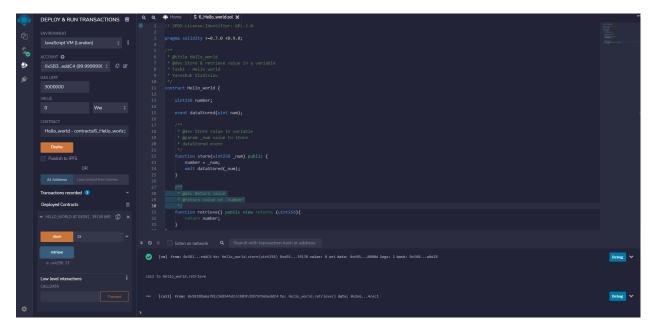
## Question 1:

1. Program a super simple "Hello World" smart contract: store an unsigned integer and then retrieve it. Please clearly comment your code. Once completed, deploy the smart contract on Remix. Include the .sol file and a screenshot of the Remix UI once deployed in your final submission pdf (more info about submission formatting below).

Code:

```
SPDX-License-Identifier: GPL-3.0
pragma solidity >=0.7.0 <0.9.0;</pre>
 * @dev Store & retrieve value in a variable
 * Task1 - Hello world
 * Yaroshuk Vladislav
contract Hello_world {
    uint256 number;
    event dataStored(uint num);
     * dataStored event
    function store(uint256 _num) public {
        number = _num;
        emit dataStored(_num);
    * @return value of 'number'
    function retrieve() public view returns (uint256){
        return number;
```



2. On the documentation page, the "Ballot" contract demonstrates a lot of features on Solidity. Read through the script and try to understand what each line of code is doing, then implement the Possible Improvements by reducing the number of transactions in the "giveRightToVote" function while maintaining the same functionality of the program.

Code:

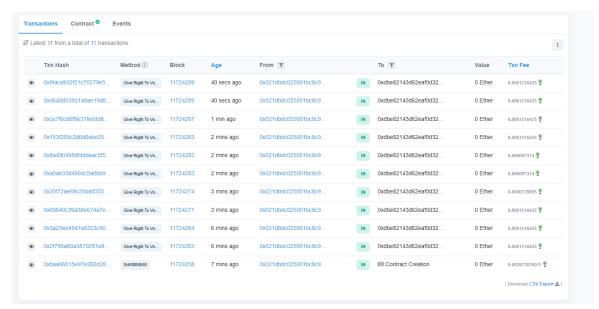
```
// SPDX-License-Identifier: GPL-3.0
pragma solidity >=0.7.0 <0.9.0;</pre>
// Task2 - Ballot upgrade
// Yaroshuk Vladislav
contract Ballot {
    // This declares a new complex type which will
    struct Voter {
        uint weight; // weight is accumulated by delegation
        bool voted; // if true, that person already voted
        address delegate; // person delegated to
        uint vote; // index of the voted proposal
    // This is a type for a single proposal.
    struct Proposal {
        bytes32 name; // short name (up to 32 bytes)
        uint voteCount; // number of accumulated votes
    address public chairperson;
    // This declares a state variable that
    mapping(address => Voter) public voters;
```

```
Proposal[] public proposals;
    constructor(bytes32[] memory proposalNames) {
        chairperson = msg.sender;
        voters[chairperson].weight = 1;
        // For each of the provided proposal names,
        for (uint i = 0; i < proposalNames.length; i++) {</pre>
            proposals.push(Proposal({
                name: proposalNames[i],
                voteCount: 0
            }));
    // Instead of taking a single address parameter in the giveRightToVote
function, take an array of addresses.
    function giveRightToVote(address[] calldata _voters) external {
       // If the first argument of `require` evaluates
        // are reverted.
        // As a second argument, you can also provide an
        // explanation about what went wrong.
           msg.sender == chairperson,
            "Only chairperson can give right to vote."
        );
each address.
        for (uint i = 0; i < _voters.length; i++) {</pre>
            require(
                !voters[ voters[i]].voted,
                "Voter already voted"
            );
            // This require cheks every voter if it has an allowance to vote
            require(voters[ voters[i]].weight == 0);
```

```
voters[_voters[i]].weight = 1;
function delegate(address to) external {
   Voter storage sender = voters[msg.sender];
   require(!sender.voted, "You already voted.");
    require(to != msg.sender, "Self-delegation is disallowed.");
   // Forward the delegation as long as
   // `to` also delegated.
   // In general, such loops are very dangerous,
   // because if they run too long, they might
   while (voters[to].delegate != address(0)) {
        to = voters[to].delegate;
        require(to != msg.sender, "Found loop in delegation.");
    sender.voted = true;
    sender.delegate = to;
    Voter storage delegate_ = voters[to];
    if (delegate_.voted) {
        proposals[delegate .vote].voteCount += sender.weight;
    } else {
        delegate_.weight += sender.weight;
function vote(uint proposal) external {
    Voter storage sender = voters[msg.sender];
    require(sender.weight != 0, "Has no right to vote");
    require(!sender.voted, "Already voted.");
    sender.voted = true;
    sender.vote = proposal;
```

```
proposals[proposal].voteCount += sender.weight;
function winningProposal() public view
        returns (uint winningProposal_)
    uint winningVoteCount = 0;
    for (uint p = 0; p < proposals.length; p++) {</pre>
        if (proposals[p].voteCount > winningVoteCount) {
            winningVoteCount = proposals[p].voteCount;
            winningProposal_ = p;
    }
// of the winner contained in the proposals array and then
function winnerName() external view
        returns (bytes32 winnerName_)
    winnerName_ = proposals[winningProposal()].name;
```

