

Machine-independent Optimization

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Global common subexpression elimination

Suppose, we have three basic blocks

▶ B_1 :

▶ $j = j - 1$

▶ $t1 = 4 * j$

▶ $t2 = a[t1]$

▶ if $t2 > v$ goto B_1

▶ B_2 :

▶ if $i \geq j$ goto B_3

▶ B_3 :

▶ $t3 = 4 * i$

▶ $x = a[t3]$

▶ $t4 = 4 * i$

▶ $t5 = 4 * j$

▶ $t6 = a[t5]$

▶ $a[t4] = t6$

▶ $t7 = 4 * j$

▶ $a[t7] = x$

Global common subexpression elimination

After local subexpression elimination, the code becomes

- ▶ B_1 :
 - ▶ $j = j - 1$
 - ▶ $t_1 = 4 * j$
 - ▶ $t_2 = a[t_1]$
 - ▶ if $t_2 > v$ goto B_1
- ▶ B_2 :
 - ▶ if $i \geq j$ goto B_4
- ▶ B_3 :
 - ▶ $t_3 = 4 * i$
 - ▶ $x = a[t_3]$
 - ▶ $t_5 = 4 * j$
 - ▶ $t_6 = a[t_5]$
 - ▶ $a[t_3] = t_6$
 - ▶ $a[t_5] = x$

t_4 and t_7 are eliminated

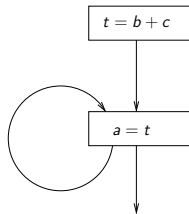
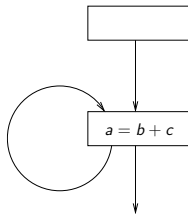
Global common subexpression elimination

After global subexpression elimination, the code becomes

- ▶ B_1 :
 - ▶ $j = j - 1$
 - ▶ $t_1 = 4 * j$
 - ▶ $t_2 = a[t_1]$
 - ▶ if $t_2 > v$ goto B_1
- ▶ B_2 :
 - ▶ if $i \geq j$ goto B_4
- ▶ B_3 :
 - ▶ $t_3 = 4 * i$
 - ▶ $x = a[t_3]$
 - ▶ $t_6 = a[t_1]$
 - ▶ $a[t_3] = t_6$
 - ▶ $a[t_1] = x$

t_5 is eliminated

Loop-invariant code motion



Reaching definition data-flow analysis

This helps for loop
invariant code motion.

- ▶ Most common and useful data-flow schemas
- ▶ It tells whether a variable x is defined when control reaches a point p
- ▶ A definition of variable x assigns or may assign values to x
- ▶ A definition is *killed* if there is any other definition of x along that path
- ▶ A definition d reaches a point p if there is a path from the point immediately following d to p , such d is not *killed* along that path
- ▶ Now debugger has the information whether x can be an undefined variable

Reaching definition

- ▶ Consider a definition $d : u = v + w$
- ▶ This statement “generates” a definition d of variable u and kills all the other definitions in the program that define variable u
- ▶ The **gen** set contains all the definitions that are visible immediately after the block It contains the last definition of each variable inside the current block.
- ▶ The gen set for the basic block
 $d_1 : a = 3$
 $d_2 : a = 4$
is d_2 and the *kill* set contains both d_1 and d_2 as they kill each other In while loop, d1 can kill d2 & d2 can kill d1 too.

Control-flow equation and transfer function

- ▶ **Control-flow equation**

$$IN[B] = \bigcup_{P \text{ a predecessor of } B} OUT[P]$$

- ▶ **Transfer function**

$$OUT[B] = gen_B \cup (IN[B] - kill_B)$$

Algorithm

Input : A flow graph for which $kill_B$ and gen_B have been computed for each block B

Output : $IN[B]$ and $OUT[B]$, the set of definitions reaching the entry and exit of each block B of the flow graph

$OUT[ENTRY] = \phi$

for(each basic block B other than $ENTRY$) $OUT[B] = \phi$

while(changes to any OUT occur)

