Randomized iterative methods for linear systems and inverting matrices

Robert Gower joint work with Peter Richtárik



Optimization and Big Data 2015, 7th of May, Edinburgh.

- Gower, Robert M., Richtárik, Peter, April 2015.
 Randomized Iterative Methods for Linear Systems (in progress)
- Gower, Robert M., Richtárik, Peter, April 2015.
 Randomized Iterative Methods for Inverting Matrices (in progress)

The Problem

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Solve a consistent linear system $Ax_* = b$, where $A \in \mathbb{R}^{m \times n}$, $m \ge n$.

Solve with an iterative method

$$x_{k+1} = \mathsf{update_formula}(A, x_k)$$
 such that $x_k \to x_*.$

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$$\begin{bmatrix} ----A_{1:} --- \\ -----A_{2:} --- \\ \vdots \\ \vdots \\ -------- \\ -------- \end{bmatrix} \begin{bmatrix} x_*^1 \\ \vdots \\ x_*^n \end{bmatrix} = \begin{bmatrix} b^1 \\ b^2 \\ \vdots \\ \vdots \\ b^{m-1} \\ b^m \end{bmatrix}$$

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Notation: Let $||x||_B^2 := x^T B x$ for B > 0.

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New Gaussian methods

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Iteratively Inverting matrices

Future work: Randomized Preconditioning?

The Return of old methods

- ▶ Old methods (Kaczmarz 1937, Guass-Seidel 1823) make a randomized return, why?
- Often suitable for Big Data problems (short recurrence, low memory,...etc)
- Easy to implement
- Easy to analyse, good complexity
- Often fits in parallel architecture

Kaczmarz method

Choose the *i*th row then iterate

$$x_{k+1} = \arg\min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{x} - \mathbf{x}_k\|_2^2$$
 subject to $A_{i:} \mathbf{x} = b^i$.

$$x_{k+1} = x_k - \frac{A_{i:}x_k - b^i}{\|A_{i:}\|_2^2} A_{i:}^T$$

- Developed in 1937 Kaczmarz
- ▶ Implemented in the first CT scanner 1972





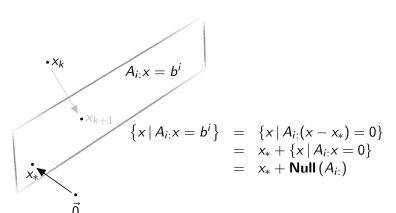


¹G.N. Hounsfield. Computerized transverse axial scanning (tomography): Part I. description of the system. British Journal Radiology. 1973

The Kaczmarz method (Stochastic gradient)

Kaczmarz Interpretation

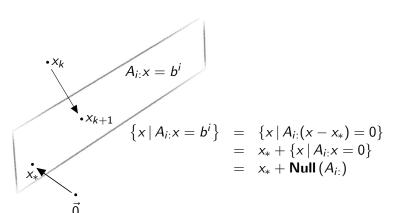
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How to choose i

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 subject to $A_{i:x} = b^i$.

- ▶ Traditional Kaczmarz: Cycle i = 1, 2, ..., m. Slow in practice + difficult to interpret complexity
- ▶ Pick *i* with probability $p_i = 1/m$. Better in practice + difficult to interpret complexity
- ▶ Break-Through (Strohmer & Vershynin, 2009): pick i with probability $p_i = ||A_{i:}||_2^2 / ||A||_F^2$.

$$\mathbf{E}\left[\|x_k - x_*\|_2^2\right] \le \left(1 - \frac{\lambda_{\min}(A^T A)}{\|A\|_F^2}\right)^k \|x_0 - x_*\|_2^2.$$

$$\lambda_{\min}(A^T A) / \|A\|_F^2 = 1 / \|A\|_F^2 \|A^{\dagger}\|_2^2$$

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Coordinate Descent (Gauss-Seidel)

