Graph Theory

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Topics

Fundamentals

Trees

Graph databases

Algorithms

Fundamentals

Seven Bridges of Königsberg

Is it possible to walk through the city crossing each of the seven bridges once and only once?

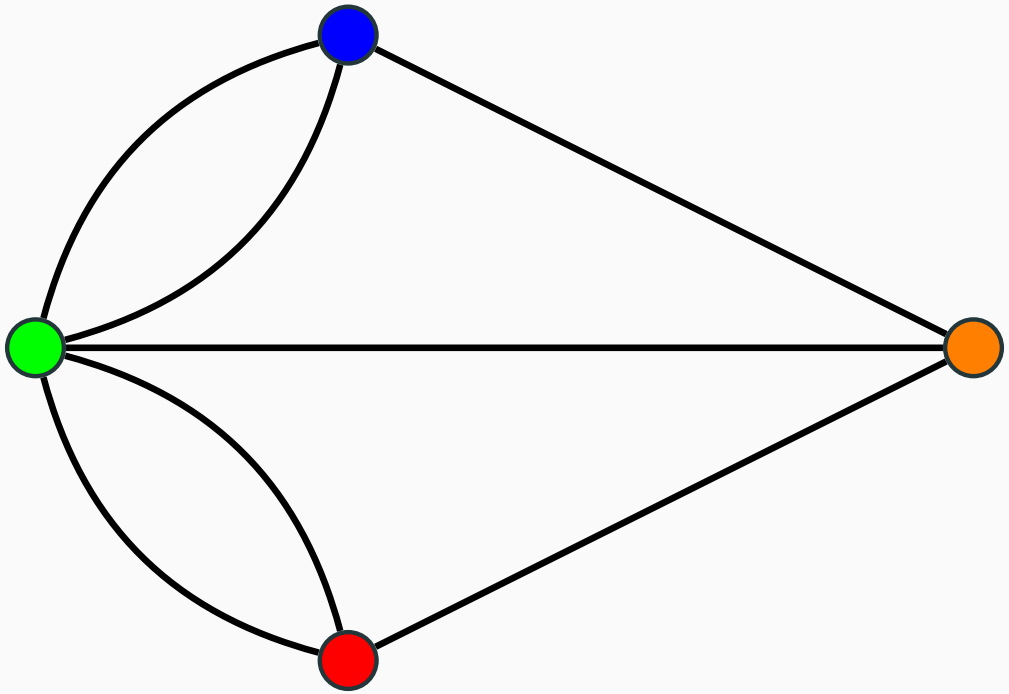
www.nature.com/nbt/journal/v29/n11

Leonhard Euler

- Born 1707 in Basel, Switzerland.
- Euler's identity: $e^{i\pi} + 1 = 0$.
- Solved the Bridges of Königsberg problem.
- It's not possible to cross all bridges once and once only.

https://en.wikipedia.org/wiki/Leonhard_Euler

Graph of Königsberg



Graph definition

Definition

A *graph* consists of a finite set V and a set E of 2-subsets of V.

Vertices – the elements of the set V are called vertices.

Edges – the elements of E are called edges.

$G = (V, E)$ – this is the way we write the graph G consists of the vertex set V and the edge set E.

Sets of Königsberg

$$V = \{Green, Blue, Orange, Red\}$$

$$\begin{aligned} E = \{ & \\ & \{Green, Blue\}, \{Green, Blue\}, \{Green, Red\}, \\ & \{Green, Red\}, \{Blue, Orange\}, \{Green, Orange\}, \\ & \{Red, Orange\} \\ \} & \end{aligned}$$

Adjanceny list

Green	Blue	Orange	Red
Blue	Green	Blue	Green
Orange	Orange	Green	Orange
Red		Red	

Defining different types of graphs

Our definition of a graph

The definition given above for a graph is not consistent with looped edges, directed edges or repeated edges. We only need to make small changes to the definition of a graph to allow for directed edges and repeated edges.

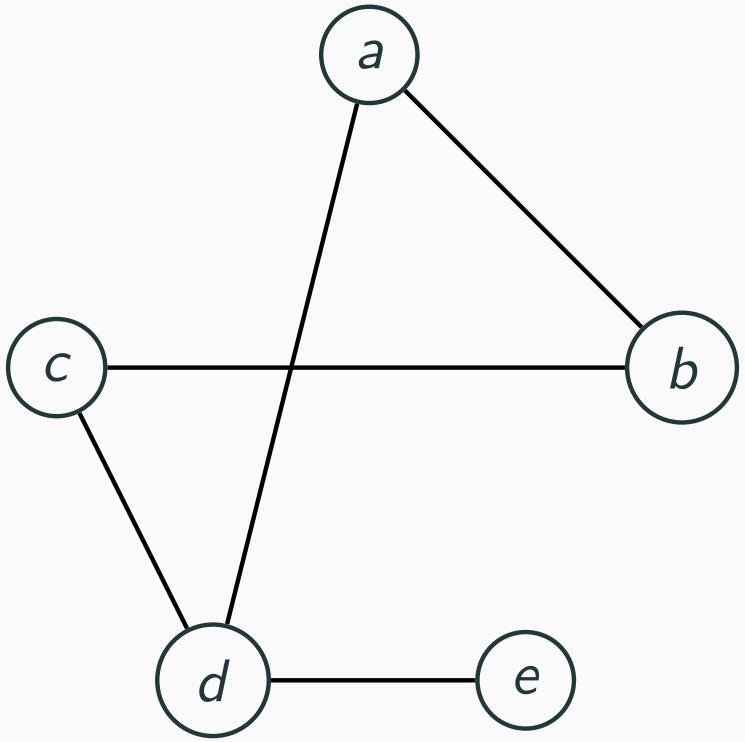
Repeated edges are edges that start and end at the same vertices.