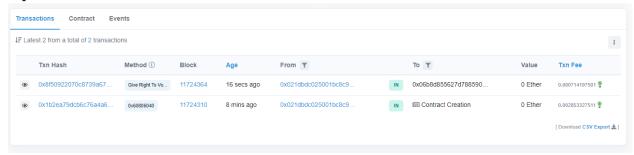
3. Deploy your script on Remix and compare the difference in gas fees between the original script and the improved script when giving 10 voters the right to vote. Once completed, submit (1) your improved version of the contract as an .sol file with comments describing the changes you made, and (2) screenshots (before and after) of the gas fees for the transaction(s) to give 10 voters the right to vote. Giving 10 voters right to vote:



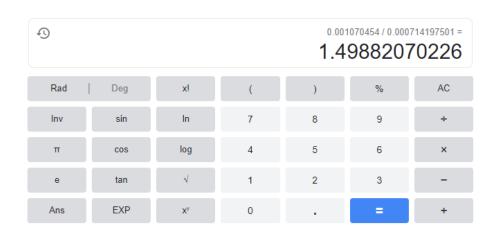
### Fees in Ropsten:



#### Updated contract:



#### Difference



## Q2:

1. Briefly describe signals, templates, components, constraints, and why we use Circom for ZK.

The arithmetic circuits built using circom operate on **signals**, which contain field elements in Z/pZ. **Signals** can be named with an identifier or can be stored in arrays and declared using the keyword signal. **Signals** can be defined as input or output, and are considered intermediate signals otherwise.

The mechanism to create generic circuits in Circom is the so-called **templates**. They are normally parametric on some values that must be instantiated when the **template** is used. The instantiation of a **template** is a new circuit object, which can be used to compose other circuits, so as part of larger circuits.

Circom allows programmers to define the **constraints** that define the arithmetic circuit. The equations that describe the circuit are called **constraints**.

With circom, you design your own circuits with your own constraints, and the compiler outputs the R1CS representation that you will need for your zero-knowledge proof. The nice thing about circuits, is that although most zero-knowledge protocols have an inherent complexity that can be overwhelming for many developers, the design of arithmetic circuits is clear and neat.

