

CSE 2320 - Homework 5

NAME: _____

Total points: 100 Topics: Recurrences , solved with methods: Master Theorem, Tree, Substitution (induction)

Convention: $\lceil \quad \rceil$ means rounded up and $\lfloor \quad \rfloor$ means rounded down.

P1. (23 points) Use the tree and table method to compute the Θ time complexity for $T(N) = 5T(\lfloor N/4 \rfloor) + 2N^3$.

Assume $T(0) = 1$ and $T(1) = 1$. Fill in the table below and finish the computations outside of it:

Level	Argument/ Problem size	Cost of one node	Nodes per level	Cost of whole level
0				
1				
2				
i				
$k=$ Leaf level. Write k as a function of N .				

Total tree cost calculation:

$T(N) = \Theta(\dots\dots\dots)$

Draw the tree. Show **levels 0,1,2** and the **leaves level**. Show the problem size $T(\dots)$ as a label next to the node and inside the node show the local cost (cost of one node) as done in class. For the leaf level and level 2 it suffices to show a few nodes.

P2. (23 points) Use the tree and table method to compute the Θ time complexity for $T(N) = 4T(N - 5) + 7$. Assume $T(N) = 1$ for all $0 \leq N \leq 4$. Assume N is a convenient value for your computations.

Fill in the table below and finish the computations outside of it:

Level	Argument/ Problem size	Cost of one node	Nodes per level	Cost of whole level
<i>0</i>				
<i>1</i>				
<i>2</i>				
<i>i</i>				
<i>k=</i> <i>Leaf level.</i> <i>Write k as a</i> <i>function of N.</i>				

Total tree cost calculation:

$T(N) = \Theta(\dots)$

Draw the tree. Show **levels 0,1,2** and the **leaves level**. Show the problem size $T(\dots)$ as a label next to the node and inside the node show the local cost (cost of one node) as done in class. For the leaf level and level 2 it suffices to show a few nodes.

P3. (36 points) Can you use the Master Theorem to solve the recurrences below? If yes, solve it with this method, if no, show why you cannot use it.

- a) $T(N) = 5T(\lfloor N/4 \rfloor) + 2N^3$. Assume $T(0) = 1$ and $T(1) = 1$.
- b) $T(N) = 4T(\lceil N/4 \rceil) + d$, for some constant $d > 0$. Assume $T(0) = 1$ and $T(1) = 1$.
- c) $T(N) = 6T(N/6) + 5N$, Assume $T(0) = 1$ and $T(1) = 1$.
- d) $T(N) = 8T(N/2) + cN^3 \lg N$, Assume $T(0) = 1$, and $T(1) = 1$.

P4. (4 points) Go to the Wikipedia webpage [https://en.wikipedia.org/wiki/Master_theorem_\(analysis_of_algorithms\)](https://en.wikipedia.org/wiki/Master_theorem_(analysis_of_algorithms)). See section "Inadmissible equations" and list the equation and the reason why it does not satisfy the Master Theorem requirements.

P5. (14 points) Show that $T(N) = 5T(\lfloor N/4 \rfloor) + 2N^3 = \Theta(N^3)$ by showing that it is $O(N^3)$ and also $\Omega(N^3)$. Assume $T(0) = 1$ and $T(1) = 1$

- a) (9 points) Use the induction method to show $O(N^3)$. As done in class, start with the inductive step and then check and refine for enough low values of N until the inductive step can be applied (See lecture from Wed, Oct 11).
- b) (5 points) Use just the definition with c and n_0 to show that it is $\Omega(N^3)$. Assume that $T(N) \geq 0$, for all $N \geq 0$. You should not need to use induction. (It is easier without induction.)

Write your answers in a document called **2320_H5.pdf**. It can be hand-written and scanned, but it must be uploaded electronically. Submit just the 2320_H5.pdf.

Remember to include your name at the top.