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Section 002

Homework 1

Q1 a) Time Complexity

```
for(i = 1; i <= N; i=i+1)
    for(k = 1; k <= i; k = 2*k)
        for(t = 1; t <= N; t = 2*t)
            printf("C");
```

Let's assume that  $N=2^p$  and  $\lg N = \log_2 N$

Iteration	Values of i	Values of k	Iterations count for k	Values of t	Iteration count for t	total repetitions of k and t for one i
1	1	1	1	1, 2, 4, 8, ..., $2^p$ (or) $2^0, 2^1, 2^2, 2^3, \dots, 2^p$	$\lg N$	$\lg N$
2	2	1	2	1, 2, 4, 8, ..., $2^p$	$\lg N$	2 $\lg N$
		2		1, 2, 4, 8, ..., $2^p$	$\lg N$	
3	3	1	2	1, 2, 4, 8, ..., $2^p$	$\lg N$	2 $\lg N$
		2		1, 2, 4, 8, ..., $2^p$	$\lg N$	
4	4	1	3	1, 2, 4, 8, ..., $2^p$	$\lg N$	3 $\lg N$
		2		1, 2, 4, 8, ..., $2^p$	$\lg N$	
		4		1, 2, 4, 8, ..., $2^p$	$\lg N$	
.	.	.	.	.	.	.
.	.	.	.	.	.	.
i	i	1	i	1, 2, 4, 8, ..., $2^p$	$\lg N$	j $\lg N$
		2		1, 2, 4, 8, ..., $2^p$	$\lg N$	
		4		1, 2, 4, 8, ..., $2^p$	$\lg N$	
		.		.	.	

		$2^j$		$1, 2, 4, 8, \dots, 2^p$	$\lg N$	
.	.	.	.	.	.	.
.	.	.	.	.	.	.
.	.	.	.	.	.	.
N	N	1	N	$1, 2, 4, 8, \dots, 2^p$	$\lg N$	$p \lg N$
		2		$1, 2, 4, 8, \dots, 2^p$	$\lg N$	
		4		$1, 2, 4, 8, \dots, 2^p$	$\lg N$	
		.		.	.	
		.		.	.	
		$2^p$		$1, 2, 4, 8, \dots, 2^p$	$\lg N$	

**Summation :  $T(N) = \lg N + 2\lg N + 2\lg N + 3\lg N + \dots + j\lg N + \dots + p\lg N$**

where j and p are the number of times k loop has executed for one i.

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Q1 b) Summation, Closed form, dominant term and  $\Theta$

```
for(i = 1; i <= N; i=i+1)
    for(k = 1; k <= S; k++)
        for(t = 1; t <= i; t++)
            printf("D");
```

Iteration	Values of i	Values of k	Iterations count for k	Values of t	Iteration count for t	total repetitions due to k for k and for t loops for one i
1	1	1 2 3 . . S	S	1 1 1 . . 1	1 1 1 . . 1	S
2	2	1 2 3 . . S	S	1,2 1,2 1,2 . . 1,2	2 2 2 . . 2	2S

.	.	.	.	.	.	.
.	.	.	.	.	.	.
i	i	1 2 3 . . S	S	1, 2, 3, ..., i 1, 2, 3, ..., i 1, 2, 3, ..., i . . 1, 2, 3, ..., i	i i i . . i	iS
.	.		.	.	.	.
.	.		.	.	.	.
N	N	1 2 3 . . S	S	1, 2, 3, ..., N 1, 2, 3, ..., N 1, 2, 3, ..., N . . 1, 2, 3, ..., N	N N N . . N	NS

$$\text{Summation: } T(N, S) = S + 2S + 3S + \dots + iS + \dots + NS$$

$$= S(1 + 2 + 3 + \dots + i + \dots + N)$$

$$= S[N(N+1)/2]$$

→ Closed Form

$$= \mathbf{S[(N^2 + N)/2]}$$

$$\text{Dominant Term} : \mathbf{S \text{ and } N^2}$$

$$\text{Theta} : \mathbf{\Theta(SN^2)}$$

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Q1 c) Summation, Closed form, dominant term and  $\Theta$

```
for(i = 1; i <= N; i=i+1){
    for(k = 1; k <= N; k++)
        for(t = 1; t <= k; t++)
            printf("E");
    for(k = 1; k <= M; k++)
        for(t = 1; t <= k; t++)
            printf("F");
}
```

