

# Graph Theory

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Graphs

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Trees

# Graphs

# Seven Bridges of Königsberg



Is it possible to walk through the city crossing each of the seven bridges once and only once?

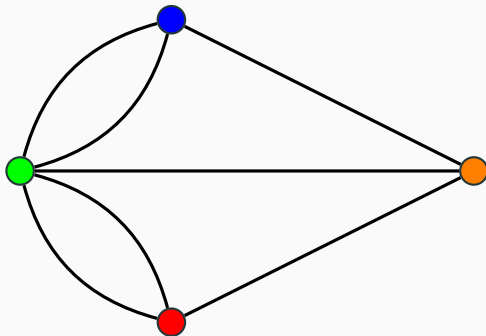
[www.nature.com/nbt/journal/v29/n11](http://www.nature.com/nbt/journal/v29/n11)



- Born 1707 in Basel, Switzerland.
- Euler's identity:  $e^{i\pi} + 1 = 0$ .
- Solved the Bridges of Königsberg problem.
- It's not possible to cross all bridges once and once only.

[https://en.wikipedia.org/wiki/Leonhard\\_Euler](https://en.wikipedia.org/wiki/Leonhard_Euler)

## Graph of Königsberg



## Definition

A *graph* consists of a finite set  $V$  and a set  $E$  of 2-subsets of  $V$ .

**Vertices** – the elements of the set  $V$  are called vertices.

**Edges** – the elements of  $E$  are called edges.

$G = (V, E)$  – this is the way we write the graph  $G$  consists of the vertex set  $V$  and the edge set  $E$ .

$$V = \{Green, Blue, Orange, Red\}$$

$$E = \{ \\ \{Green, Blue\}, \{Green, Blue\}, \{Green, Red\}, \\ \{Green, Red\}, \{Blue, Orange\}, \{Green, Orange\}, \\ \{Red, Orange\} \\ \}$$



## Adjacency list

Green	Blue	Orange	Red
Blue	Green	Blue	Green
Orange	Orange	Green	Orange
Red		Red	

# Defining different types of graphs

## Our definition of a graph

The definition given above for a graph is not consistent with looped edges, directed edges or repeated edges. We only need to make small changes to the definition of a graph to allow for directed edges and repeated edges.

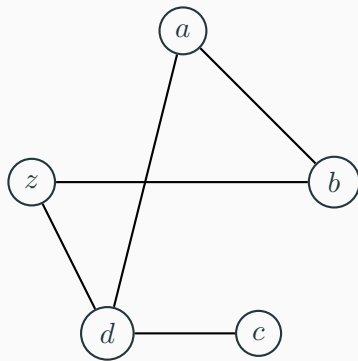
**Repeated edges** are edges that start and end at the same vertices.

**Directed edges** are edges where a direction is added.

**Looped edges** begin and end at the same vertex.

The application will determine the definition we want to use.

## A better example



### Exercise

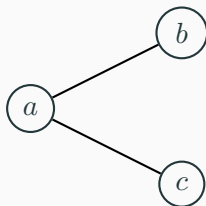
Determine the vertex set, edge set and adjacency list of this graph.

[global.oup.com/booksites/content/9780198507185/](http://global.oup.com/booksites/content/9780198507185/)

# Degree of a vertex

## Definition

The degree of a vertex is the number of edges that contain it.



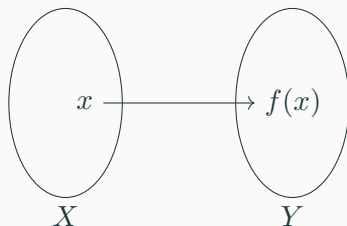
The degree of the vertex  $a$  is 2.

## Exercise

For each of the vertices on the previous slide, determine its degree.

## Definition

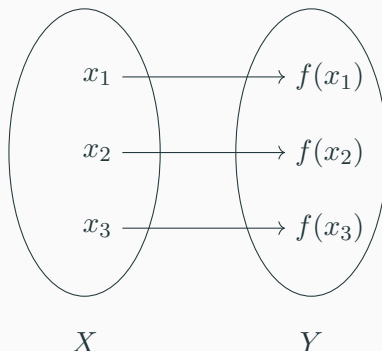
Suppose that  $X$  and  $Y$  are sets. We say we have a function  $f$  from  $X$  to  $Y$  if for each  $x$  in  $X$  we can specify a unique element in  $Y$ , which we denote by  $f(x)$ .