Choose the *i*th coordinate then

$$x_{k+1} = \arg\min_{x \in \mathbb{R}^n} \lVert Ax - b
Vert_2^2 \quad ext{ subject to } \quad x = x_k + te_i, \quad t \in \mathbb{R}.$$

$$x_{k+1} = x_k - \frac{(A_{:i})^T (Ax_k - b)}{\|A_{:i}\|_2^2} e_i$$

Note that
$$||Ax - b||_2^2 = ||A(x - x_*)||_2^2 = ||x - x_*||_{A^T A}^2$$

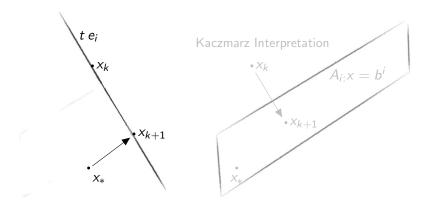
Convergence (Leventhal & Lewis, 2010)

Pick *i* with probability $p_i = ||A_{:i}||_2^2/||A||_F^2$.

$$\mathbf{E}\left[\|x_k - x_*\|_{A^T A}^2\right] \le \left(1 - \frac{\lambda_{\min}(A^T A)}{\|A\|_F^2}\right)^k \|x_0 - x_*\|_{A^T A}^2.$$

Coordinate Descent Interpretation

$$x_{k+1} = \arg\min_{x \in \mathbb{R}^n} ||x - x_*||_{A^T A}^2$$
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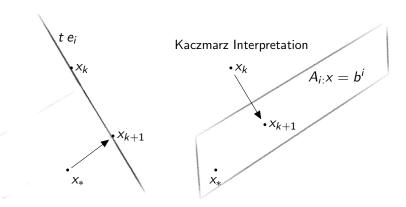


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Framework for designing randomized methods

Choose $B \succ 0 \in \mathbb{R}^{n \times n}$ and a random matrix S independently drawn at each iteration k. **Two** viewpoints of the **same** method.

(I)
$$x_{k+1} = \arg\min_{x \in \mathbb{R}^n} \|x - x_k\|_B^2$$
 s. t. $S^T A x = S^T b$,
(II) $x_{k+1} = \arg\min_{x \in \mathbb{R}^n} \|x - x_*\|_B^2$ s. t. $x \in x_k + B^{-1} \mathbf{Range} \left(A^T S\right)$

(I): Project x_k onto a randomly compacted system.

$$\begin{bmatrix} S^T \\ A \end{bmatrix} = \begin{bmatrix} S^T A \\ X \end{bmatrix}$$

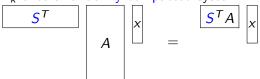
Kaczmarz fits nicely with B = I and $S = e_i$. **Block Kaczmarz** choose B = I and $S = I_{:C}$ a subset of columns of identity.

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Kaczmarz fits nicely with B = I and $S = e_i$. **Block Kaczmarz** choose B = I and $S = I_{:C}$ a subset of columns of identity.

Coordinate descent methods fit (II)

$$\textbf{(II)} \quad \textit{x}_{k+1} = \arg\min_{x \in \mathbb{R}^n} \lVert x - \textit{x}_* \rVert_B^2 \quad \text{ subject to } \quad \textit{x} \in \textit{x}_k + \mathsf{Range}\left(B^{-1}A^TS\right)$$

▶ Least-Squares Coord. Desc: With $B = A^T A$ and $S = Ae_i = A_{:i}$

$$x_{k+1} = \arg\min_{x \in \mathbb{R}^n} \|Ax - b\|_2^2$$
 subject to $x = x_k + t e_i$, $t \in \mathbb{R}$.

Stochastic Newton (SDNA² Method 1) Let $S = Al_{:C} = A_{:C}$ subset of columns of A

$$x_{k+1} = \arg\min_{x \in \mathbb{R}^n} ||Ax - b||_2^2$$
 subject to $x = x_k + t I_{:C}, t \in \mathbb{R}^{|C|}$.

▶ Positive Definite Coord. Desc: When $A \succ 0$, B = A and $S = I_{:C}$ then

$$x_{k+1} = \arg\min_{\mathbf{x} \in \mathbb{R}^n} \underbrace{\frac{1}{2} \mathbf{x}^T A \mathbf{x} - \mathbf{x}^T b}_{=\|\mathbf{x} - \mathbf{x}_*\|_A^2} \quad \text{ subject to } \quad \mathbf{x} = \mathbf{x}_k + t \, I_{:C}, \quad t \in \mathbb{R}^{|C|},$$

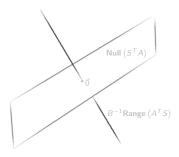
²Qu, Z., Richtárik, P., Takáč, M., & Fercoq, O. (2015). SDNA: Stochastic Dual Newton Ascent for Empirical Risk Minimization.

