

Knowledge Graphs

Lecture 4 - Knowledge Representation with Ontologies
Excursion 5: Description Logics

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Knowledge Graphs

Lecture 4: Knowledge Representation with Ontologies

4.1 A Brief History of Ontologies

4.2 Why we do need Logic

Excursion 4: A Brief Recap of Essential Logics

Excursion 5: Description Logics

4.3 First Steps in OWL

4.4 More OWL

4.5 OWL and beyond

4.6 How to Design your own Ontology

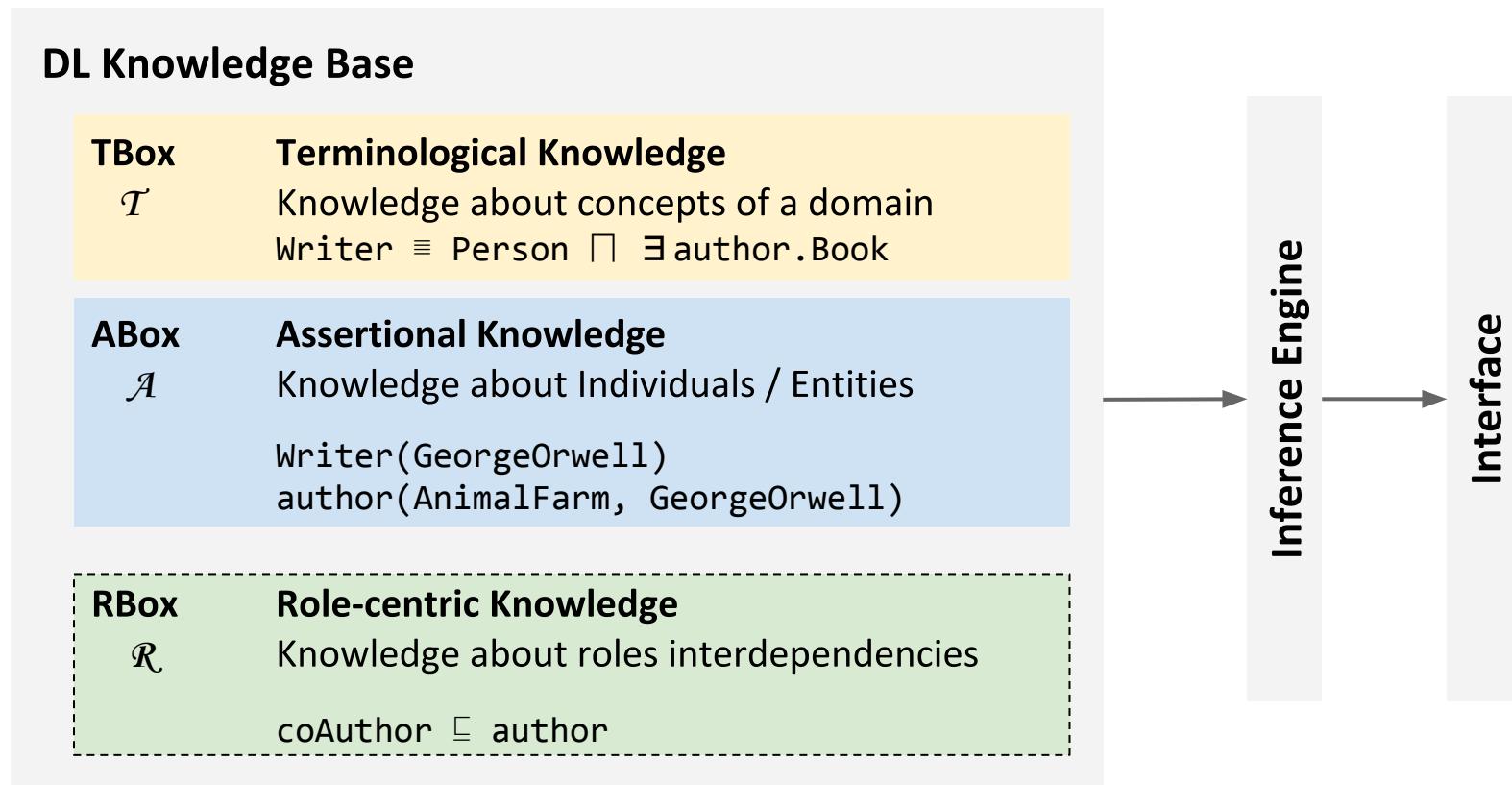
Knowledge Representation with FOL

- Why not simply take FOL for Ontologies?
- With FOL you can do everything, but...
 - *you could also program everything with assembler instead of higher programming languages*
- FOL has high **expressivity**
- FOL is too **bulky** for modelling
- FOL is not appropriate to find consensus in modelling
- FOL proof theoretically is very **complex** (semi-decidable)
- *FOL is of course not a Markup Language for the Web*
- **Basic idea:**
 - Look for an **appropriate fragment of FOL**....
 - ...and then make it a vocabulary for RDF(S)

Description Logics - DLs

- DLs are fragments of FOL (*compromise of expressivity and scalability*)
- A DL models **concepts**, **roles** and **individuals**, and their relationships.
- In DL from *simple descriptions* more *complex descriptions* are created with the help of **constructors**.
- DLs differ in the applied constructors (Expressivity)
- DLs have been developed from „semantic networks“
- DLs are **decidable** (most times)
- DLs possess **sufficient expressivity** (most times)
- DLs are related to modal logics
- Example for a DL:
W3C Standard OWL 2 DL is based on description logics $SROIQ(D)$

General DL Architecture



Description Logics - DLs

- DLs are a **family** of logic-based formalisms applied for knowledge representation
- ***ALC (Attribute Language with Complement)*** is the smallest deductively complete DL
 - **Conjunction, Disjunction, Negation** are class constructors, denoted as \sqcap , \sqcup , \neg
 - Quantifiers restrict domain and range of roles

