Intermediate Code Generation

Sudakshina Dutta

a = 5;

a = 5; L :

P.code

Non-term	inal	Semantic rules
P -> S		S.next = new label() P.code = S.code S.next

S.code

S.next:

Types of statement

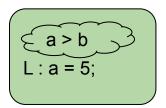
- Assignment
- If-statement
- If-else statement
- While statement
- Sequence of statements

a = 5;

a = 5;

S.code

Non-terminal	Semantic rules
S -> assign	S.code = assign.code



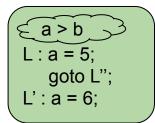
L':
$$b = 6$$
;

Non-terminal	Semantic rules
S -> if (B) S ₁	B.true = new label(); B.false = S ₁ .next = S.next S.code = B.code label(B.true) S ₁ .code

S.code

B.false:

B.code	
B.true :	
S ₁ .code	



L":

Non-terminal	Semantic rules
S -> if (B) S ₁ else S ₂	B.true = newlabel() B.false = newlabel() S_1 .next = S_2 .next = S.next S.code = B.code label(B.true) S_1 .code gen('goto' S.next) label(B.false) S_2 .code

S.code:

B.code:		
B.true : S ₁ .code		
	goto S.next:	
B.false:		
S ₂ .code		

S.next:

while (a < b) a = a + 5; begin : a < b

L: a = a + 5 goto begin begin: B.code

S.code:

B.true:

S₁.code

goto begin

L':

Non-terminal	Semantic rules
S -> while (B) S ₁	begin = new label() B.true = new label() B.false = S.next S ₁ .next = begin S.code = label(begin) B.code label(B.true) S ₁ .code gen('goto' begin)

S.next: (B.false)

while (a < b) a = a + 5; a = a + 7 begin : (a < b)

L: a = a + 5 goto begin

L': a = a + 7L': S.code

S₁.code

S₁.next: S₂.code

S₂.next (S.next)

Non-terminal	Semantic rules
S -> S ₁ S ₂	S_1 .next = new label() S_2 .next = S.next S.code = S_1 .code label(S_1 .next) S_2 .code

Non-terminal	Semantic rules
B -> B ₁ B ₂	B_1 .true = B.true B_1 .false = newlabel() B_2 .true = B.true B_2 .false = B.false $B.code = B_1.code $ B_2 .code

B ₁ .code	
B ₁ .false : B ₂ .code	

B.code

Non-terminal	Semantic rules
B -> B ₁ && B ₂	B_1 .true = newlabel() B_1 .false = B.false B_2 .true = B.true B_2 .false = B.false $B.code = B_1.code $ $abel(B_1.true) B_2.code$

B ₁ .code	
B ₁ .true: B ₂ .code	

B.code

Non-terminal	Semantic rules
B -> ! B ₁	B_1 .true = B.false B_1 .false = B.true B.code = B_1 .code

B.code

B ₁ .code	
B ₁ .true (B.false): B ₁ .false (B.true):	

Non-terminal	Semantic rules
B -> E1 rel E2	B.code = E1.code E2.code gen('if' E1.addr rel.op E2.addr 'goto' B.true) gen('goto' B.false)

b.code		
E ₁ .code		
E ₂ .code		
	ʻif' E1.addr rel.op E2.addr ʻgoto' B.true	
	'goto' B.false	

B.code

Non-terminal	Semantic rules
B -> true	B.code = gen('goto' B.true)

'goto' B.true

B.code

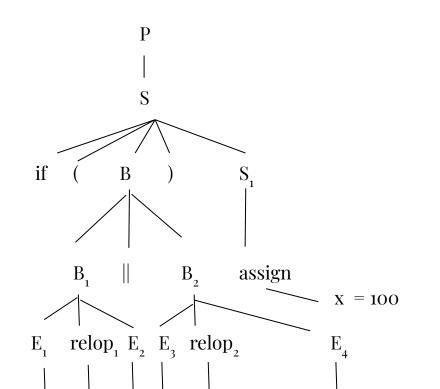
'goto' B.false

Non-terminal	Semantic rules
B -> false	B.code = gen('goto' B.false)

Example

Consider the following string

• if (x < 100 || x > 200) x = 100;

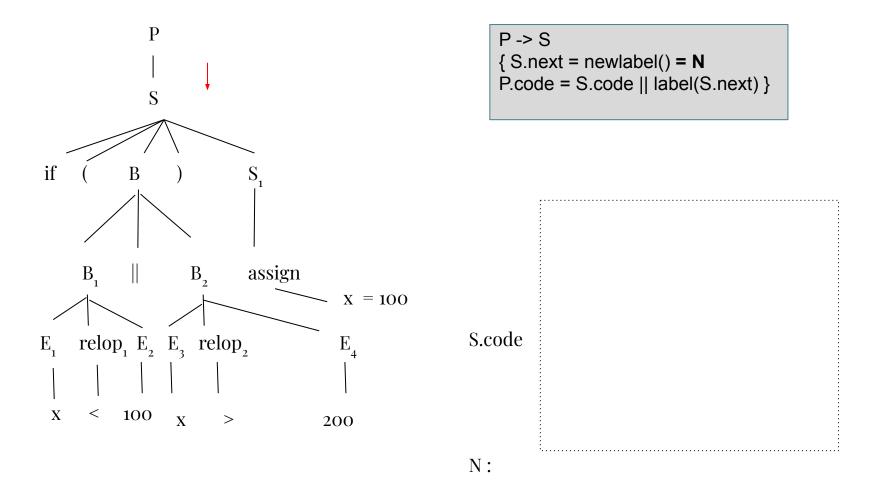


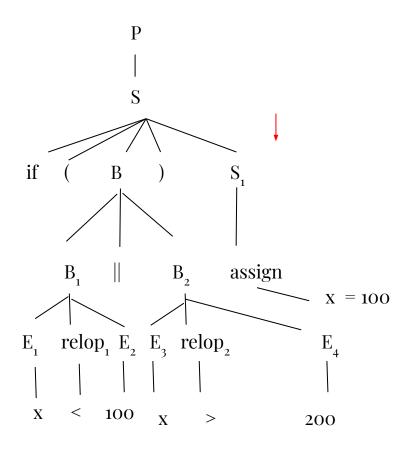
200

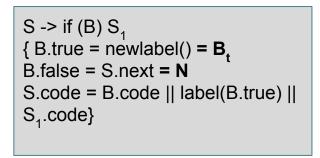
< 100

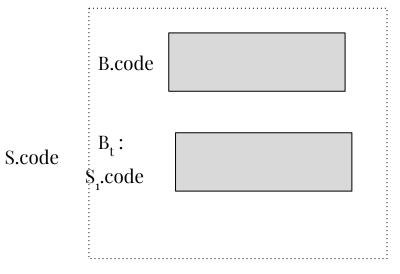
X

X

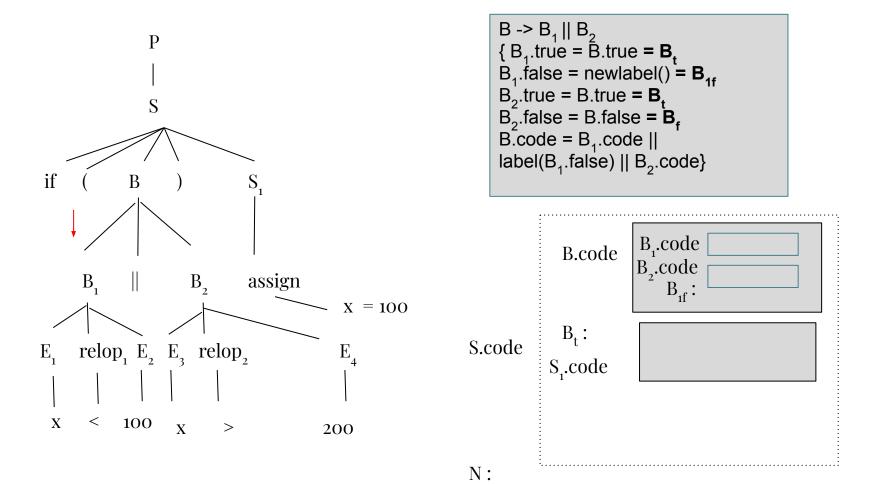


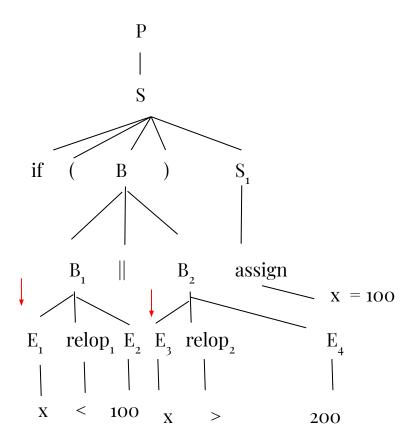






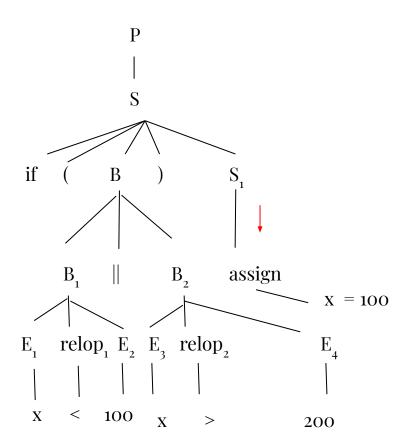
N:





 $B \rightarrow E_1 \text{ relop } E_2$ { B.code = E_1 .code || E_2 .code || gen('if' E_1 .addr relop E_2 .addr 'goto' B.true)} || gen('goto' B.false)}

	B.code	B ₁ .code	if x < 100 goto B _t goto B _{1f}	
S.code		B_2 .code B_{if} :	if x > 200 goto B _t goto N	
	B _t : S ₁ .code			
	S ₁ .code			



S -> assign {S.code = assign.code}

B.code B_1 .code if x < 100 goto B_t goto B_{1f} B₂.code S.code B_{1f} : if x > 200 goto B_t goto Nx = 0

N:

```
if x < 100 goto B<sub>t</sub>
    goto B<sub>if</sub>

B<sub>if</sub>: if x > 200 goto B<sub>t</sub>
    goto N

B<sub>t</sub>: x = 0
N:
```

Non-terminal	Semantic rules
B -> B ₁ M B ₂	{backpatch(B ₁ .falselist, M.instr); B.truelist = merge(B ₁ .truelist, B ₂ .truelist); B.falselist = B ₂ .falselist;}

Non-terminal	Semantic rules
B -> B ₁ && M B ₂	{backpatch(B ₁ .truelist, M.instr); B.truelist = B ₂ .truelist; B.falselist = merge(B ₁ .falselist, B ₂ .falselist);}

Non-terminal	Semantic rules
B -> ! B ₁	{B.truelist = B ₁ .falselist; B.falselist = B ₁ .truelist;}

Non-terminal	Semantic rules
B→(B ₁)	{B.truelist=B ₁ .truelist; B.falselist=B ₁ .falselist;}

Non-terminal	Semantic rules
B→E ₁ rel E ₂	{B.truelist = makelist(nextinstr); B.falselist = makelist(nextinstr+ 1); gen(if E ₁ .addr rel.op E ₂ .addr goto _); gen(goto _);}

Non-terminal	Semantic rules
B→true	{B.truelist = makelist(nextinstr); gen(goto _);}

Non-terminal	Semantic rules
B→false	{B.falselist = makelist(nextinstr); gen(goto _);}

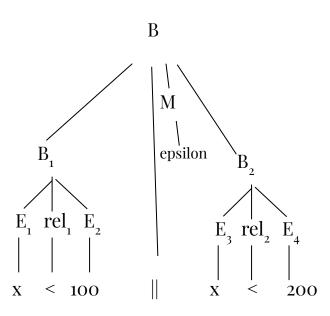
Non-terminal	Semantic rules
M -> epsilon	{M.instr = nextstr;}

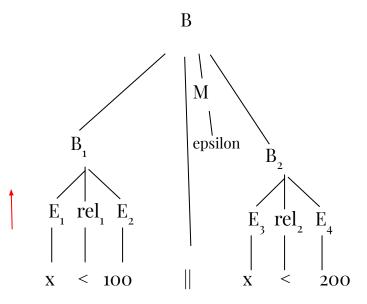
Example

Consider the following boolean expression

• x < 100 || x < 200

Parse Tree





```
B_1 \rightarrow E_1 \text{ rel}_1 E_2

\{B_1.\text{truelist} = \text{makelist}(\text{nextinstr}) = \{100\};

B_1.\text{falselist} = \text{makelist}(\text{nextinstr} + 1) = \{101\};

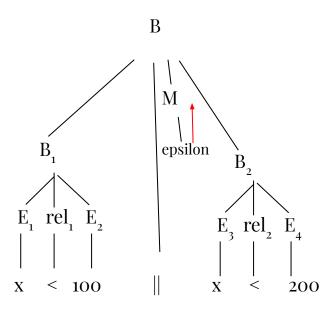
\text{gen}(\text{if } E_1.\text{addr rel.op } E_2.\text{addr goto } \_);

\text{gen}(\text{goto } \_);
```

100: if x < 100 goto _

101 : goto _

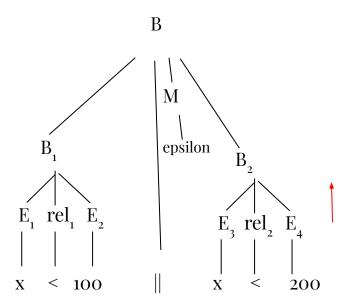
M -> epsilon {M.instr = nextinstr = {102}; }



100: if x < 100 goto _

101: goto _

IR generation



```
B_2 \rightarrow E_3 \text{ rel}_2 E_4

\{B_2.\text{truelist} = \text{makelist}(\text{nextinstr}) = \{102\};

B_2.\text{falselist} = \text{makelist}(\text{nextinstr} + 1) = \{103\};

\text{gen}(\text{if } E_3.\text{addr rel.op } E_4.\text{addr goto } \_);

\text{gen}(\text{goto } \_);
```

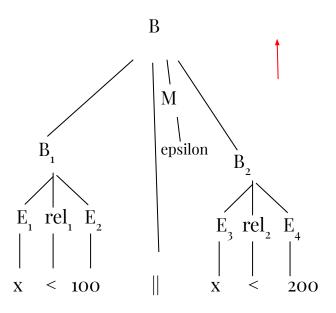
100: if x < 100 goto _

101 : goto _

102: if x < 200 goto _

103 : goto

IR generation



```
B \rightarrow B_1 \parallel M B_2
{backpatch(B_1.falselist, M.instr);
B.truelist = merge(B_1.truelist, B_2.truelist) = {100, 102};
<math>B.falselist = B_2.falselist = {103};
```

100: if x < 100 goto _

101 : goto <u>102</u>

102: if x < 100 goto _

103: goto _

Non-terminal	Semantic rules
S→if(B)M S ₁	{backpatch(B.truelist, M.instr); S.nextlist = merge(B.falselist, S ₁ .nextlist);}

Non-terminal	Semantic rules
$S\rightarrow if(B) M_1 S_1 N else M_2 S_2$	{backpatch(B.truelist, M ₁ .instr); backpatch(B.falselist, M ₂ .instr); temp = merge(S ₁ .nextlist, N.nextlist); S.nextlist = merge(temp, S ₂ .nextlist);}

Non-terminal	Semantic rules
S→while M ₁ (B)M ₂ S ₁	{backpatch(S ₁ .nextlist, M ₁ .instr); backpatch(B.truelist, M ₂ .instr); S.nextlist = B.falselist; gen(goto M ₁ .instr)}

Non-terminal	Semantic rules
S→{L}	{S.nextlist = L.nextlist;}

Non-terminal	Semantic rules
S→A	{S.nextlist = null;}

Non-terminal	Semantic rules
M -> epsilon	{M.instr = nextinstr;}

Non-terminal	Semantic rules
N -> epsilon	{N.nextlist=makelist(nextinstr); gen(goto _);}

Non-terminal	Semantic rules
L→L ₁ M S	{backpatch(L ₁ .nextlist,M.instr); L.nextlist = S.nextlist;}