## Definition

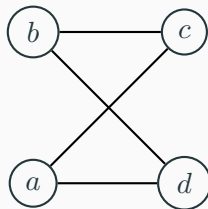
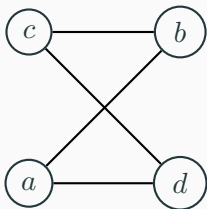
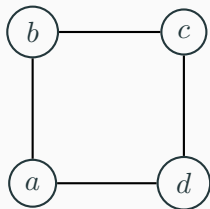
A bijection is function  $f$  from a set  $X$  to a set  $Y$  where both of the following are true:

- every  $y$  in  $Y$  is a value  $f(x)$  for at most one  $x$  in  $X$ .
- every  $y$  in  $Y$  is a value  $f(x)$  for at least one  $x$  in  $X$ .



## Definition

Two graphs  $G_1$  and  $G_2$  are said to be isomorphic when there is a bijection  $\alpha$  for the vertex set  $V_1$  of  $G_1$  to the vertex set  $V_2$  of  $G_2$  such that  $\{\alpha(x), \alpha(y)\}$  is an edge of  $G_2$  if and only if  $(x, y)$  is an edge of  $G_1$ .



# Sum of degrees

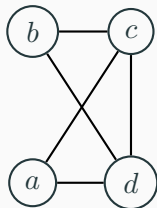
## Theorem

*The sum of the degrees of the vertices of a graph  $G = (V, E)$  is equal to twice the number of edges:*

$$\sum_{v \in V} \delta(v) = 2|E|$$

## Proof.

The degree  $\delta(v)$  of a vertex  $v$  is equal to the number of edges incident on it. Every edge is incident on two vertices. So every edge contributes 1 to the degrees of two distinct vertices. Therefore every edge contributes 2 to the sum total of the degrees of all the vertices.  $\square$

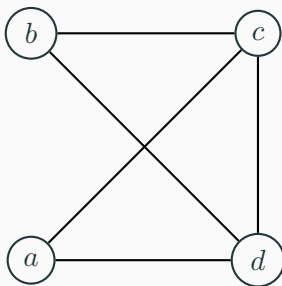




# Handshaking lemma

## Definition

A vertex is an odd vertex if its degree is odd, and it is an even vertex if its degree is even. The set of all odd vertices is denoted  $V_o$  and the set of all even vertices is denoted  $V_e$ .



## Exercise

Which of the above vertices are even, and which are odd?

## Lemma

*The number of odd vertices  $|V_o|$  in a graph is even.*

## Proof.

The sets  $V_o$  and  $V_e$  are disjoint (i.e. they don't have any elements in common.) Also, every vertex is either in  $V_o$  or  $V_e$ . Therefore  $V = V_o \cup V_e$  and  $|V| = |V_o| + |V_e|$ .

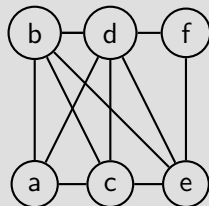
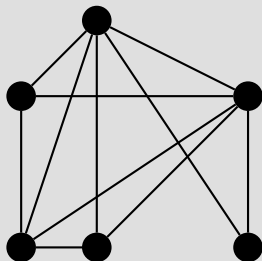
Furthermore:

$$\sum_{v \in V_o} \delta(v) + \sum_{v \in V_e} \delta(v) = 2|E|$$

Both  $2|E|$  and  $\sum_{v \in V_e} \delta(v)$  are even, so  $\sum_{v \in V_o} \delta(v)$  must be. Since  $\delta(v)$  is odd for every  $v$  in  $V_o$ , this must mean that  $|V_o|$  is even.  $\square$

## Exercise

Determine if these two graphs are isomorphic.



