

# Optimizing corridor placement using simulated annealing

Michael D. Catchen<sup>1,2</sup>

<sup>1</sup> McGill University; <sup>2</sup> Québec Centre for Biodiversity Sciences

**Correspondance to:**

Michael D. Catchen — michael.catchen@mail.mcgill.ca

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how many different chapter ones will i have hmmmm

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## Introduction

Human activity has rapidly reshaped the face of Earth's surface, leaving fragments of patchy habitat. Although there is no shortage of debate as to the effects of fragmentation *per se* on biodiversity and ecosystem function (**cite?**), it is generally accepted that the combination of habitat and ensuing subdivision produce negative outcomes for ecosystem function and services (**resasco?** review).

In order to mitigate the consequences of landscape change on ecosystems, developing landscape *corridors* has seen much attention in the last several decades. Bit more evidence for corridors here. But still, the spatter of fragments in a landscape, where should ecologists choose to use their limit resources to build a corridor?

Here we propose to answer that question by proposing an algorithm to estimate the landscape modification that results in optimizing a specific ecosystem process (in this paper maximizing the time until extinction of a metapopulation, although the algorithm and associated software can be generally applied to any process-based model with a quantifiable target state).

Although algorithms have been proposed for this (**peterman?** etc), they are focused on finding the where the paths of least existance for a given species is given data on that species dispersal.

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## An algorithm for optimizing corridor placement

Start with some definitions and notation.

Define the set of possible landscape modifications,  $\mathfrak{M}$ , in optimization language called the *search-space*. Introduce uncountability argument of this space.

Because we cannot test every possible modification in  $\mathfrak{M}$ , we use simulated annealing, a method for estimating the global optimum of functions with NP search-spaces.

**2.1. Proposing landscape modifications** This is really important. We propose (no pun intended) several algorithms for generating landscape modifications. Some of the details here might have to go in a supplement/appendix.

**2.1.1 Graph-based** Consider only modifications that consist of connecting nodes. TODO: only choose topologies from the minimum spanning tree of the nodes.

**2.1.1.1 The two stage approach** Stage-one: accept a new topology of connected nodes with probability in proportion to chain temperature (see next section).

Stage-two: modify way that the connection for a given topological structure is chosen. Because we are working in a 2D raster, all distances between points are Manhattan distances, and any link between points is composed of  $x$  horizontal steps and  $y$  vertical steps. There are thus  $2^{\min(x,y)}$  ways to connect two nodes that far apart.

**2.1.2 Not graph-based** The reason to avoid this is because the search-space grows much faster with lattice size and budget. That being said, we can use some simply heuristics to weight proposals using “common-sense.”

**2.2. Simulated annealing to explore the space of landscape modifications** The transition probability function,  $q$ , which gives the probability of moving from one modification  $i \in \mathbb{M}$  to a new proposed state  $j \in \mathbb{M}$ , as a function of a chains temperature.

Here we define  $q(i, j)$  using a logistic function,

$$q(i, j, \alpha) = \frac{1}{1 + e^{\alpha(s(j)-s(i))}}$$

$s(i)$  is the function that gives the score of a proposed modification. Here, the mean time to extinction.

Simulated annealing can be written described as the following.

A markov-chain, denoted  $\pi_\alpha$

**Figure 1: concept fig**

**2.3. Process-based optimization (using occupancy dynamics)** Here we use occupancy dynamics as the process, although we emphasize that this method works for arbitrary process models and is instead limited only by the computational demands of a given process model.

Compute new resistance surface which gives pairwise potential values for each pair of points.

**2.3.1 Occupancy model** This pairwise potential value becomes normalized dispersal potential in spatially-explicit metapopulation dynamics model (**ovask2003?**). Done using `MetacommunityDynamics.jl`.

There might be attempts at analytic stuff but maybe not here?

**Figure 2: MTE versus epoch fig:** shows the chains move toward higher \*extinction times over time, i.e. it works.

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## Simulation of data for testing the algorithm

In this section we describe the generation of simulated land cover and sites for testing the algorithm.

### 3.1. Simulation of landscapes

**3.1.1 Generation of landcover maps** DiamondSquare with high autocorrelatoin (0.7). Binned into  $N_{cat}$  land cover categories.

**3.1.2 Generation of points** Random Poisson process rounded to be integer coordinates. Padding around the edges because real data doesn't have points on the edges, 10% on each side.

**3.1.3 Resistance values assigned to each land cover type** Each simulated land cover has  $N_{cat}$  categories. The values of resistance

*Some type of performance fig vs. raster size and budget figure*

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**Actual data St. Lawrence lowlands**

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**Discussion**