

ECE 602: LUMPED LINEAR SYSTEMS

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Reachability of Discrete-Time (DT) Linear
Time-Invariant (LTI) Systems

Reachability of discrete-time (DT) linear time-invariant (LTI) systems

- **Objective:** Introduce notion of reachability of DT LTI systems modeled as

$$\mathbf{x}[k+1] = \mathbf{A}\mathbf{x}[k] + \mathbf{B}\mathbf{u}[k], \quad \mathbf{x}[0] = \mathbf{x}_0,$$

where $\mathbf{A} \in \mathbb{R}^{n \times n}$ and $\mathbf{B} \in \mathbb{R}^{n \times m}$

- First, obtain a solution of the system
- Note that

$$\mathbf{x}[1] = \mathbf{A}\mathbf{x}[0] + \mathbf{B}\mathbf{u}[0]$$

and

$$\begin{aligned}\mathbf{x}[2] &= \mathbf{A}\mathbf{x}[1] + \mathbf{B}\mathbf{u}[1] \\ &= \mathbf{A}(\mathbf{A}\mathbf{x}[0] + \mathbf{B}\mathbf{u}[0]) + \mathbf{B}\mathbf{u}[1] \\ &= \mathbf{A}^2\mathbf{x}[0] + \mathbf{A}\mathbf{B}\mathbf{u}[0] + \mathbf{B}\mathbf{u}[1]\end{aligned}$$

Solving DT LTI system modeling equation

- We have

$$\mathbf{x}[2] = \mathbf{A}^2 \mathbf{x}[0] + \mathbf{A} \mathbf{B} \mathbf{u}[0] + \mathbf{B} \mathbf{u}[1]$$

- Iterate to obtain

$$\begin{aligned} \mathbf{x}[3] &= \mathbf{A} \mathbf{x}[2] + \mathbf{B} \mathbf{u}[2] \\ &= \mathbf{A} (\mathbf{A}^2 \mathbf{x}[0] + \mathbf{A} \mathbf{B} \mathbf{u}[0] + \mathbf{B} \mathbf{u}[1]) + \mathbf{B} \mathbf{u}[2] \\ &= \mathbf{A}^3 \mathbf{x}[0] + \mathbf{A}^2 \mathbf{B} \mathbf{u}[0] + \mathbf{A} \mathbf{B} \mathbf{u}[1] + \mathbf{B} \mathbf{u}[2] \\ &= \mathbf{A}^3 \mathbf{x}[0] + \begin{bmatrix} \mathbf{B} & \mathbf{A} \mathbf{B} & \mathbf{A}^2 \mathbf{B} \end{bmatrix} \begin{bmatrix} \mathbf{u}[2] \\ \mathbf{u}[1] \\ \mathbf{u}[0] \end{bmatrix} \end{aligned}$$

Solving DT LTI system modeling equation—Contd

- We have

$$\mathbf{x}[3] = \mathbf{A}^3 \mathbf{x}[0] + \begin{bmatrix} \mathbf{B} & \mathbf{A}\mathbf{B} & \mathbf{A}^2\mathbf{B} \end{bmatrix} \begin{bmatrix} \mathbf{u}[2] \\ \mathbf{u}[1] \\ \mathbf{u}[0] \end{bmatrix}$$

- In general

$$\mathbf{x}[i] = \mathbf{A}^i \mathbf{x}[0] + \begin{bmatrix} \mathbf{B} & \mathbf{A}\mathbf{B} & \cdots & \mathbf{A}^{i-1}\mathbf{B} \end{bmatrix} \begin{bmatrix} \mathbf{u}[i-1] \\ \vdots \\ \mathbf{u}[1] \\ \mathbf{u}[0] \end{bmatrix}$$

Properties of $\begin{bmatrix} B & AB & \dots & A^{i-1}B \end{bmatrix}$

- Let

$$U_i = \begin{bmatrix} B & AB & \dots & A^{i-1}B \end{bmatrix}$$

- Note that

$$U_i \in \mathbb{R}^{n \times mi}$$

- If

$$\text{rank } U_i = \text{rank } U_{i+1}$$

then all the columns of $A^i B$ are linearly dependent on those of U_i

- This, in turn, implies that all the columns of $A^{i+1}B, A^{i+2}B, \dots$ must also be linearly dependent on those of U_i
- Hence,

$$\text{rank } U_i = \text{rank } U_{i+1} = \text{rank } U_{i+2} = \dots$$

Properties of $[B \ AB \ \dots \ A^{i-1}B]$ —Contd

- If

$$\text{rank } \mathbf{U}_I = \text{rank } \mathbf{U}_{I+1}$$

then

$$\text{rank } \mathbf{U}_I = \text{rank } \mathbf{U}_{I+1} = \text{rank } \mathbf{U}_{I+2} = \dots$$

- Indeed, if $\text{rank } \mathbf{U}_I = \text{rank } \mathbf{U}_{I+1}$, then all the columns of $\mathbf{A}'\mathbf{B}$ are linearly dependent on those of \mathbf{U}_I
- Hence,
$$\mathbf{A}^{I+1}\mathbf{B} = \mathbf{A}(\mathbf{A}'\mathbf{B}) = \mathbf{A} * (\text{linear combo of columns of } \mathbf{U}_I)$$
- But since $\text{rank } \mathbf{U}_I = \text{rank } \mathbf{U}_{I+1}$, every column of $\mathbf{A}^{I+1}\mathbf{B}$ must be a linear combo of the columns of \mathbf{U}_I
- Therefore, $\text{rank } \mathbf{U}_I = \text{rank } \mathbf{U}_{I+1} = \text{rank } \mathbf{U}_{I+2} = \dots$

