

### UMD DATA605 - Big Data Systems

## **Apache HBase**

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- Content in slides
  - Web
    - 2006, BigTable paper
    - https://hbase.apache.org/
    - https://github.com/apache/hbase
  - Good overview:
    - Seven Databases in Seven Weeks, 2e





# (Apache) HBase

- HBase = Hadoop DataBase
  - Supports large tables on commodity hardware clusters
  - Column-oriented DB
  - Part of Apache Hadoop ecosystem
  - Uses Hadoop filesystem (HDFS)
    - HDFS modeled after Google File System (GFS)
    - HBase based on Google BigTable
    - Google BigTable runs on GFS, HBase runs on HDFS
  - Used by Google, Airbnb, eBay
- When to use HBase
  - For large DBs (e.g., many 100 GBs or TBs)
  - With at least 5 nodes in production
- Applications
  - Large-scale online analytics
  - Heavy-duty logging
  - Search systems (e.g., Internet search)
  - Facebook Messages (based on Cassandra)
  - Twitter metrics monitoring





#### **HBase:** Features

- Data versioning
  - Store versions of data
- Data compression
  - Compress and decompress on-the-fly
- Garbage collection for expired data
- In-memory tables
- · Atomicity at row level
- Strong consistency guarantees
- Fault tolerant for machines and network
  - Write-ahead logging
    - Write data to in-memory log before disk
  - Distributed configuration
    - Nodes rely on each other, not centralized source



#### From HDFS to HBase

- Different types of workloads for DB backends
  - OLTP (On-Line Transactional Processing)
    - Read and write individual data items in a large table
    - E.g., update inventory and price as orders come in
  - OLAP (On-Line Analytical Processing)
    - Read large data amounts and process it
    - E.g., analyze item purchases over time
- Hadoop FileSystem (HDFS) supports OLAP workloads
  - Provide a filesystem with large files
  - Read data sequentially, end-to-end
  - Rarely updated
- HBase supports OLTP interactions
  - Built on HDFS
  - Use additional storage and memory to organize tables
  - Write tables back to HDFS as needed



#### **HBase Data Model**

- Warning: HBase uses names similar to relational DB concepts, but with different meanings
- A database consists of multiple tables
- Each table consists of multiple rows, sorted by row key
- Each row contains a row key and one or more column families
- Each column family
  - Contains multiple columns (family:column)
  - · Defined when the table is created
- A cell
  - Uniquely identified by (table, row, family:column)
  - Contains metadata (e.g., timestamp) and an uninterpreted array of bytes (blob)
- Versioning
  - New values don't overwrite old ones
  - put() and get() allow specifying a



SCIENCtimestamp (otherwise uses current time)

```
\# HBase Database: from table name to Table.

Database = Dict[str, Table]

\# HBase Table = {

# Row key

'rou': {

# (column family:column) → value

'cfi:col1': 'value1',

'cfi:col2': 'value2',

'cf2:col1': 'value2',

'cf2:col1': 'value3'

},

'rou2': {

... # More row data

}
}
database = {'table1': table}

\# Querwing data.
```

(value, metadata) = \
table['row1']['cf1:col1']

## **Example 1: Colors and Shape**

- Table with:
  - 2 column families: "color" and "shape"
  - 2 rows: "first" and "second"
- Row "first":
  - 3 columns in "color": "red", "blue", "yellow"
  - 1 column in "shape": shape = 4
- Row "second":
  - No columns in "color"
  - 2 columns in "shape"
  - Access data using row key and



column (family:qualifier)

	row keys	column family "color"	column family "shape"
LOM	"first"	"red": "#F00" "blue": "#00F" "yellow": "#FF0"	"square": "4"
10M	"second"		"triangle": "3" "square": "4"
\			



# Why All This Convoluted Stuff?

- A row in HBase is like a mini-database
  - A cell has many values
  - Data stored sparsely
- Rows in HBase are "deeper" than in relational DBs
  - Relational DBs: rows have many column values (fixed array with types)
  - HBase: rows like a two-level nested dictionary with metadata (e.g., timestamp)
- Applications
  - Store versioned website data
  - Store a wiki

	row keys	column family "color"	column family "shape"
10M	"first"	"red": "#F00" "blue": "#00F" "yellow": "#FF0"	"square": "4"
row	"second"		"triangle": "3" "square": "4"