Non-terminal	Semantic rules
L→S	{L.nextlist = S.nextlist;}

Input string

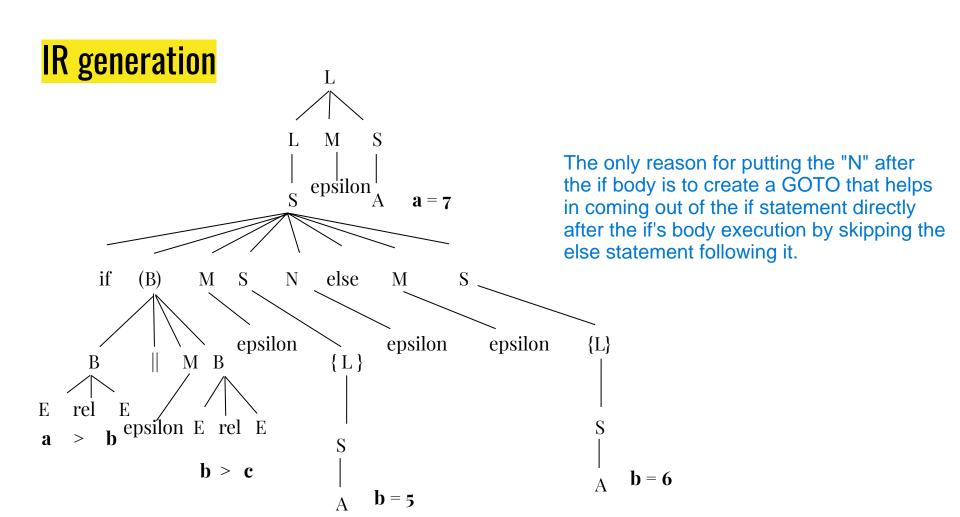
```
if(a > b \parallel b > c)
  b = 5;
else
  b = 6:
a = 7;
```

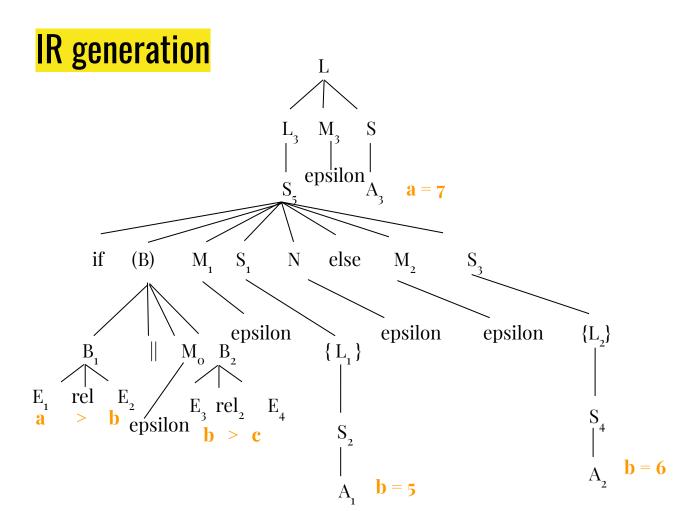
Boolean expressions have TRUE & FALSE lists. Statements have NEXT lists.

We are using backpatching operations to fill the unknown jump targets using known information.

We do bottom up parsing.

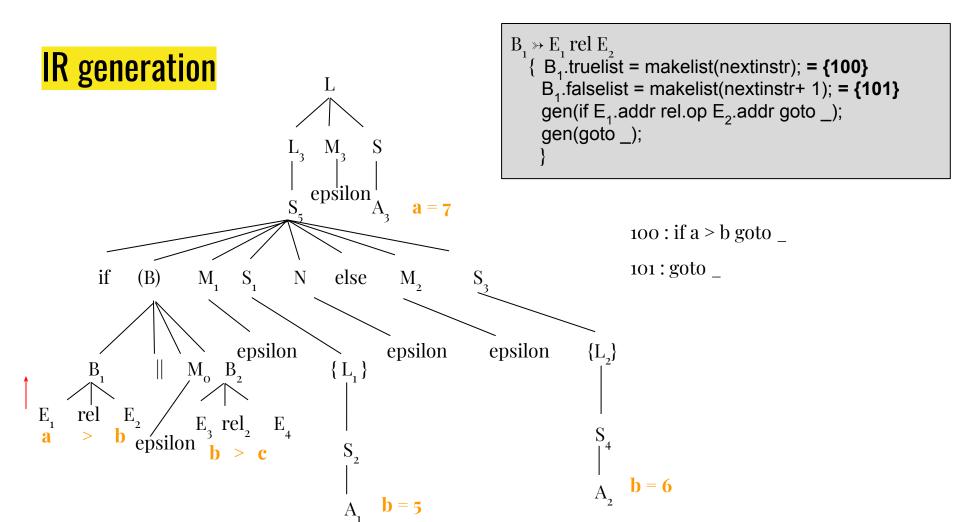
We start with left most & bottom & parse the things & group up &move a bit up if possible, etc..

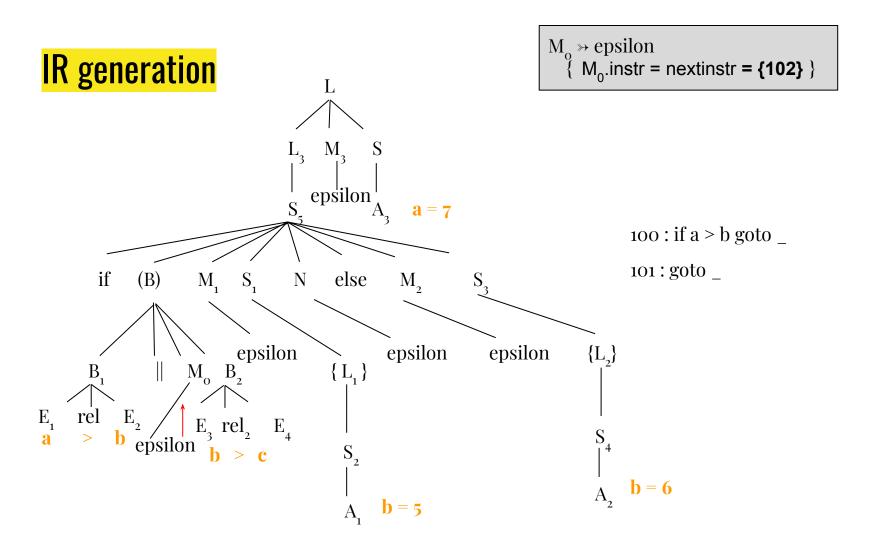


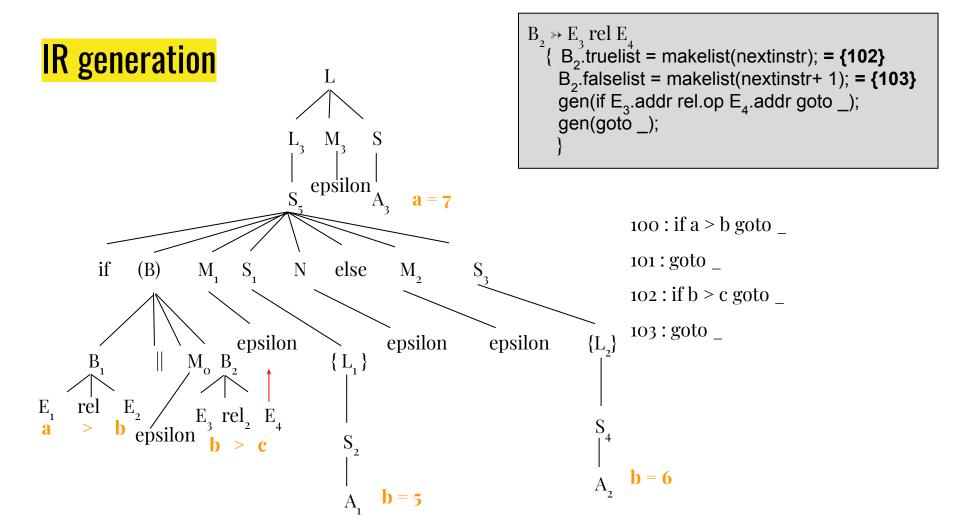


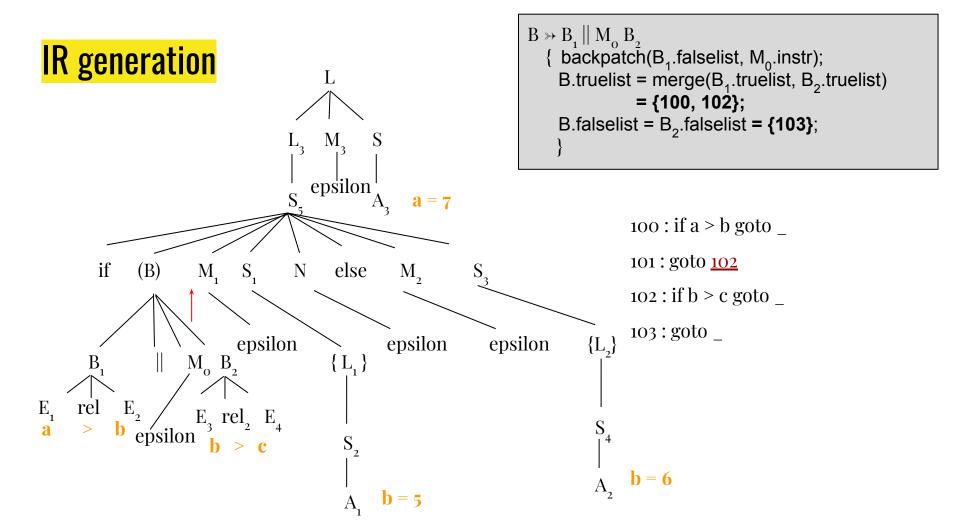
Note that all Ls, all Ms, all Ss are individually same.