Directed edges are edges where a direction is added.

Looped edges begin and end at the same vertex.

The application will determine the definition we want to use.

A better example

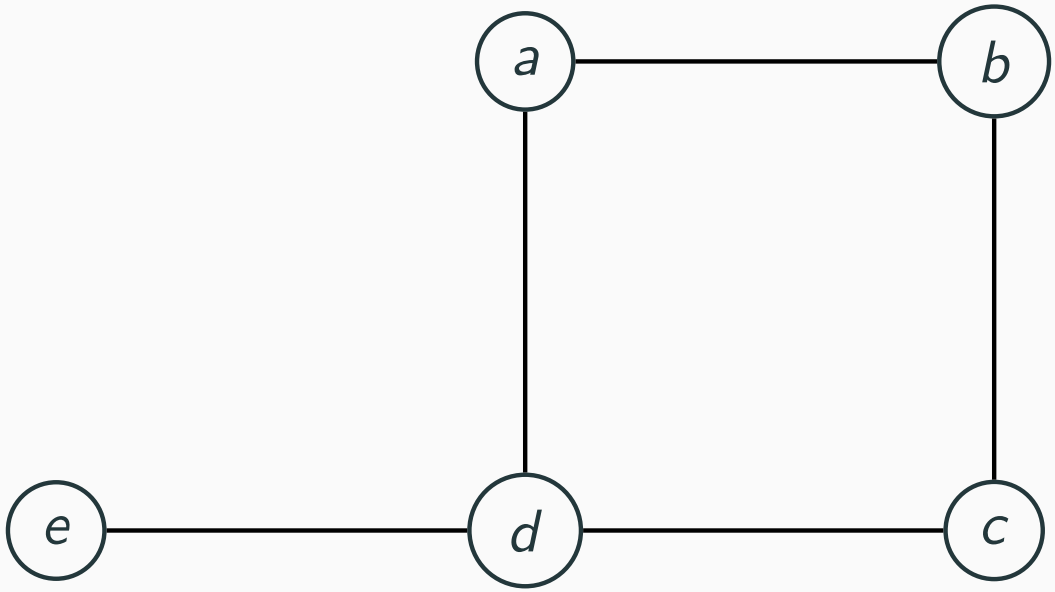


Exercise

Determine the vertex set, edge set and adacency list of this graph.

global.oup.com/booksites/content/9780198507185/

Another better example



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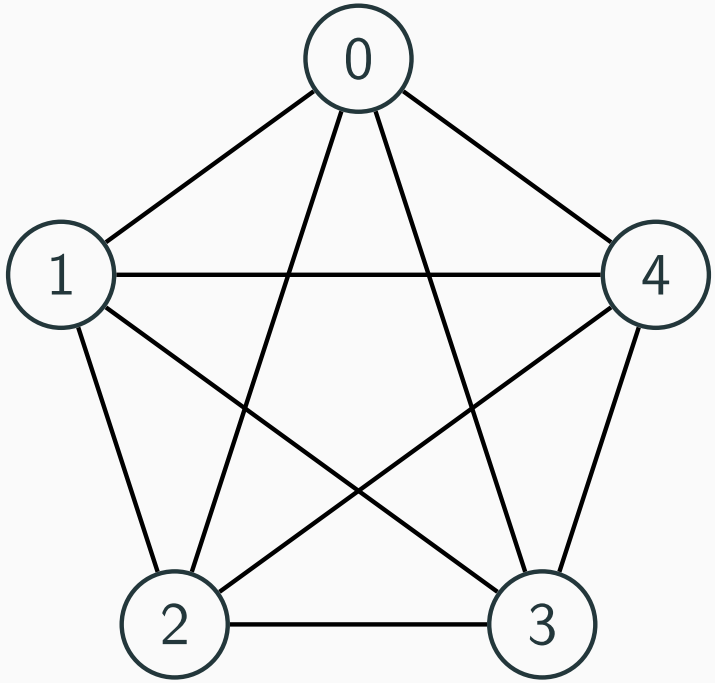
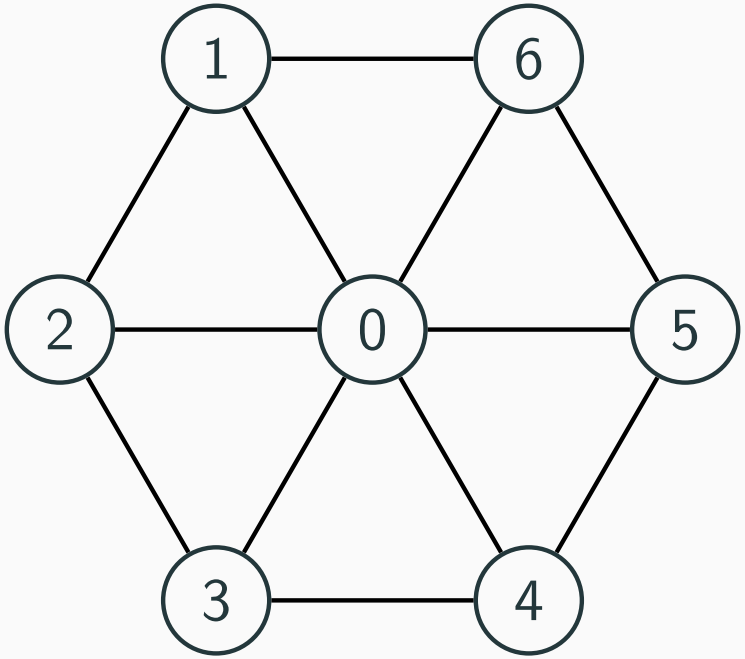
Complete graph K_n



