isn't it?

*“Four former bankers have been **jailed** for their role in attempting to manipulate the interest rate benchmark Libor”*

What's a long running transaction?

- 1. Check Credit rating**
- 2. Do some manual checks (2 days)**
- 3. Decide whether to extend a loan or not**

3 steps *Business Process*

Is this always consistent?

- 1. Check Credit rating**
- 2. Do some manual checks (2 days)**
- 3. Decide whether to extend a loan or not**

What if ***Credit Rating*** changes while we are in ***Step 2***

We check the credit rating...

- 1. Check Credit rating**
- 2. Do some manual checks (2 days)**
- 3. Decide whether to extend a loan or not**

What if *Credit Rating* changes while we are in *Step 2*

Time passes...

1. **Check Credit rating**
2. **Do some manual checks (2 days)**
3. **Decide whether to extend a loan or not**

What if *Credit Rating* changes while we are in *Step 2*

And our initial Guess was wrong...

1. **Check Credit rating**
2. **Do some manual checks (2 days)**
3. **Decide whether to extend a loan or not**

What if *Credit Rating* changes while we are in *Step 2*

KA-BOOM

What do we do now?

- 1. Check Credit rating**
- 2. Do some manual checks (2 days)**
- 3. Decide whether to extend a loan or not**

But **Business people** have already sorted this out – basing on initial Credit Check they **promise you a loan, on condition** that your rating does not change and **then re-check**

There are different sorts of “apologies”...

- 1. Check Credit rating**
- 2. Do some manual checks (2 days)**
- 3. Decide whether to extend a loan or not**

Another Business Level decision might be to just go-ahead and acknowledge the risk from potential default.

But consequences may vary...

1. **Check Prohibited Country Status** *So...*
2. **Do some manual checks (2 days)**
3. **Decide whether to onboard or not**

And we may not always be able to ...

1. Check **Prohibited Country Status**
2. Do some manual checks (2 days)
3. **Decide whether to onboard or not**



... apologise

1. **Check Prohibited Country Status**
2. **Do some manual checks (2 days)**
3. **Decide whether to onboard or not**

Wrong decision can have your license revoked – be careful...

You are not always free to just take the risk...

Is this problem visible in computer systems too? In distributed systems? In databases?

YES, we've known that for quite a long time...

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RFC 677 (1975) Maintenance of duplicate databases

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Sagas Hector Garcia-Molina (1987) Long running transactions

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RFC 677 (1975) Maintenance of duplicate databases

Sagas Hector Garcia-Molina **(1987)** Long running transactions

A Critique of ANSI SQL Isolation Levels
Berenson et al. (1995)

Let me show you how this works in a traditional database

<https://blog.acolyer.org/2016/02/24/a-critique-of-ansi-sql-isolation-levels/>

[illegible]

What are the common defaults?

Highly Available Transactions: Virtues

and Limitations – Peter Bailis et al.

<http://www.bailis.org/papers/hat-vldb2014.pdf>

Database	Default	Maximum
Action Ingres 10.0/10S	S	S
Aerospike	RC	RC
Akiban Persistit	SI	SI
Clustrix CLX 4100	RR	RR
Greenplum 4.1	RC	S
IBM DB2 10 for z/OS	CS	S
IBM Informix 11.50	Depends	S
MySQL 5.6	RR	S
MemSQL 1b	RC	RC
MS SQL Server 2012	RC	S
NuoDB	CR	CR
Oracle 11g	RC	SI
Oracle Berkeley DB	S	S
Oracle Berkeley DB JE	RR	S
Postgres 9.2.2	RC	S
SAP HANA	RC	SI
ScaleDB 1.02	RC	RC
VoltDB	S	S
RC: read committed, RR: repeatable read, SI: snapshot isolation, S: serializability, CS: cursor stability, CR: consistent read		

What are the common defaults?

Highly Available Transactions: Virtues

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READ COMMITTED – allows Lost Update – T1 reads $x = 100$, T2 reads $x = 100$, T1 finishes committing and writes $x = 110$, T2 uses old *stale* value of $x = 100$, increases it by 30 and commits $x = 130$.

Why Serializable is not default??? – it impacts performance...

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MySQL 5.6	RR	S
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Serialize all the things?

A Critique of ANSI SQL Isolation Levels, H. Berenson et al.

Why **SERIALIZABLE** is not the default setting??? – (asked my Boss once... 😊) it severely impacts performance... and availability...

But is still an option...