Framework for randomized methods

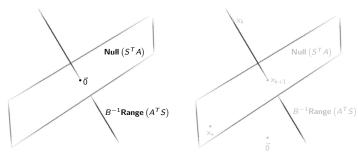
Geometry and Duality

Let
$$\langle x, y \rangle = x^T B y$$
 in \mathbb{R}^n . $\{ x \mid S^T A x = S^T b \} = x_* + \{ x \mid S^T A x = 0 \}$ and

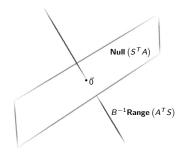
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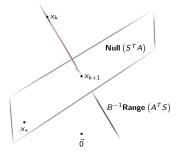


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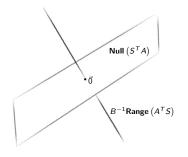


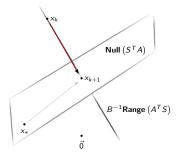


I Project x_k onto $x_* + \text{Null}(S^T A)$

$$x_{k+1} = \arg\min_{\mathbf{x} \in \mathbb{R}^n} \lVert \mathbf{x} - \mathbf{x}_k \rVert_B^2 \quad \text{ subject to } \quad S^T A \mathbf{x} = S^T b$$

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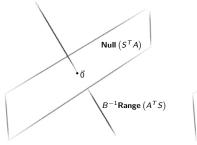
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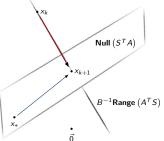
$$x_{k+1} = \arg\min_{x \in \mathbb{R}^n} ||x - x_k||_B^2$$
 subject to $S^T A x = S^T b$

II Project x_* onto $x_k + B^{-1}$ Range $(A^T S)$

$$x_{k+1} = \arg\min_{\mathbf{x} \in \mathbb{R}^n} ||\mathbf{x} - \mathbf{x}_*||_B^2$$
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Let
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The Solution

Assuming A^TS has full column rank \Rightarrow closed form solution

II Project
$$x_*$$
 onto $x_k + B^{-1}$ Range (A^TS)

$$x_{k+1} = x_k + \operatorname{proj}_{B^{-1}\text{Range}(A^TS)}(x_* - x_k)$$

$$= x_k + B^{-1}A^TS(S^TAB^{-1}A^TS)^{-1}S^TA(x_k - x_*)$$

$$= x_k + B^{-1}A^TS\underbrace{(S^TAB^{-1}A^TS)^{-1}}_{\text{Solve small system.}} S^T(Ax_k - b)$$

Project x_k onto $x_* + \text{Null}(S^T A)$

$$x_{k+1} = x_* + \text{proj}_{\text{Null}(A^T S)}(x_k - x_*)$$

$$= x_* + (I - B^{-1} \underbrace{A^T S(S^T A B^{-1} A^T S)^{-1} S^T A}_{Z})(x_k - x_*)$$

$$= x_* + (I - B^{-1} Z)(x_k - x^*).$$

The Solution

Assuming $A^T S$ has full column rank \Rightarrow closed form solution

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$$= x_k + B^{-1}A^T S (S^T A B^{-1}A^T S)^{-1} S^T A (x_k - x_*)$$

$$= x_k + B^{-1}A^T S (S^T A B^{-1}A^T S)^{-1} S^T (A x_k - b)$$
Solve small system.

