

**Lemma 1.** Given a timestamp list  $T = [t_1, t_2, \dots, t_n]$  ( $n \geq 1$ ) and an integer  $t \geq 1$ , let  $T' \subseteq T$  be a sublist no shorter than  $t$  whose frequency is equal to the  $t$ -frequency  $F_t(T)$ , there exists an integer  $i \in [1, n - t + 1]$  such that  $T' = [t_i, t_{i+1}, \dots, t_{i+|T'|-1}]$ .

*Proof.* Given a  $T' = [t_j, \dots, t_k]$  which is a sublist of  $T$ , and the elements in  $T'$  do not have consecutive subscripts in  $T$ , we can calculate that the local frequency derived from  $T'$  is  $\frac{|T'|}{t_k - t_j + 1}$ , however, if we insert missing elements between  $t_k$  and  $t_j$  into  $T$ , denoted by  $T''$ , such that the subscripts of the elements are consecutive in  $T$ . Then, the local frequency derived from  $T''$  is  $\frac{|T''|}{t_k - t_j + 1}$  which is obviously greater because  $T''$  have more elements.  $\square$

**Lemma 2.** Given a timestamp list  $T = [t_1, t_2, \dots, t_n]$  ( $n \geq 1$ ) and an integer  $t \geq 1$ ,  $F_t(T)$  is the reciprocal of the minimum slope of all  $t$ -lines between  $p'_i = (i - 1, t_i - 1) \in P_{start}$  and  $p_j = (j, t_j) \in P_{end}$  with  $j \geq i$ .

*Proof.* Based on Lemma 1, the elements in  $T'$  must have continuous subscripts in  $T$ , so  $F_t(u, v) = \{\max(\frac{|T'|}{t_k - t_j + 1}) | |T'| \geq t\} = \{\max(\frac{j-i+1}{t_j - t_i + 1}) | j - i + 1 \geq t\} = \{\frac{1}{\min(\frac{t_j - (t_i - 1)}{j - (i - 1)})} | j - i + 1 \geq t\} = \{\frac{1}{\min(\text{slp}(p'_i, p_j))} | j - i + 1 \geq t\}$ .  $\square$

**Lemma 3.** For an end point  $p_j$  and an integer  $t$ , any start point  $p'_i$  with  $1 \leq i \leq j - t + 1$  such that  $\text{slp}(p'_i, p_j) = \text{mslp}(\cdot, p_j)$  must satisfy that  $p'_i \in \text{UCC}_j$ .

*Proof.* For the  $p'_j$  which is not in the upper convex curve and the two points next to it  $p'_i$  and  $p'_k$  which are in the curve, based on  $\text{slp}(p'_i, p'_j) \leq \text{slp}(p'_j, p'_k)$  we consider three cases for the end point  $p_j$ . Firstly, if  $\text{slp}(p'_j, p_j) \geq \text{slp}(p'_j, p'_k)$ , then  $\text{slp}(p'_i, p_j) \leq \text{slp}(p'_j, p_j)$ . Secondly, if  $\text{slp}(p'_i, p'_j) \leq \text{slp}(p'_j, p_j) < \text{slp}(p'_j, p'_k)$ , then both  $\text{slp}(p'_i, p_j)$  and  $\text{slp}(p'_k, p_j)$  are less than  $\text{slp}(p'_j, p_j)$ . Third, if  $\text{slp}(p'_j, p_j) < \text{slp}(p'_i, p'_j)$ , then  $\text{slp}(p'_k, p_j) < \text{slp}(p'_j, p_j)$ . In either case,  $\text{slp}(p'_j, p_j)$  is not the minimum.  $\square$

**Lemma 4.** The slopes of the  $t$ -lines from points in  $\text{UCC}_j$  to endpoint  $p_j$  decreases monotonically first and then increases. (In boundary cases, it may be monotonically decreasing or increasing).

*Proof.* The slope of  $p_j$  to  $\text{UCC}_j$  can be denoted by  $k = \frac{y_j - y}{x_j - x}$ , its derivative is  $k' = \frac{k - \frac{dy}{dx}}{(x_j - x)^3} = \{\frac{k - \text{slp}(p'_i, p'_{i+1})}{(x_j - x)^3} | 1 \leq i < |T|\}$ . Consider three cases, firstly, if the initial  $k$  is greater than  $\text{slp}(p'_i, p'_{i+1})$ , i.e.,  $\text{slp}(p'_1, p_j) > \text{slp}(p'_1, p'_2)$ , then  $k$  is monotonically increasing because  $\text{slp}(p'_i, p'_{i+1})$  is decreasing which means  $k'$  is always greater than 0. In this case,  $p'_1$  is OSP. Secondly, if  $\text{slp}(p'_1, p_j) \leq \text{slp}(p'_1, p'_2)$ , thus  $k$  decreases initially until it reaches a certain point  $p'_i$  where  $\text{slp}(p'_i, p_j) > \text{slp}(p'_i, p'_{i+1})$ , and after that  $k$  starts to increase. In this case,  $p'_i$  is OSP. Third, if there is no point  $p'_i$  such that  $\text{slp}(p'_i, p_j) > \text{slp}(p'_i, p'_{i+1})$ , then  $k$  is monotonically decreasing, and  $p'_{|T|}$  is OSP.  $\square$

**Lemma 5.** Those start points other than  $p'_{j-t+1}$  that do not exist in  $\text{UCC}_{j-1}$  cannot exist in  $\text{UCC}_j$ .

