

# Graph Theory

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Fundamentals

Trees

Graph databases

Algorithms

# Fundamentals

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# Seven Bridges of Königsberg



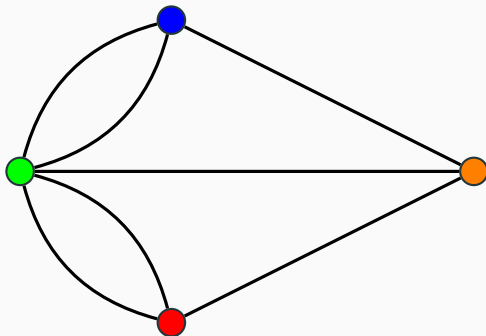
Is it possible to walk through the city crossing each of the seven bridges once and only once?

[www.nature.com/nbt/journal/v29/n11](http://www.nature.com/nbt/journal/v29/n11)



- Born 1707 in Basel, Switzerland.
- Euler's identity:  $e^{i\pi} + 1 = 0$ .
- Solved the Bridges of Königsberg problem.
- It's not possible to cross all bridges once and once only.

# Graph of Königsberg



## Definition

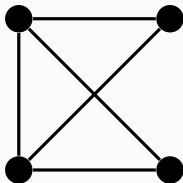
A *graph* consists of a finite set  $V$  and a set  $E$  of 2-subsets of  $V$ .

**Vertices** – the elements of the set  $V$  are called vertices.

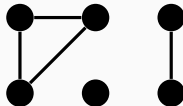
**Edges** – the elements of  $E$  are called edges.

$G = (V, E)$  – this is the way we write the graph  $G$  consists of the vertex set  $V$  and the edge set  $E$ .

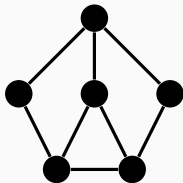
## Example graphs



4 vertices, 5 edges



6 vertices, 4 edges



6 vertices, 8 edges



1 vertex, 0 edges



$$V = \{Green, Blue, Orange, Red\}$$

$$E = \{\{Green, Blue\}, \{Green, Blue\}, \{Green, Red\}, \\ \{Green, Red\}, \{Blue, Orange\}, \{Green, Orange\}, \\ \{Red, Orange\}\}$$

## Not a graph by our definition

Note that the Bridges of Königsberg graph above is not a graph, due to the repeated edges. It's a multi-graph.

## Adjacency list

Green	Blue	Orange	Red
Blue	Green	Blue	Green
Orange	Orange	Green	Orange
Red		Red	

# Defining different types of graphs

## Our definition of a graph

The definition given above for a graph is not consistent with looped edges, directed edges or repeated edges. We only need to make small changes to the definition of a graph to allow for directed edges and repeated edges.

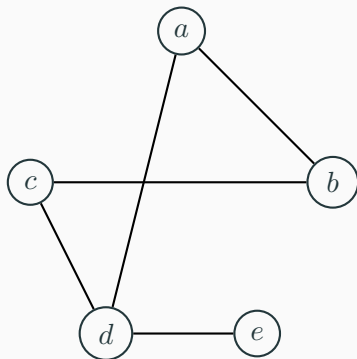
**Repeated edges** are edges that start and end at the same vertices.

**Directed edges** are edges where a direction is added.

**Looped edges** begin and end at the same vertex.

The application will determine the definition we want to use.

## A better example

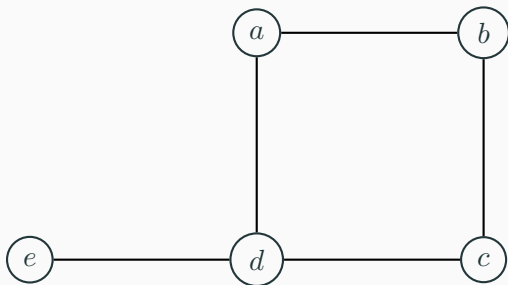


### Exercise

Determine the vertex set, edge set and adjacency list of this graph.

