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Q2. Ans: Let 'A' is array of size 5, and elements are 8, 5, 7, 3, 2.

A =

0	1	2	3	4
8	5	7	3	2

, n = 5

② Bubble sort

For 1st pass

8	5	5	5	5
5	8	7	7	7
7	7	8	3	3
3	3	3	8	2
2	2	2	2	8

 } 4 comp & 4 swaps
 8 at sorted position.

For 2nd pass

5	5	5	5
7	7	3	3
3	3	7	2
2	2	2	7
8	8	8	8

 } 3 comp & 3 swaps
 7 & 8 are present at sorted position.

For 3rd pass

5	3	3
3	5	2
2	2	5
7	7	7
8	8	8

 } 2 Comparisons & 2 swaps.
 5, 7, & 8 are present at sorted position.

For 4th pass

3	2
2	3
5	5
7	7
8	8

 } 1 Comparison & 1 swap
 all elements are present at sorted position.

No of passes = 4 i.e. (n-1) passes

No of comparisons = $1 + 2 + 3 + \dots + (n-1) = \frac{n(n-1)}{2} = O(n^2)$

Max/No of swaps = $1 + 2 + 3 + 4 \rightarrow 1 + 2 + \dots + (n-1) = \frac{n(n-1)}{2} = O(n^2)$

hence time complexity of Bubble sort is $O(n^2)$.

⑥ Quick Sort:

Let A is array of size 9, and elements are 7, 6, 10, 5, 9, 2, 1, 15, 7.

	l	1	2	3	4	5	6	7	h
A	0	1	2	3	4	5	6	7	8
	7	6	10	5	9	2	1	15	7

Sort given array by using partition method
Pivot = A[l] = 7

7	6	10	5	9	2	1	15	7
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$i \leftarrow$ swap equal $\rightarrow j$

If any element greater than pivot then i stop &
If any element less than pivot then j will be stop.

7	6	7	5	9	2	1	15	10
---	---	---	---	---	---	---	----	----

i j

7	6	7	5	9	2	1	15	10
---	---	---	---	---	---	---	----	----

i swap j

7	6	7	5	1	2	9	15	10
---	---	---	---	---	---	---	----	----

i j

7	6	7	5	1	2	9	15	10
---	---	---	---	---	---	---	----	----

j i

here $i > j$ then swap pivot with A[j],

2	6	7	5	1	7	9	15	10
---	---	---	---	---	---	---	----	----

all element in left side of pivot is small & equal to pivot

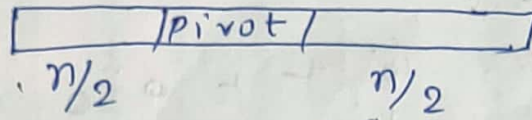
all element in right side of pivot is greater than pivot.

Again use partition of left of pivot & right of pivot.
Recursively.

⇒ Time complexity of Quick sort:

① Best case time analysis of Quick sort.

It occurs when list is divided into 2 part after placement of pivot at its proper location.



2 sub problems exist of size $n/2$

Recurrence Relation

$$T(n) = \begin{cases} c & \text{if } n=1, \text{ b/c single element are always sorted.} \\ 2T(n/2) + cn & \text{if } n > 1 \end{cases}$$

no of sub problems

solving above equation using substitution method

$$T(n) = 2T(n/2) + cn \rightarrow (i)$$

$$T(n/2) = 2T(n/4) + cn/2$$

$$T(n/4) = 2T(n/8) + cn/4$$

$$T(n) = 8T(n/8) + cn + 2cn$$

3rd term

$$T(n) = 2^3 T\left(\frac{n}{2^3}\right) + 3cn$$

kth terms

$$T(n) = 2^k T\left(\frac{n}{2^k}\right) + Kcn$$

$$\text{let } n = 2^k$$

$$\log_2 n = \log_2 2^k$$

$$k = \log_2 n$$

$$T(n) = nT(1) + cn \log_2 n$$

$$T(n) = n(n \log_2 n)$$

Best case time complexity of Quick sort.
→ omega is use b/c this is best case

② worst case time analysis of quick sort

It occurs when the list is **already sorted** in Ascending order.

Let A is a array of size 5 & elements are

	0	1	2	3	4
A	5	4	3	2	1
	i			J	

If any element greater & equal to pivot then i should be stop & If any element less than pivot then J should be stop

Pivot = 5

5	4	3	2	1
			J	i

n comparisons

here $i > J$ then swap pivot with $A[J]$

1	4	3	2	5
J	i			

$(n-1)$ element ie $(n-1)$ comparisons

here again $i > J$ then swap pivot with $A[J]$

1	4	3	2	5

$(n-2)$ element ie $(n-2)$ comparison

1	2	3	4	5

$(n-3)$ element ie $(n-3)$ comparison

1	2	3	4	5

$(n-4)$ element ie $(n-4)$ comparison

the total no of comparisons

$$T(n) = n + (n-1) + (n-2) + \dots + 1$$

$$= \frac{n(n+1)}{2}$$

$$T(n) = O(n^2)$$

• worst case time complexity of quick sort.

Average case time complexity of Quick sort is also $\theta(n \log_2 n)$.

⇒ category of both bubble sort & Quick sort is in-place algorithm because it does not take extra space.

⇒ Comparing bubble sort & Quick sort with insertion sort & merge sort.

class	Bubble sort	Quick sort	insertion sort	merge sort
Worst-case performance	$O(n^2)$ comp $O(n^2)$ swaps	$O(n^2)$	$O(n^2)$ comp $O(n^2)$ swaps	$O(n \log n)$
Best-case performance	$O(n)$ comp $O(1)$ swaps	$O(n \log n)$ (simple partition) or $O(n)$ (three-way partition & equal key)	$O(n)$ comp $O(1)$ swaps	$O(n \log n)$ typical, $O(n)$ natural variant
Average performance	$O(n^2)$ comp $O(n^2)$ swaps	$O(n \log n)$	$O(n^2)$ comp $O(n^2)$ swaps	$O(n \log n)$

Some important points about above four sorting algorithms.

- (i) Bubble sort, insertion sort and merge sort all are stable but Quick sort is not stable.
- (ii) only Bubble sort and insertion sort are Adaptive out of all sorting algorithms.
- (iii) Bubble sort, Quick sort & insertion sort are in-place sorting algorithms but merge sort is outplace sorting method is more efficient in term of time complexity.