2. For the following truth table, what is the constraint system? Hint: your answer should be a simple equation relating x1 and x2 to y.

x1	x2	у
0	0	0
0	1	0
1	0	0
1	1	1

Answer:  $y = x_1 \cdot x_2$ 

- 3. Using open source resources like circomlib, **define**, **compile**, **and prove** a circuit that checks if a **number is less than 10**.
  - 1) Submit a screenshot of the terminal outputs (to demonstrate your code effectively runs) along with (a well-commented) .circom file. (You will just copy and paste the contents of this file into your final pdf.)
  - 2) Concisely describe what each command in the compilation and proving steps are doing under the hood. If you submit a screenshot of terminal outputs, you could write this in a separate file. If you write a bash file you can add comments above each line and then submit the bash file.

3)

## Writing circuits

```
pragma circom 2.0.0;
include "./circim/circomlib/circuits/comparators.circom";
// We will be using LessThan function from circomlib
/* Logic:
  It's convinient to take a = 15, it can be represented as [1 \ 1 \ 1].
  The LessThan circuit attempts to build the bit representation of the minimum
integer that requires
  one (1) additional bit, which is 16 (represented with 5 bits: 1 0 0 0 0)
  The LessThan circuit will then add the result of (a - 10) to the 5-bit number.
  If a - 10 < 0, then the most significant bit of 5-bit number will become 0.
  If a - 10 > 0, then the most significant bit of 5-bit number will remain 1.
  Signal output = 1 will give either 0 (meaning a \geq= 10) or 1 (meaning a \leq 10).
*/
// N is the number of bits the input have.
template ex2(n) {
  // Declaration of signals.
  signal input a;
  signal output out;
  var x = 10;
  component lessThan = LessThan(n);
  // Constraints.
  lessThan.in[0] < --a;
  lessThan.in[1] \leq-- x;
  out <== lessThan.out;
// I will assign n to be 32-bit, it should be enough.
component main = ex2(32);
```

## Compiling our circuit

After we write our arithmetic circuit using **circom**, we should save it in a file with the **.circom** extension.

We create a file called **comparison.circom.** Now is time to compile the circuit to get a system of arithmetic equations representing it. As a result of the compilation

we will also obtain programs to compute the witness. We can compile the circuit with the following command:

\$ circom comparison.circom --r1cs --wasm --sym --c

```
ubuntu@ip-172-31-9-53:-$ circom comparison.circom --rlcs --wasm --sym --c
template instances: 3
non-linear constraints: 33
linear constraints: 0
public inputs: 0
public outputs: 1
private inputs: 1
private outputs: 0
wires: 36
labels: 40
Written successfully: ./comparison.rlcs
Written successfully: ./comparison.sym
Written successfully: ./comparison.cpp and ./comparison.cpp/comparison.dat
Written successfully: ./comparison.cpp/main.cpp, circom.hpp, calcwit.hpp, calcwit.cpp, fr.hpp, fr.cpp, fr.asm and Makefile
Written successfully: ./comparison_js/comparison.wasm
Everything went okay, circom safe
ubuntu@ip-172-31-9-53:-$
```

# Computing our witness

Before creating the proof, we need to calculate all the signals of the circuit that match all the constraints of the circuit. For that, we will use the **Wasm** module generated by **circom** that helps to do this job.

We need to create a file named **input.json** containing the inputs written in the standard **json** format.

```
ubuntu@ip-172-31-9-53:~$ sudo nano input.json
ubuntu@ip-172-31-9-53:~$

GNU nano 4.8
{"a": 15}
```

# Computing the witness with WebAssembly

Enter in the directory **comparison\_js**, add the input in a file **input.json** and execute:

\$ node generate\_witness.js comparison.wasm input.json witness.wtns

```
ubuntu@ip-172-31-9-53:~/comparison_js$ node generate_witness.js comparison.wasm input.json witness.wtns ubuntu@ip-172-31-9-53:~/comparison_js$
```

It will generate the same **witness.wtns** file. This file is encoded in a binary format compatible with **snarkjs**, which is the tool that we use to create the actual proofs.

