

Terminating sequences of Bunny Trainer Transforms

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In this note, we prove a necessary and sufficient condition for the termination of Bunny Trainer Transform sequences.

Definition. A *Bunny Trainer Transform* (*BTT*) is the following function on pairs of nonnegative integers (m, n) :

$$BTT(m, n) = \begin{cases} BTT(n, m) & m > n \\ (2m, n - m) & 3m \leq n \\ (n - m, 2m) & m \leq n < 3m \end{cases}$$

BTT represents the outcome of a match where the bunny trainer corresponding to the smaller integer (e.g. m in (m, n)) bets the entire position value m and wins. BTT is defined in such a way that if $(x, y) = BTT(m, n)$ then $x \leq y$.

Definition. A BTT sequence *terminates* or is *terminating* when $BTT^k(m, n) = (0, p)$ for some nonnegative integers k and p . If no such k and p exist, then the BTT sequence is *nonterminating*.

That is to say, the BTT sequence terminates when successive matches lead to one of the bunny trainers losing everything. We have $BTT(0, p) = (0, p)$, so the BTT sequence terminates if it eventually reaches a fixed point where the first integer in the pair is zero.

Definition. A pair of nonnegative integers (m, n) satisfies the *terminating condition with exponent k* (*TC- k*), if for nonnegative integer k we have

$$m + n = 2^k \gcd(m, n).$$

(We take $\gcd(0, 0) = 0$ by definition.) More generally the pair (m, n) is said to satisfy the *terminating condition* (*TC*) if for some nonnegative integer k it satisfies TC- k .

Proposition 1. *If (m, n) satisfies TC- k for some positive integer k , then $BTT(m, n)$ satisfies TC- $(k - 1)$.*

Proof. $(0, 0)$ satisfies TC- k for all nonnegative integers k , and $BTT(0, 0) = (0, 0)$, so the proposition is true trivially in this case.

$(0, n)$ for a positive integer n only satisfies TC-0 and hence does not meet the condition of the proposition.

So without loss of generality, we can assume that $0 < m \leq n$. Let $d = \gcd(m, n)$, and let $d' = \gcd(2m, n - m)$.

We claim $d' = 2d$. Certainly $2d \mid 2m$ since $d \mid m$, and since $d \mid n$ also then $d \mid (n - m)$. But

$$\frac{m}{d} + \frac{n}{d} = 2^k$$

is even as $k > 0$, and so both integers $m/d, n/d$ are odd. (They can't both be even integers because then d would not be the *greatest* common divisor of m and n .) It follows that

$$\frac{n - m}{d} = \frac{n}{d} - \frac{m}{d}$$

is an even integer, and so $2d \mid (n - m)$. Thus $2d \mid d'$, being a divisor both of $2m$ and $n - m$.

Now suppose $d' = 2dq$. Then

$$\frac{2m}{d'} = \frac{m}{dq}$$

is an integer, which contradicts $d = \gcd(m, n)$ unless $q = 1$.

Finally since $BTT(m, n) = (2m, n - m)$ or $(n - m, 2m)$, and

$$2m + (n - m) = m + n = 2^k d = 2^{k-1} d',$$

the proposition is established. \square

Proposition 2. *If (m, n) satisfies TC- k for some nonnegative integer k , then every point in the preimage $BTT^{-1}(m, n)$ satisfies*

1. TC-0 or TC-1, if $k = 0$,
 - and if the preimage $BTT^{-1}(m, n)$ is nonempty, it contains points satisfying TC-0;
2. TC- $(k + 1)$, if $k > 0$.

Proof. $BTT^{-1}(0, 0) = \{(0, 0)\}$, so both claims are trivially true for $(0, 0)$ as it satisfies TC- k for all nonnegative integers k .

Also $BTT^{-1}(m, n)$ is the empty set when $m > n$, so both claims are trivially true in this case also.

Case 1. TC-0:

So suppose (m, n) satisfies TC-0. That is, $m + n = \gcd(m, n)$. If both m and n are positive this is impossible, since then $m + n > \max(m, n)$ whereas $\gcd(m, n) \leq \min(m, n)$. Hence $m = 0$ or $n = 0$. Without loss of generality we may assume that $m = 0, n > 0$.

If n is odd, then $BTT^{-1}(0, n) = \{(0, n), (n, 0)\}$, and the first claim is satisfied.

If $n = 2p$ is even, then $BTT^{-1}(0, n) = \{(0, n), (n, 0), (p, p)\}$. Since $p > 0$, (p, p) satisfies TC-1, and again the first claim is established.