More Analysis of U_i

- The rank of U_i increases by at least one when i is increased by one, until the maximum value of rank U_i is attained
- The maximum value of the rank of U_i is guaranteed to be achieved for $i = n$
- Indeed, by the Cayley-Hamilton theorem the matrix A satisfies its own characteristic equation
- Hence

$$A^n = -a_{n-1}A^{n-1} - \dots - a_1A - a_0I_n.$$

- Therefore

$$\text{rank} \begin{bmatrix} B & \dots & A^{n-1}B \end{bmatrix} = \text{rank} \begin{bmatrix} B & \dots & A^{n-1}B & A^nB \end{bmatrix}$$

The controllability matrix

- By the Cayley-Hamilton (C-H) theorem, the maximum value of the rank of U_i is guaranteed to be achieved for $i = n$
- Indeed, by C-H, $A^n = -a_{n-1}A^{n-1} - \dots - a_1A - a_0I_n$
- Hence

$$\begin{aligned}A^{n+1} &= A(A^n) = A(-a_{n-1}A^{n-1} - \dots - a_1A - a_0I_n) \\&= -a_{n-1}A^n - \dots - a_1A^2 - a_0A \\&= -a_{n-1}(-a_{n-1}A^{n-1} - \dots - a_0I_n) - \dots - a_0A\end{aligned}$$

- So we always have $\text{rank } U_n = \text{rank } U_{n+1} = \text{rank } U_{n+2} = \dots$
- The matrix

$$U_n = \begin{bmatrix} B & AB & \dots & A^{n-1}B \end{bmatrix} \in \mathbb{R}^{n \times mn}$$

is called the **controllability matrix**

Reachability

Definition

We say that the system $\mathbf{x}[k+1] = \mathbf{A}\mathbf{x}[k] + \mathbf{B}\mathbf{u}[k]$ is reachable if for any \mathbf{x}_f there exists a finite integer $q > 0$ and a control sequence, $\{\mathbf{u}[i] : i = 0, 1, \dots, q-1\}$, that transfers $\mathbf{x}[0] = \mathbf{0}$ to $\mathbf{x}[q] = \mathbf{x}_f$.

Theorem

The system $\mathbf{x}[k+1] = \mathbf{A}\mathbf{x}[k] + \mathbf{B}\mathbf{u}[k]$ is reachable if and only if

$$\text{rank} \begin{bmatrix} \mathbf{B} & \mathbf{AB} & \dots & \mathbf{A}^{n-1}\mathbf{B} \end{bmatrix} = n.$$

Necessary and Sufficient Condition for Reachability

- For $\mathbf{x}[0] = \mathbf{0}$ and $\mathbf{x}[q] = \mathbf{x}_f$,

$$\begin{aligned}\mathbf{x}_f &= \sum_{k=0}^{q-1} \mathbf{A}^{q-k-1} \mathbf{B} \mathbf{u}[k] \\ &= \begin{bmatrix} \mathbf{B} & \mathbf{AB} & \dots & \mathbf{A}^{q-1} \mathbf{B} \end{bmatrix} \begin{bmatrix} \mathbf{u}[q-1] \\ \vdots \\ \mathbf{u}[0] \end{bmatrix}\end{aligned}$$

- Thus for any \mathbf{x}_f there exists a control sequence $\{\mathbf{u}[i] : i = 0, 1, \dots, q-1\}$ that transfers $\mathbf{x}[0] = \mathbf{0}$ to $\mathbf{x}[q] = \mathbf{x}_f$ if and only if

$$\text{rank} \begin{bmatrix} \mathbf{B} & \mathbf{AB} & \dots & \mathbf{A}^{q-1} \mathbf{B} \end{bmatrix} = \text{rank } \mathbf{U}_q = n.$$

Example of Nonreachable DT System

- A DT system,

$$\begin{aligned}\mathbf{x}[k+1] &= \mathbf{A}\mathbf{x}[k] + \mathbf{b}u[k] \\ &= \begin{bmatrix} 0 & a \\ 0 & 0 \end{bmatrix} \mathbf{x}[k] + \begin{bmatrix} 1 \\ 0 \end{bmatrix} u[k], \quad a \in \mathbb{R}\end{aligned}$$

- The solution for $\mathbf{x}[0] = \mathbf{0}$ and arbitrary \mathbf{x}_f

$$\mathbf{x}_f = \begin{bmatrix} \mathbf{b} & \mathbf{A}\mathbf{b} \end{bmatrix} \begin{bmatrix} u[1] \\ u[0] \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} u[1] \\ u[0] \end{bmatrix}$$

- Take the state $\mathbf{x}_f = \begin{bmatrix} 0 & 1 \end{bmatrix}^\top$
- Then, this state is reachable \iff there are $u[0]$ and $u[1]$ such that

$$\begin{bmatrix} 0 \\ 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} u[1] \\ u[0] \end{bmatrix}$$

So Why Nonreachable?

- The state $\mathbf{x}_f = \begin{bmatrix} 0 & 1 \end{bmatrix}^\top$ is not reachable because there are no $u[0], u[1]$ such that

$$\begin{bmatrix} 0 \\ 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} u[1] \\ u[0] \end{bmatrix}$$

- Equivalently, the necessary and sufficient condition for reachability is NOT satisfied,

$$\text{rank} \begin{bmatrix} \mathbf{b} & \mathbf{A}\mathbf{b} \end{bmatrix} = \text{rank} \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} = 1 < 2.$$

- Note that nonreachable states are

$$\left\{ c \begin{bmatrix} 0 \\ 1 \end{bmatrix}, c \in \mathbb{R} \right\}$$