Table 2. Degrees of Consistency and Locking Isolation Levels defined in terms of locks.		
Consistency Level = Locking Isolation Level	Read Locks on Data Items and Predicates (the same unless noted)	Write Locks on Data Items and Predicates (always the same)
Degree 0	none required	Well-formed Writes
Degree 1 = Locking READ UNCOMMITTED	none required	Well-formed Writes Long duration Write locks
Degree 2 = Locking READ COMMITTED	Well-formed Reads Short duration Read locks (both)	Well-formed Writes, Long duration Write locks
Cursor Stability (see Section 4.1)	Well-formed Reads Read locks held on current of cursor Short duration Read Predicate locks	Well-formed Writes, Long duration Write locks
Locking REPEATABLE READ	Well-formed Reads Long duration data-item Read locks Short duration Read Predicate locks	Well-formed Writes, Long duration Write locks
Degree 3 = Locking SERIALIZABLE	Well-formed Reads Long duration Read locks (both)	Well-formed Writes, Long duration Write locks

So why didn't all of us heard about this problem already?

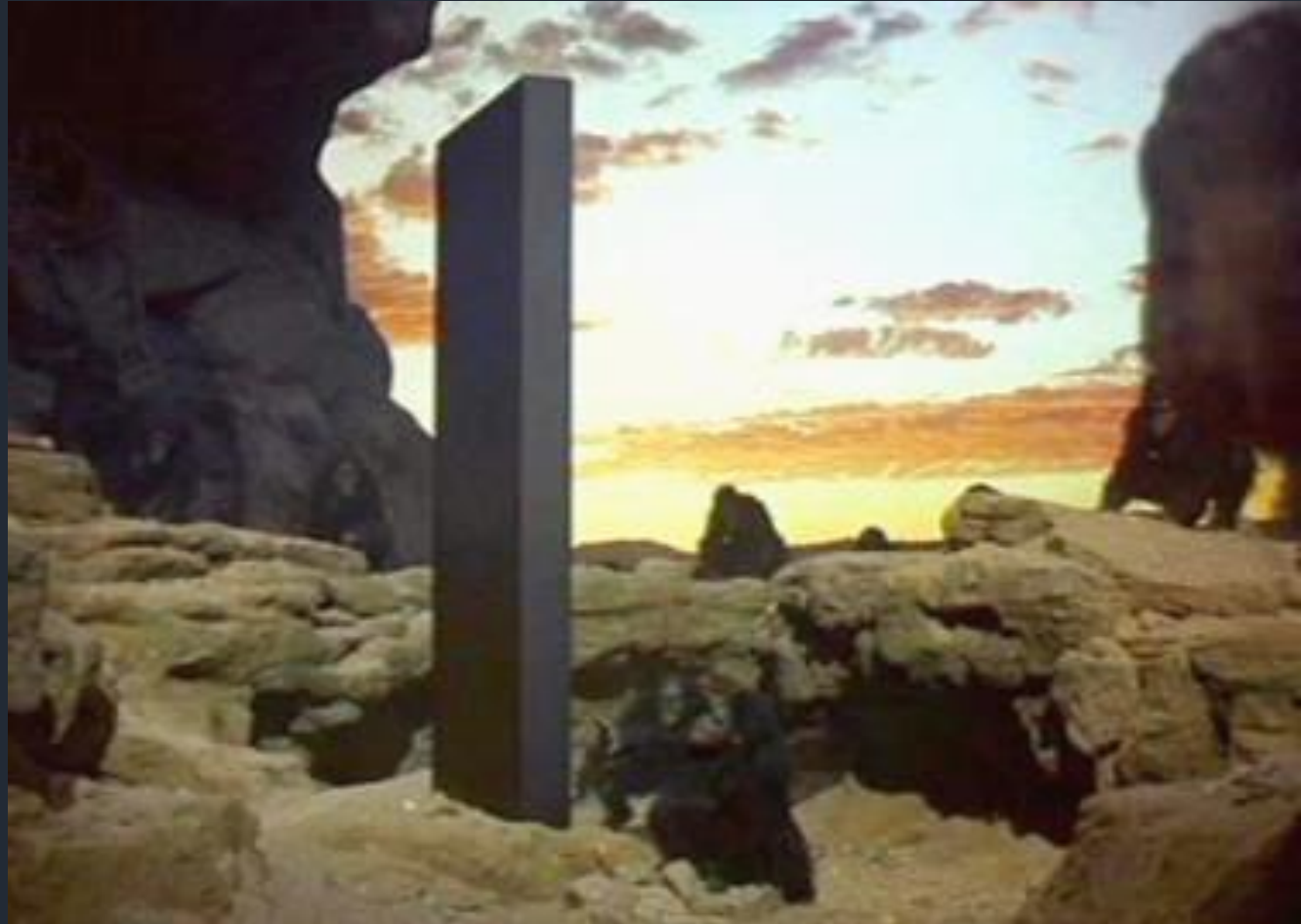
Database transactions have been very fast

because

All changes were LOCAL

So why didn't all of us heard about this problem already?

DB MONOLITH



By Source, Fair use,
<https://en.wikipedia.org/w/index.php?curid=31738209>

When did we start noticing?

We tried moving this successful DB-based model to “the cloud”.

When did we start noticing?

We tried moving this successful DB-based model to “the cloud”.

