	Groce Miguel HW2
	I pledge my honor that I've abided by the Stevens
	Honor system.
	1942
4	a) Mystery(n) n=4
	S-0 (1-n1xP= L(1-n)x6 16 = Cn)x
	S=1 This algorithm computes
	. S= 1+ 4=5 the sym of squares up to
	S=15+9=14 and including n.
	5= 14+16= 30
	b) Basic operation is setting S = itself + i2
	$(2-0) \in \mathcal{F}(n) \times \mathbb{R}$
	c) Operation is executed in times
	Hank?
	d) O(n)
	Jane 1
	e) [D(n+)(2n+1)] new algorithm w/ time complexity.
	6 · of O(1)
	$((l-n)-n)$ $\in (n)_{\times}$
1)	a) x(n)=x(n-1)+5 for n>1 x(1)=0
	Step 1 Step 4
	$x(n-1)^{2}, x(n-2) + 5$ SICF-1 $n-i=0$
	x(n) = x(n-2) + 10
174 (1)	1-1)+ (1+1+1)+ (+1)x 1-1+ (1-1)x = (+1)x
	Step 2 Step 5
	x(n-2)=x(n-3)+5 $x(n-(n+1))+5(n-1)$
	x(n) = x(n-3) + 15
	5n-5
	Step 3
	x(n-i+1) + 5i

b) 
$$x(n) = 3x(n-1)$$
 for  $n>1$   $x(1)=4$ 
 $x(n-1) = 3x(n-2)$ 
 $x(n) = 3[3x(n-2)] = 9x(n-2)$ 
 $x(n) = 9[3x(n-3)] = 27x(n-3)$ 
 $x(n) = 9[3x(n-3)] = 27x(n-3)$ 
 $x(n) = 3[(n-i)]$ 
 $x(n) = 3$ 

d) x(n)=x(2) to for no! x(1)=1 solve for n=2K Step 1 = x(2k) + 2k -> x(2k-1) + 2k  $\frac{5+cp\ 2}{1\times(2^{k-1})} = \times(2^{k-2}) + 2^{k-1}$   $\times(2^{k})^{2} \times(2^{k-2}) + 2^{k-1} + 2^{k}$ x(2k-2) = x(2k-3) + 2k-2 x(2k) = x(2k-3) + 2k-2 + 2k-1 + 2k Step 3 General form: X(2k-i) + 2k-i + 2k-i+1 + 2k-i+2 ... 2k Stop4 IC = x(1)=1 2 = 1 2 = 20 K-1=0 [i=K] Step5 x(20 K) + 2 K-K+1 K-K+1 K-K+2 K [x(n) = 2n = 1)

x(n)=x(n/3)+1 n>1 x(1)=1 solve for n=3k x(3K) = x(3K-1)+1 Step 1 = x(3k-2)+1" x(3k) = x(3k-2)+2 Step2 = (3K-3) +1 x(3K-i) +i -Step 4

IC = K-3 = 1 3 = 3 Step5 X(3KK)+K. n=312 log3n=K S(n-1) +2 b) This has the same number of multiplications as the S(n) = S(n-2)+4 recursive but the nonrecusive S(n-2) -S(n-3) +2 will take less storage and SCn)= SCn-3)+6 +me. s(n) = S(n-i) + 2i s k-1=1 i=n-1 S(n-(n-1)) +2(n-1) (x(n)=2(n-1))