# Paths and Cycles

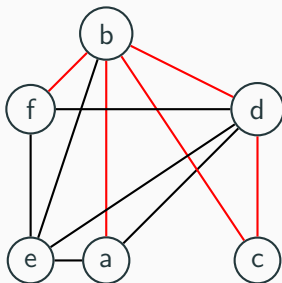
## Definition

A *walk* in a graph is a sequence of vertices

$$v_1, v_2, \dots, v_k$$

such that  $v_i$  and  $v_{i+1}$  are adjacent for  $1 \leq i < k$ .

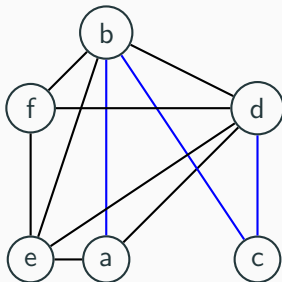
For instance, in the following graph  $a, b, c, d, b, f$  is a walk.



## Definition

A *path* is a walk where each vertex is distinct.

For instance, in the following graph  $a, b, c, d$  is a path.

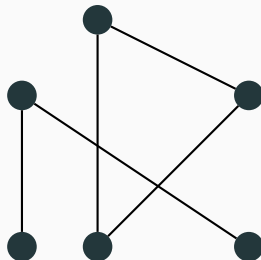
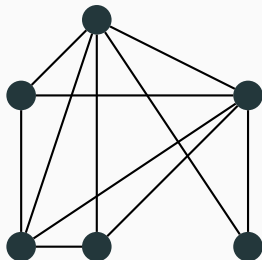


# Connected graphs

## Definition

A graph is *connected* if there is a path between each pair of vertices.

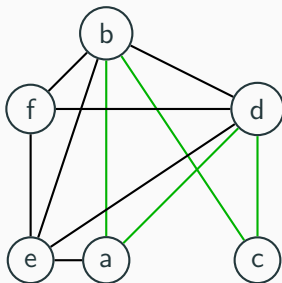
For instance, the graph on the left below is connected, but the one on the right is not.



## Definition

A *cycle* is a walk where each vertex is distinct, except for the start and end vertices being equal.

For instance, in the following graph  $a, b, c, d, a$  is a cycle.



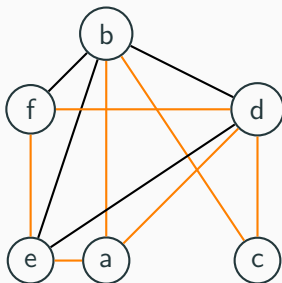


# Hamiltonian cycles

## Definition

A *Hamiltonian cycle* is a cycle that contains all of vertices of the graph.

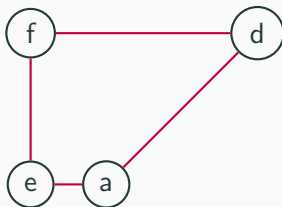
For instance, in the following graph  $a, b, c, d, f, e, a$  is a Hamiltonian cycle.



## Definition

An *Eulerian walk* is a walk which uses each edge of the graph exactly once.

For instance, in the following graph  $a, d, f, e, a$  is an Eulerian walk.



# Graph databases

## Definition

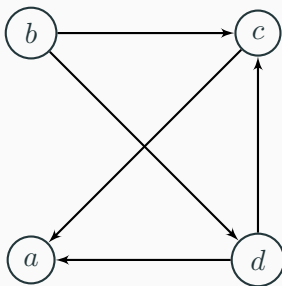
A *digraph* (short for directional graph) consists of a finite set  $V$  and a set  $E$  of ordered pairs of elements of  $V$ .

**Degrees** of vertices can now be split into in-degrees and out-degrees.

**Walks, paths, cycles** must be redefined.

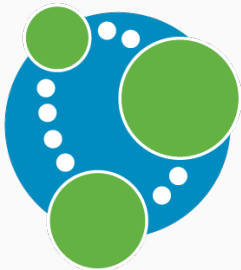
**Loops** are allowed in the above definition, unless we rule them out.

## Digraph example





- Neo4j is an open-source NoSQL graph database implemented in Java and Scala.
- Development started in 2003, it has been publicly available since 2007
- Available on GitHub.
- A graph is composed of two elements: a node and a relationship.