Cloud actually means distributed system.

Cloud

**So my Database is
now no longer a
Monolith, it is a
System. Changes are
no longer LOCAL
CAP theorem applies**

By Photo by John Kerstholt, original upload by Solitude (From English Wikipedia) [GFDL (<http://www.gnu.org/copyleft/fdl.html>) or CC-BY-SA-3.0 (<http://creativecommons.org/licenses/by-sa/3.0/>)], via Wikimedia Commons



Distributed consistency...

<https://aphyr.com/posts/313-strong-consistency-models>

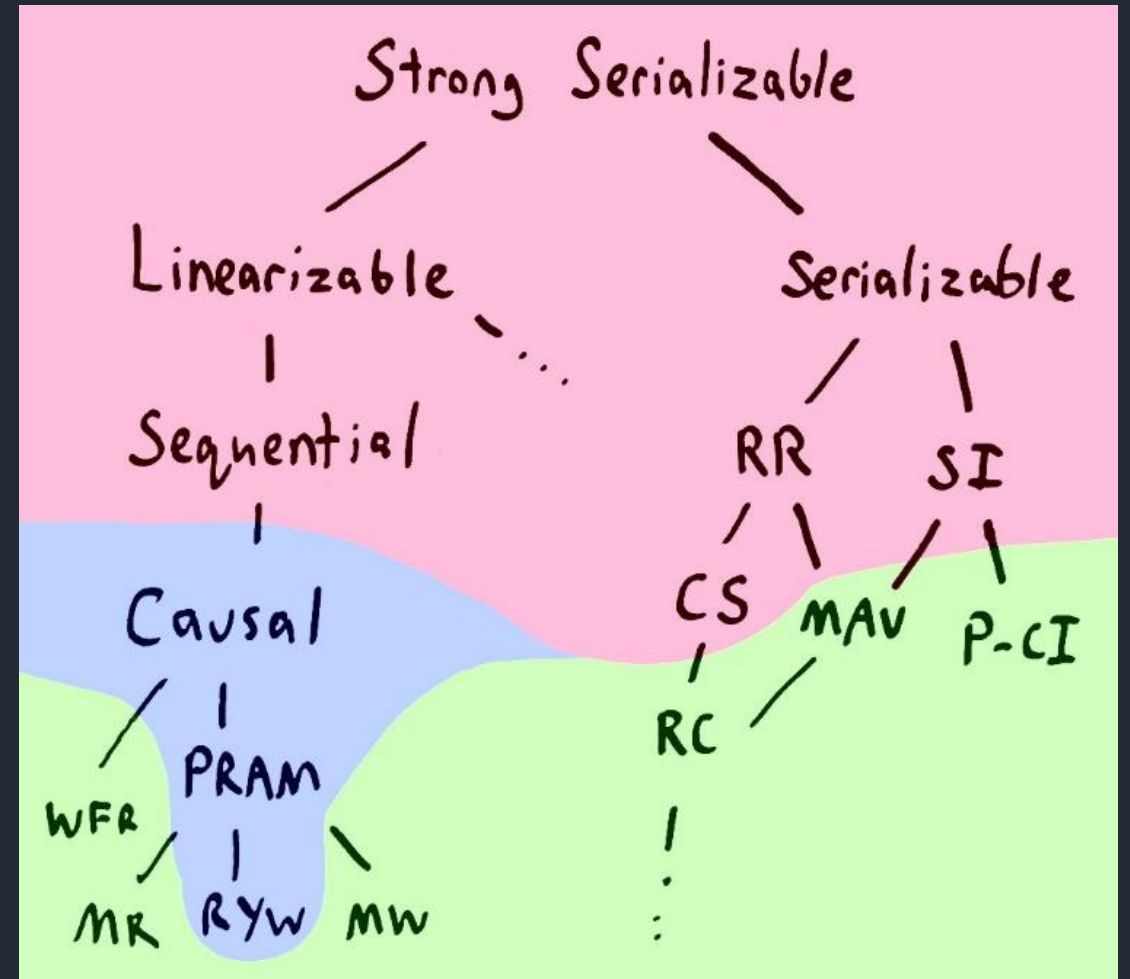
<http://www.bailis.org/blog/linearizability-versus-serializability/>