## Proving circuits

After compiling the circuit and running the witness calculator with an appropriate input, we will have a file with extension .wtns that contains all the computed signals and, a file with extension .r1cs that contains the constraints describing the circuit. Both files will be used to create our proof.

```
ubuntu@ip-172-31-9-53:~/comparison_js$ ls comparison.wasm generate_witness.js input.json witness.wtns witness_calculator.js ubuntu@ip-172-31-9-53:~/comparison_js$
```

Now, we will use the **snarkjs** tool to generate and validate a proof for our input.

We are going to use the **Groth16 zk-SNARK protocol**. To use this protocol, you will need to generate a trusted setup. **Groth16** requires a per circuit trusted setup. In more detail, the trusted setup consists of 2 parts:

- The powers of tau, which is independent of the circuit.
- The phase 2, which depends on the circuit.

Next, we provide a very basic ceremony for creating the trusted setup and we also provide the basic commands to create and verify **Groth16** proofs..

### Powers of Tau

First, we start a new "powers of tau" ceremony:

\$ snarkjs powersoftau new bn128 12 pot12\_0000.ptau -v

```
ubuntu@ip-172-31-9-53:~/comparison_js$ snarkjs powersoftau new bn128 12 pot12_0000.ptau -v
[DEBUG] snarkJS: Calculating First Challenge Hash
[DEBUG] snarkJS: Calculate Initial Hash: tauG1
[DEBUG] snarkJS: Calculate Initial Hash: tauG2
[DEBUG] snarkJS: Calculate Initial Hash: alphaTauG1
[DEBUG] snarkJS: Calculate Initial Hash: betaTauG1
[DEBUG] snarkJS: Blank Contribution Hash:
               786a02f7 42015903 c6c6fd85 2552d272
               912f4740 e1584761 8a86e217 f71f5419
               d25e1031 afee5853 13896444 934eb04b
               903a685b 1448b755 d56f701a fe9be2ce
[INFO] snarkJS: First Contribution Hash:
               9e63a5f6 2b96538d aaed2372 481920d1
               a40b9195 9ea38ef9 f5f6a303 3b886516
               0710d067 c09d0961 5f928ea5 17bcdf49
               ad75abd2 c8340b40 0e3b18e9 68b4ffef
```

Then, we contribute to the ceremony:

\$ snarkjs powersoftau contribute pot12\_0000.ptau pot12\_0001.ptau --name="First contribution" -v

Now, we have the contributions to the powers of tau in the file pot12\_0001.ptau and we can proceed with the Phase 2.

### Phase 2

The phase 2 is circuit-specific. Execute the following command to start the generation of this phase:

snarkjs powersoftau prepare phase2 pot12\_0001.ptau pot12\_final.ptau -v

```
ubuntu@ip-172-31-9-53:~$ snarkjs powersoftau prepare phase2 pot12_0001.ptau pot12_final.ptau -v
[DEBUG] snarkJS: Starting section: tauG1
[DEBUG] snarkJS: tauG1: fft 0 mix start: 0/1
[DEBUG] snarkJS: tauG1: fft 0 mix end: 0/1
[DEBUG] snarkJS: tauG1: fft 1 mix start: 0/1
[DEBUG] snarkJS: tauG1: fft 1 mix end: 0/1
[DEBUG] snarkJS: tauG1: fft 2 mix start: 0/1
[DEBUG] snarkJS: tauG1: fft 2 mix end: 0/1
[DEBUG] snarkJS: tauG1: fft 3 mix start: 0/1
[DEBUG] snarkJS: tauG1: fft 3 mix end: 0/1
[DEBUG] snarkJS: tauG1: fft 4 mix start: 0/1
[DEBUG] snarkJS: tauG1: fft 4 mix end: 0/1
[DEBUG] snarkJS: tauG1: fft 5 mix start: 0/1
[DEBUG] snarkJS: tauG1: fft 5 mix end: 0/1
[DEBUG] snarkJS: tauG1: fft 6 mix start: 0/1
[DEBUG] snarkJS: tauG1: fft 6 mix end: 0/1
[DEBUG] snarkJS: tauG1: fft 7 mix start: 0/1
[DEBUG] snarkJS: tauG1: fft 7 mix end: 0/1
[DEBUG] snarkJS: tauG1: fft 8 mix start: 0/1
[DEBUG] snarkJS: tauG1: fft 8 mix end: 0/1
[DEBUG] snarkJS: tauG1: fft 9 mix start: 0/1
[DEBUG] snarkJS: tauG1: fft 9 mix end: 0/1
[DEBUG] snarkJS: tauG1: fft 10 mix start: 0/1
[DEBUG] snarkJS: tauG1: fft 10 mix end: 0/1
[DEBUG] snarkJS: tauG1: fft 11 mix start: 0/1
[DEBUG] snarkJS: tauG1: fft 11 mix end: 0/1
[DEBUG] snarkJS: tauG1: fft 12 mix start: 0/1
[DEBUG] snarkJS: tauG1: fft 12 mix end: 0/1
```