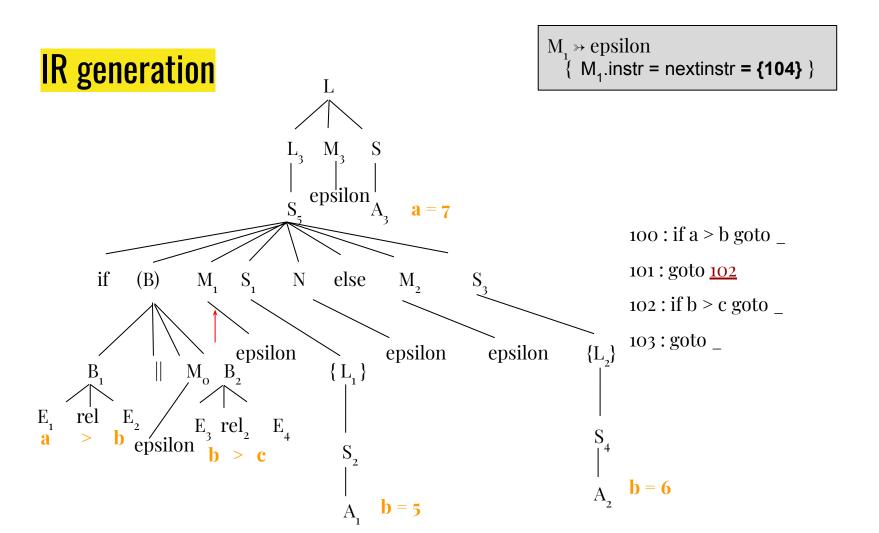
But to make it precise we are indexing them like: S1, S2, S3,...
M1, M2, M3...