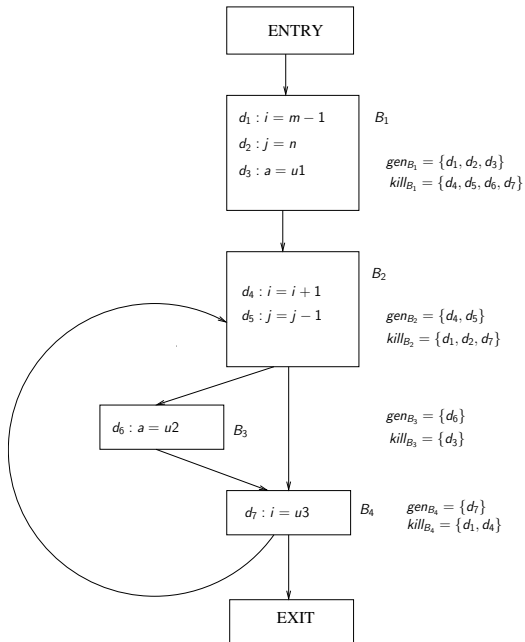
for(each basic block B other than $ENTRY$) {

$IN[B] = \bigcup_{P \text{ a predecessor of } B} OUT[P]$

$OUT[B] = gen_B \cup (IN[B] - kill_B)$

}

Example



Example

- ▶ $IN[B_2]^1 = OUT[B_1]^1 \cup OUT[B_4]^0$
 - ▶ $= gen_{B_1} \cup (IN[B_1] - kill_{B_1})$
 - ▶ $= \{d_1, d_2, d_3\} \cup \phi = \{d_1, d_2, d_3\}$
- ▶ $OUT[B_2]^1 = gen_{B_2} \cup (IN[B_2]^1 - kill_{B_2})$
 - ▶ $= \{d_4, d_5\} \cup (\{d_1, d_2, d_3\} - \{d_1, d_2, d_7\})$
 - ▶ $= \{d_3, d_4, d_5\}$
- ▶ $IN[B_3]^1 = OUT[B_2]^1$
- ▶ $OUT[B_3]^1 = gen_{B_3} \cup (IN[B_3]^1 - kill_{B_3})$
 - ▶ $\{d_6\} \cup (\{d_3, d_4, d_5\} - \{d_3\})$
 - ▶ $\{d_4, d_5, d_6\}$
- ▶ $IN[B_4]^1 = OUT[B_3]^1 \cup OUT[B_2]^1$
 - ▶ $= \{d_4, d_5, d_6\} \cup \{d_3, d_4, d_5\}$
 - ▶ $= \{d_3, d_4, d_5, d_6\}$

Example

- ▶ $OUT[B_4]^1 = gen_{B_4} \cup (IN[B_4]^1 - kill_{B_4})$
 - ▶ $= \{d_7\} \cup (\{d_3, d_4, d_5, d_6\} - \{d_1, d_4\})$
 - ▶ $= \{d_7\} \cup \{d_3, d_5, d_6\}$
 - ▶ $= \{d_3, d_5, d_6, d_7\}$
- ▶ $IN[B_2]^2 = OUT[B_1]^1 \cup OUT[B_4]^1$
 - ▶ $= \{d_1, d_2, d_3\} \cup \{d_3, d_5, d_6, d_7\}$
 - ▶ $= \{d_1, d_2, d_3, d_5, d_6, d_7\}$
- ▶ $OUT[B_2]^2 = gen_{B_2} \cup (IN[B_2]^2 - kill_{B_2})$
 - ▶ $= \{d_4, d_5\} \cup (\{d_1, d_2, d_3, d_5, d_6, d_7\} - \{d_1, d_2, d_7\})$
 - ▶ $= \{d_4, d_5\} \cup \{d_3, d_5, d_6\}$
 - ▶ $= \{d_3, d_4, d_5, d_6\}$
- ▶ $IN[B_3]^2 = OUT[B_2]^2$
 - ▶ $= \{d_3, d_4, d_5, d_6\}$

Example

- ▶ $OUT[B_3]^2 = gen_{B_3} \cup (IN[B_3]^2 - kill_{B_3})$
 - ▶ $= \{d_6\} \cup (\{d_3, d_4, d_5, d_6\} - \{d_3\})$
 - ▶ $= \{d_4, d_5, d_6\}$
- ▶ $IN[B_4]^2 = OUT[B_2]^2 \cup OUT[B_3]^2$
 - ▶ $= \{d_3, d_4, d_5, d_6\} \cup \{d_4, d_5, d_6\}$
 - ▶ $= \{d_3, d_4, d_5, d_6\}$
- ▶ $OUT[B_4]^2 = gen_{B_4} \cup (IN[B_4]^2 - kill_{B_4})$
 - ▶ $= \{d_7\} \cup (\{d_3, d_4, d_5, d_6\} - \{d_1, d_4\})$
 - ▶ $= \{d_3, d_5, d_6, d_7\}$