Solving Polynomial Systems Using Subresultants

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28 April 2024

In this paper, we will develop mathematical theories and algorithms for the following problem.

Input: $f \in \mathbb{C}[x_1, \dots, x_n]^n$

Output: S = V(f), the set of all complex solutions of f.

1 Theory

Notation 1 Let

$$f = a_p x^p + \dots + a_0 x^0$$
$$g = b_q x^q + \dots + b_0 x^0$$

where x is a variable and a_i and b_i are coefficients which might be again polynomials in other variables.

Definition 2 (Subresultant) The k-th subresultant of f and g with respect to x, written as $R_{x,k}(f,g)$, is defined by

$$\begin{vmatrix} a_p & a_{p-1} & \cdots & a_{\epsilon+1} & a_{\epsilon} & a_{\epsilon-1} & \cdots & a_{2k-q+2} & a_{2k-q+1}x^k + \cdots + a_0x^{q-k-1} \\ 0 & a_p & \cdots & a_{\epsilon+2} & a_{\epsilon+1} & a_{\epsilon} & \cdots & a_{2k-q+3} & a_{2k-q+2}x^k + \cdots + a_0x^{q-k-2} \\ \vdots & \vdots & & \vdots & & \vdots & & \vdots \\ 0 & 0 & \cdots & a_{\delta+1} & a_{\delta} & a_{\delta-1} & \cdots & a_{2k+q+r+1} & \sum_{n=0}^{2k-q+r} a_nx^{n+q-k-r} \\ \vdots & \vdots & & \vdots & & \vdots & & \vdots \\ 0 & 0 & \cdots & a_p & a_{p-1} & a_{p-2} & \cdots & a_k & a_{k-1}x^k + \cdots + a_0x^1 \\ 0 & 0 & \cdots & 0 & a_p & a_{p-1} & \cdots & a_{k+1} & a_kx^k + \cdots + a_0x^0 \\ b_q & b_{q-1} & \cdots & b_{\epsilon+1} & b_{\epsilon} & b_{\epsilon-1} & \cdots & b_{2k-p+2} & b_{2k-p+1}x^k + \cdots + b_0x^{p-k-1} \\ 0 & b_q & \cdots & b_{\epsilon+2} & b_{\epsilon+1} & b_{\epsilon} & \cdots & b_{2k-p+3} & b_{2k-q+2}x^k + \cdots + b_0x^{p-k-2} \\ \vdots & \vdots & & \vdots & & \vdots & & \vdots \\ 0 & 0 & \cdots & b_{\delta+1} & b_{\delta} & b_{\delta-1} & \cdots & b_{3k-p-q+r+1} & \sum_{n=1}^{3k-p-q+r} b_nx^{n-2k+p+q-r} \\ \vdots & \vdots & & \vdots & & \vdots & & \vdots \\ 0 & 0 & \cdots & b_q & b_{q-1} & b_{q-2} & \cdots & b_k & b_{k-1}x^k + \cdots + b_0x^1 \\ 0 & 0 & \cdots & 0 & b_q & b_{q-1} & \cdots & b_{k+1} & b_kx^k + \cdots + b_0x^0 \\ \end{vmatrix}$$

where r is given by the row number, starting at 1 and ending at q+p-2k. The matrix defining the subresultant is in $\mathbb{R}^{q+p-2k\times q+p-2k}$.

Theorem 3 We have

1.
$$\deg_x R_{x,k} \leq k$$

2.
$$R_{x,k}(f,g) \in \langle f,g \rangle$$
.

Proof.

- 1. Obvious from Laplace expansion along the last column.
- 2. For the sake of simple presentation, we show the proof only for p = 3, q = 4 and k = 1. Note that $f = a_3x^3 + a_2x^2 + a_1x + a_0$ and $g = b_4x^4 + b_3x^3 + b_2x^2 + b_1x + b_0$ and that we can define $R_{x,1}(f,g)$ as

$$\begin{vmatrix} a_3 & a_2 & a_1 & a_0 & 0 \\ 0 & a_3 & a_2 & a_1 & a_0 x \\ 0 & 0 & a_3 & a_2 & a_1 x + a_0 \\ b_4 & b_3 & b_2 & b_1 & b_0 x \\ 0 & b_4 & b_3 & b_2 & b_1 x + b_0 \end{vmatrix}$$