## **Example 2: Storing a Wiki**

- Wiki (e.g., Wikipedia)
  - Contains pages
  - Each page has a title, article text varies over time
- HBase data model
  - Table name  $\rightarrow$  wikipedia
  - Row → entire wiki page
  - Row keys → wiki identifier (e.g., title or URL)
  - Column family  $\rightarrow$  text
  - Column  $\rightarrow$  " (empty)
  - Cell value  $\rightarrow$  article text

	row keys (wiki page titles)	column family "text"
(bage)	"first page's title"	"": "Text of first page"
(bage)	"second page's title"	"": "Text of second page"

```
wikipedia table = {
  # wiki i.d.
  'Home': {
    # Column family:column $\to$ value
    ':text': 'Welcome to the wiki!'.
  },
  'Welcome page': {
    ... # More row data
Database = Dict[str, Table]
database: Database = {'wikipedia':
wiki table}
(article, metadata) = \
wiki_table['Home']['text']
```

## **Example 2: Storing a Wiki**

#### Add data

- Columns don't need to be predefined when creating a table
- The column is defined as text

```
> put 'wikipedia', 'Home', 'text',
'Welcome!'
```

- Query data
- Specify the table name, the row key, and optionally a list of columns

```
> get 'wikipedia', 'Home', 'text'
text: timestamp=1295774833226,
value=Welcome!
```

 HBase returns the timestamp (ms since the epoch 01-01-1970 UTC)

```
row keys
(wiki page titles)

COMMEN "first page's title"

COMMEN "second page's title"

COMMEN "second page's title"

COMMEN "second page's title"

COMMEN "second page's title"

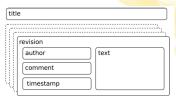
COMMEN "Text of second page"
```

```
wikipedia table = {
  # wiki i.d.
  'Home': {
    # Column family, column → value
    'text': 'Welcome to the wiki!',
  },
  'Welcome page': {
    ... # More row data
Database = Dict[str, Table]
database: Database = {'wikipedia':
wiki table}
(queried value, metadata) = \
wiki_table['Home']['text']
```



## **Example 2: Improved Wiki**

- Improved wiki using versioning
- A page
  - Uniquely identified by title
  - Can have multiple revisions
- A revision
  - Made by an author
  - · Optionally contains a commit comment
  - Identified by timestamp
  - Contains text
- HBase data model
- Add family column "revision" with multiple columns
  - E.g., author, comment,
- Timestamp automatically binds article text and metadata
- Title not part of revision
  - Fixed and uniquely identifies page (like a primary key)
  - To change title, re-write entire row



	keys (title)	family "text"	family "revision"
(bage)	"first page"	"": ""	"author": "" "comment": ""
(Dage)	"second page"		"author": "" "comment": ""



#### **Data in Tabular Form**

	Name		Home		Office	
Key	First	Last	Phone	Email	Phone	Email
101	Florian	Krepsbach	555-1212	florian@wobegon.org	666-1212	fk@phc.com
102	Marilyn	Tollerud	555-1213		666-1213	
103	Pastor	Inqvist			555-1214	inqvist@wel.org

#### Fundamental operations

- CREATE table, families
- PUT table, rowid, family:column, value
- PUT table, rowid, whole-row
- GET table, rowid
- SCAN table (WITH filters)
- DROP table