Exercise

Determine the vertex set, edge set and adacency list of K_5.

Wheel graph W_n



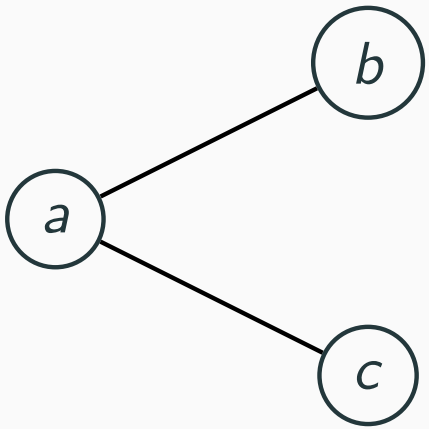
Exercise

Determine the vertex set, edge set and adacency list of W_6.

Degree of a vertex

Definition

The degree of a vertex is the number of edges that contain it.



The degree of the vertex a is 2.

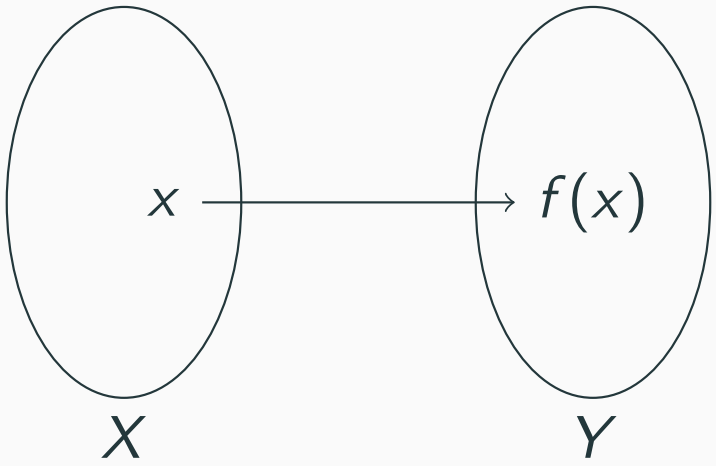
Exercise

For each of the vertices on the previous slide, determine its degree.

Functions

Definition

Suppose that X and Y are sets. We say we have a function f from X to Y if for each x in X we can specify a unique element in Y, which we denote by $f(x)$.

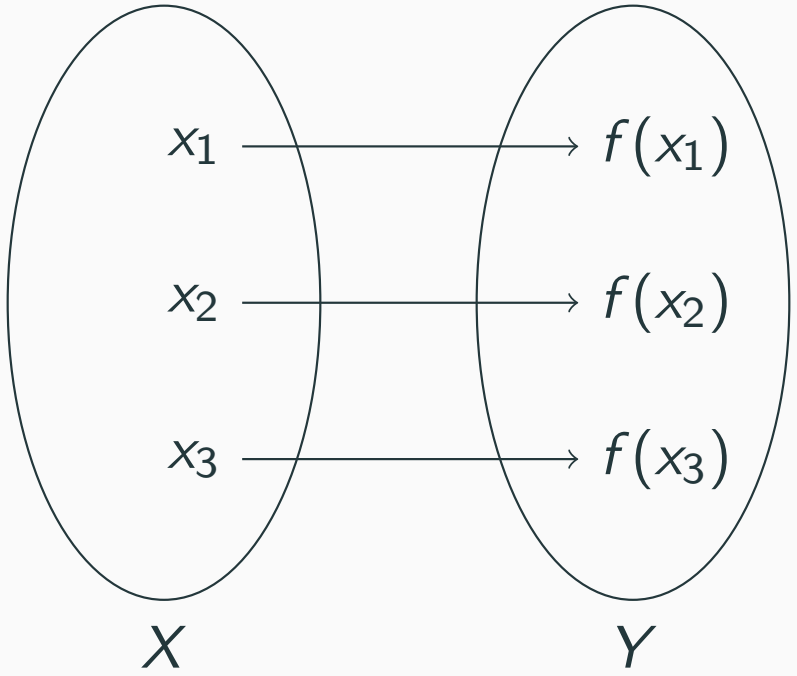


Bijections

Definition

A bijection is function f from a set X to a set Y where both of the following are true:

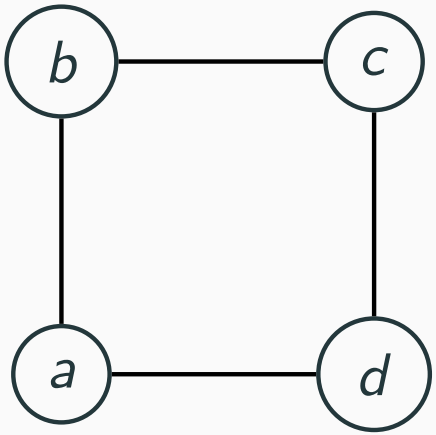
- every y in Y is a value $f(x)$ for at most one x in X.
- every y in Y is a value $f(x)$ for at least one x in X.

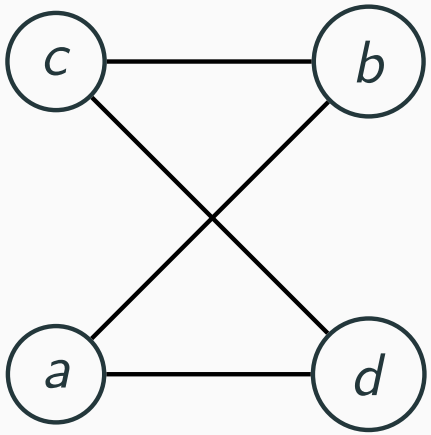


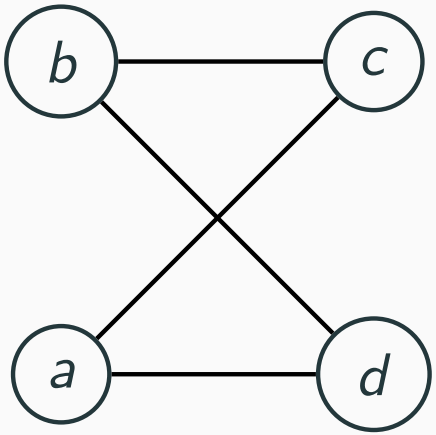
Isomorphism

Definition

Two graphs G_1 and G_2 are said to be isomorphic when there is a bijection α for the vertex set V_1 of G_1 to the vertex set V_2 of G_2 such that $\{\alpha(x), \alpha(y)\}$ is an edge of G_2 if and only if (x, y) is an edge of G_1.







Sum of degrees

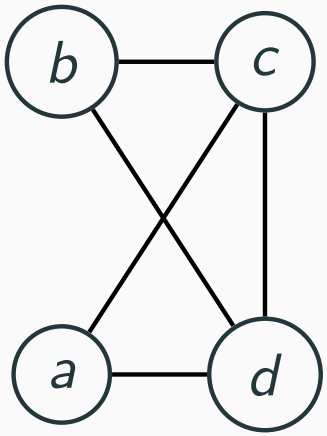
Theorem

The sum of the degrees of the vertices of a graph $G = (V, E)$ is equal to twice the number of edges:

$$\sum_{\mathrm{v} \in V} \delta(\mathrm{v}) = 2|E|$$

Proof.

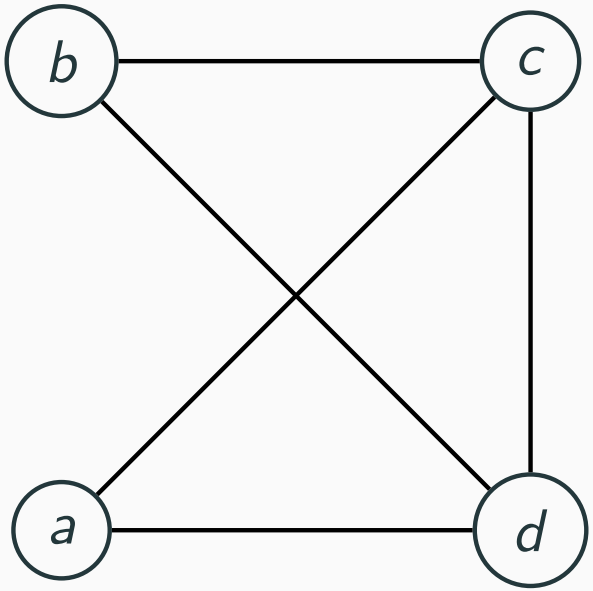
The degree $\delta(\mathrm{v})$ of a vertex v is equal to the number of edges incident on it. Every edge is incident on two vertices. So every edge contributes 1 to the degrees of two distinct vertices. Therefore every edge contributes 2 to the sum total of the degrees of all the vertices. □



Handshaking lemma

Definition

A vertex is an odd vertex if its degree is odd, and it is an even vertex if its degree is even. The set of all odd vertices is denoted V_o and the set of all even vertices is denoted V_e.