Linearizability: single-operation, single-object, real-time order



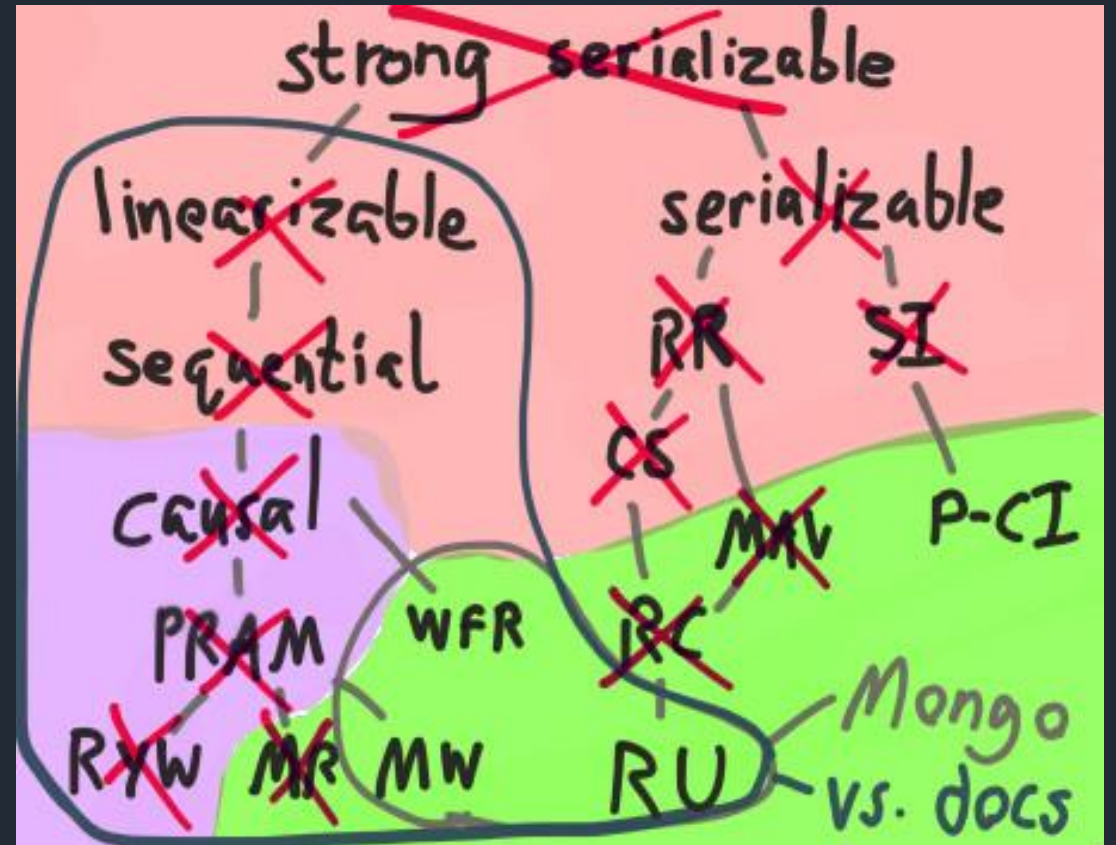
Serializability: multi-operation, multi-object, arbitrary total order