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$$= x_* + (I - B^{-1} Z)(x_k - x^*).$$

All randomness is in the range space projection $B^{-1}Z$

$$\mathbf{Z} \stackrel{\text{def}}{=} A^T S (S^T A B^{-1} A^T S)^{-1} S^T A.$$

For analysis, fixed point form

$$x_{k+1}-x_*=(I-B^{-1}Z)(x_k-x_*).$$

All randomness is in the range space projection $B^{-1}Z$

$$\mathbf{Z} \stackrel{\text{def}}{=} \mathbf{A}^T \mathbf{S} (\mathbf{S}^T \mathbf{A} \mathbf{B}^{-1} \mathbf{A}^T \mathbf{S})^{-1} \mathbf{S}^T \mathbf{A}.$$

For analysis, fixed point form

$$\mathbf{E}[x_{k+1} - x_* | x_k] = (I - B^{-1}\mathbf{E}[Z])(x_k - x_*).$$

$$\mathbf{E}[\mathbf{E}[x_{k+1} - x_* | x_k]] = \mathbf{E}[x_{k+1} - x_*]$$

$$= \mathbf{E}[(I - B^{-1}\mathbf{E}[Z])(x_k - x_*)]$$

$$= (I - B^{-1}\mathbf{E}[Z])\mathbf{E}[x_k - x_*].$$

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Convergence Theorems

$$\|\mathbf{E}[x_k] - x_*\| \le \left(1 - \lambda_{\min}(B^{-1/2}\mathbf{E}[Z]B^{-1/2})\right)^k \|x_0 - x_*\|$$

and when $\mathbf{E}[Z]$ nonsingular

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$$\mathbf{E}[\|x_k - x_*\|] \le \left(1 - \lambda_{\min}(B^{-1/2}\mathbf{E}[Z]B^{-1/2})\right)^{\kappa} \|x_0 - x_*\|$$

Theorem (General S)

Let S be a random matrix such that A^TS full column rank. Then for all $k \ge 0$,

$$\|\mathbf{E}[x_k] - x_*\| \le \rho^k \|x_0 - x_*\|,$$

where

$$\rho = 1 - \lambda_{\min}(B^{-1/2}\mathbf{E}[Z]B^{-1/2})$$
 and $0 \le \rho \le 1$.

Proof.

Taking conditional expectation with respect to x_k , we get

$$\mathbf{E}[x_{k+1} - x_* \mid x_k] = (I - B^{-1}\mathbf{E}[Z])(x_k - x_*). \tag{1}$$

Taking full expectation, we get

$$\mathbf{E}[x_{k+1} - x_*] = \mathbf{E}[\mathbf{E}[x_{k+1} - x_* \mid x_k]]
\stackrel{(1)}{=} \mathbf{E}[(I - B^{-1}\mathbf{E}[Z])(x_k - x_*)]
= (I - B^{-1}\mathbf{E}[Z])\mathbf{E}[x_k - x_*].$$

Now unroll the recurrence and apply the operator norm. As $B^{-1}Z$ is a projection, by Jensen's inequality with the convex functions λ_{\max} and $-\lambda_{\min}$, we have

$$0 \le \lambda_{\max}(B^{-1}\mathbf{E}[Z]) \le \lambda_{\max}(B^{-1}Z) \le 1.$$

Unifying previous methods & analysis

Theorem (Discrete random vector)

Let S be discrete r.v. such $S = s_i \in \mathbb{R}^n$ (for concreteness, think of $s_i = e_i$) with probability $p_i > 0$, for i = 1, ..., m, and let

$$S = [s_1, \ldots, s_m]$$
.

Then

$$x_{k+1} = x_k + \frac{s_i^T (Ax_k - b)}{s_i^T A B^{-1} A^T s_i} B^{-1} A^T s_i, \quad \text{with prob } p_i.$$

If we choose

$$p_i = rac{\mathbf{s}_i^T A B^{-1} A^T \mathbf{s}_i}{\|B^{-1/2} A^T \mathbf{S}\|_F^2}, \quad \text{for } i = 1, \dots, m,$$

then

$$\mathbf{E}\left[Z\right] = \frac{A^T \mathbf{S} \mathbf{S}^T A}{\|B^{-1/2} A^T \mathbf{S}\|_F^2} \quad \text{and} \quad \rho = 1 - \frac{\lambda_{\min}\left(B^{-1/2} A^T \mathbf{S} \mathbf{S}^T A B^{-1/2}\right)}{\|B^{-1/2} A^T \mathbf{S}\|_F^2}.$$

Furthermore, if S^TA has full column rank then $\rho < 1$.