[global.oup.com/booksites/content/9780198507185/](http://global.oup.com/booksites/content/9780198507185/)

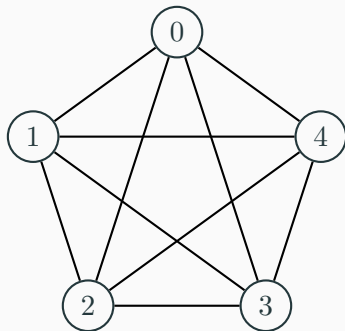
## Another better example



### Exercise

Determine the vertex set, edge set and adjacency list of this graph.

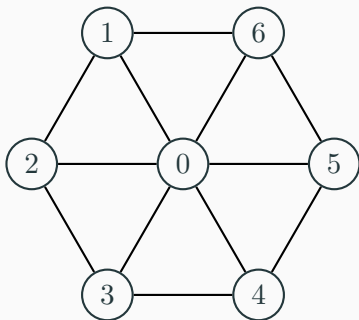
## Complete graph $K_n$



### Exercise

Determine the vertex set, edge set and adjacency list of  $K_5$ .

## Wheel graph $W_n$



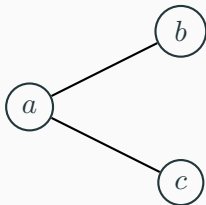
### Exercise

Determine the vertex set, edge set and adjacency list of  $W_6$ .

# Degree of a vertex

## Definition

The degree of a vertex is the number of edges that contain it.



The degree of the vertex  $a$  is 2.

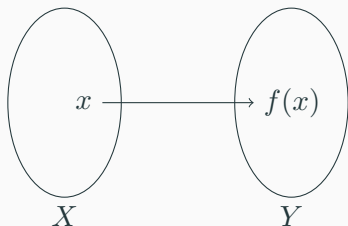
## Exercise

For each of the vertices on the previous slide, determine its degree.



## Definition

Suppose that  $X$  and  $Y$  are sets. We say we have a function  $f$  from  $X$  to  $Y$  if for each  $x$  in  $X$  we can specify a unique element in  $Y$ , which we denote by  $f(x)$ .

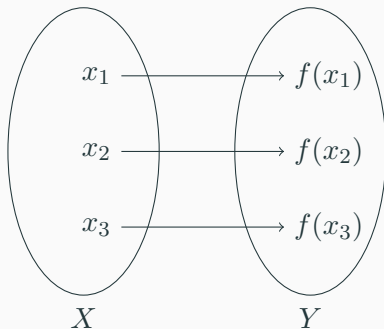


# Bijections

## Definition

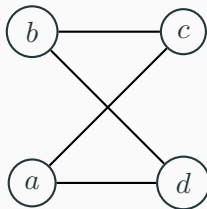
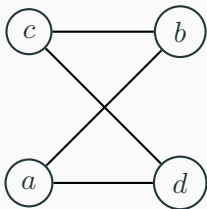
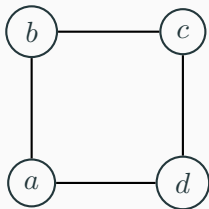
A bijection is function  $f$  from a set  $X$  to a set  $Y$  where both of the following are true:

- every  $y$  in  $Y$  is a value  $f(x)$  for at most one  $x$  in  $X$ .
- every  $y$  in  $Y$  is a value  $f(x)$  for at least one  $x$  in  $X$ .

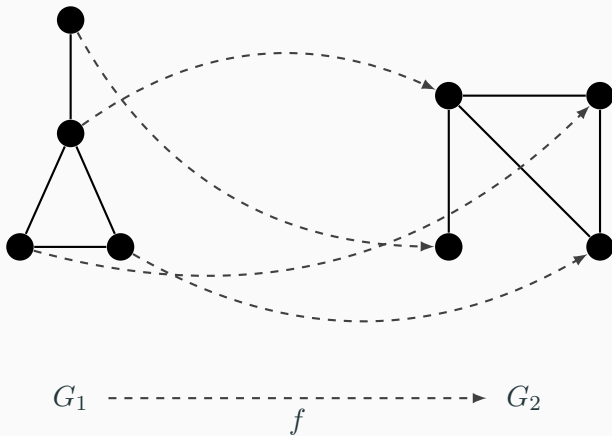


## Definition

Two graphs  $G_1$  and  $G_2$  are said to be isomorphic when there is a bijection  $\alpha$  for the vertex set  $V_1$  of  $G_1$  to the vertex set  $V_2$  of  $G_2$  such that  $\{\alpha(x), \alpha(y)\}$  is an edge of  $G_2$  if and only if  $(x, y)$  is an edge of  $G_1$ .