This matrix is 5×5 and we will define its columns as c_1, c_2, c_3, c_4 , and c_5 respectively. Using determinant rules, we can redefine c_5 as $c_5 + c_4x^2 + c_3x^3 + c_2x^4 + c_1x^5$ without changing the determinant of the matrix. Thus

$$\begin{vmatrix} a_3 & a_2 & a_1 & a_0 & 0 \\ 0 & a_3 & a_2 & a_1 & a_0x \\ 0 & 0 & a_3 & a_2 & a_1x + a_0 \\ b_4 & b_3 & b_2 & b_1 & b_0x \\ 0 & b_4 & b_3 & b_2 & b_1x + b_0 \end{vmatrix} = \begin{vmatrix} a_3 & a_2 & a_1 & a_0 & a_0x^2 + a_1x^3 + a_2x^4 + a_3x^5 \\ 0 & a_3 & a_2 & a_1 & a_0x + a_1x^2 + a_2x^3 + a_3x^4 \\ 0 & 0 & a_3 & a_2 & a_0 + a_1x + a_2x^2 + a_3x^3 \\ b_4 & b_3 & b_2 & b_1 & b_0x + b_1x^2 + b_2x^3 + b_3x^4 + b_4x^5 \\ 0 & b_4 & b_3 & b_2 & b_0 + b_1x + b_2x^2 + b_3x^3 + b_4x^4 \end{vmatrix}$$

We can then use determinant rules to factor out f and g from c_5 , giving

If we set
$$U(x) = \begin{pmatrix} a_3 & a_2 & a_1 & a_0 & x^2 \\ 0 & a_3 & a_2 & a_1 & x \\ 0 & 0 & a_3 & a_2 & 1 \\ b_4 & b_3 & b_2 & b_1 & 0 \\ 0 & b_4 & b_3 & b_2 & 0 \end{pmatrix}$$
 and $V(x) = \begin{pmatrix} a_3 & a_2 & a_1 & a_0 & 0 \\ 0 & a_3 & a_2 & a_1 & 0 \\ 0 & 0 & a_3 & a_2 & 0 \\ b_4 & b_3 & b_2 & b_1 & x \\ 0 & b_4 & b_3 & b_2 & 1 \end{pmatrix}$, then we can say that $R_{x,1}(f,g) = U(x)f + V(x)g$, thus $R_{x,1}(f,g) \in \langle f,g \rangle$

Definition 4 (Triangularization) Let

$$f = (f_1, \ldots, f_n) \in \mathbb{C} [x_1, \ldots, x_n]^n$$

Then the trianglarization of f, denoted as

$$\tilde{f} = \left(\tilde{f}_1, \dots, \tilde{f}_n\right) \in \mathbb{C}\left[x_1, \dots, x_n\right]^n$$

is defined by the following process.

(For the sake of simple presentation, we will show the n=4 case only. The generalization to arbitrary n is straightforward).

1. Repeated 0-th order subresultant (resultant):

2. 1-st order subresultant:

$$\tilde{f}_1 = R_{x_1,1} (f_1, f_2)
\tilde{f}_2 = R_{x_2,1} (f_{12}, f_{13})
\tilde{f}_3 = R_{x_3,1} (f_{123}, f_{124})
\tilde{f}_4 = f_{1234}$$

Theorem 5 $V(\tilde{f}) \supseteq V(f)$

Proof. For the sake of simple presentation, we show the proof only for n=4. We will first show that $\tilde{f}_1, \tilde{f}_2, \tilde{f}_3, \tilde{f}_4 \in \langle f_1, f_2, f_3, f_4 \rangle$, then we will use that to show that for any $\alpha \in V(f)$, $\alpha \in V(\tilde{f})$ as well. Note the following:

$$\begin{split} \tilde{f}_{1} &= R_{x_{1},1}(f_{1},f_{2}) \in \langle f_{1},f_{2} \rangle \subset \langle f_{1},f_{2},f_{3},f_{4} \rangle \\ \tilde{f}_{2} &= R_{x_{2},1}\left(f_{12},f_{13}\right) \in \langle f_{12},f_{13} \rangle = \langle R_{x_{1},0}\left(f_{1},f_{2}\right),R_{x_{1},0}\left(f_{1},f_{3}\right) \rangle \\ &\subset \langle \langle f_{1},f_{2} \rangle, \langle f_{1},f_{3} \rangle \rangle \subset \langle f_{1},f_{2},f_{3} \rangle \subset \langle f_{1},f_{2},f_{3},f_{4} \rangle \\ \tilde{f}_{3} &= R_{x_{3},1}\left(f_{123},f_{124}\right) \in \langle f_{123},f_{123} \rangle = \langle R_{x_{2},0}(f_{12},f_{13}),R_{x_{2},0}(f_{12},f_{14}) \rangle \\ &\subset \langle \langle f_{12},f_{13} \rangle, \langle f_{12},f_{14} \rangle \rangle \subset \langle f_{12},f_{13},f_{14} \rangle = \langle R_{x_{1},0}(f_{1},f_{2}),R_{x_{1},0}(f_{1},f_{3}),R_{x_{1},0}(f_{1},f_{4}) \rangle \\ &\subset \langle \langle f_{1},f_{2} \rangle, \langle f_{1},f_{3} \rangle, \langle f_{1},f_{4} \rangle \rangle \subset \langle f_{1},f_{2},f_{3},f_{4} \rangle \\ \\ \tilde{f}_{4} &= R_{x_{3},0}\left(f_{123},f_{124}\right) \in \langle f_{123},f_{124} \rangle \subset \langle f_{1},f_{2},f_{3},f_{4} \rangle \end{split}$$

Thus $\tilde{f} \in \langle f \rangle$. Because of this, we can define \tilde{f} as follows:

$$\begin{split} \tilde{f}_1 &= u_{11}f_1 + u_{12}f_2 + u_{13}f_3 + u_{14}f_4 \\ \tilde{f}_2 &= u_{21}f_1 + u_{22}f_2 + u_{23}f_3 + u_{24}f_4 \\ \tilde{f}_3 &= u_{31}f_1 + u_{32}f_2 + u_{33}f_3 + u_{34}f_4 \\ \tilde{f}_4 &= u_{41}f_1 + u_{42}f_2 + u_{43}f_3 + u_{44}f_4 \end{split}$$

for some set of u's in \mathbb{R} . Let α solve f, that is $f_1(\alpha) = f_2(\alpha) = f_3(\alpha) = f_4(\alpha) = 0$. From this, we can say that $\tilde{f}_1(\alpha) = u_{11} \cdot 0 + u_{12} \cdot 0 + u_{13} \cdot 0 + u_{14} \cdot 0 = 0$, which is similarly true for \tilde{f}_2 , \tilde{f}_3 , and \tilde{f}_4 . Because $\tilde{f}_1(\alpha) = \tilde{f}_2(\alpha) = \tilde{f}_3(\alpha) = \tilde{f}_4(\alpha) = 0$, we know that $\alpha \in V(\tilde{f})$. Therefore it must be true that $V(f) \subseteq F(\tilde{f})$.

2 Algorithms

Algorithm 6 (Triangularize)

- 1. $h \leftarrow f$
- 2. $\tilde{f} \leftarrow 0_{n \times 1}$
- 3. For i = 0, ..., n 1 $\tilde{f}_i \leftarrow R_{x_i,1}(h_i, h_{i+1})$ where for j = i + 1, ..., n 1 $h_j \leftarrow R_{x_i,0}(h_i, h_j)$
- 4. $\tilde{f}_n \leftarrow h_{n-1}$

Algorithm 7 (BackSubstitute)

- 1. $rs \leftarrow All \ complex \ roots \ of \ \tilde{f}_{n-1}$
- 2. $T \leftarrow [\]$, an empty list
- 3. For a root r in rs,

$$t \leftarrow 0_{n \times 1}$$

$$t_{n-1} \leftarrow r$$

where for
$$i = n - 2, n - 1, \dots, 1, 0$$

where for
$$j = i + 1, \dots, n - 1$$

$$h \leftarrow \tilde{f}_i(t_j)$$

 $t_i \leftarrow -\frac{\operatorname{coeff}(h, x_i, 0)}{\operatorname{coeff}(h, x_i, 1)}$, where the coefficients are evaluated on h at x_i^0 and x_i^1 Append t to T

4. $T \leftarrow All \ possible \ complex \ solutions \ to \ \tilde{f}$

Algorithm 8 (ChooseSolution)

- 1. $S \leftarrow [\]$, an empty list
- 2. For a solution t in T,

$$h \leftarrow f(t)$$

and if
$$||h||_2 \leq \epsilon$$
, append h to S

3. $S \leftarrow All \ possible \ complex \ solutions \ to \ f$