Independent Project EE3025

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## Chapter 1

# Introduction to Eda Playground

### 1.1 Blink Program

Implementation of Blink program on Eda Playground (online simulator) can be found here and GitHub link here

A short video explanation of working with the simulator and code can be found here.

### 1.2 Seven Segment display counter

Implementation of Seven segment display counter on Eda Playground can be found here and GithuB link here

A short video explanation of working with the simulator and code can be found here.



### Chapter 2

## FFT Algorithm

#### 2.1 Background

Let x[n] be a series of complex signals, for  $0 \le n \le N-1$  and X[k] the Discrete Fourier Transform of  $0 \le k \le N-1$ . The frequency domain signals X[k] can be obtained according to equation 2.2 or in its matrix form as can be seen in equation 2.3, where

$$W_N^k n = exp(\frac{2ikn\pi}{N}) \tag{2.1}$$

$$X[k] = \sum_{n=0}^{N-1} X[n] W_N^k n$$
 (2.2)

$$\begin{bmatrix} X[0] \\ \vdots \\ X[k] \end{bmatrix} = \begin{bmatrix} W_N^{0*0} & \dots & W_N^{0*N-1} \\ \vdots & \ddots & \vdots \\ W_N^{n-1*0} & \dots & w_N^{N-1*N-1} \end{bmatrix} * \begin{bmatrix} x[0] \\ \vdots \\ x[k] \end{bmatrix}$$
(2.3)

The computation of X[k] has complexity of  $O(n^2)$ . However, the expression of the equation 2.2 may be split in two terms according to equation 2.4.

$$X[k] = \sum_{n=0}^{N-1} X[2n]W_N^{2kn} + \sum_{n=0}^{N-1} X[2n+1]W_N^{2kn+1}$$
 (2.4)

Note that applying the properties  $W_N^{2kn}=W_{N/2}^{kn}$  and  $W_N^{k+N/2}=-W_N^{kn}$  in the above equation, we get

$$X[K] = Xe[k] + W_N^k Xo[k]$$

$$X[K + N/2] = Xe[k] - W_N^k Xo[k]$$
(2.5)

where

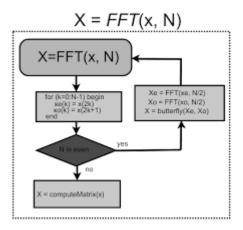
$$Xe[k] = \sum_{n=0}^{\frac{N}{2}-1} X[n] W_{\frac{N}{2}}^{kn}$$

$$Xo[k] = \sum_{n=0}^{\frac{N}{2}-1} X[2n+1] W_{\frac{N}{2}}^{kn}$$

According to equation 2.5, if N is a power of 2, the computation of X[k] has complexity of  $O(N*log_2(N))$ .

### 2.2 Recursive Algorithm

A very simple algorithm to compute the FFT can be defined taking advantage of the recursive nature of the FFT, as can be seen in



if the size N of the FFT is even then call two FFT of order N/2, one to compute the Fourier Transform of the signals with even index (x[2n]) and other to compute the signals with odd index (x[2n+1]).

The Fourier Transform will be scaled with the twiddle factor  $W_N{}^k$ . if N is odd

$$X_{1}(k) \xrightarrow{1} X(k) = X_{1}(k) + W_{N}^{k} X_{2}(k)$$

$$X_{2}(k) \xrightarrow{W_{N}^{k}} -1 X(\frac{N}{2} + k) = X_{1}(k) - W_{N}^{k} X_{2}(k)$$

then the FFT will be slowly calculated using the Fourier matrix of equation 2.3(Matrix multiplication).

If the length N of the FFT is of the form  $N=m*2^p$ , the the complexity of the algorithm of Figure 1 will be  $O(m^2)\times O(P*\log_2(p))$ . Sudo code of the above algorithm considering general case m=1 is

#### FFT Algorithm

```
NOTE: Procedure FFT is presented here in pseudo-code,
for a generic field F in which it is possible to define \omega,
a primitive n-th root of unity.
procedure FFT (A, n, w)
    # Preconditions:
       A is a Vector of length n;
        n is a power of 2;
        w is a primitive n-th root of unity.
    # The Vector A represents the polynomial
       a(z) = A[1] + A[2]*z + ... + A[n]*z^{(n-1)}.
    # The value returned is a Vector of the values
    # [ a(1), a(w), a(w^2), ..., a(w^{n-1}) ] # computed via a recursive FFT algorithm.
        return A
    else
        Aeven <-- Vector( [A[1], A[3], ..., A[n-1]] )
Aodd <-- Vector( [A[2], A[4], ..., A[n]] )
         Veven <-- FFT( Aeven, n/2, w^2)
         Vodd <-- FFT( Aodd, n/2, w^2)
         V <-- Vector(n) # Define a Vector of length n
         for i from 1 to n/2 do

V[i] <-- Veven[i] + w^(i-1)*Vodd[i]
             V[n/2 + i] <-- Veven[i] - w^(i-1)*Vodd[i]
         end do
         return V
    end if
end procedure
```

#### 2.3 Code and Links

SystemVerilog code implemented in Eda Playground can be found here and corresponding GitHub link is here.

A short video explaining the recursive fft code using SystemVerilog can be found here.

A faster algorithm iterative FFT is implemented in c code; Github link for c is here

Github Link for all the codes in project is here.