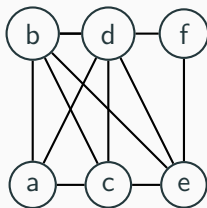
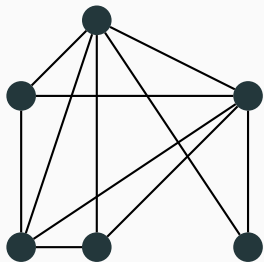


## Isomorphism example



## Exercise

Determine if these two graphs are isomorphic.



# Sum of degrees

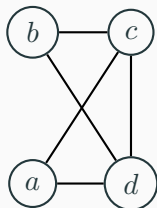
## Theorem

*The sum of the degrees of the vertices of a graph  $G = (V, E)$  is equal to twice the number of edges:*

$$\sum_{v \in V} \delta(v) = 2|E|$$

## Proof.

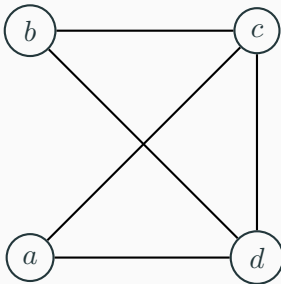
The degree  $\delta(v)$  of a vertex  $v$  is equal to the number of edges incident on it. Every edge is incident on two vertices. So every edge contributes 1 to the degrees of two distinct vertices. Therefore every edge contributes 2 to the sum total of the degrees of all the vertices.  $\square$



# Handshaking lemma

## Definition

A vertex is an odd vertex if its degree is odd, and it is an even vertex if its degree is even. The set of all odd vertices is denoted  $V_o$  and the set of all even vertices is denoted  $V_e$ .



## Exercise

Which of the above vertices are even, and which are odd?

# Handshaking lemma

## Lemma

*The number of odd vertices  $|V_o|$  in a graph is even.*

## Proof.

The sets  $V_o$  and  $V_e$  are disjoint (i.e. they don't have any elements in common.) Also, every vertex is either in  $V_o$  or  $V_e$ . Therefore  $V = V_o \cup V_e$  and  $|V| = |V_o| + |V_e|$ .

Furthermore:

$$\sum_{v \in V_o} \delta(v) + \sum_{v \in V_e} \delta(v) = 2|E|$$

Both  $2|E|$  and  $\sum_{v \in V_e} \delta(v)$  are even, so  $\sum_{v \in V_o} \delta(v)$  must be. Since  $\delta(v)$  is odd for every  $v$  in  $V_o$ , this must mean that  $|V_o|$  is even.  $\square$



# Directed graph definition

## Definition

A *directed graph* consists of a finite set  $V$  and a set  $E$  of 2-tuples (ordered pairs of elements) from  $V$ .

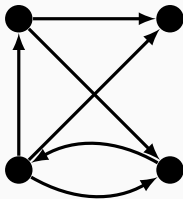
**Looped edges** are allowed in this definition. A single one per vertex.

**Multiple edges** between the same start and end vertices are not allowed. However, two edges are allowed between every pair of vertices so long as they have opposite directions.

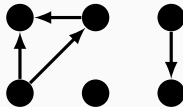
**Direct edges** use round brackets rather than curly braces:

$$E = \{(a, b) \mid a, b \in V\}.$$

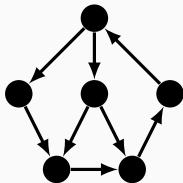
## Example directed graphs



4 vertices, 6 directed edges



6 vertices, 4 directed edges



6 vertices, 8 directed edges



1 vertex, 0 directed edges

# Multigraph definition

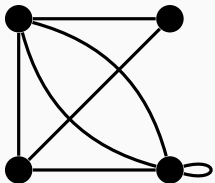
## Definition

A *multigraph* consists of a finite set  $V$  and a multiset  $E$  of 2-multisubsets from  $V$ .