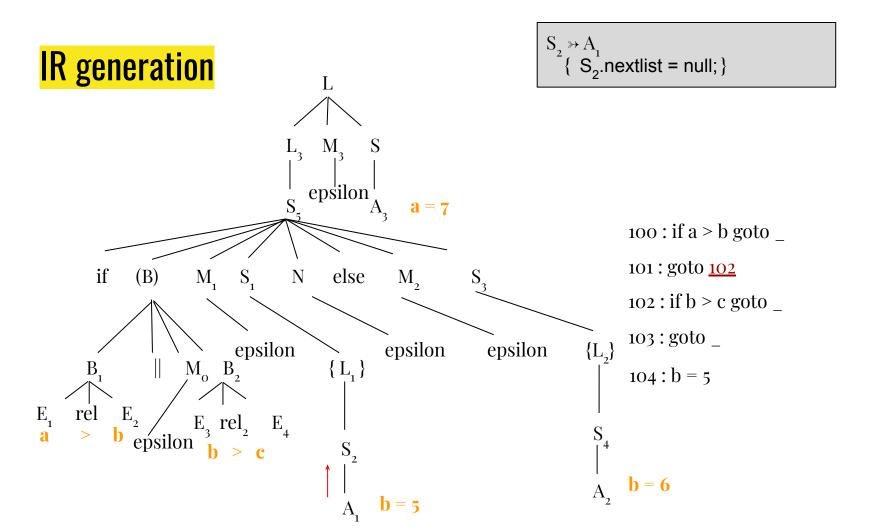


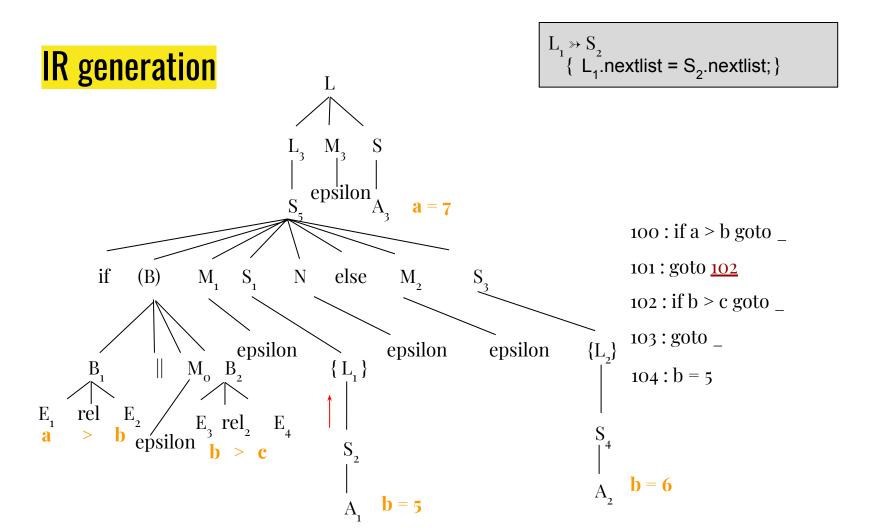


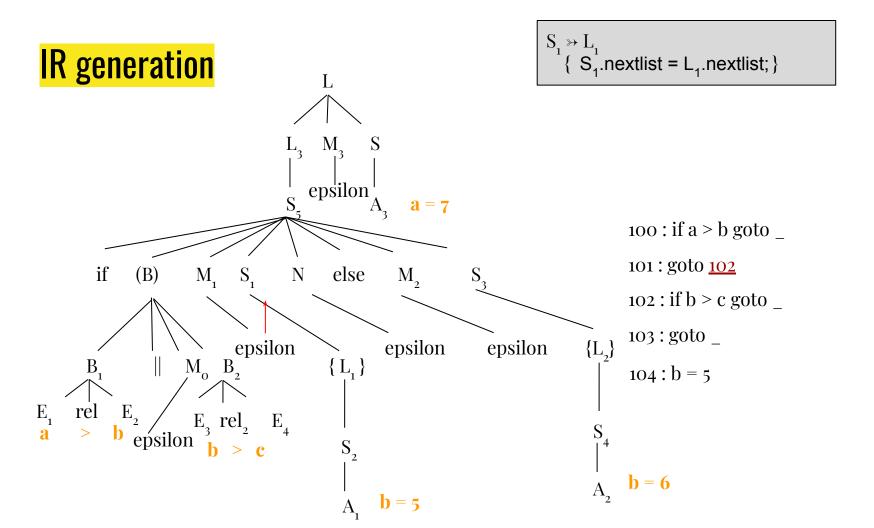


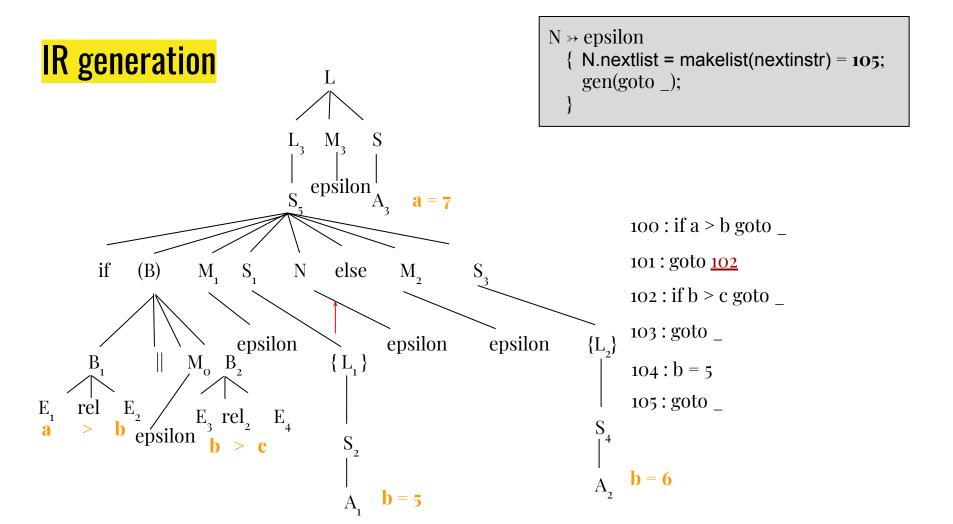


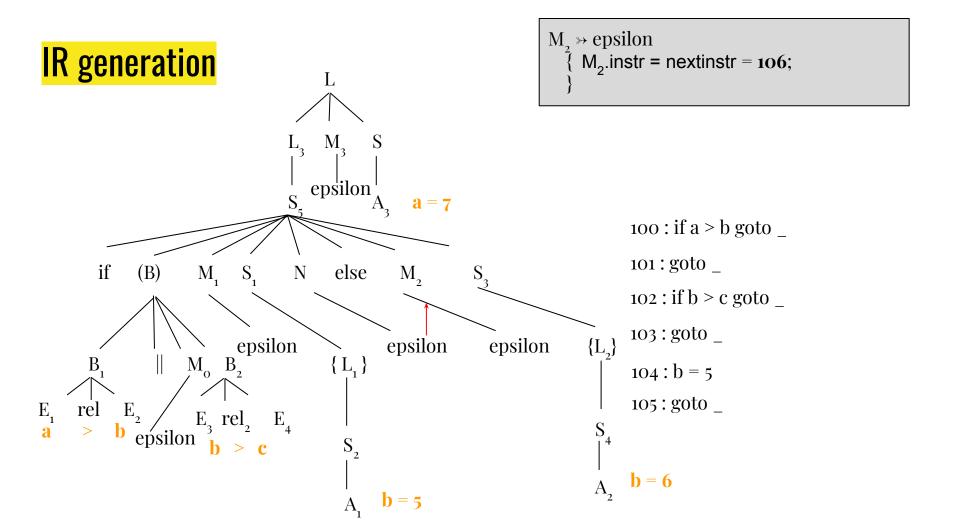


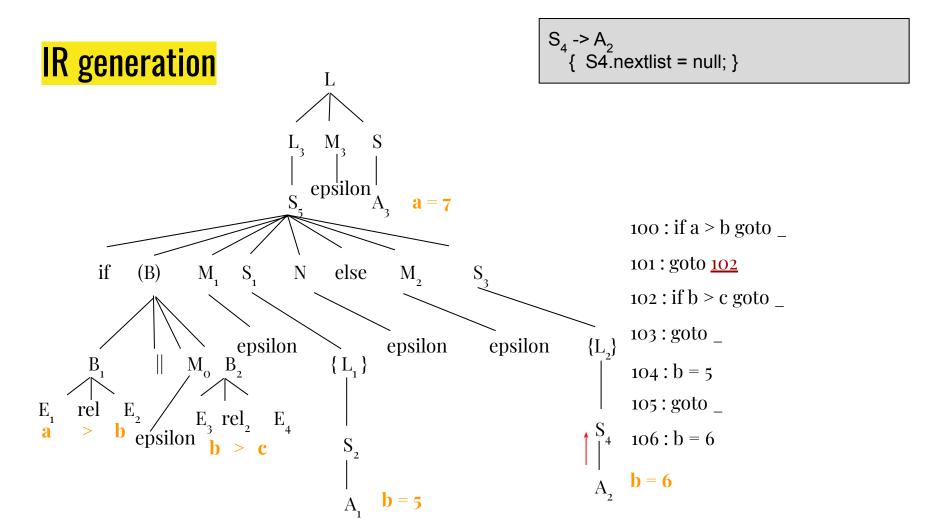


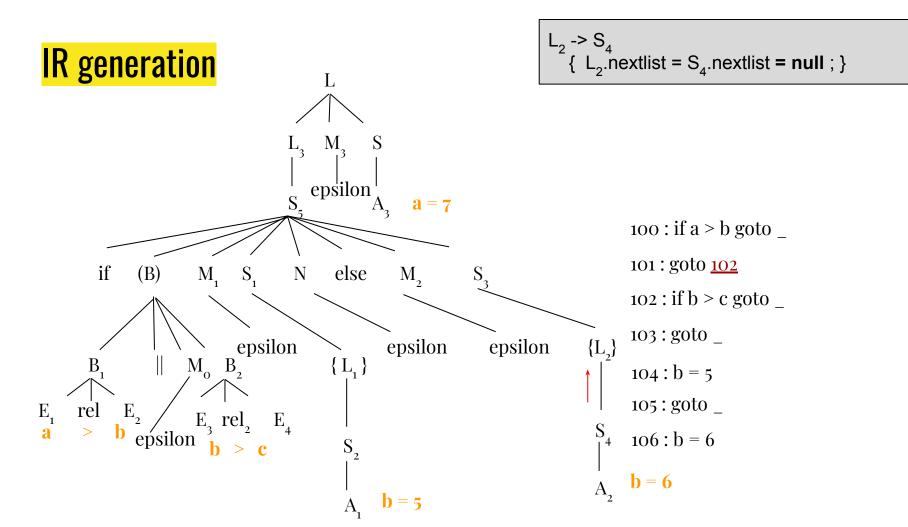


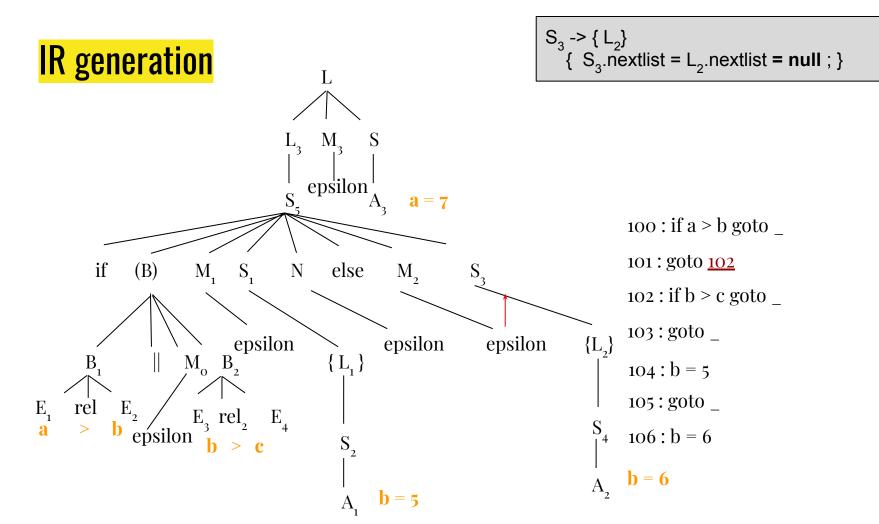


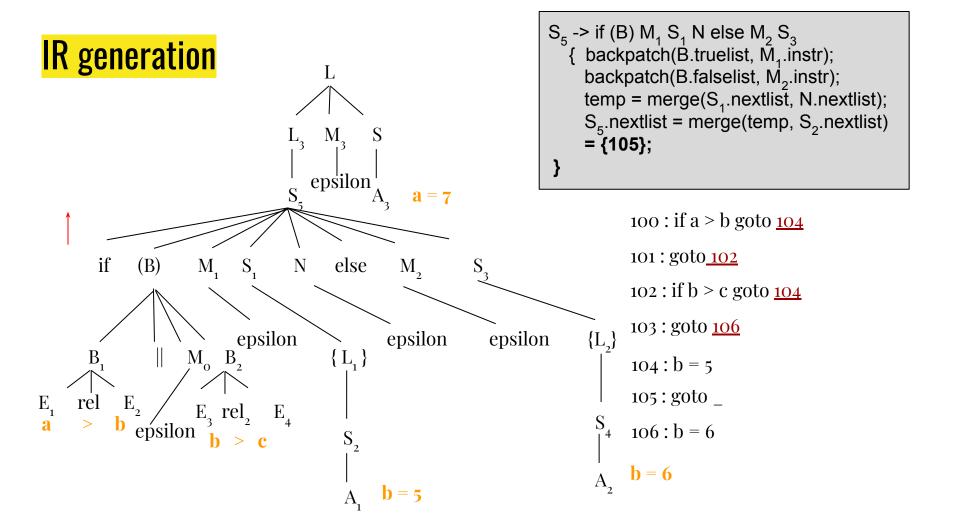


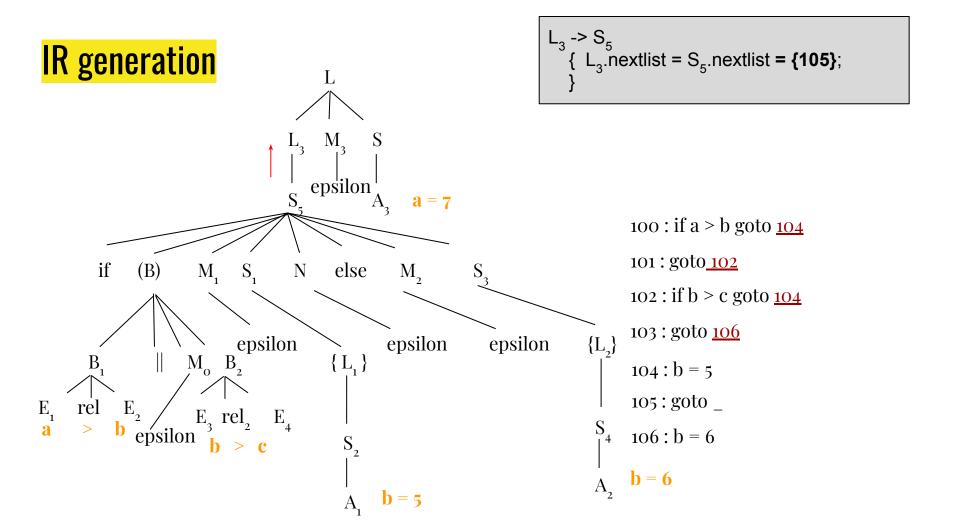


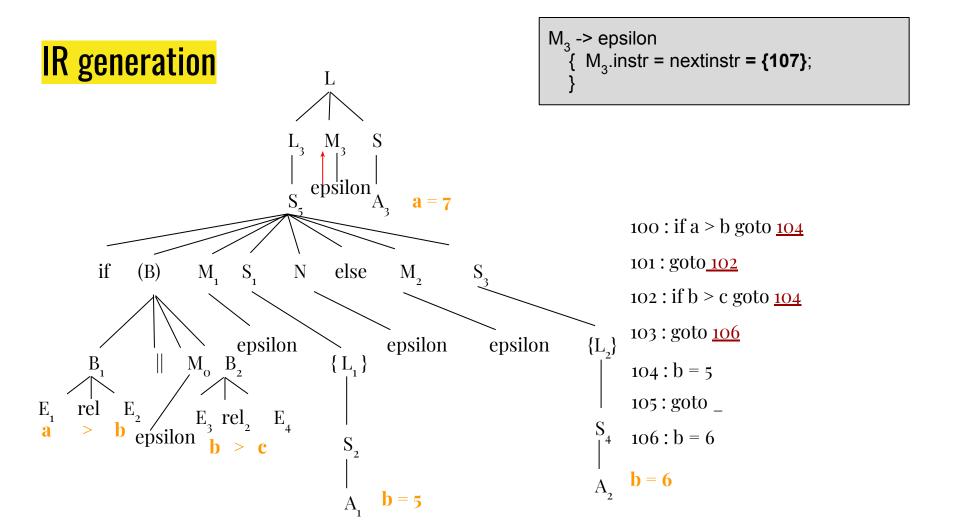


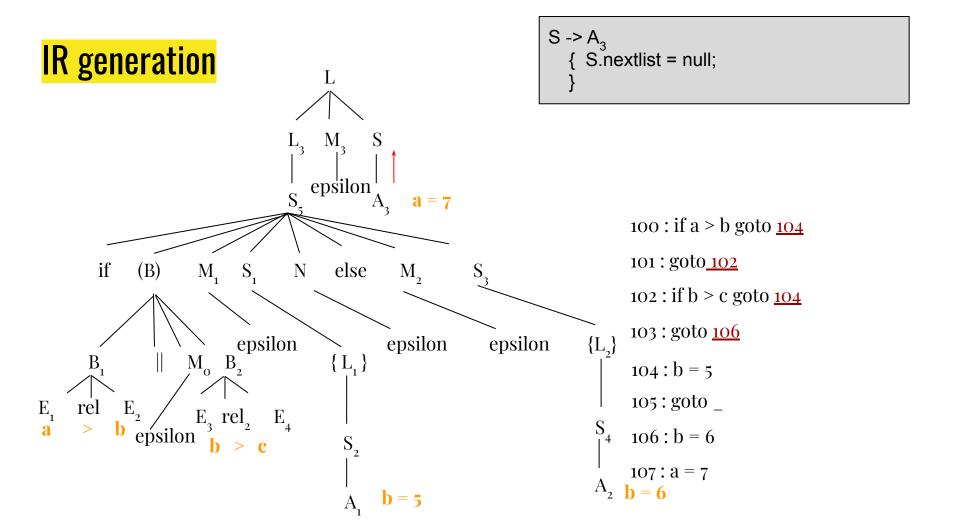


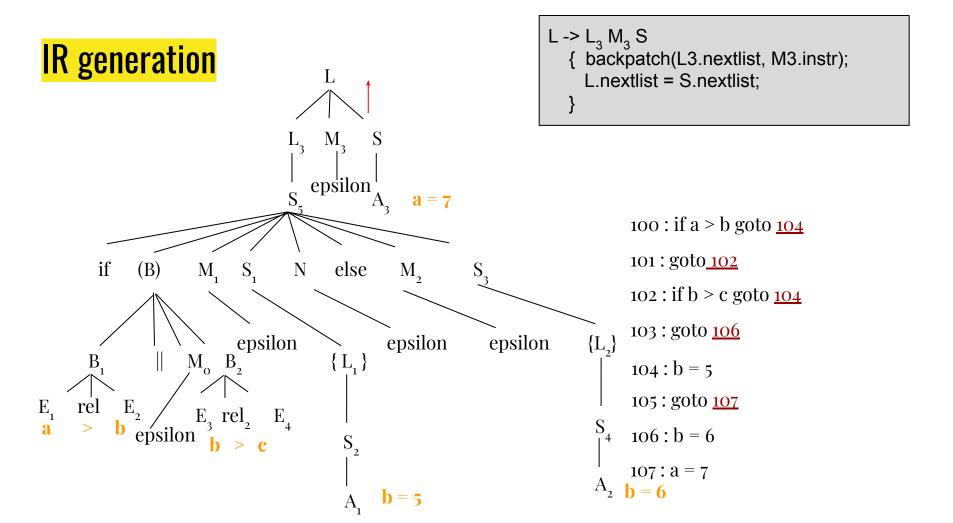










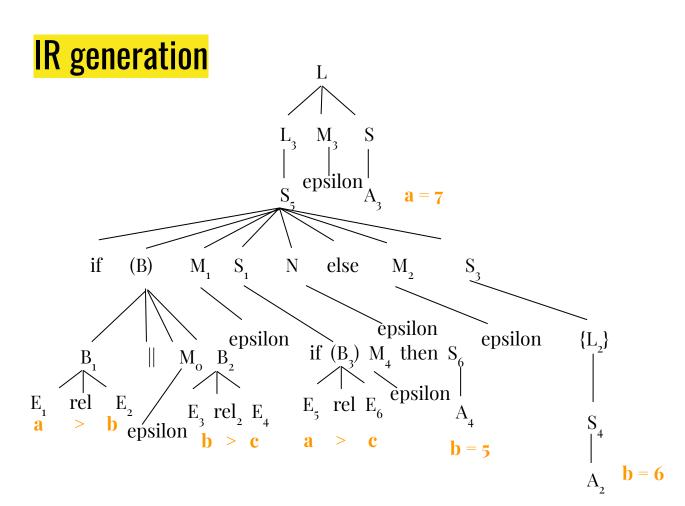


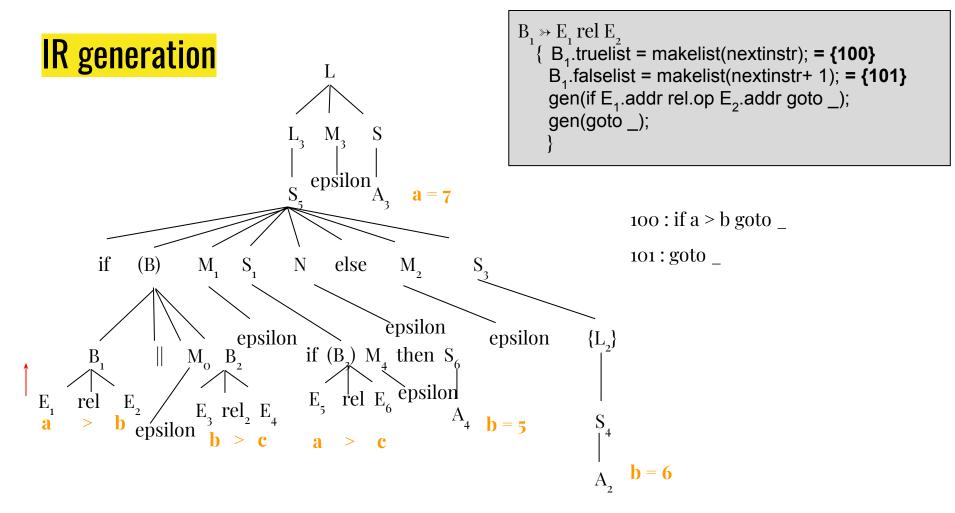