Case 2. TC- k for $k > 0$:

Suppose (m, n) satisfies TC- k , and without loss suppose $0 < m \leq n$. If one of m, n is odd and one is even, then $m + n$ is odd, which is impossible since $2^k \gcd(m, n)$ is even when $k > 0$.

If (m, n) is in the image of BTT then one of m, n must be even and hence both must be even, since (m, n) satisfies TC- k . Therefore if both m, n are odd, $BTT^{-1}(m, n)$ is the empty set and the second claim is satisfied.

So now assume that $m = 2p$ and $n = 2q$ for some positive integers $p \leq q$.

Suppose that $(u, v) \in BTT^{-1}(m, n)$, with $u \leq v$. Then we have $(m, n) \in \{(2u, v - u), (v - u, 2u)\}$. It follows that $(u, v) = (p, 2q + p)$ or $(u, v) = (q, 2p + q)$. (So in fact $u < v$.) Let $d = \gcd(m, n)$. d is certainly even, and we know from the proof of Proposition 1 that $2p/d$ and $2q/d$ are both odd.

Consider when $(u, v) = (p, 2q + p)$. Let $d' = \gcd(p, 2q + p)$. We claim that $d = 2d'$.

Certainly $d'|p$ and $d'|(2q + p)$ together imply that $d'|2p$ and $d'|2q$. Hence $d'|d$. Then we can write $d = sd'$ for some positive integer s .

Since d is even, write $d = 2r$ for some positive integer r . Then $d|2p$ and $d|2q$ together imply that $r|p$ and $r|q$ and so $r|(2q + p)$. From this it follows that $r|d'$. So write $d' = tr$ for some positive integer t .

Putting the previous statements together, we see $2r = d = str$, and so $st = 2$. If $s = 1$, $t = 2$ then $d = d'$. But then $d'|p$ means that $d|p$ and so $2p/d$ must be even, a contradiction. The only remaining possibility is that $s = 2$, $t = 1$ and so $d = 2d'$.

Now we can conclude that

$$p + (2q + p) = m + n = 2^k d = 2^{k+1} d',$$

and so (u, v) satisfies TC- $(k + 1)$, which was to be shown.

We have not used the fact that $p \leq q$, and so symmetry allows us to make a completely analogous argument to show that $(u, v) = (q, 2p + q)$ satisfies TC- $(k + 1)$ also.

If instead $u > v$, we can make the same arguments above with u and v interchanged, since $BTT(u, v) = BTT(v, u)$ by definition. We've shown, by construction, that when $m \leq n$ are both even, positive and (m, n) satisfies TC- k , then $BTT^{-1}(m, n)$ is nonempty and every point in it satisfies TC- $(k + 1)$, completing the proof of the second claim.

□

Theorem. *A pair of nonnegative integers (m, n) has a terminating BTT sequence if and only if the pair satisfies the terminating condition.*

Proof. If $m = 0$ then $(0, n) = BTT^k(0, n)$ for any nonnegative k . If $n = 0$ and $m > 0$ then $BTT^k(m, 0) = (0, m)$ for any positive k . Thus (m, n) has a terminating BTT sequence when $m = 0$ or $n = 0$, and we can see that such (m, n) also satisfy TC-0. This establishes the theorem when $m = 0$ or $n = 0$.

Now suppose that m, n are positive integers such that $BTT(m, n) = (0, p)$. p must be positive and even, since $BTT^{-1}(0, 0) = \{(0, 0)\}$ and $BTT^{-1}(0, p) = \{(0, p), (p, 0)\}$ when p is positive and odd, as we've seen in the proof of Proposition 2, in Case 1. So also from there we know that we must have $m = n = p/2$. $(p/2, p/2)$ satisfies TC-1.

If m, n are positive integers such that $BTT^k(m, n) = (0, p)$ for some $k > 1$, with k the smallest positive integer for which the equation is true, then we will write $BTT(u, v) = (0, p)$ where $(u, v) = BTT^{k-1}(m, n)$. Now u, v are both positive integers (otherwise k is not the smallest positive integer for which $BTT^k(m, n) = (0, p)$ holds), and from the previous paragraph we know that (u, v) satisfies TC-1. But (m, n) is in the $(k-1)$ -fold BTT preimage of (u, v) , i.e. $(m, n) \in BTT^{-(k-1)}(u, v)$. Applying Proposition 2(2), inductively $k-1$ times, we find that (m, n) satisfies TC- k .

We've shown the *only if* part of the equivalence, and now must show the *if* case.

Suppose (m, n) are positive integers that satisfy TC- k for some, necessarily positive, integer k . Applying Proposition 1 k times shows that $BTT^k(m, n)$ satisfies TC-0. But from Proposition 2 in Case 1, this means that $BTT^k(m, n) = (0, p)$ for some positive integer p . \square