Typical cloud database - MongoDB

<https://aphyr.com/posts/322-jepsen-mongodb-stale-reads>

NOT GOOD



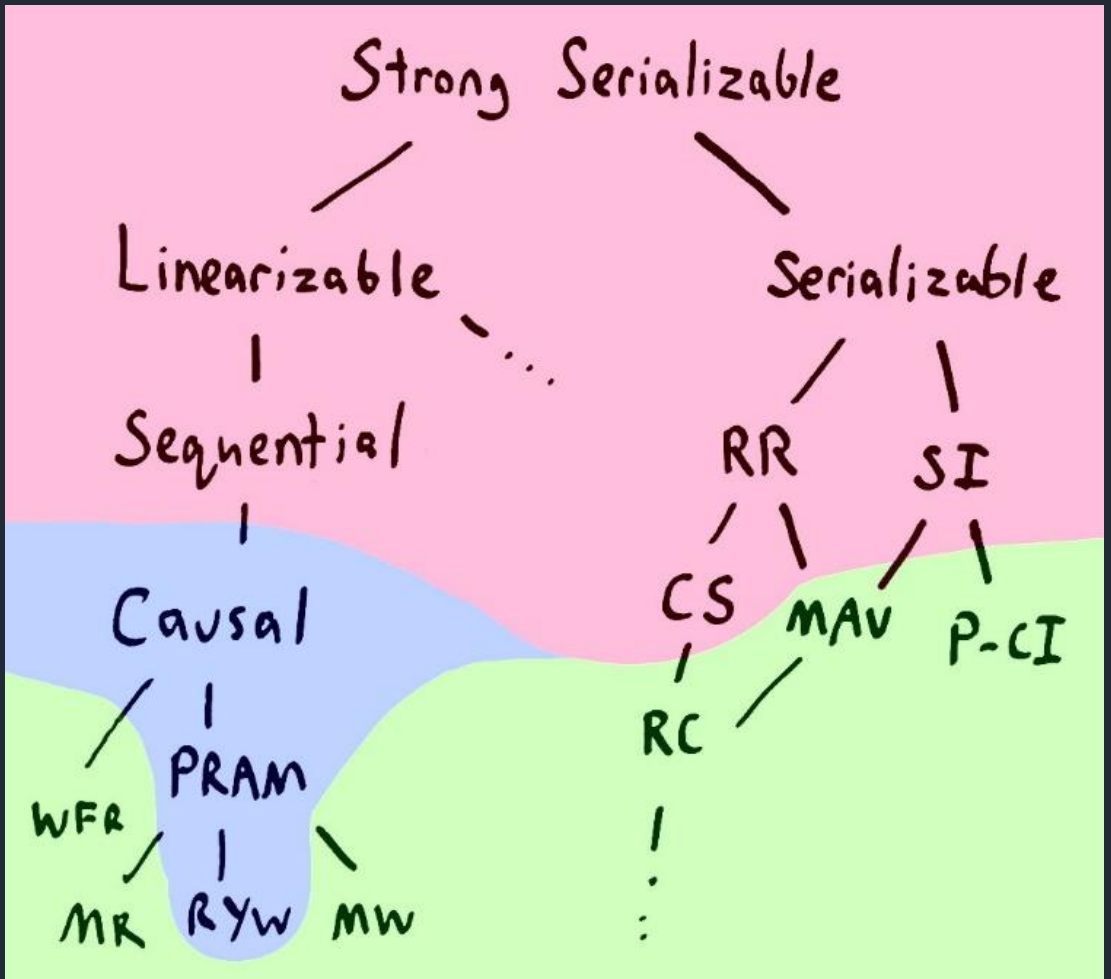
Why? – the HAT paper

Bailis et al. 2014

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GREEN ones do not impact availability



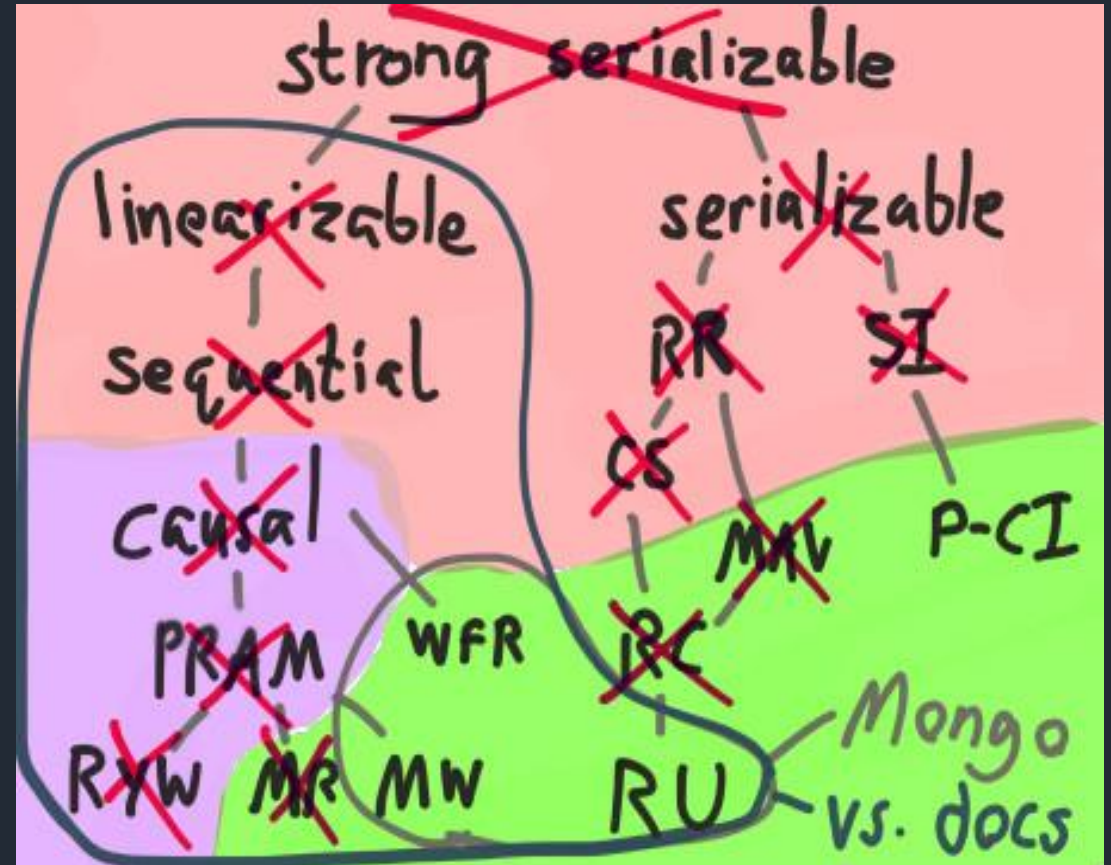
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Bailis et al. 2014

