

MST / (Minimum Spanning Tree) : $G(V, E)$ & for each edge $(u, v) \in E$, we have a weight $w(u, v)$ specifying the cost to connect u and v . We want to find an acyclic subset $T \subseteq E$ that connects all the vertices & whose total weight $w(T) = \sum_{(u, v) \in T} w(u, v)$ is minimized.

MST-KRUSKAL(G, w)

- $O(1) \rightarrow$ 1 $A = \emptyset \rightarrow$ is a forest whose vertices are all those of given graph.
 2 for each vertex $v \in G.V$ $A: \rightarrow \textcircled{a} \quad \textcircled{b} \quad \textcircled{c} \quad \dots \quad \textcircled{h} \quad \textcircled{g} \quad \textcircled{i}$
 3 MAKE-SET(v)
 $O(E \log E) \rightarrow$ 4 sort the edges of $G.E$ into nondecreasing order by weight
 5 for each edge $(u, v) \in G.E$, taken in nondecreasing order by weight
 6 if FIND-SET(u) \neq FIND-SET(v)
 7 $A = A \cup \{(u, v)\}$
 8 UNION(u, v)
 9 return A

$O(E) \rightarrow$ FIND-SET & UNION
 \downarrow no. of MAKE-SET
 $O((V+E) \log(V)) \rightarrow$ very slow growing fun \rightarrow $\{h, g\}$
 $O((V+E) \log E) = O(E \log E) \leftarrow \log E. \{i, c\}$
 $= O(E \log V)$
 $\left\{ \begin{array}{l} G \text{ is connected} \\ |E| \geq |V| - 1, \\ |E| \leq |V|^2 \\ \log E = 2 \log V \\ = O(\log V) \end{array} \right.$