Next, we generate a **.zkey** file that will contain the proving and verification keys together with all phase 2 contributions. Execute the following command to start a new zkey:

\$ snarkjs groth16 setup comparison.r1cs pot12\_final.ptau comparison\_0000.zkey

Contribute to the phase 2 of the ceremony:

\$ snarkjs zkey contribute comparison\_0000.zkey comparison\_0001.zkey -- name="1st Contributor Name" -v

Export the verification key:

\$ snarkjs zkey export verificationkey comparison 0001.zkey verification key.json

```
ubuntu@ip-172-31-9-53:~$ snarkjs zkey export verificationkey comparison_0001.zkey verification_key.json ubuntu@ip-172-31-9-53:~$ ■
```

## Generating a Proof

\$ snarkjs groth16 prove comparison\_0001.zkey witness.wtns proof.json public.json

```
ubuntu@ip-172-31-9-53:-$ snarkjs groth16 prove comparison_0001.zkey comparison_js/witness.wtns proof.json public.json
ubuntu@ip-172-31-9-53:-$ ls
circom comparison.rlcs comparison_0000.zkey comparison_comparison pot12_0001.ptau proof.json verification_key.json
comparison.circom comparison.sym comparison_0001.zkey comparison_js pot12_0000.ptau pot12_final.ptau public.json
```

This command generates a Groth16 proof and outputs two files:

- proof.json: it contains the proof.
- public.json: it contains the values of the public inputs and outputs.

Verifying a Proof

To verify the proof, execute the following command:

\$ snarkjs groth16 verify verification\_key.json public.json proof.json

The command uses the files **verification\_key.json** we exported earlier, **proof.json** and **public.json** to check if the proof is valid. If the proof is valid, the command outputs an OK.

```
ubuntu@ip-172-31-9-53:~$ snarkjs groth16 verify verification_key.json public.json proof.json [INFO] snarkJS: OK! ubuntu@ip-172-31-9-53:~$ ■
```

We were checking if 15 < 10, result is **False**.

\$ nano public.json

```
GNU nano 4.8
[
"""
]
```

## Verifying from a Smart Contract

It is also possible to generate a Solidity verifier that allows verifying proofs on Ethereum blockchain.