#### **Data in Tabular Form**

**CADEMY** 

```
| Name | | | Home | | Office | | Social |
                                                                     - | | Key | First | | Last | Phone | Email | Phone | Email | FacebookID
                                                                      — | 101 | Florian | Garfield | Krepsbach | 555-1212 |
florian@wobegon.org | 666-1212 | fk@phc.com | - | | 102 | Marilyn | - | Tollerud | 555-1213 | | 666-1213 | | - | | 103 | Pastor | - | Inqvist | | |
555-1214 | inqvist@wel.org | - | :::columns ::::{.column width=20%}
:::: ::::{.column width=20%}
New columns can be
added at runtime
:::: ::::{.column width=50%}
:::: ::::{.column width=20%}
Column families cannot
be added at runtime
Table People(Name, Home, Office)
    101: {
        Timestamp: T403:
        Name: {First="Florian", Middle="Garfield", Last="Krepsbach"},
        Home: {Phone="555-1212", Email="florian@wobegon.org"},
        Office: {Phone="666-1212", Email="fk@phc.com"}
    ٦.
    102: {
        Timestamp: T593;
        Name: {First="Marilyn", Last="Tollerud"},
        Home: {Phone="555-1213"},
    CIEdNGE {Phone="666-1213"}
```

### **Nested Data Representation**

	Name		Home		Office	
Key	First	Last	Phone	Email	Phone	Email
101 102	Florian Marilyn	Krepsbach Tollerud	555-1212 555-1213	florian@wobegon.org	666-1212 666-1213	fk@phc.com
103	Pastor	Inqvist	555 1215		555-1214	inqvist@wel.org

```
GET People:101
{
    Timestamp: T403;
    Name: {First="Florian", Last="Krepsbach"},
    Home: {Pone="555-1212", Email="florian@wobegon.org"},
    Office: {Phone="666-1212", Email="fk@phc.com"}}
}
GET People:101:Name
{First="Florian", Last="Krepsbach"}
GET People:101:Name:First
"Florian"
```



# Column Family vs Column

- Adding a column
  - Cheap
  - Done at run-time
- Adding a column family
  - Not at run-time
  - Requires table copy (expensive)
  - Indicates data storage method
    - Easy to add: map
    - · Hard to add: static array
    - E.g., mongoDB document vs Relational DB column
- Why differentiate column families vs columns?
  - Why not store all row data in one column family?
  - Each column family configured independently, e.g.,
    - Compression
    - Performance tuning
    - Stored together in files
  - Designed for specific data types
    - E.g., timestamped web data for search engine



## **Consistency Model**

#### Atomicity

- Update entire rows atomically or not at all
- Independent of column count

#### Consistency

- GET returns a complete row from the table's history
  - Weak/eventual consistency
  - · Check timestamp for certainty
- SCAN
  - Includes all data written before scan
  - May include updates since start

#### Isolation

- Concurrent vs sequential semantics
- Not guaranteed beyond a single row
- Row is the atom of information
- Durability
  - Successful writes are durable on disk



# **Checking for Row or Column Existence**

- HBase uses Bloom filters to check row or column existence
  - Acts like a cache for keys
  - Track presence without querying
- Hashset complexity
  - Unbounded space for data storage
  - No false positives
  - O(1) average/amortized time
- Bloom filter implementation
  - Probabilistic hash set
  - Array of bits initially set to 0
  - Hash new data, set bits to 1
  - To test data presence, compute hash, check bits
    - All bits 0: data not seen
    - All bits 1: likely seen (false positive possible)
- Bloom filter complexity
  - Constant space usage
  - False positives possible (no false negatives)
  - O(1) time complexity



# Write-Ahead Log (WAL)

- Technique used by DBs
  - Provide atomicity and durability
  - Protect against node failures
  - Equivalent to journaling in file systems
- HBase and Postgres use WAL
- WAL mechanics
- Updated state of tables:
  - Not written to disk immediately
  - Buffered in memory
  - Written to disk as checkpoints periodically
- Problem
  - Server crash during this period loses state
- Solution
- Use append-only disk-resident data structure SCIENCE Log operations since last checkpoint in WAL ACADEM Clear WAL when tables stored to disk

## **Storing Variable-Length Data in Dbs**

#### **SQL** Table

People(ID: Integer, FirstName: CHAR[20], LastName: CHAR[20], Phone: CHAR[8])
UPDATE People SET Phone="555-3434" WHERE ID=403;