- Cypher is a declarative graph query language.
- What to retrieve from a graph, not on how to retrieve it.
- Allows for expressive and efficient querying and updating of the graph store.
- Cypher borrows its structure from SQL.

Create a node with the label User, and two properties:

---

```
1 CREATE (user:User { Id: 123, Name: "Jim" });
```

---

Find the node(s) with label User and their Id property being 123:

---

```
1 MATCH (user:User)
2 WHERE user.Id = 123
3 RETURN user;
```

---



Create a relationship with label FOLLOWS from user(s) with Id 123 to user(s) with Id 456:

---

```
1 MATCH (user1:User), (user2:User)
2 WHERE user1.Id = 123 AND user2.Id = 456
3 CREATE user1-[:FOLLOWS]->user2;
```

---

Create a relationship with label INVITED from user(s) with Id 123 to a new user with Id 789 and Name Jack:

---

```
1 MATCH (invitee:User)
2 WHERE invitee.Id = 123
3 CREATE invitee-[:INVITED]->(invited:User {Id: 789,
4                                     Name: "Jack"});
```

---

Delete all nodes:

---

```
1 MATCH (x)
2 DELETE x;
```

---

- Suppose we have nodes representing people.
- We give them the label People.
- We also want to identify each person as either Male or Female.
- Should we use Male and Female labels, or a Gender property?
- If you are going to use the person's gender in a lot of queries, a normal property will be relatively slow, so you should use a label.
- However, you can also index some of your properties to high-light them as important.

Find the minimum number of hops between two nodes.

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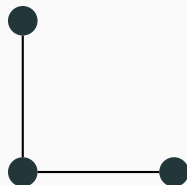
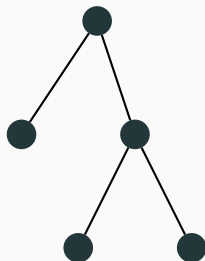
```
1 MATCH p=shortestPath(  
2   (a:Actor {id: 1})-[*]-(b:Actor {id: 10})  
3   )  
4 RETURN p;
```

---

# Trees

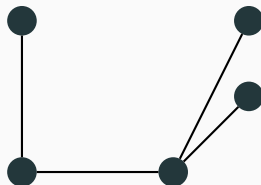
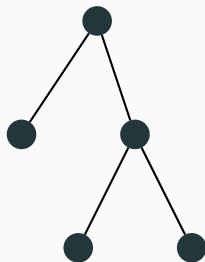
## Tree

A *tree* is a graph where every pair of nodes has a path between them, and there are no cycles.



## Tree

A *tree* is a graph where every pair of vertices has a path between them, and there are no cycles.



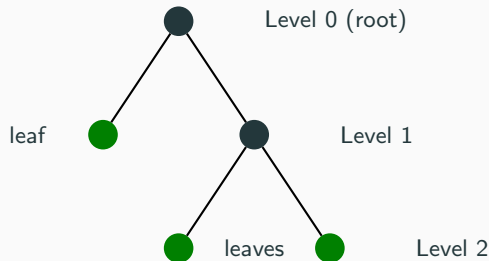


**Root** Specified vertex for some purpose.

**Levels** Root is at level 0, neighbours of the root are at level 1, their other neighbours at level 2, and so on.

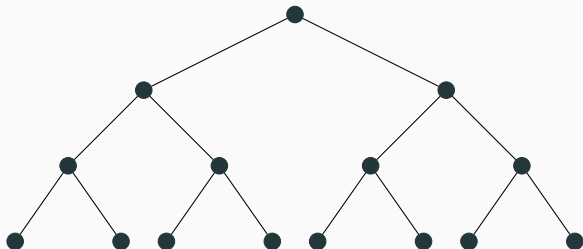
**Leaf** Vertex at level  $i$  with no neighbours at level  $i + 1$ .

**Internal vertex** A non-leaf.



## Definition

When a vertex at level  $i$  is connected to a vertex at level  $i + 1$  it's common to call the former the *father* and the latter the *son*. A rooted tree is  $m$ -ary if every father has the same number of sons. A 2-ary rooted tree is called a *binary tree*.



## Theorem

*The height of an  $m$ -ary rooted tree with  $l$  leaves is at least  $\log_m l$ .*