Proof.

$$\mathbf{E}[Z] = \sum_{i=1}^{m} A^{T} s_{i} (s_{i}^{T} A B^{-1} A^{T} s_{i})^{-1} s_{i}^{T} A \rho_{i}$$

$$= \frac{1}{\|B^{-1/2} A^{T} \mathbf{S}\|_{F}^{2}} \sum_{i=1}^{m} A^{T} s_{i} s_{i}^{T} A$$

$$= \frac{1}{\|B^{-1/2} A^{T} \mathbf{S}\|_{F}^{2}} A^{T} \mathbf{S} \mathbf{S}^{T} A.$$

Thus the ρ is given by

$$\rho = 1 - \lambda_{\min} \left(B^{-1/2} \mathbf{E} \left[Z \right] B^{-1/2} \right) = 1 - \frac{\lambda_{\min} \left(B^{-1/2} A^T \mathbf{S} \mathbf{S}^T A B^{-1/2} \right)}{\| B^{-1/2} A^T \mathbf{S} \|_F^2}.$$

As $\mathbf{S}^T A$ has full column rank, $\mathbf{E}[Z]$ is positive definite and $\rho < 1$.

Unifying previous methods & analysis

$$p_i = \frac{s_i^T A B^{-1} A^T s_i}{\|B^{-1/2} A^T \mathbf{S}\|_F^2} \quad \text{with} \quad \rho = 1 - \frac{\lambda_{\min} \left(B^{-1/2} A^T \mathbf{S} \mathbf{S}^T A B^{-1/2}\right)}{\|B^{-1/2} A^T \mathbf{S}\|_F^2}.$$

Name	В	S	S	p _i	1- ho
Kaczmarz	I	ei	1	$ A_{i:} _2^2/ A _F^2$	$\lambda_{\min}(A^TA)/\ A\ _F^2$
CD $ Ax - b _2^2$	A^TA	$A_{:i}$	Α	$ A_{:i} _2^2/ A _F^2$	$\lambda_{min}(A^TA)/\ A\ _F^2$
$CD x^T Ax/2 - x^T b$	Α	e _i	1	$A_{ii}/\mathbf{Tr}\left(A\right)$	$\lambda_{min}\left(\mathit{A}\right)/Tr\left(\mathit{A}\right)$

New possibilities suggested:

ightharpoonup Covers new cases, e.g., $S = \alpha_i e_i + \alpha_j e_j$

Unifying previous methods & analysis

$$p_i = \frac{s_i^T A B^{-1} A^T s_i}{\|B^{-1/2} A^T \mathbf{S}\|_F^2} \quad \text{with} \quad \rho = 1 - \frac{\lambda_{\min} \left(B^{-1/2} A^T \mathbf{S} \mathbf{S}^T A B^{-1/2}\right)}{\|B^{-1/2} A^T \mathbf{S}\|_F^2}.$$

Name	В	-	S	"	1- ho
Kaczmarz	I	ei	1	$ A_{i:} _2^2/ A _F^2$	$\lambda_{\min}(A^TA)/\ A\ _F^2$
$CD \ \ Ax - b\ _2^2$	A^TA	$A_{:i}$	Α	$ A_{:i} _2^2/ A _F^2$	$\lambda_{min}(A^TA)/\ A\ _F^2$
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Guassian based sampling

Why not make S a continuous random matrix? Sample $S = \xi \sim N(0, \Sigma)$ a normal random variable then

$$Z = A^{T} S (S^{T} A B^{-1} A^{T} S)^{-1} S^{T} A = \frac{A^{T} \xi \xi^{T} A}{\xi^{T} A B^{-1} A^{T} \xi}.$$
$$x_{k+1} = x_{k} - \frac{\xi^{T} (A x_{k} - b)}{\xi^{T} A B^{-1} A^{T} \xi} B^{-1} A^{T} \xi.$$

Iteration cost $O(\text{product } A^T \cdot \xi)$.

The convergence rate determined by

$$\rho = 1 - \lambda_{\min}(B^{-1/2}\mathbf{E}[Z]B^{-1/2}) = 1 - \lambda_{\min}\left(\mathbf{E}\left[\frac{\bar{\xi}\bar{\xi}^T}{\bar{\xi}^T\bar{\xi}}\right]\right),$$

where $\bar{\xi} = B^{-1/2}A^T\xi \sim N(0,\Omega)$, and $\Omega = B^{-1/2}A\Sigma A^TB^{-1/2}$.