Exercise

Which of the above vertices are even, and which are odd?

Handshaking lemma

Lemma

The number of odd vertices $|V_o|$ in a graph is even.

Proof.

The sets V_o and V_e are disjoint (i.e. they don't have any elements in common.) Also, every vertex is either in V_o or V_e. Therefore $V = V_o \cup V_e$ and $|V| = |V_o| + |V_e|$.

Furthermore:

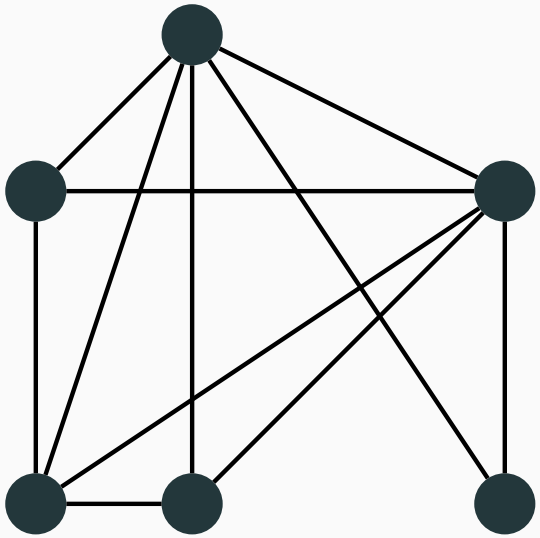
$$\sum_{\mathrm{v} \in V_o} \delta(\mathrm{v}) + \sum_{\mathrm{v} \in V_e} \delta(\mathrm{v}) = 2|E|$$

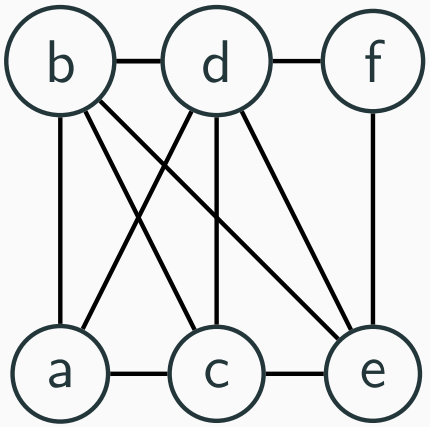
Both $2|E|$ and $\sum_{\mathrm{v} \in V_e} \delta(\mathrm{v})$ are even, so $\sum_{\mathrm{v} \in V_o} \delta(\mathrm{v})$ must be. Since $\delta(\mathrm{v})$ is odd for every v in V_o, this must mean that $|V_o|$ is even. □

Isomorphism: degrees

Exercise

Determine if these two graphs are isomorphic.



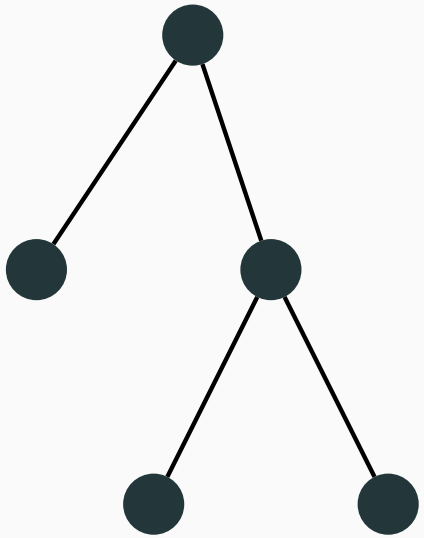


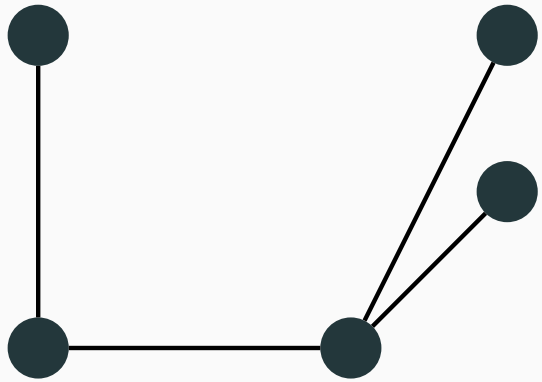
Trees

Definition

Tree

A *tree* is a graph where every pair of vertices has a path between them, and there are no cycles.





