## Why BST ?

-> Search -> Delete. -> Closest.

-> Insert -> Traverse -> Rank
-> Sortes.

	Search	Insert	Debeles	closest vale.	Sorted Traversal
Arroy (unsorted)	0(4)	000	٥(٢)	٥٢٤)	0(n 10gh)
(Sonces)	0(1092)	0(n)	0(n)	0(logh)	0(n)
Linked List	0(n)	0(5)	0(1)/ 0(4)	0(u)	o (hlogh)
Mash Table	6(1)	٥(١)	6(1)	0(h)	o(n logh)
BST	O(cogn)	o Clogn)	o(1094)	O (logh)	OCN)

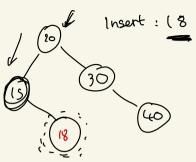
reys on the left subtree 2 root rightmost note of the left.

When in an interiew/ Lozing sets, maps, ordered sets unu Itimaps. How are these implemented? Balanced BSTs 3 -> AVL.

Search: Iteratively - O(h) a stack space.

(next: 20,15,30,40 bod search (Note\* root, int se) {

Insert: (8 Note\* cur = root)



Note\* cur = root;

while (cur = NOLL) {

if (cur = NoLL) {

if (cur = NoLL) {

refurn true;

if (cur > Note = x)

cur = cur > left;

else

cur = cur > right;

return false;

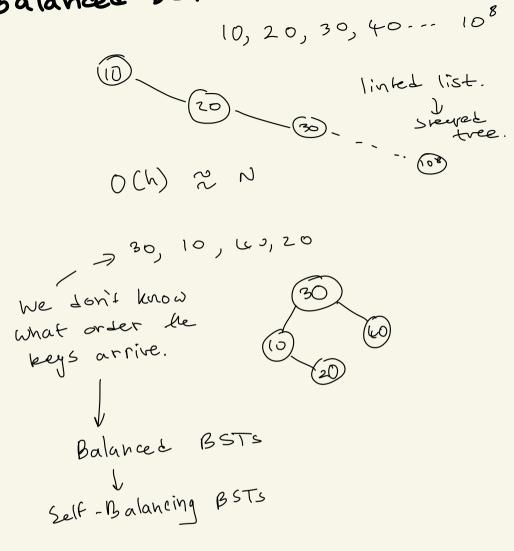
HW: Change for insert Deletion: No children 70 80 70 βÒ JUST MAKE IT NULL NULL Ore child. to pavent 60 Two children E:Ker choose 80 rightmost in left or leftmost in right. Inorder predecessor/successor.

HW: Implement these iteratively

Q: Given x output the greatest evenent in the tree that is  $\leq \infty$ . Note\* floor ( Note\* root, int x) { Note\* ans = NULL; Note\* cur = root; while (cur != NULL) & if ( cur -> key == 2 ) { ans = curj if (cur > kmy > 2) & cur = cur > left;

ans = cur; cur = right; & élse &

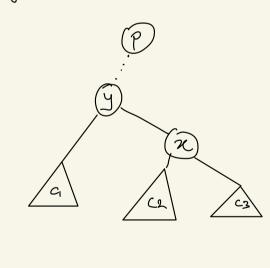
HW: Implement ceil (n). Balanced BST

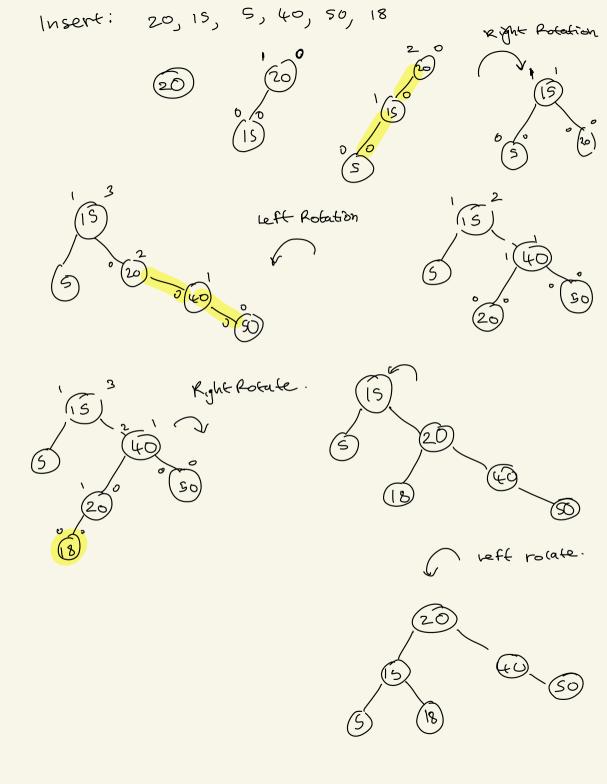


For each hode. 1h = height of left subtree n right rh =

AUL Property -| lh - rh | = 1.

Rotations 3/5-3/=2 Right Rotation. Valid BST y L x L all c3 all c, Lre all c2 L x //





LL -> Right Rotate

RR -> Left Rotate. 2 Single Rotation ? Double Rotation LR -> L, R RL -> R,L Coding -> HW Insert }

## Further Reading

Problems on BSTs > 4-5 Assignment.,

> Try solving without IDE.

-> policy based ds.

Extra ordered - set.

> cosing rounds. ( know how to use, syntax).

optional Rel - Black Trees.