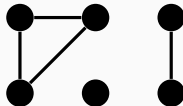
**Multisets** are like sets, but the same element can be in the set more than once.

**Directed multigraphs** are similar, but  $E$  is a set of 2-tuples of elements in  $V$ .

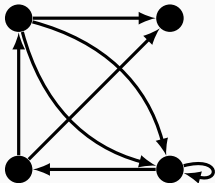
# Example multigraphs



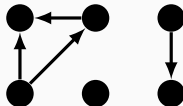
Multigraph



Still a multigraph



Directed multigraph



Still a directed multigraph

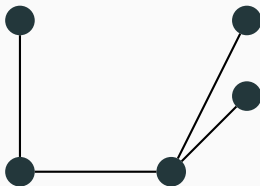
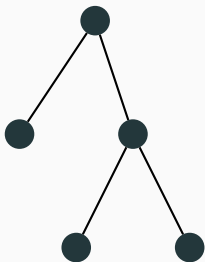
# Trees

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# Definition

## Tree

A *tree* is a graph where every pair of vertices has a path between them, and there are no cycles.



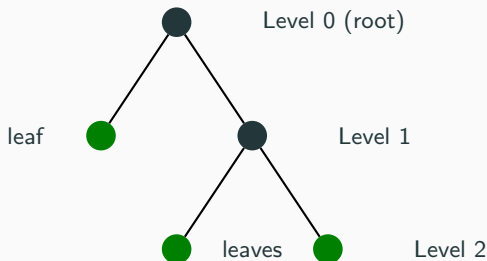
# Rooted trees

**Any vertex** of a tree can be called its root.

**Levels** Root is at level 0, neighbours of the root are at level 1, their other neighbours at level 2, and so on.

**Height** of a tree is  $h$ , where there's vertex at level  $h$  but not at level  $h + 1$ .

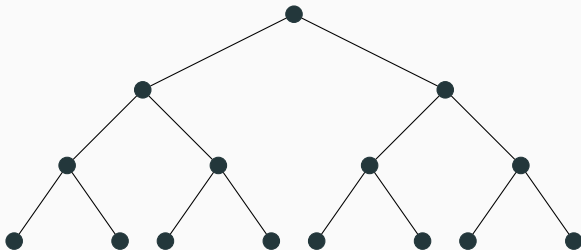
**Leaf** Vertex at level  $i$  not connected to a vertex at level  $i + 1$ .



# *m*-ary Rooted Tree

## Definition

When a vertex at level  $i$  is connected to a vertex at level  $i + 1$  it's common to call the former the *parent* and the latter the *child*. A rooted tree is *m*-ary if every parent has the same number of children. A 2-ary rooted tree is called a *binary tree*.

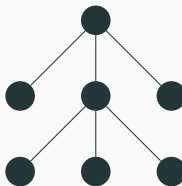
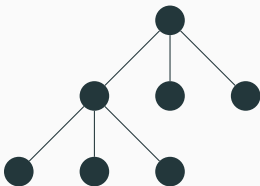




# Isomorphic Rooted Trees

## Definition

Two rooted trees are said to be *isomorphic* if there is a graph isomorphism between them which takes the root of one tree to the root of the other.



We define  $\log$  in the following way:

$$m^h = l \Leftrightarrow \log_m l = h$$

## What does *log* mean?

Suppose we have two numbers  $m$  and  $h$  and we ask the question “what is  $m$  to the power of  $h$ ?” Let’s call the answer  $l$ , so  $l = m^h$ .

The  $\log$  function asks the inverse question: “what do we need to raise  $m$  to the power of to get  $l$ ?” The answer is  $h$ .

For example,  $10^2 = 100$  so  $\log_{10} 100 = 2$ . The subscript 10 is called the *base*.