Highly Available Transactions: Virtues and Limitations – Peter

*When partition occurs,
Mongo chooses **Availability**
over **Consistency***

*Besides, being **Consistent** in
Distributed Land is **HARD** and
VERY SLOW*

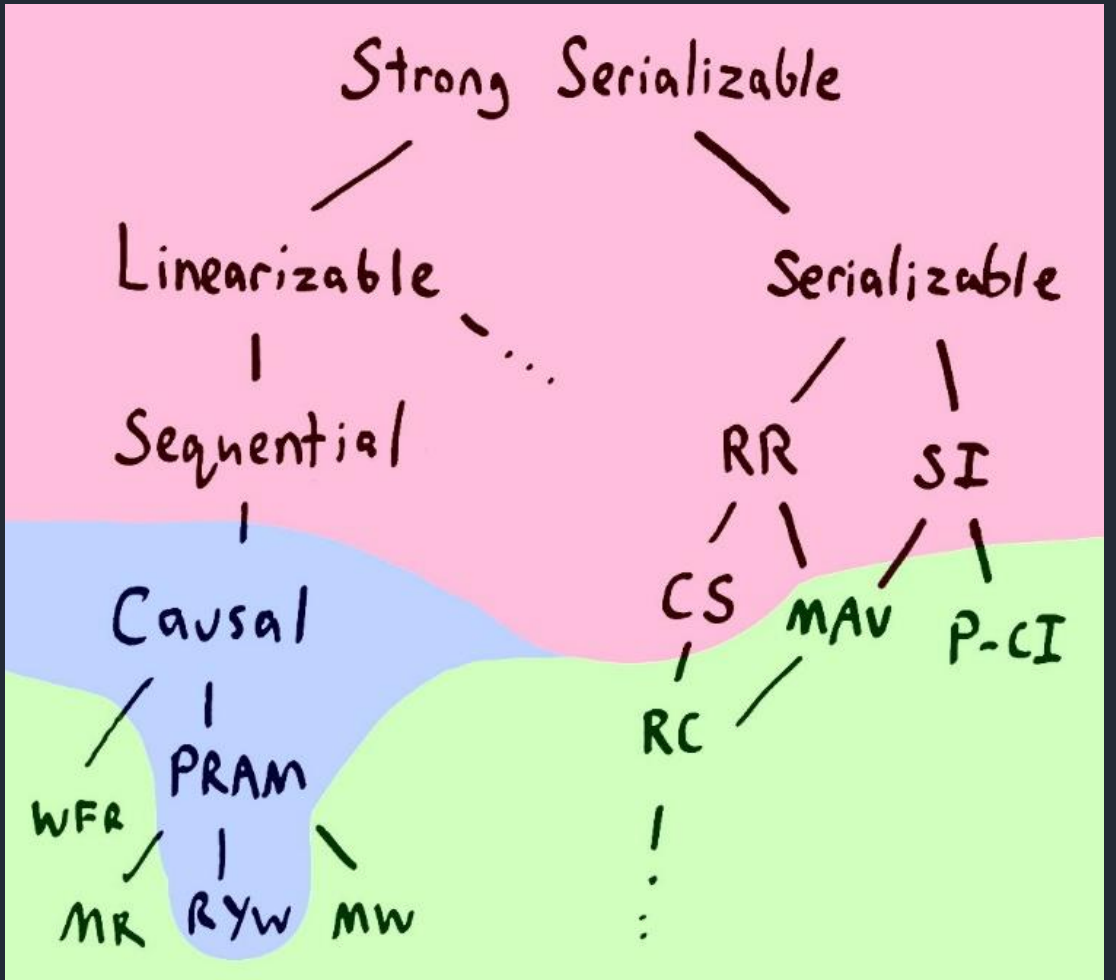


Why?