Summation table for first 4 lines of code (first 3 nested loops):

Assumed value of N	Iteration	Values of i	Values of k	Iterations count for k	Values of t	Iteration count for t	total repetitions due to k for k and for t loops for one i
1	1	1	1	1	1	1	1
2	1	1	1	2	1, 2	1	6
	2	2	1	2	1	1	
			2	2	1, 2	2	
3	1	1	1	3	1	1	18
			2		1, 2	2	
			3		1, 2, 3	3	
	2	2	1	3	1	1	
			2		1, 2	2	
			3		1, 2, 3	3	
	3	3	1	3	1	1	
			2		1, 2	2	
			3		1, 2, 3	3	
N	1	1	1	N	1	1	N (1+2+...+N)
			2		1, 2	2	
			.		.	.	
			.		.	.	
			N		1, 2, ..., N	N	
	2		1	N	1	1	
			2		1, 2	2	
			.		.	.	
			.		.	.	
			N		1, 2, ..., N	N	
	3		1		1	1	
			2		1, 2	2	
			.		.	.	
			.		.	.	
			.		.	.	



			3	4 5 1 2 3 4 5	5     5	1, 2, 3, 4 1, 2, 3, 4, 5 1 1, 2 1, 2, 3 1, 2, 3, 4 1, 2, 3, 4, 5	4 5 1 2 3 4 5	
N	M	1  2  3  . . N	1  2  3  . . N	1 2 . . M 1 2 . . M 1 2 . . M . . 1 2 . . M	    M    M    M . .   M	1 1, 2 . . 1, 2, ..., M 1 1, 2 . . 1, 2, ..., M 1 1, 2 . . 1, 2, ..., M . . 1 1, 2 . . 1, 2, ..., M	1 2 . . M 1 2 . . M . . 1 2 . . M	N (1+2+3 +...+M)

Summation :  $T_2(N, M) = N(1+2+3+\dots+M)$

$$= N[M(M+1)/2]$$

$$= N[(M^2+M)/2]$$

$$= (NM^2+NM)/2$$

Summation for the entire code:  $T_1(N) + T_2(N, M)$

$$= [(N^3+N^2)/2] + [(NM^2+NM)/2]$$

$$= \frac{1}{2}(N^3+N^2+NM^2+NM) \quad \rightarrow \text{Closed Form}$$

Dominant terms :  $N^3$  and  $NM^2$

Theta :  $\Theta(N^3+NM^2)$

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Q2) for(i=1; i<=N; i=i+7) ;

Let's assume number of values of i=p

Iteration	Values of i	Total no.of loop iterations
1	1	1
2	8	1
3	15	1
4	22	1
.	.	.
.	.	.
t	i	1
.	.	.
.	.	.
p	N	1

For 4<sup>th</sup> iteration, N=22, total no. of for loop iterations so far is 4 times which is  $22/7=4$ (approx.)

For 5<sup>th</sup> iteration, N=29, total no. of for loop iterations so far is 5 times which is  $29/7=5$ (approx.)

For p<sup>th</sup> iteration, N, total no. of for loop iterations so far is  $N/7$  times.

$$\begin{aligned}\text{Summation} \quad : T(N) &= 1+1+1+\dots+1 \\ &= (N/7)1 \\ &= \mathbf{N/7}\end{aligned}$$

Dominant term : **N/7**

Theta :  **$\Theta(N/7)$**

=====

Q3) Theta time complexity

```
i = 0;
while (i<=N){
    for(t=0, k=1; k<N; t=t+1, k=2*k)
        printf("G");
    i=i+t;
}
```

Assumed value of N	Iteration of i	Values of i	Values of k	Iterations count for k	Values of t inside and outside for loop	total repetitions due to k for t loop for one i
1	1	0	1	0	0	0

Let's assume the value of **N=1**.