Levels

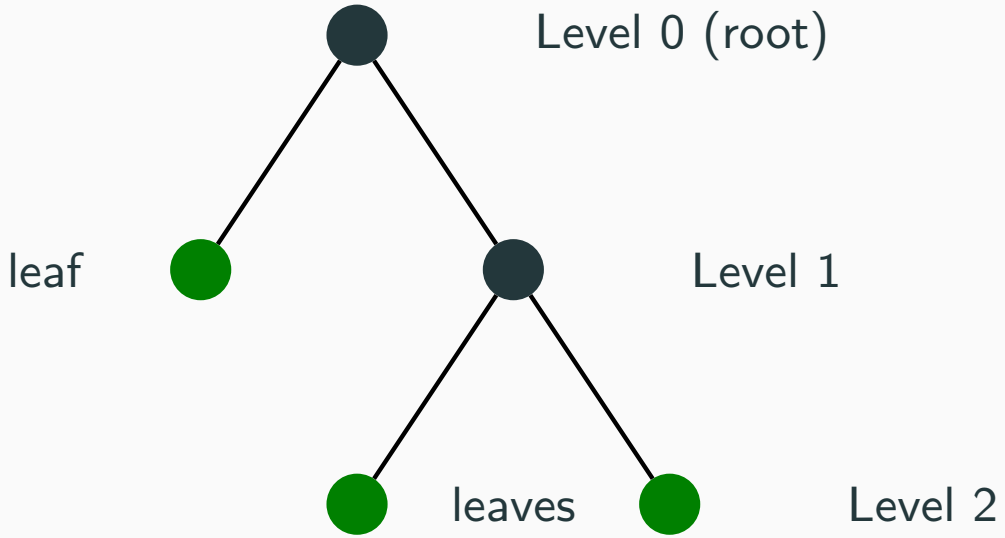
Root Specified vertex for some purpose.

Levels Root is at level 0, neighbours of the root are at level 1, their other neighbours at level 2, and so on.

Height h, where vertex at level h but not at level $h+1$.

Leaf Vertex at level i with no neighbours at level $i+1$.

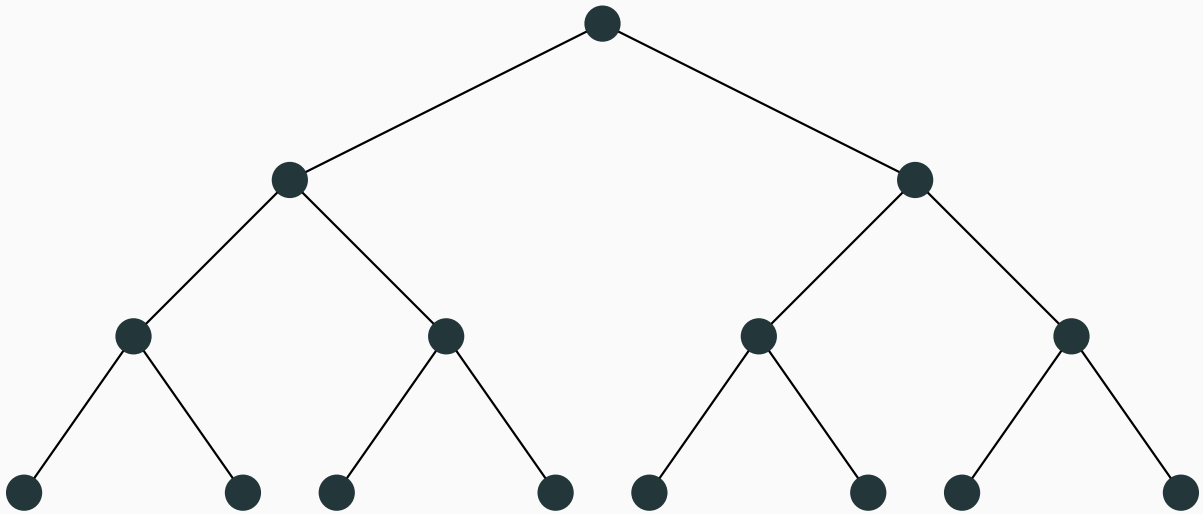
Internal vertex A non-leaf.



m-ary Rooted Tree

Definition

When a vertex at level i is connected to a vertex at level $i+1$ it's common to call the former the *father* and the latter the *son*. A rooted tree is m-ary if every father has the same number of sons. A 2-ary rooted tree is called a *binary tree*.



Isomorphic Rooted Trees

Definition

Two rooted trees are said to be *isomorphic* if there is an isomorphism between them which takes the root of one tree to the root of the other.

Height of an m-ary rooted tree

Theorem

The height h of an m-ary rooted tree with l leaves is at least $\log_m l$. That is:

$$h \geq \log_m l$$

Proof.

First note that: $h \geq log_m l \Leftrightarrow m^h \geq m^{log_m l} \Leftrightarrow m^h \geq l$.