Attributive Language with Complements - \mathcal{ALC}

Basic Building Blocks:

- Classes
- Roles / Properties
- Individuals
- **Person(GeorgeOrwell)**
Individual GeorgeOrwell is of class Person
- **Book(NineteenEightyfour)**
Individual NineteenEightyfour is of class Book
- **author(NineteenEightyfour, GeorgeOrwell)**
The book NineteenEightyfour has the author GeorgeOrwell

\mathcal{ALC} - Building Blocks

- **\mathcal{ALC} Atomic Types**
 - Concept names A, B, \dots
 - Special concepts
 - T - Top (universal concept)
 - \perp - Bottom concept
 - Role names R, S, \dots
- **\mathcal{ALC} Constructors**
 - Negation: $\neg C$
 - Conjunction: $C \sqcap D$
 - Disjunction: $C \sqcup D$
 - Existential quantifier: $\exists R.C$
 - Universal quantifier: $\forall R.C$

\mathcal{ALC} - Building Blocks

- **Class Inclusion**
 - $\text{Novel} \sqsubseteq \text{Book}$
 - *every novel is also a book*
 - equals FOL $(\forall x) (\text{Novel}(x) \rightarrow \text{Book}(x))$

- **Class Equivalence**
 - $\text{Novel} \equiv \text{Prose}$
 - *all Prose are exactly Novels*
 - equals FOL $(\forall x) (\text{Novel}(x) \leftrightarrow \text{Prose}(x))$

\mathcal{ALC} - Complex Class Relations

- **Conjunction** \sqcap
- **Disjunction** \sqcup
- **Negation** \neg

Novel \sqsubseteq (Book \sqcap Fiction) \sqcup (Paperback \sqcap \neg Poetry)

\mathcal{ALC}

$$(\forall x) (\text{Novel}(x) \rightarrow ((\text{Book}(x) \wedge \text{Fiction}(x)) \vee (\text{Paperback}(x) \wedge \neg \text{Poetry}(x)))$$

FOL

\mathcal{ALC} - Quantifiers on Roles

- **Strict Binding** of the Range of a Role to a Class

- $\text{Book} \sqsubseteq \forall \text{author}.\text{Writer}$
 - *A Book must be authored by a Writer*
 - $(\forall x) (\text{Book}(x) \rightarrow (\forall y) (\text{author}(x, y) \rightarrow \text{Writer}(y)))$

- **Open Binding** of the Range of a Role to a Class

- $\text{Book} \sqsubseteq \exists \text{author}.\text{Person}$
 - *Every Book has at least one author (who is a person)*
 - $(\forall x) (\text{Book}(x) \rightarrow (\exists y) (\text{author}(x, y) \wedge \text{Person}(y)))$