Since **i** and **t** will always be **0**, we will run into an **infinite loop** where **G** will not be printed even once.

If  $N > 1$ , 'G' will not go into infinite loop. But we are going into an infinite loop for  $N=1$ , hence we are not discussing that.

If  $N > 1$ , then below is the summation table. Let  $\lg N = \log_2 N$

Assumed value of N	Iteration of i	Values of i	Values of k	Iterations count for k	Values of t inside for loop	Values of t outside for loop/inside while loop	total repetitions due to k for k loop and t loop for one i
2	1	0	1	1	0	1	3
	2	1	1	1	0	1	
	3	2	1	2	0	1	
3	1	0	1 2	1 2	0 1	 2	4
	2	2	1 2	1 2	0 1	 2	
$N > 1$	1	0	1	$\lg N$	0	1	N  N  . N
	2	$0+t$	1, 2, 4, ..., $N-1$	$\lg N$	0, 1, 2, ..., $N-1$	1, 2, ..., N	
	.	.	.	.	.	.	
	.	$i+t$	1, 2, 4, ..., $N-1$	$\lg N$	0, 1, 2, ..., $N-1$	1, 2, ..., N	
	.	.	.	.	.	1, 2, ..., N	
	.	.	.	$\lg N$	.	.	



	N	N	1, 2, 3, ..., N-1	. lgN	0, 1, 2, ..., N-1	1, 2, ..., N	. N
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Value of i depends on the value of t ( $i=i+t$ ) and t depends on k (t iterates until k iterates), whereas k depends on N ( $k < N$ ).

Summation:  $T(N) = N \lg N$

For the value of N, there is a lot of ambiguity in the values of i, k and t. Hence this code is highly susceptible to a lot of errors.

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**Q4)  $T(n) = 1+3^1+3^2+3^3+ \dots +3^n$**

Code for the above time complexity is

```
for(i = 1; i <= N; i=i*3)
    for (k = 1; k <= i; k++);
```

Let's assume N is a multiple of 3

For the  $m^{th}$  iteration : Value of  $i=i$

For the  $t^{th}$  iteration : Value of  $i=N$

Iteration of i	Values of i	Values of k	total repetitions due to k for k loop for one i
1	1	1	$1=3^0$
2	3	1, 2, 3	$3=3^1$
3	9	1, 2, 3, 4, 5, 6, 7, 8, 9	$9=3^2$
.	.	.	.
.	.	.	.
m	i	1, 2, 3, ..., i	$3^{m-1}$
.	.	.	.
.	.	.	.
t	N	1, 2, 3, ..., i, ..., N	$3^{t-1}=3^n$

Summation :  $T(N) = 3^0+3^1+3^2+3^3+ \dots +3^n = 1+3^1+3^2+3^3+ \dots +3^n$

Closed Form :  $(3^{n+1}-1)/2$

Dominant Term :  $3^n/2$

Theta :  $\Theta(3^n/2)$

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Q5) Dominant term and Theta

$$N^3 + 500N^2 + NM + 106$$

Dominant term :  $N^3$  and  $NM$

Theta :  $\Theta(N^3 + NM)$

$$100N^3 + 20N^2 + 15M + 5N$$

Dominant term :  $N^3$  and  $M$

Theta :  $\Theta(N^3 + M)$

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Q6) how long they take for  $N=10$ ,  $N=100$ ,  $N=300$ . (Also try  $N=1000$ )

- the runtime effect of replacing the 'res = res+1' with the 'print...' instruction for runtime\_increment and runtime\_print.

- how the runtime depends on the value of  $N$  for runtime\_increment.

- how much faster (i.e. for smaller values of  $N$ ) the runtime\_pow becomes too slow.

```
// run for N = 10, N = 100, N = 1000
void runtime_increment(int N){
int i, k, t, res = 0;
for(i = 1; i <= N; i=i+1)
    for(k = 1; k <= i; k++)
        for(t = 1; t <= N; t++)
            res = res + 1;
}
```

**For  $N = 10$**

Runtime: It's super fast. Time taken is **less than a second** (in microseconds).