# Heights and leaves of $m$ -ary rooted trees

## Theorem

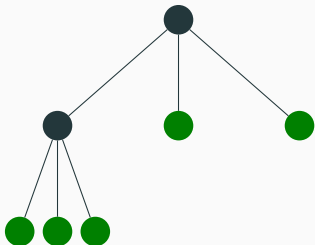
*The height  $h$  of an  $m$ -ary rooted tree with  $l$  leaves is at least  $\log_m l$ . That is:  $h \geq \log_m l$ .*

## Proof.

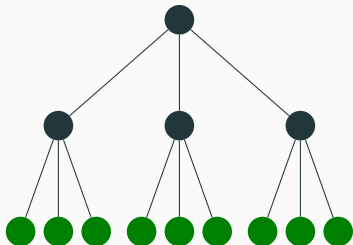
Note:  $h \geq \log_m l \Leftrightarrow m^h \geq m^{\log_m l} \Leftrightarrow m^h \geq l$ . So we'll just show  $l$  is at most  $m^h$ . For a tree of height 0,  $l = 1$  and  $m^0 = 1$  so  $m^h \geq l$ . Now assume the theorem is true for trees of height  $i - 1$ . Consider a tree of height  $i$  with  $l$  leaves. We can create  $m$  trees of height  $i - 1$  from it by deleting the root. Each of these smaller trees has at most  $m^{h-1}$  leaves by assumption. There are  $m$  of these, so the big tree has at most  $m \times m^{h-1} = m^h$  leaves. □

## Examples of heights and leaves of $m$ -ary trees

$$m = 3, h = 2, l = 5$$

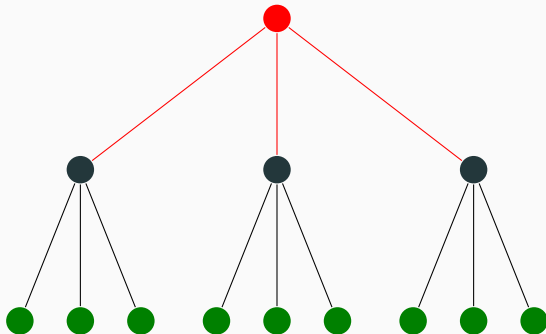


$$m = 3, h = 2, l = 9$$



## Deleting the root of an $m$ -ary tree

$m$  smaller trees of height  $h - 1$



## Subgraph

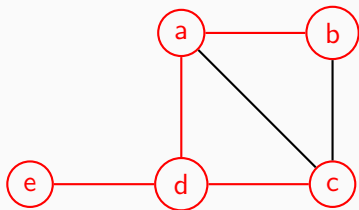
A *subgraph*  $H = (V_H, E_H)$  of a graph  $G = (V, E)$  is a graph such that  $V_H$  is a subset of  $V$ ,  $E_H$  is a subset of  $E$ , and no edge in  $E_H$  contains a vertex not in  $V_H$ .

## Spanning Tree

A *spanning tree*  $T$  of a connected graph  $G$  is a subgraph of  $G$  such that:

- the vertex set of  $T$  is the vertex set of  $G$  and
- $T$  is a tree.

## Spanning tree example



Spanning tree in red.

# Graph databases

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## Definition

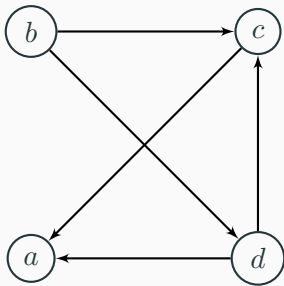
A *digraph* (short for directional graph) consists of a finite set  $V$  and a set  $E$  of ordered pairs of elements of  $V$ .

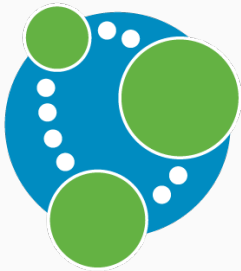
**Degrees** of vertices can now be split into in-degrees and out-degrees.

**Walks, paths, cycles** must be redefined.

**Loops** are allowed in the above definition, unless we rule them out.

## Digraph example





- Neo4j is an open-source NoSQL graph database implemented in Java and Scala.
- Development started in 2003, it has been publicly available since 2007
- Available on GitHub.
- A graph is composed of two elements: a node and a relationship.



- Cypher is a declarative graph query language.
- What to retrieve from a graph, not on how to retrieve it.
- Allows for expressive and efficient querying and updating of the graph store.
- Cypher borrows its structure from SQL.