Tutorial - 6

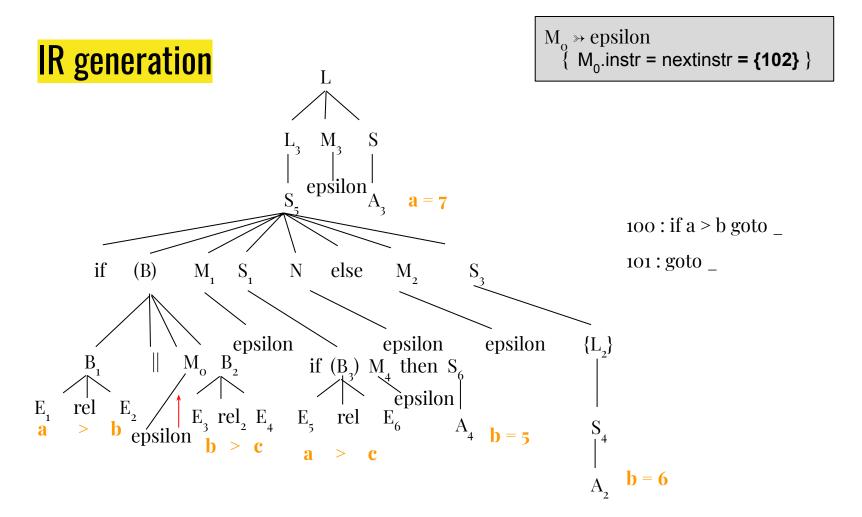
a = 7;

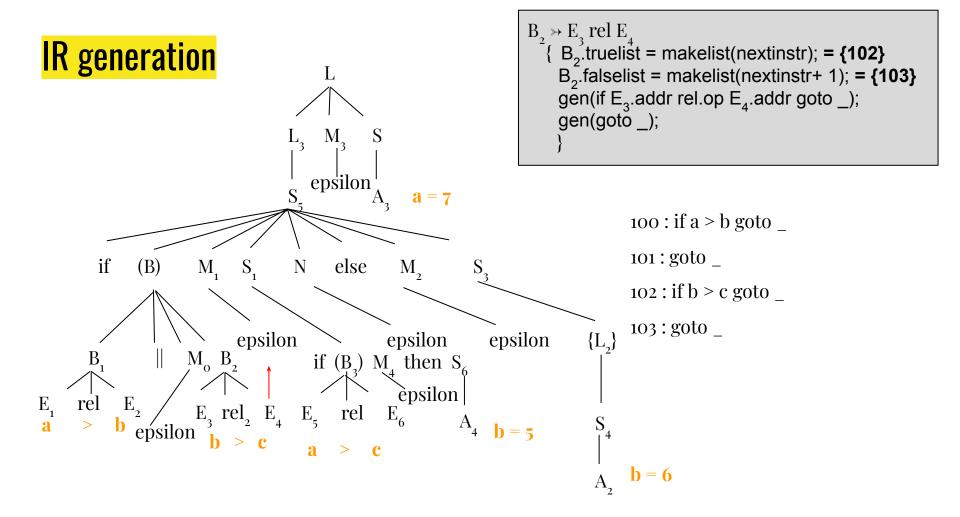
```
if(a > b \parallel b > c)
                                                             100 : if a > b goto <u>104</u>
                                                             101 : goto <u>102</u>
  if(a > c)
                                                             102 : if b > c goto <u>104</u>
                                                             103 : goto <u>108</u>
                                                             104 : if a > c goto <u>106</u>
  b = 5;
                                                             105 : goto <u>109</u>
                                                             106 : b = 5
                                                             107 : goto <u>109</u>
else
                                                             108 : b = 6
                                                             109: a = 7
   b = 6;
```

IR generation M epsilon $\mathbf{a} = \mathbf{7}$ if else (B) N M M epsilon epsilon epsilon {L} if (B) M then S M B rel E èpsilon E rel A epsilon E rel E a > \mathbf{c} a $\mathbf{b} = \mathbf{5}$ $\mathbf{b} > \mathbf{c}$ $\mathbf{b} = \mathbf{6}$