<http://www.bailis.org/papers/hat-vldb2014.pdf>

ROSE potential loss of availability

BLUE guaranteed by available system only inside a session

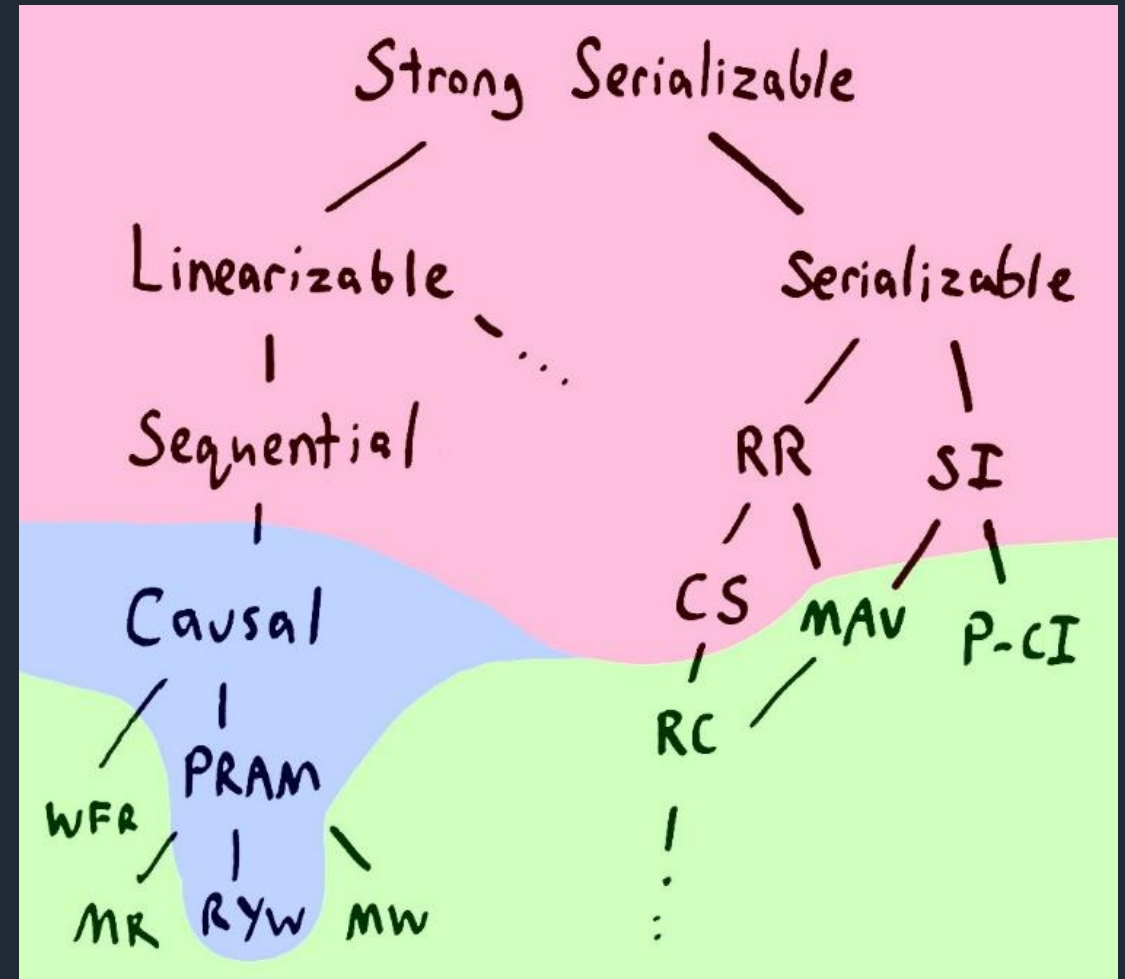


What left and right rose ones have in common?

<http://www.bailis.org/papers/hat-vldb2014.pdf>

ROSE prohibits Lost Updates

BLUE prohibits Lost Updates inside a session



But Lost Updates were not that bad in Monolith?

Because system was Local everything was fast and Lost Updates were very, very rare...

Plus you can make your system Strongly Consistent and only acceptably slower if you are Local.

Can we get back to Local?

Can we get better?

Can we get back to Local?

Microservices?

Database per Service pattern -

<http://microservices.io/patterns/data/database-per-service.html>

<http://www.lagomframework.com/blog/lagom-1-0.html>



What is the problem now?

Microservice orchestration – how to control the emergent behaviour of group of microservices performing a chain of transactions.

When partition heals, we either merge results or take back affected operations.

Merging 😊

<https://chuva-inc.com/blog/fast-tip-enable-git-rerere-right-now>

<https://git-scm.com/docs/git-rerere>

Doing merges is hard.

And requires domain knowledge.

the number of planets are

<<<<<< HEAD

nine

=====

eight

>>>>>> branch-a



Sagas – *Hector Garcia-Molina, 1987*

Also: <http://kellabyte.com/2012/05/30/clarifying-the-saga-pattern/> and “ACTA: The SAGA Continues” – Chrysanthis and Ramamritham.

Saga - one long running transaction consisting of many transactional steps - pair each transactional operation with its compensation.