# Cypher: Nodes

Create a node with the label User, and two properties:

---

```
1 CREATE (user:User { Id: 123, Name: "Jim" });
```

---

Find the node(s) with label User and their Id property being 123:

---

```
1 MATCH (user:User)
2 WHERE user.Id = 123
3 RETURN user;
```

---

Create a relationship with label FOLLOWS from user(s) with Id 123 to user(s) with Id 456:

---

```
1 MATCH (user1:User), (user2:User)
2 WHERE user1.Id = 123 AND user2.Id = 456
3 CREATE user1-[:FOLLOWS]->user2;
```

---

# Cypher: Relationships and Nodes

Create a relationship with label INVITED from user(s) with Id 123 to a new user with Id 789 and Name Jack:

---

```
1 MATCH (invitee:User)
2 WHERE invitee.Id = 123
3 CREATE invitee-[:INVITED]->(invited:User {Id: 789,
4                                     Name: "Jack"});
```

---

Delete all nodes:

---

```
1 MATCH (x)
2 DELETE x;
```

---



# Labels and properties

- Suppose we have nodes representing people.
- We give them the label People.
- We also want to identify each person as either Male or Female.
- Should we use Male and Female labels, or a Gender property?
- If you are going to use the person's gender in a lot of queries, a normal property will be relatively slow, so you should use a label.
- However, you can also index some of your properties to highlight them as important.

# Cypher: shortestPath

Find the minimum number of hops between two nodes.

---

```
1 MATCH p=shortestPath(  
2   (a:Actor {id: 1})-[*]-(b:Actor {id: 10})  
3   )  
4 RETURN p;
```

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# Algorithms

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**A decision tree** is a rooted tree where each vertex represents a decision.

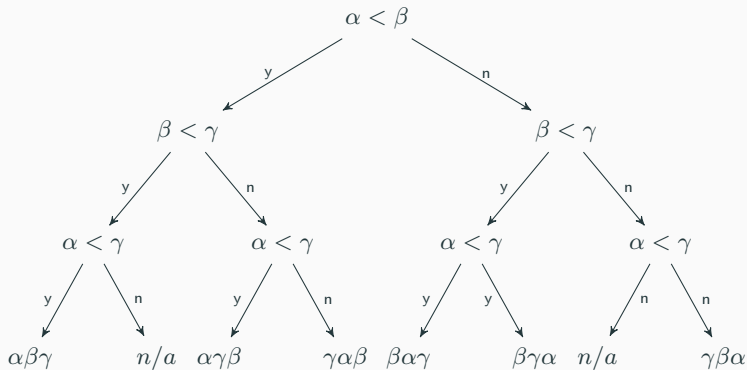
**The results** of a decision are represented by the edges connecting the vertex to the vertices at the next level down.

