Govher Sapardurdyyeva

$$|r - c_n| \leq \frac{b_0 - a_0}{2^{n+1}}$$

$$C_n = \frac{a_n + b_n}{2}$$

$$\delta_n - \alpha_n = \frac{b_n - a_n}{2^n}$$

$$a_n \leq r \leq \delta_n$$

$$|r_n - c_n| = |r - \frac{a_n + b_n}{2}|$$

$$|r - \frac{a_n + b_n}{2}| \leq \frac{b_n - a_n}{2} = \frac{b_n - a_n}{2^{n+1}}$$

$$|r - c_n| \leq \frac{b_n - a_n}{2^{n+1}}$$

14. Denote the successive intervals that arise in the bisection method by $[a_0, b_0]$, $[a_1, b_1]$, $[a_2, b_2]$, and so on. Show

a. $a_0 \le a_1 \le a_2 \le \cdots$ and $b_0 \ge b_1 \ge b_2 \ge \cdots$

b. $b_n - a_n = 2^{-n}(b_0 - a_0).$

c. $a_n b_n + a_{n-1} b_{n-1} = a_{n-1} b_n + a_n b_{n-1}$, for all n.

If f(an) = 0 then an is the root If f(an) \$ 0 then we check!

If fa) * f(an) < 0, then root lies [a,an] If flan) + f(b) <0, then root lies [amb] Then a, E [a,b] such that f(a,) = 0 The interval is [an, bn]

let the midpoint b cn, ther we

have fice, to then

f(an), f(cn) (o, in interval [ansn]

This oxives us an = an

 $b_n = C_n$

an remains the same, but by dicreases to a

In another case

f(G) * f(b,) (O and the root lies on interval

Therefore and = Ch, bn, = bn,

Con increases and bo stays some.

8. If a=0.1 and b=1.0, how many steps of the bisection method are needed to determine the root with an error of at most $\frac{1}{2}\times 10^{-8}$?

a=0.1 b=1.0

$$\frac{b-a}{2^n} \leq 8$$

$$n \geq \frac{\ln(b-\alpha) - \ln(\varepsilon)}{\ln(2)}$$

$$h \ge \frac{\ln(10-0.1) - \ln(\frac{1}{2} \cdot 10^{-8})}{\ln(2)}$$

$$W \ge \frac{|n(0.9) - (-19.11)}{0.69315}$$

 $n \ge 27.4234$

w. 18 steps

M Since and childh an increouses or Stays Same by decreases or stays same Which Shows $a_0 \le a_1 \le a_2 \le \dots = b_0 \ge b \ge b_1 \ge \dots$

next interval is either [a,, Co] or [co,bo]

$$b_1 - a_1 = \frac{b_0 - a_0}{2}$$

further iterations:

$$b_{3}-a_{3} = \frac{b_{1}-a_{1}}{2}$$

$$= \frac{b_{0}-a_{0}}{2^{2}}$$

$$b_3 - A_3 = b_3 - A_2$$

$$= b_3 - A_3$$

Which gives

$$b_{n} - a_{n} = \frac{b_{0} - a_{0}}{2^{n}} = 2^{-n} (b_{0} - a_{0})$$

The length of an interval is the difference between the interval's upper and lower endpoints

C) Using the induction theorem,
Suppose
$$a_n b_n + a_{n+} b_{n+} = a_{n+} b_n + a_n b_n + a_n$$

 $a_1b_1 + a_0b_0 = a_0b_1 + a_1b_0$, so since $c_0 = \underbrace{a_0 + b_0}$

If $f(a_0)$, $f(c_0) < 0$ then $a_1 = a_0$ and $b_1 = c_0$ So, $a_1b_1 + a_0b_0 = a_0c_0 + a_1b_0$ $= q_0b_1 + a_1b_0$

and if $f(c_0) \cdot f(b_0)(0)$ then $a_1 = c_0$, $b_1 = b_0$ then $a_1b_1 + a_0b_0 = c_0b_0 + a_0b_1$ $= a_1b_0 + a_0b_1$

Therefore both of them satisfy for n=1

anbn+ anbn- = anibnanbn- for all n

(5)

15. (Continuation) Can $a_0 = a_1 = a_2 = \cdots$ happen?

16. (Continuation) Let $c_n = (a_n + b_n)/2$. Show that

$$\lim_{n\to\infty} c_n = \lim_{n\to\infty} a_n = \lim_{n\to\infty} b_n$$

(15) $a_0 = a_1 = a_2 = \dots$ can happen if and only if $f(a_1) = a_2 = \dots$ So if $f(a_1) = a_2$ is the

 $f(a_0) = 0$, so iff a_0 is the root

(16) from previous

$$b_n - a_n = \frac{b_o - a_o}{2^n}$$

interval decreases as n-xx, so an and bn get closer

to each othe which

makes them converge to the same limit. So,

lim an = lim bn and

Suppose liman = lim bn = L

exists for L

then

$$\lim_{n\to\infty} c_n = \frac{L}{2} = L$$

lim an = lim bn = lim Cn

22. If the bisection method is applied with starting interval $[2^m, 2^{m+1}]$, where m is a positive or negative integer, how many steps should be taken to compute the root to full machine precision on a 32-bit word-length computer?

To compute number of iterations required for bisection method:

$$\frac{b-a}{2^{n}} \le \mathcal{E}$$

$$n = \log_{2}\left(\frac{b-a}{\mathcal{E}}\right)$$

$$\mathcal{E} = 1.19 \cdot 10^{-7} \cdot 2^{-23}$$

$$interval is \left[2^{m}, 2^{m+1}\right]$$

$$n = \log_{2}\left(\frac{2^{m+1}-2^{m}}{2^{-23}}\right)$$

$$= \log_{2}\left(\frac{2^{m}}{2^{-23}}\right)$$

$$= \log_{2}\left(\frac{2^{m}}{2^{-23}}\right)$$

$$= \log_{2}\left(2^{m+23}\right)$$

$$= m + 23$$

in number of steps required is n = m + 23