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─ New Gaussian methods

New Gaussian Methods

Sample
$$S = \xi \sim N(0, \Sigma)$$
. Let $\eta \sim N(0, I)$.
Gauss. Kaczmarz $B = I$ and $\Sigma = I$

$$x_{k+1} = \arg\min_{\mathbf{x} \in \mathbb{R}^n} ||\mathbf{x} - \mathbf{x}_k||_2^2$$
 subject to $\eta^T (A\mathbf{x} - b) = 0$.

Gauss Least-squares
$$B = A^T A$$
 and $\Sigma = AA^T$

$$x_{k+1} = \arg\min_{\mathbf{x} \in \mathbb{R}^n} \|A\mathbf{x} - b\|_2^2 \quad \text{ subject to } \quad \mathbf{x} = \mathbf{x}_k + t\,\eta, \quad t \in \mathbb{R}.$$

Gauss. Pos. Def. B = A and $\Sigma = I$

$$x_{k+1} = \arg\min_{x \in \mathbb{R}^n} 1/2x^T A x - x^T b$$
 subject to $x = x_k + t \eta$, $t \in \mathbb{R}$.

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Future work: Randomized Preconditioning?

Dense Overdetermined Gaussian Matrix

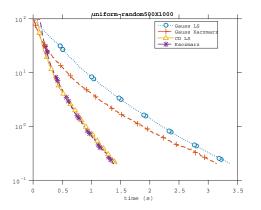


Figure: $m \times n = 500X1000$, A = randn(m, n)

Dense matrix \Rightarrow High iteration cost of Gaussian methods $O(A \cdot \eta)$.

Sparse Square Gaussian Matrix

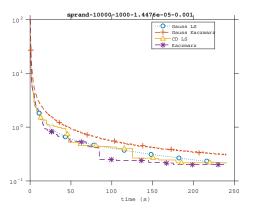
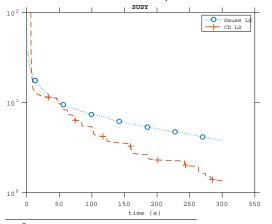


Figure : $m \times n = 10000 \times 1000$, density = $1/\sqrt{m} = 1\%$; $\kappa = \sqrt{n}$; A = sprandsym(n,density,rc)

Sparse matrices \Rightarrow Guass methods become competitive.

Regression SUSY

The SUSY³ Classification problem



Solving least-squares regression $\min \|Ax - b\|_2^2$ with $m = 5 \cdot 10^6$ and n = 18

 $^{^3}$ Baldi, P., P. Sadowski, and D. Whiteson. Searching for Exotic Particles in High-energy Physics with Deep Learning, Nature Communications 5 (July 2, 2014)

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Why iteratively invert a matrix $A \in \mathbb{R}^{n \times n}$?

- Needed to calculate Schur complements, a orojection operator...etc
- ▶ Iterative is good when we can tolerate an error
- ▶ Iterative is good when we have an initial guess $X_0 \approx A^{-1}$.
- Staging for randomized variable metric methods and randomized Preconditioning.

New context: $A \in \mathbb{R}^{n \times n}$ non-singular.

Framework

- \triangleright Assume we observe S^TA where S is random.
- ▶ Given $X_k \approx A \in \mathbb{R}^{n \times n}$, we want to iteratively calculate

$$X_{k+1} = \mathsf{update_formula}(S^T A, X_k)$$

such that $X_{k+1} \to A^{-1}$.

$$X_{k+1} = \arg\min_{X \in \mathbb{R}^{n \times n}} ||X - X_k||^2_{\mathsf{Frobenius}(B)}$$
 s.t. $S^T A X = X$,

The solution

$$X_{k+1} = X_k + \text{proj}_{B^{-1}\text{Range}(A^TS)}(A^{-1} - X_k)$$

$$= X_k + B^{-1}A^TS(S^TAB^{-1}A^TS)^{-1}S^TA(A^{-1} - X_k)$$

$$= X_k + B^{-1}A^TS(S^TAB^{-1}A^TS)^{-1}S^T(I - AX_k).$$

What about the symmetric case $A^T = A$?