ID	FirstName	LastName	Phone
101	Florian	Krepsbach	555-3434
102	Marilyn	Tollerud	555-1213
103	Pastor	Ingvist	555-1214

- Each row: 52 bytes
- Move to next row: fseek(file,+52)
- Get to Row 401: fseek(file, 401\*52)
- Overwrite data in place

#### **HBase Table**

Need to use pointers

## **HBase Implementation**

- How to store the web on disk?
- HBase is backed by HDFS
  - Store each table (e.g., Wikipedia) in one file
  - "One file" means one gigantic file stored in HDFS
    - HDFS splits/replicates file into blocks on different servers
- Idea in several steps:
  - Idea 1: Put entire table in one file
    - Overwrite file with any cell change
    - Too slow
  - Idea 2: One file + WAL
    - Better, but doesn't scale to large data
  - Idea 3: One file per column family + WAL
    - · Getting better!
  - Idea 4: Partition table into regions by key
    - Region = chunk of rows [a, b)
    - Regions never overlap



# Idea 1: Put the Table in a Single File

- How do we do the following operations?
  - CREATE, DELETE (easy / fast)
  - SCAN (easy / fast)
  - GET, PUT (difficult / slow)

```
Table People(Name, Home, Office) { 101: { Timestamp: T403; Name: {First="Florian", Middle="Garfield", Last="Krepsbach"}, Home: {Phone="555-1212", Email="florian@wobegon.org"}, Office: {Phone="666-1212", Email="fk@phc.com"} }, 102: { Timestamp: T593; Name: {First="Marilyn", Last="Tollerud"}, Home: {Phone="555-1213"}, Office: {Phone="666-1213"} }, \dots }
```

File "People"



#### Idea 2: One File + WAL

#### Table People(Name, Home, Office)

PUT 101:Office:Phone = "555-3434" PUT 102:Home:Email = mt@yahoo.com

#### **WAL** for Table People

- Changes are applied only to the log file
- The resulting record is cached in memory
- · Reads must consult both memory and disk

#### **Memory Cache for Table People**

101

102

**GET People:101** 

**GET People:103** 

PUT People:101:Office:Phone = "555-3434"



# Idea 2 Requires Periodic Table Update

```
101: {Timestamp: T403;Name: {First="Florian", Middle="Garfield",
Last="Krepsbach"},Home: {Phone="555-1212",
Email="florian@wobegon.org"},Office: {Phone="666-1212",
Email="fk@phc.com"}}, 102: {Timestamp: T593;Name: { First="Marilyn",
Last="Tollerud"}, Home: { Phone="555-1213" }, Office: { Phone="666-1213"
}}, . . .
Table for People on Disk (Old)
PUT 101:Office:Phone = "555-3434" PUT 102:Home:Email = mt@yahoo.com
WAL for Table People:
101: {Timestamp: T403;Name: {First="Florian", Middle="Garfield",
Last="Krepsbach"},Home: {Phone="555-1212",
Email="florian@wobegon.org"},Office: {Phone="555-3434",
```

Email="fk@phc.com"}},102: {Timestamp: T593;Name: { First="Marilyn", Last="Tollerud"},Home: { Phone="555-1213", Email="my@yahoo.com" },

SCIENCE TADAE HORY People on Disk (New)

# Idea 3: Partition by Column Family

Data for Column Family Name

#### Tables for People on Disk (Old)

PUT 101:Office:Phone = "555-3434" PUT 102:Home:Email = mt@yahoo.com . . .

#### WAL for Table People

#### Tables for People on Disk (New)

- Write out a new copy of the table, with all of the changes applied
- Delete the log and memory cache
- Start over

Data for Column Family Home

Data for Column Family Office

Data for Column Family Home (Changed)

Data for Column Family Office (Changed)



# Idea 4: Split Into Regions

Region 1: Keys 100-200

Region 2: Keys 100-200

Region 3: Keys 100-200

Region 4: Keys 100-200

Region Server

Region Master

Region Server

Region Server

Region Server

Transaction Log

Memory Cache

Table



# Final HBase Data Layout