**For  $N = 100$**

Runtime: It's also super fast. Time taken is **less than a second** (in microseconds).

**For  $N = 1000$**

Runtime: **Approximately 0.293 seconds.** Less than a second.

*As the value of  $N$  increases, the **runtime** of the loop also **increases** for runtime\_increment. But the rate of increase of runtime is slow when compared to increase in value of  $N$ .*

When 'res = res+1' is replaced with the 'print...', then

```
// run for N = 10, N = 100, N = 300, N = 1000
void runtime_print(int N){
int i, k, t;
for(i = 1; i <= N; i=i+1)
    for(k = 1; k <= i; k++)
        for(t = 1; t <= N; t++)
            printf("A");
}
```

**For N = 10**

Runtime: **Approximately 0.007 seconds.** Time taken is less than a second. A gets printed 550 times.

**For N = 100**

Runtime: **Approximately 2.5 to 3 seconds.** Time taken is just little more than 3 seconds. A gets printed 505000 times.

**For N = 300**

Runtime: **Approximately 83 to 85 seconds.** Time taken is little less than a minute and a half. A gets printed 13545000 times.

**For N = 1000**

Runtime: **Approximately 1 hour 35 mins** (5695.366 seconds). A gets printed 500500000 times.

It took too much time to give the result. But I was patient enough for the program to complete execution.

*By replacing 'res = res+1' with the 'print...', we infer that the time taken for 'print...' is slower or longer than the execution time of 'res = res+1'. **The runtime of 'res = res+1' is much faster than runtime of 'print...'.***

```
// run for N = 10, N = 15, N = 20, N = 25, N = 30
void runtime_pow(int N){
int i, res = 0;
for(i = 1; i <= pow(2.0, (double)N); i=i+1)
res = res + 1;
}
```

**For N = 10**

Runtime: It's super fast. Time taken is **less than a second**(in few microseconds).

**For N = 15**

Runtime: It's also fast. Time taken is less than a second.  
**Approximately 0.004 seconds.**

**For N = 20**

Runtime: It's also fast. Time taken is less than a second.  
**Approximately 0.015 seconds.**

**For N = 25**

Runtime: It's also fast. Time taken is less than a second.  
**Approximately 0.069 seconds.**

**For N = 30**

Runtime: It's also fast. Time taken is a little more than a second. **Approximately 2 seconds.**

*As the value of **N increases**, the **runtime** of the loop also **increases** for runtime\_pow. But the rate of increase of runtime is fast for higher values of N.*

Example: For N=30, it took two seconds. For N=40, I had to wait for a very long time, so I stopped execution.

Q6 b) Time complexity 'closer' to that of the runtime\_rec

```
#include <stdio.h>
#include <stdlib.h>
#include <math.h>

void runtime_rec(int N, char * str){

    if (N==0) {
        //printf("%s\n", str);
        return;
    }

    str[N-1] = 'L';
    runtime_rec(N-1, str);
    str[N-1] = 'R';
    runtime_rec(N-1, str);

}

int main(int argc, char** argv) {

    int N = 0;
    char ch;
    char str[100];
    printf("run for: N = ");
    scanf("%d", &N);
    str[N] = '\0'; //to use it as a string of length N.
    printf("runtime_rec(%d)\n", N);
    runtime_rec(N, str);

}
```

**For N=10**

Runtime: It is very fast. Time taken is less than a second.  
**Approximately 0.001 seconds.**

**For N=15**

Runtime: It is also fast. Time taken is less than a second.  
**Approximately 0.222 seconds.**

**For N=20**

Runtime: It took some time. **Approximately 7.299 seconds.**

**For N=25**

Runtime: It took a long time. **Approximately 250 seconds.**

As the value of **N increases**, the **runtime** of the program also **increases**.

Moreover, we are printing the values of `runtime_rec`. From question 6a, we infer that printing values(`runtime_print`) take longer time than just calculating(`runtime_increment`).

Rate of increase of runtime is drastically fast for higher values of N.

Hence, **time complexity of `runtime_print` is closer to that of `runtime_rec`**.