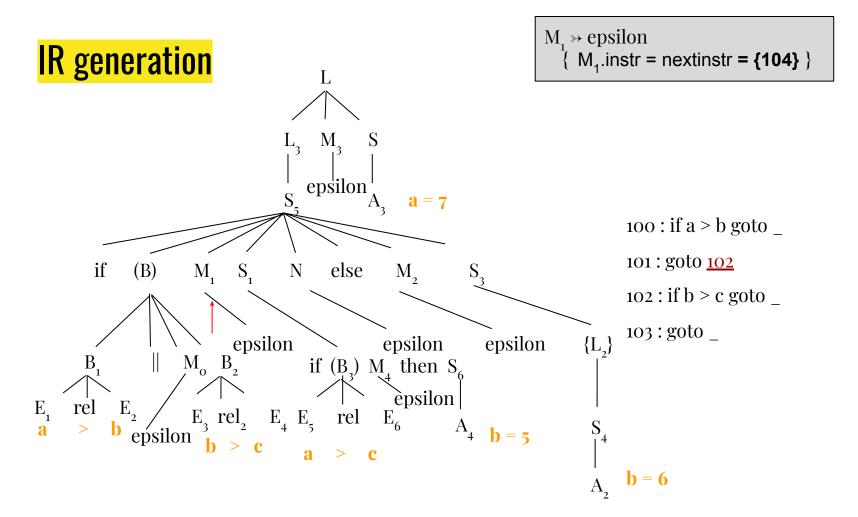


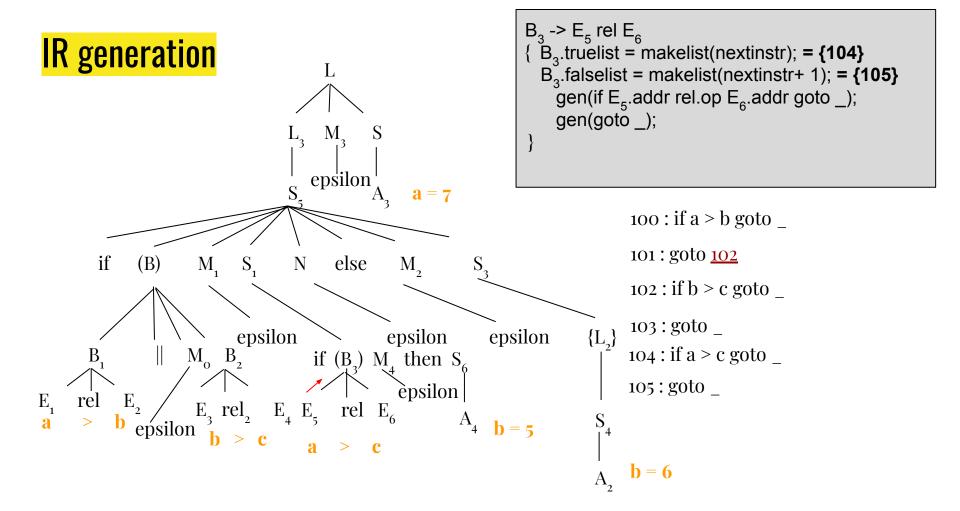




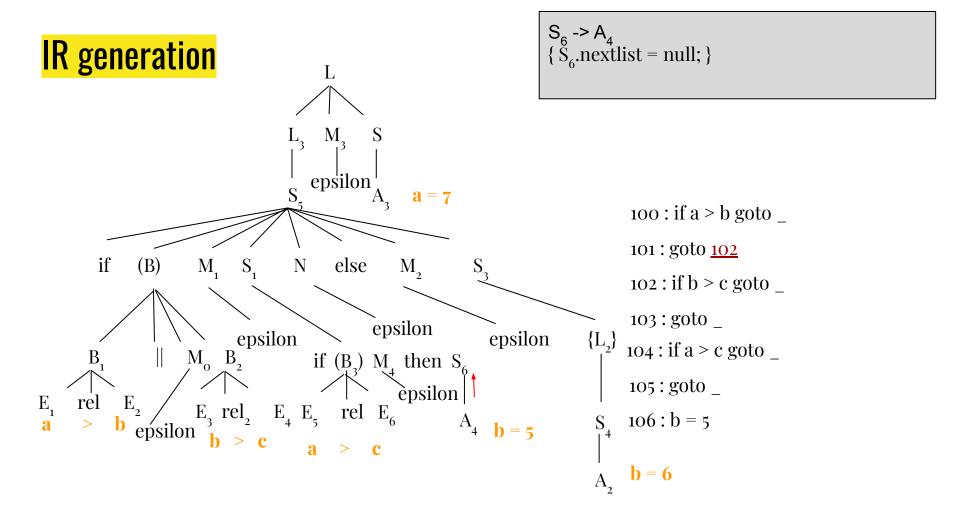


$B \gg B_1 \parallel M_0 B_2$ IR generation { backpatch(B₁.falselist, M₀.instr); B.truelist = merge(B₁.truelist, B₂.truelist) **=** {100, 102}; B.falselist = B_2 .falselist = {103}; L_{3} M_{2} epsilon $\mathbf{a} = \mathbf{7}$ 100: if a > b goto 101: goto <u>102</u> if (B) $M_{_{1}}$ S N else M_{2} 102: if b > c goto 103: goto _ epsilon epsilon epsilon $M_0 B_2$ if (B₂) M₄ then S₆ èpsilon rel epsilon $\mathbf{b} = \mathbf{6}$





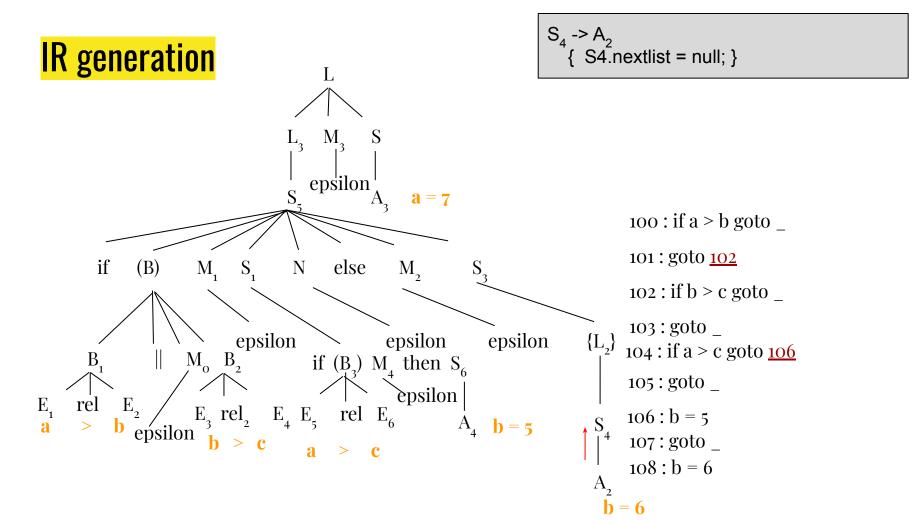
M₄ -> epsilon IR generation $\{ M_A.instr = nextinstr = \{106\}; \}$ $L_3 M_3$ epsilon 100: if a > b goto 101: goto <u>102</u> if $S_{_{1}}$ (B) $M_{_{1}}$ N else M_{2} 102: if b > c goto 103 : goto _ epsilon epsilon epsilon M_{o} 104: if a > c goto _ if (B₂) M₄ then S₆ epsilon 105 : goto _ rel epsilon $\mathbf{b} = \mathbf{6}$



$S_1 \rightarrow if (B_3) M_4 then S_6$ IR generation {backpatch(B₃.truelist, M₄.instr); S_1 .nextlist = $merge(B_3$.falselist, S_6 .nextlist) = **{105}**;} L_{3} M_{2} epsilon 100: if a > b goto 101 : goto <u>102</u> if (B) else M_{2} 102: if b > c goto_ 103 : goto _ M_o B, epsilon èpsilon epsilon $104 : \text{if } a > c \text{ goto } \underline{106}$ if (B₂) M₁ then S₂ 105 : goto _ èpsilon 106 : b = 5

N » epsilon IR generation { N.nextlist = makelist(nextinstr) = 107; gen(goto); L_3 M_3 epsilon 100: if a > b goto 101 : goto <u>102</u> if $S_{_{1}}$ (B) else M_{2} $M_{_{1}}$ 102: if b > c goto 103 : goto _ epsilon epsilon epsilon $\{L_2\}$ 104: if a > c goto 106 M_{o} if (B₂) M₄ then S₆ 105 : goto èpsilon rel 106 : b = 5epsilon 107 : goto _

M₂ » epsilon IR generation M_2 .instr = nextinstr = 108; $L_3 M_3$ epsilon 100 : if a > b goto101 : goto <u>102</u> if (B) N else M_{2} 102: if b > c goto 103 : goto _ epsilon epsilon epsilon $104 : \text{if } a > c \text{ goto } \frac{106}{}$ M_{o} if (B₂) M₄ then S₆ 105 : goto _ èpsilon rel 106 : b = 5epsilon 107: goto _

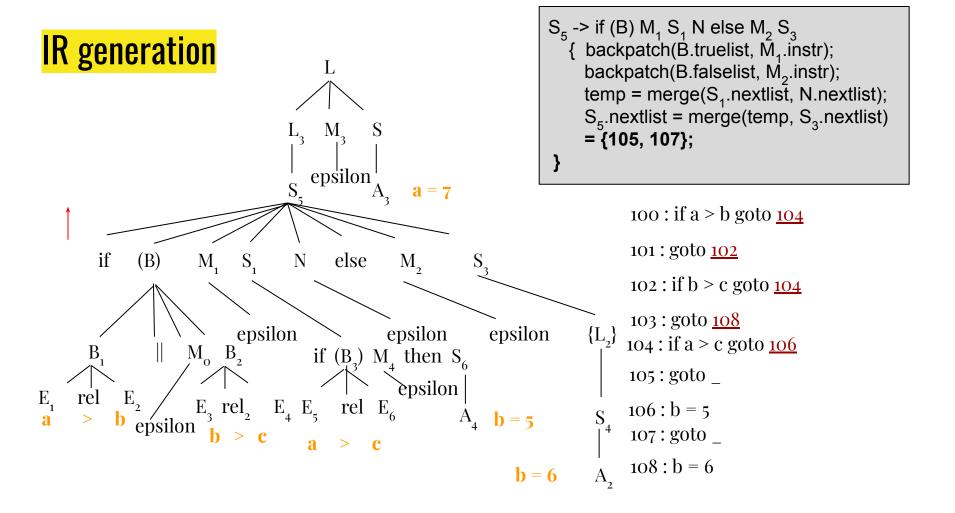


IR generation L_3 M_3 epsilon if (B) $S_{_{1}}$ N else M_{2} $M_{_{1}}$ epsilon epsilon epsilon M_{o} if (B_3) M_4 then S_6 epsilon rel

epsilon

{ L₂.nextlist = S₄.nextlist = **null**; } 100: if a > b goto 101 : goto <u>102</u> 102: if b > c goto 103 : goto _ $\{L_{2}\}$ 104: if a > c goto 106 105: goto _ 106 : b = 5107 : goto _ 108 : b = 6 $\mathbf{b} = \mathbf{6}$

$S_3 -> \{L_2\}$ IR generation { S₃.nextlist = L₂.nextlist = **null**; } L_3 M_3 epsilon 100 : if a > b goto101 : goto <u>102</u> if (B) else M_{2} 102: if b > c goto 103 : goto _ epsilon epsilon epsilon 104: if a > c goto 106 M_{o} if (B₃) M₄ then S₆ 105 : goto _ èpsilon 106 : b = 5epsilon 107 : goto _



IR generation { L_3 .nextlist = S_5 .nextlist = {105, 107}; M_{2} epsilon 100 : if a > b goto 104101 : goto <u>102</u> if (B) N else M_{2} $102 : \text{if } b > c \text{ goto } \underline{104}$ 103 : goto <u>108</u> epsilon epsilon epsilon $104 : \text{if } a > c \text{ goto } \frac{106}{}$ M_{o} if (B_3) M_4 then S_6 105 : goto _ epsilon rel 106 : b = 5 A_4 b = 5epsilon 107 : goto _ 108 : b = 6

