Ucycles and Gucycles of Integer Compositions and Partitions

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In this particular Gucycle, we consider windows of size 3, and read the number of edges between the third and first vertex, the number of edges between the first and second vertex, and the number of edges between the second and third vertex, to be the three numerical values of a particular string being represented.

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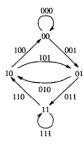
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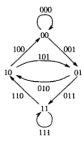
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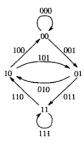


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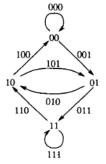


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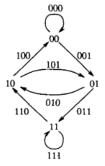
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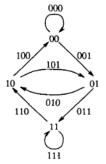
The transition digraph of this DeBruijn Sequence is generated as follows: let the vertices be strings of size n-1. If the last n-2 elements of a string match the first n-2 elements of another string, there is a directed edge from the first vertex to the second. From this we get our edge labels, which are the intersection of the two vertices. As such, they are length n.



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11101000,00010111

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Graph Representation

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Shifting the window to the right, we have that (1,3,1,2) can connect to (3,1,2,1), or (3,1,3) in the transition digraph. Shifting to the left, we have that (1,3,1,2) can come from (1,1,3,1,1), or (2,3,1,1).





Note that the edge labels are compositions of 8 ((2,3,1,2)) and (1,1,3,1,2) from the left, (1,3,1,3) and (1,3,1,2,1) to the right), so any Eulerian Circuits will be Gucycles of Compositions of (n+1).



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Theorem

Let G_C be the transition digraph of integer compositions of any n. G_C is balanced, with each vertex having an in-degree and out-degree of 2. Additionally, G_C is strongly connected. Thus, G_C is Eulerian, meaning Gucycles of Integer Compositions exist for any n.

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$$11101000 \implies (4) \rightarrow (3,1) \rightarrow (2,2) \rightarrow (1,2,1) \rightarrow (2,1,1) \rightarrow (1,1,1,1) \rightarrow (1,1,2) \rightarrow (1,3) \rightarrow (4)$$

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Therefore, since we already know DeBruijn Sequences exist for binary strings of any length, Ucycles must exist for compositions of any $n \in \mathbb{N}$.

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From this definition, two things are evident:

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- Choose whether to connect the new vertex to any of the existing parts of p, or let it remain a lone vertex. Recall that d is the number of distinct parts of p. Since there is one way to leave a lone vertex, and d distinct ways to connect a vertex to one of the parts of p, there are (d+1) ways to do this.

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- Delete any vertex. For any p(n, t, d), there are d distinct ways to delete a vertex from any part. However, after Step 2 we have multiple different partitions of (n+1) that could potentially have different values of d. Thus, we don't have a closed form for this.

The best we can do is this:

Theorem

Let p be an integer partition graph with d(p) distinct parts, and let α be the set of partitions of (n+1) that are created after Step 2. We have then that the number of partitions that p is connected to is equal to

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By our definition of the edges, we have that each vertex has at least one loop to itself. However, since most of these loops would have the same labels as other loops, and loops don't really affect Eulerian Circuits, we choose to ignore all loops that don't have a unique edge label.

'Thank you!



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