\mathcal{ALC} - Formal Syntax

- Production rules for creating classes in \mathcal{ALC} :
(A is an atomic class, C and D are complex Classes and R a Role)

○ $C, D ::= A \mid \top \mid \perp \mid \neg C \mid C \sqcap D \mid C \sqcup D \mid \exists R.C \mid \forall R.C$

- An \mathcal{ALC} TBox contains assertions of the form $C \sqsubseteq D$ and $C \equiv D$, where C, D are complex classes.
- An \mathcal{ALC} ABox contains assertions of the form $C(a)$ and $R(a, b)$, where C is a complex Class, R a Role and a, b Individuals.
- An \mathcal{ALC} -Knowledge Base contains an ABox and a TBox.

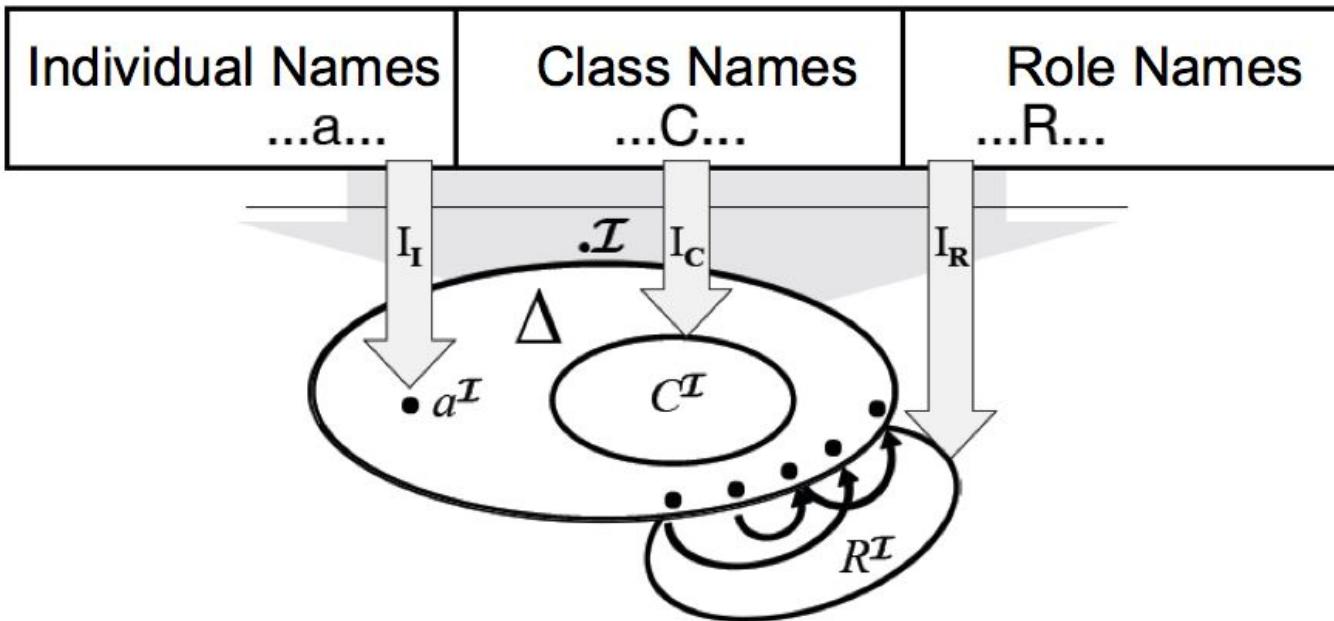
Syntax = symbols without meaning, defines rules, how to construct well-formed and valid sequences of symbols (strings)

\mathcal{ALC} - Model-theoretic Semantics

- The model-theoretic semantic for \mathcal{ALC} is defined via interpretations
- An **Interpretation** $I = (\Delta^I, .^I)$ contains
 - a set Δ^I (**Domain**) of individuals and
 - an **interpretation function** $.^I$ that maps
 - Individual names a
to domain elements $a^I \in \Delta^I$
 - Class names C
to a set of domain elements $C^I \subseteq \Delta^I$
 - Role names R
to a set of pairs of domain elements $R^I \subseteq \Delta^I \times \Delta^I$