So, we just need to show the number of leaves is at most m^h. For a tree of depth 0, there is only one vertex and $m^0 \leq 1$. Suppose we know that the theorem is true for trees of height $h-1$. Consider a tree of height h, with l leaves. We can create m trees of height $h-1$ from it by deleting the root. Each of these smaller trees has at most m^{h-1} leaves, and there are m of them. So the big tree has at most $m \times m^{h-1} = m^h$ leaves.

□

Graph databases

Digraph definition

Definition

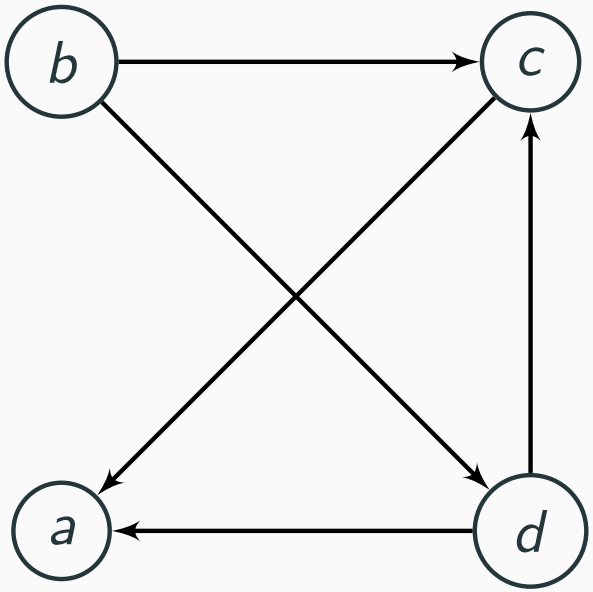
A *digraph* (short for directional graph) consists of a finite set V and a set E of ordered pairs of elements of V.

Degrees of vertices can now be split into in-degrees and out-degrees.

Walks, paths, cycles must be redefined.

Loops are allowed in the above definition, unless we rule them out.

Digraph example



Neo4j

- Neo4j is an open-source NoSQL graph database implemented in Java and Scala.
- Development started in 2003, it has been publicly available since 2007
- Available on GitHub.
- A graph is composed of two elements: a node and a relationship.

http://neo4j.com/developer/graph-database/

Cypher

- Cypher is a declarative graph query language.
- What to retrieve from a graph, not on how to retrieve it.
- Allows for expressive and efficient querying and updating of the graph store.
- Cypher borrows its structure from SQL.

http://neo4j.com/docs/stable/cypher-introduction.html

Cypher: Nodes

Create a node with the label User, and two properties:

```
1 CREATE (user:User { Id: 123, Name: "Jim" });
```

Find the node(s) with label User and their Id property being 123:

```
1 MATCH (user:User)
2 WHERE user.Id = 123
3 RETURN user;
```

neo4j.com/docs/stable/cypher-query-lang.html

Cypher: Relationships

Create a relationship with label FOLLOWS from user(s) with Id 123 to user(s) with Id 456:

```
1 MATCH (user1:User), (user2:User)
2 WHERE user1.Id = 123 AND user2.Id = 456
3 CREATE user1-[:FOLLOWS]->user2;
```

neo4j.com/docs/stable/cypher-query-lang.html

Cypher: Relationships and Nodes

Create a relationship with label INVITED from user(s) with Id 123 to a new user with Id 789 and Name Jack:

```
1 MATCH (invitee:User)
2 WHERE invitee.Id = 123
3 CREATE invitee-[:INVITED]->(invited:User {Id: 789,
4                                           Name: "Jack"});
```

neo4j.com/docs/stable/cypher-query-lang.html

Cypher: DELETE

Delete all nodes:

```
1 MATCH (x)
2 DELETE x;
```

neo4j.com/docs/stable/cypher-query-lang.html

Labels and properties

- Suppose we have nodes representing people.
- We give them the label People.
- We also want to identify each person as either Male or Female.
- Should we use Male and Female labels, or a Gender property?
- If you are going to use the person's gender in a lot of queries, a normal property will be relatively slow, so you should use a label.
- However, you can also index some of your properties to high-light them as important.

neo4j.com/docs/stable/cypher-query-lang.html

Cypher: shortestPath

Find the minimum number of hops betweem two nodes.

```
1 MATCH p=shortestPath(
2   (a:Actor {id: 1})-[*]-(b:Actor {id: 10})
3   )
4 RETURN p;
```

neo4j.com/docs/stable/cypher-query-lang.html

Algorithms

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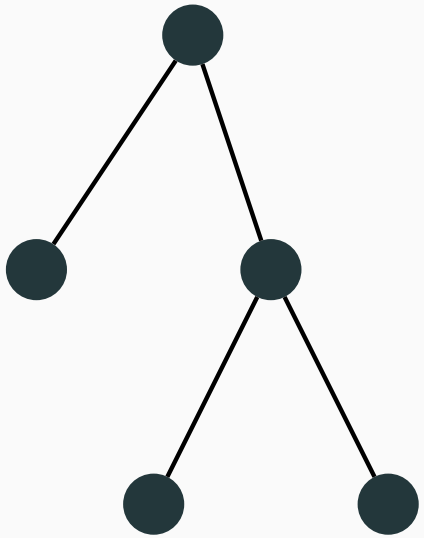