**Decisions** can connect to other decisions further down the tree.

**Final outcomes** of the procedure are represented by the leaves.

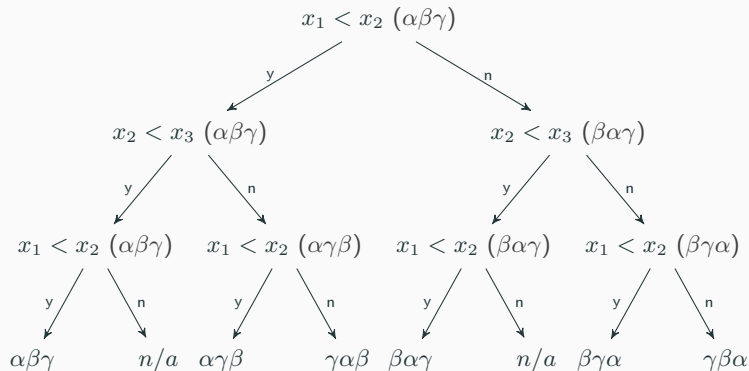
# Decision tree for sorting three items

Three items –  $(\alpha, \beta, \gamma)$



# Decision tree for bubble sort with three items

Three items –  $(\alpha, \beta, \gamma)$



# Heap sort

**Heap sort** uses a tree as part of the algorithm.

**This tree** is not a decision trees, it is for another purpose.

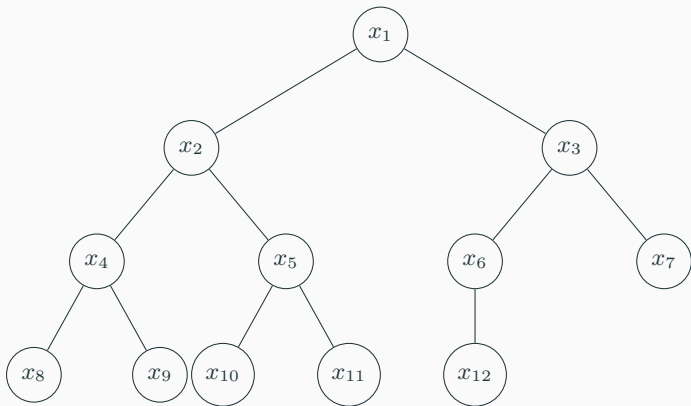
**Different** steps in the algorithm manipulate the tree.

**Worst case** – performs better than quick sort.

**They have** the same average performance though.

## Heapsort initial tree

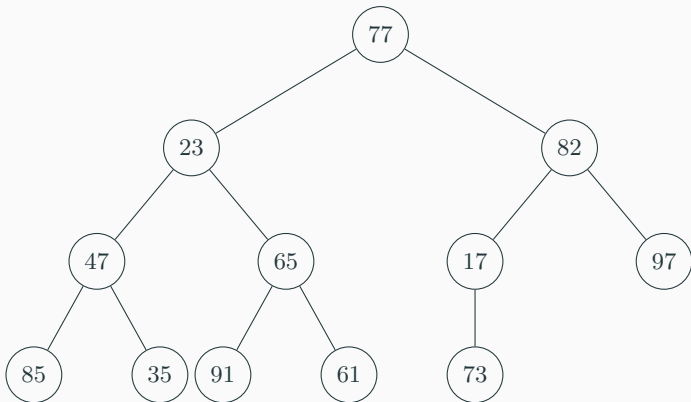
$(x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9, x_{10}, x_{11}, x_{12})$





## Heapsort initial tree

(77, 23, 82, 47, 65, 17, 97, 85, 35, 91, 61, 73)



## Transforming to a heap

**Suppose** the trees at  $x_{2r}$  and  $x_{2r+1}$  are already heaps.

**The vertex**  $x_r$  is their parent.

**Compare**  $x_r$  to  $x_{2r}$  and  $x_{2r+1}$ .

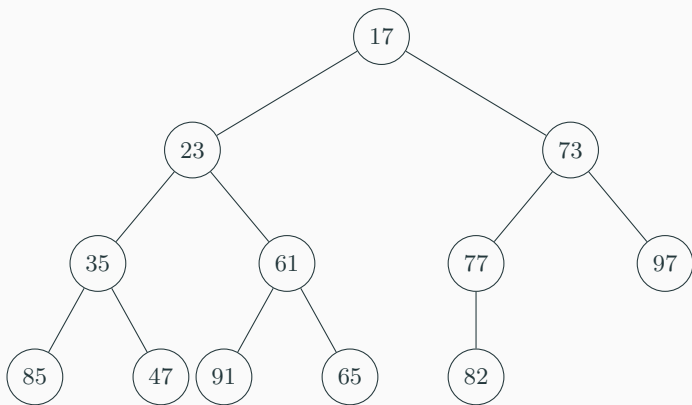
**If**  $x_r$  is smaller, do nothing.

**Otherwise** replace  $x_r$  with the smaller of it's children and fill the new vacancy with the smaller of  $x_r$  and the two children, and repeat if necessary.

**Start** this procedure at the last parent, and move backwards through the other parents.

## Heapsort - the heap

(17, 23, 73, 35, 61, 77, 97, 85, 47, 91, 65, 82)



$$x_r < x_{2r} \text{ and } x_r < x_{2r+1}$$

## Transforming to a sorted list

**Start** with a new empty list.

**Place** the root of the heap at the end of the list.

**Remove** the last leaf and place it at the root.

**Transform** the tree to a heap again. This is relatively easy since the subtrees at  $x_2$  and  $x_3$  are already heaps.

**Repeat** from step 2.

(17, 23, 35, 47, 61, 65, 73, 77, 82, 85, 91, 97)