L: T1 → T2 → T3 → T4

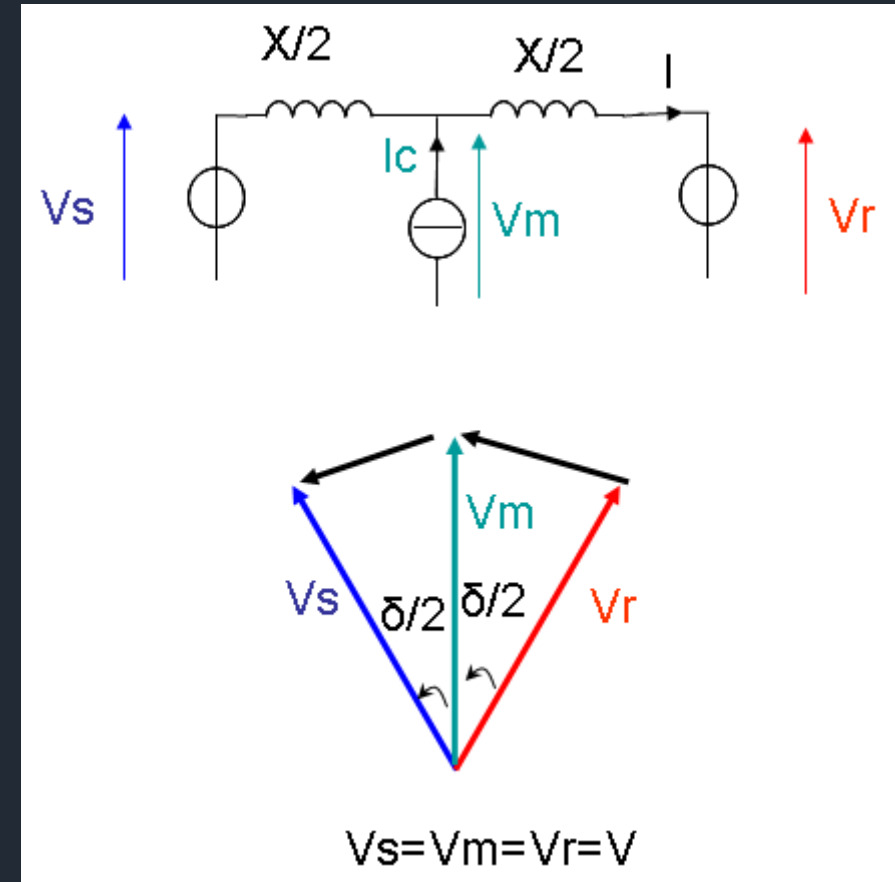
C1 ← C2 ← C3 ← C4

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**We just need to
define what should
the compensation
steps comprise of.**

Shunts compensation - By Hermione12 - Own work, Public Domain,
<https://commons.wikimedia.org/w/index.php?curid=2509621>



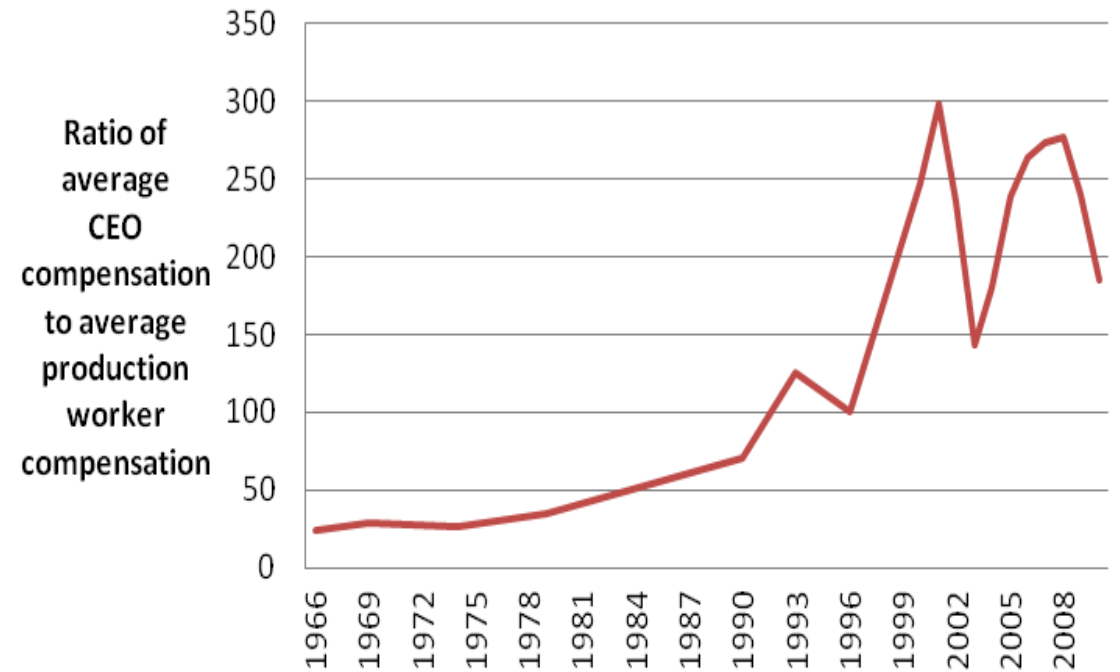
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**What is a good
compensation?**

**Is the compensation
possible at all?**

Growth of CEO pay in America



Exercise: What is the right level and manner of compensation?

Cases:

- **Missing transactions in the Web Banking app**
- **Knightmare on Wall Street**
 - <https://dougseven.com/2014/04/17/knightmare-a-devops-cautionary-tale/>
- **Multiplicated Visa Charges (up to 15x)**
 - <http://prawo.vagla.pl/node/7663>

Exercise: What is the right level of compensation?

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Final words

**When we allow inconsistency it is always a
Business level decision what to do.**

So it is a Functional Requirement

We need to communicate this clearly.

Thank you!

Questions?