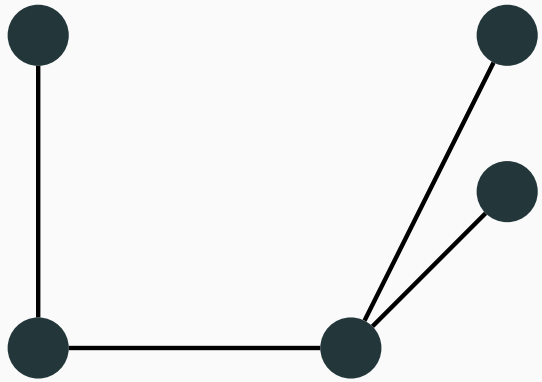

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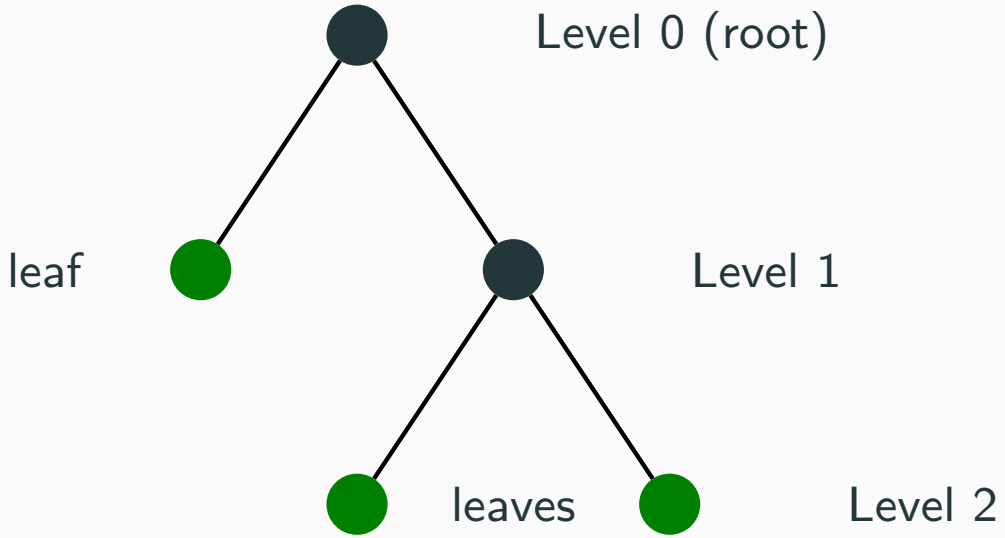

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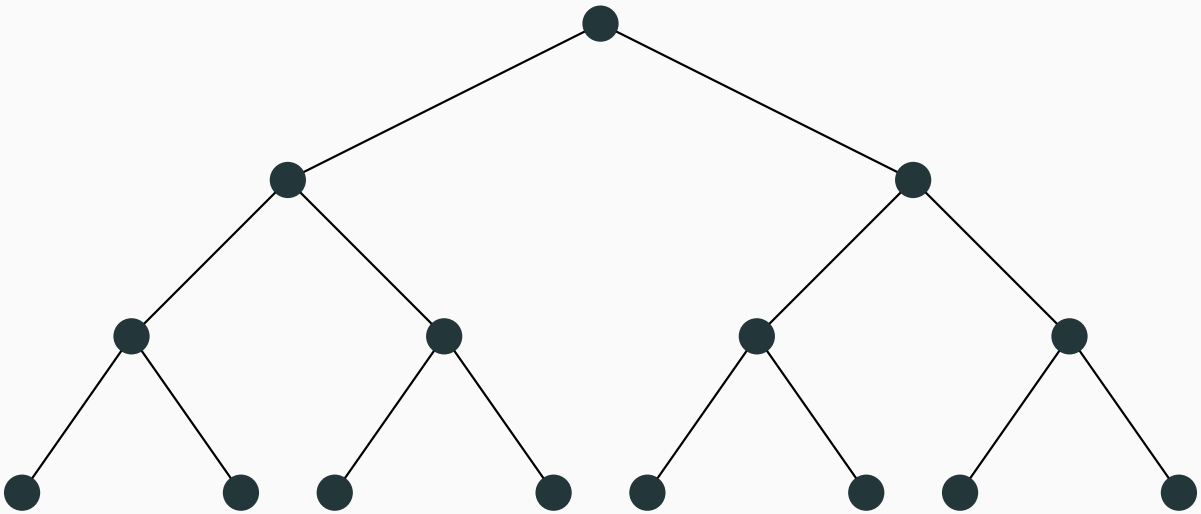

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□

Decision tree

Definition

A decision tree is a rooted tree where each vertex represents a decision. The possible results of each decision are represented by the edges connecting the vertex to the vertices at the next level down. Final outcomes of the procedure are represented by the leaves.