Model-theoretic Semantics performs the semantic interpretation of artificial and natural languages by „*identifying meaning with an exact and formally defined interpretation with a model*“

\mathcal{ALC} - Model-theoretic Semantics



\mathcal{ALC} - Model-theoretic Semantics

- Extension (of Interpretation I) for complex classes:

- $\top^I = \Delta^I$ and $\perp^I = \emptyset$
- $(C \sqcup D)^I = C^I \cup D^I$ and $(C \sqcap D)^I = C^I \cap D^I$
- $(\neg C)^I = \Delta^I \setminus C^I$
- $\forall R.C = \{a \in \Delta^I \mid (\forall b \in \Delta^I)((a, b) \in R^I \rightarrow b \in C^I)\}$
- $\exists R.C = \{a \in \Delta^I \mid (\exists b \in \Delta^I)((a, b) \in R^I \wedge b \in C^I)\}$

\mathcal{ALC} - Model-theoretic Semantics

- Extension (of Interpretation I) for axioms:

- $C(a)$ holds, iff $a^I \in C^I$
- $R(a, b)$ holds, iff $(a^I, b^I) \in R^I$
- $C \sqsubseteq D$ holds, iff $C^I \subseteq D^I$
- $C \equiv D$ holds, iff $C^I = D^I$

Beyond \mathcal{ALC} - More DL Constructors

- **Number restrictions** for roles: $\geq 1 \text{ hasChild}$, $\leq 1 \text{ hasMother}$
- **Qualified number restrictions** for roles:
 $\geq 2 \text{ hasChild.Female}$, $\leq 1 \text{ hasParent.Male}$
- **Nominals** (definition by extension):
 $\{\text{CarbonDioxide}, \text{Methane}, \text{NitrousOxide}, \text{Ozone}, \text{WaterVapour}, \text{Chlorofluorocarbons}, \text{Hydrofluorocarbons}\}$
- **Concrete domains** (datatypes): $\text{hasAge.}(\geq 21)$
- **Inverse roles**: $\text{hasChild}^- \equiv \text{hasParent}$
- **Transitive roles**: $\text{hasAncestor} \sqsubseteq^+ \text{hasAncestor}$
- **Role composition**: $\text{hasParent}.\text{hasBrother(uncle)}$

Operator/Constructor	Syntax	Language
Conjunction	$A \sqcap B$	\mathcal{FL}
Value Restriction	$\forall R.C$	
Existential Quantifier	$\exists R$	
Top	\top	\mathcal{AL}^*
Bottom	\perp	
Negation	$\neg A$	
Disjunction	$A \sqcup B$	
Existential Restriction	$\exists R.C$	
Number Restriction	$(\leq nR) \ (\geq nR)$	
Set of Individuals	$\{a_1, \dots, a_n\}$	
Role Hierarchy	$R \sqsubseteq S$	\mathcal{H}
Inverse Role	R^{-1}	I
Qualified Number Restriction	$(\leq nR.C) \ (\geq nR.C)$	Q

Description Logics Family

- \mathcal{ALC} : Attribute Language with Complement
- \mathcal{S} : \mathcal{ALC} + Transitivity of Roles
- \mathcal{H} : Role Hierarchies
- \mathcal{O} : Nominals
- \mathcal{I} : Inverse Roles
- \mathcal{N} : Number restrictions $\leq nR$ etc.
- \mathcal{Q} : Qualified number restrictions $\leq nR.C$ etc.
- (\mathcal{D}) : Datatypes
- \mathcal{F} : Functional Roles
- \mathcal{R} : Role Constructors

- **OWL 2 DL** is $\mathcal{SROIQ}(\mathcal{D})$



First Steps in OWL

Next Lecture...

Picture References:

- [1] Archibald Thorburn, Little Owl and Scops Owl, 1925 [Public Domain]
[https://commons.wikimedia.org/wiki/Category:Athene_noctua_\(illustrations\)#/media/File:Archibald_Thorburn_Little_Owl_and_Scops_Owl.jpg](https://commons.wikimedia.org/wiki/Category:Athene_noctua_(illustrations)#/media/File:Archibald_Thorburn_Little_Owl_and_Scops_Owl.jpg)