First, we need to generate the Solidity code using the command:

\$ snarkjs zkey export solidityverifier comparison\_0001.zkey verifier.sol

#### **Output:**

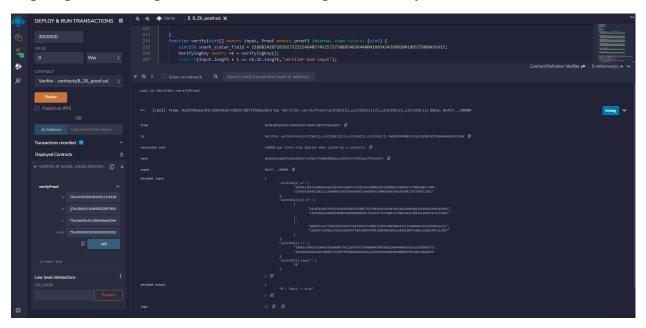
["0x247fe335cb9165311c6158374c8d56e932a92e05c1a9f92ea0fb23e4de49afa4", "0x0e098b69511178f32b43ef04851ef41be09a643db144232e5fa8e22299fcbf93"],["0x2861b740d005029678932be7158563a5eede0fd9ddefece4742536b295514561"

"0x1dbf0ed262589aa148c67e9a43e446527828881352c4cd673e5f3a64b3fc9dbb"], ["0x1406bba9c45c1b4c6067f4847db86f050b060e76167e5b0c0fb4d4b194af32a2", "0x195b1839c2bd9cdb271fc86ec212c93774898ddd6e501b58a38144bfe2b67701"] ], ["0x2be93b44136b946ab526e4c6475e175bdeccfcaa66593e00d46c32e1deafefe5"]

<sup>&</sup>quot;0x21980e7f82a5c4ae707663cda31ae9dca0f2adc8df1546c0dd64a278bf6b1621"], [

Cut and paste the output of the command to the parameters field of the **verifyProof** method in Remix. If everything works fine, this method should return TRUE. You can see contract code in my github.

https://github.com/grGred/Zero-Knowledge-University



# Question 3: Thinking with ZK

1. Describe an application you believe would benefit from ZK. Explain why ZK can benefit the application. Then formulate the use of ZK in terms of a statement and witness. Be as precise as possible in your use of variables to represent which data should be public and private, as well as the exact relations between them.

It would great to see zk-rollups in Uniswap for example. Right now there are millions of arbitrage bots and frontrunners that can easily spoil your mood. In order to prevent such types of attacks, we can have public function (swap), private input (tx calldata), public output (balance change).

When we making swap, we will receive private and public input, when smart contract will want to know our balance (and to make further transaction), we will give him a prove (private input) that we really own this funds, before that swap everyone will now that you possibly own some amount of this token.

2. List the key difference between an interactive and a non-interactive ZK proof. What type has more relevance to crypto?

SNARKs are short for succinct non-interactive arguments of knowledge. In this general setting of so-called interactive protocols, there is a prover and a verifier and the prover wants to convince the verifier about a statement (e.g. that f(x) = y) by exchanging messages. The generally desired properties are that no prover can convince the verifier about a wrong statement (soundness) and there is a certain strategy for the prover to convince the verifier about any true statement (completeness).

Non-interactive: there is no or only little interaction. For zkSNARKs, there is usually a setup phase and after that a single message from the prover to the verifier. Furthermore, SNARKs often have the so-called "public verifier" property meaning that anyone can verify without interacting anew, which is important for blockchains.

Non-interactive ZK proof is more secure and decentralized, less vulnerable to the economic attacks.

3. Beyond privacy, ZK helps with scalability. In a paragraph or less, describe what ZK-rollups are, how they work, and why people are excited about them.

Rollups perform transaction execution outside layer 1 and then the data is posted to layer 1 where consensus is reached. As transaction data is included in layer 1 blocks, this allows rollups to be secured by native Ethereum security. The ZK-rollup smart contract maintains the state of all transfers on layer 2, and this state can only be updated with a validity proof. This means that ZK-rollups only need the validity proof instead of all transaction data. With a ZK-rollup, validating a block is quicker and cheaper because less data is included. With a ZK-rollup, there are no delays when moving funds from layer 2 to layer 1 because a validity proof accepted by the ZK-rollup contract has already verified the funds. Being on layer 2, ZK-rollups can be optimised to reduce transaction size further. For instance, an account is represented by an index rather than an address, which reduces a transaction from 32 bytes to just 4 bytes. Transactions are also written to Ethereum as *calldata*, reducing gas.