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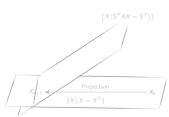
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What about the symmetric case $A^T = A$?

Symmetric matrices

$$X_{k+1} = \arg\min_{X \in \mathbb{R}^{n \times n}} ||X - X_k||^2_{\mathsf{Frobenius}(B)}$$

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 $X = X^T$



⁴Gower and Gondzio 2014

Symmetric matrices

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$$\text{s.t.} \quad S^T A X = X$$

$$X = X^T$$

$$X_{k+1} = \arg\min_{X \in \mathbb{R}^{n \times n}} \|X - X_k\|^2_{\mathsf{Frojection}}$$

$$X_k = X^T$$

Solution:⁴

$$\begin{aligned} X_{k+1} &= X_k + \operatorname{proj}_{B^{-1}\mathsf{Range}(AS)}(X_k - A^{-1})\operatorname{proj}_{B^{-1}\mathsf{Range}(AS)} \\ &- (X_k - A^{-1})\operatorname{proj}_{B^{-1}\mathsf{Range}(AS)} - \operatorname{proj}_{B^{-1}\mathsf{Range}(AS)}(X_k - A^{-1}) \end{aligned}$$

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$$X_{k+1} = \max_{X \in \mathbb{R}^{n \times n}} \|X - X_k\|^2_{\mathsf{Projection}}$$

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Solution:4

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Theorem (Convergence)

Let S be equal to a column of a full rank matrix $\mathbf{S} := [s_1, \dots, s_n]$ with probability $\|B^{-1/2}As_i\|^2/\|B^{-1/2}A\mathbf{S}\|_F^2$. Then from a given $X_0 \in \mathbb{R}^{n \times n}$, the iteration

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converges with

$$\mathbf{E}\left[\|X_k - A^{-1}\|^2_{Frob(B)}\right] = \left(I - \frac{1}{\kappa_F^2(B^{-1/2}A\mathbf{S})}\right)^k \|X_0 - A^{-1}\|^2_{Frob(B)},$$

where $\kappa_F(B^{-1/2}A\mathbf{S}) = \|B^{-1/2}A\mathbf{S}\|_F \|\mathbf{S}^{-1}A^{-1}B^{1/2}\|_F$.

Self-preconditioning Method: This suggests that $\mathbf{S} \approx A^{-1}$. But $X_k \approx A^{-1}$ so try $S = \text{sample columns of } X_k$.

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But $X_k \approx A^{-1}$ so try S = sample columns of X_k .

Initial experiments A positive definite

Newton-Schulz: $X_0 = A^T/(0.99||A^TA||_2)$, $X_{k+1} = 2X_k - X_kAX_k$. Self-preconditioning Method: B = A, $X_0 = I$, $X_{k+1} = \text{proj}_S + (I - \text{proj}_S A)X_k(I - A\text{proj}_S)$,

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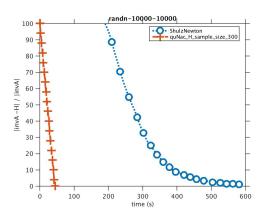


Figure: n = 10'000, with $nnz = 10^8$, A = randn(n, n), A = (A')*A; 32/35

Towards Randomized Preconditioning

```
Initialize X_0 \in \mathbb{R}^{n \times n} and x_0 \in \mathbb{R}^n.

While (stopping_criteria)

S_k = \text{sample\_function}(A, X_k)

x_{k+1} = \arg\min_{x \in \mathbb{R}^n} \|x - x_k\|_B^2 \text{ s.t } S_k^T A x = S_k^T b

X_{k+1} = \arg\min_{X \in \mathbb{R}^{n \times n}} \|X - X_k\|_{\text{Frobenius}(B)} \text{ s.t } S_k^T A X = S_k^T, X = X^T.

k = k+1
```

end

What if $(A_k)_k$ and (b_k) are slowly changing (like a Hessian matrix)?

Conclusion

- ► A natural framework for designing and analysing randomized iterative methods
- Analyse previous methods through one Theorem.
- New Gaussian methods, with potential on sparse problems
- New randomized matrix inversion methods.
- Paving a path towards randomized preconditioning.

Thank you for your attention!

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