HELLO DEVDAY

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GMU BCS '06

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Ruby '07 GMU BCS '06 Viget Labs Justin Marney

Viget '08 Ruby '07 GMU BCS '06 Viget Labs





DISTRIBUTING YOUR DATA

WHY?

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web applications are judged by their level of availability

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increase throughput

ability to manage availability during node failure

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increase durability

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web applications are judged by their level of availability

increase scalability

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ability to continue operating during failure scenarios

"I can add twice as much X and get twice as much Y."

X = processor, RAM, disks, servers, bandwidth

Y = throughput, storage space, uptime

scalability is a ratio.

2:2 = linear scalability ratio

scalability ratio allows you to predict how much it will cost you to grow.

UP/DOWN/VERTICAL/HORIZONTAL/L/R/L/R/A/B/START



UP grow your infrastructure multiple data centers higher bandwidth faster machines

DOWN
shrink your infrastructure
mobile
set-top
laptop

VERTICAL
add to a single node
CPU
RAM
RAID

HORIZONTAL
add more nodes
distribute the load
commodity cost
limited only by capital

@gary_hustwit: Dear Twitter: when a World Cup match is at the 90th minute, you might want to turn on a few more servers.

A distributed transaction is bound by availability of all nodes.

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$$(.99^1) = .99$$

$$(.99^2) = .98$$

$$(.99^3) = .97$$

Asynchronous systems operate without the concept of global state.

The concurrency model more accurately reflects the real world.

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What about my ACID!?



Atomic

Series of database operations either all occur, or nothing occurs.

Consistent

Transaction does not violate any integrity constraints during execution.

solated

Cannot access data that is modified during an incomplete transaction.

Durable

Transactions that have committed will survive permanently.



Defines a set of characteristics that aim to ensure consistency.

What happens when we realize that in order scale we need to distribute our data and handle asynchronous operations?

ACID

Without global state, no Atomicity.

Without a linear timeline, no transactions and no Isolation.

The canonical location of data might not exist, therefore no D.

Without A, I, or D, Consistency in terms of entity integrity is no longer guaranteed.

CAP Theorem

Eric Brewer @ 2000 Principles of Distributed Computing (PODC).

Seth Gilbert and Nancy Lynch published a formal proof in 2002.

CAP Acronym

Consistency: Multiple values for the same piece of data are not allowed.

Availability: If a non-failing node can be reached the system functions.

Partition-Tolerance: Regardless of packet loss, if a non-failing node is reached the system functions.

CAP Theorem

Consistency, Availability, Partition-Tolerance: Choose One...

CAP Theorem

Single node systems bound by CAP.

100% Partition-tolerant 100% Consistent No Availability Guarantee

CAP Theorem

Multi-node systems bound by CAP.

CA: DT, 2PC, ACID

CP: Quorum, distributed databases

AP: Dynamo, no ACID

CAP Theorem

CAP doesn't say AP systems are the solution to your problem.

Not an absolute decision.

Most systems are a hybrid of CA, CP, & AP.

CAP Theorem

Understand the trade-offs and use that understanding to build a system that fails predictably.

Enables you to build a system that degrades gracefully during a failure.

BASE

Dan Pritchett

BASE: An ACID Alternative

Associate for Computing Machinery Queue, 2008

BASE

BASE: An ACID Alternative

Basically Available
Soft State
Eventually Consistent

BASE

BASE: An ACID Alternative

Basically Available Soft State

Eventually Consistent

Eventually Consistent

Rename to Managed Consistency.

Does not mean probable or hopeful or indefinite time in the future.

Describes what happens during a failure.

Eventually Consistent

During certain scenarios a decision must be made to either return inconsistent data or deny a request.

EC allows you control the level of consistency vs. availability in your application.

Eventually Consistent

In order to achieve availability in an asynchronous system, accept that failures are going to happen.

Understand failure points and know what you are willing to give up in order to achieve availability.

How can we model the operations we perform on our data to be asynchronous & EC?

Model system as a network of independent components.

Partition components along functional boundaries.

Don't interact with your data as one big global state.

This doesn't meant every part of your system must operate this way!

Use ACID 2.0 to help identify and architect components than can.

ACID 2.0

Associative

Order of operations does not change the result.

Commutative

Operations can be aggregated in any order.

Idempotent

Operation can be applied multiple times without changing the result.

Distributed

Operations are distributed and processed asynchronously.

OPS BROS

Incremental scalability

Homogeneous node responsibilities

Heterogeneous node capabilities

LINKS

Base: An ACID Alternative

Into the Clouds on New Acid

Brewer's CAP theorem

Embracing Concurrency At Scale

Amazon's Dynamo



http://sorescode.com

http://github.com/gotascii

http://spkr8.com/s/1

@vigemarn