*Proof.* Considering end points  $p_j$  and  $p_{j-1}$ , since timestamps in  $T$  are sorted incrementally, thus  $p_j$  is to the upper right of  $p_{j-1}$ , which means the proof in Lemma 3 of  $p_{j-1}$  can be applied to  $p_j$ .  $\square$

**Lemma 6.** When iterating to a new end point  $p_j$ , points before its OSP can not be the start point to obtain the minimum slope, which means that they should be removed from  $\text{UCC}$ .

*Proof.* Suppose that the OSP of  $p_j$  is  $p'_l$ , and the OSP of a new end point  $p_{j'}$  (after  $p_j$ ) is  $p'_{l'}$ , while the  $x$ -value of  $p'_{l'}$  is less than that of  $p'_l$ . According to the property of optimal start point and the definition of upper convex curve, there is  $\text{slp}(p'_{l'}, p_{j'}) > \text{slp}(p'_{l'}, p'_{l'+1}) > \text{slp}(p'_{l-1}, p'_l) > \text{slp}(p'_l, p_j)$ , which means  $p'_{l'}$  can not be the start point to obtain the minimum slope.  $\square$

**Lemma 7.** Given a temporal graph  $\mathcal{G}$ ,  $\mathcal{C}_{t,f'}^k \subseteq \mathcal{C}_{t,f}^k$  if  $f' > f$ .

*Proof.* Assume a vertex  $u \in \mathcal{C}_{t,f'}^k$ , we have  $|\mathcal{N}_{t,f'}(u)| \geq |\mathcal{N}_{t,f}(u)| \geq k$  if  $f' > f$ . Thus,  $u$  must also be present in  $\mathcal{C}_{t,f}^k$ , namely,  $\mathcal{C}_{t,f'}^k$  contains all vertices in  $\mathcal{C}_{t,f}^k$ .  $\square$

**Lemma 8.** Given a vertex  $v$  and an integer  $k$ , we have  $\text{cf}(v, k, t) \geq \text{cf}(v, k, t')$  if  $t < t'$ .

*Proof.* We prove by contradiction. Supposing that  $t < t'$  and there exists a vertex  $v$  satisfying  $\text{cf}(v, k, t) < \text{cf}(v, k, t')$ . We have  $v \in \mathcal{C}_{t',\text{cf}(v,k,t')}^k$  and  $v \notin \mathcal{C}_{t,\text{cf}(v,k,t')}^k$  since  $\text{cf}(v, k, t) < \text{cf}(v, k, t')$ . However,  $\mathcal{C}_{t',\text{cf}(v,k,t')}^k \subseteq \mathcal{C}_{t,\text{cf}(v,k,t')}^k$  due to  $t < t'$  which is a contradiction.  $\square$

**Lemma 9.** Given  $k$  and  $t$ , let  $N_t^k(u)$  be the set of neighbors of a vertex  $u$  such that for each  $v \in N_t^k(u)$  we have  $\text{cf}(v, k, t) \geq \text{cf}(u, k, t - 1)$  and  $F_t(u, v) \geq \text{cf}(u, k, t - 1)$ . Consequently,  $\text{cf}(u, k, t) < \text{cf}(u, k, t - 1)$  if and only if  $|N_t^k(u)| < k$ .

*Proof.* ( $\Rightarrow$ ) We prove by contradiction, suppose  $\text{cf}(u, k, t) < \text{cf}(u, k, t - 1)$  and  $|N_t^k(u)| \geq k$ , then  $u$  have more than  $k$  neighbors  $v$  that satisfying  $\text{cf}(v, k, t) \geq \text{cf}(u, k, t - 1)$  and  $F_t(u, v) \geq \text{cf}(u, k, t - 1)$ , which means  $u \in \mathcal{C}_{t,\text{cf}(u,k,t-1)}^k$ . Based on the definition of core frequency, we have  $\text{cf}(u, k, t - 1) \leq \text{cf}(u, k, t)$ , which contradicts to the assumption.

( $\Leftarrow$ ) We also prove by contradiction, suppose that  $|N_t^k(u)| < k$  and  $\text{cf}(u, k, t) \geq \text{cf}(u, k, t - 1)$ . Firstly,  $u \in \mathcal{C}_{t,\text{cf}(u,k,t-1)}^k$  because  $\text{cf}(u, k, t) \geq \text{cf}(u, k, t - 1)$ . However, vertex  $u$  have less than  $k$  neighbors who have more than  $k(t, \text{cf}(u, k, t - 1))$ -neighbors due to  $|N_t^k(u)| < k$ , which is a contradiction.  $\square$

**Lemma 10.** Given  $t$  and  $k$ ,  $\text{cf}(u, k, t)$  equals to the maximum float  $f \leq 1$  such that the number of neighbor vertices (denoted by  $v$ ) of  $u$ , which satisfy  $\text{cf}(v, k, t) \geq f$  and  $F_t(u, v) \geq f$ , is no less than  $k$ .

*Proof.* We consider the value of  $f'$  in two cases. First, if  $f' > f$ , then  $u$  does not exist in  $\mathcal{C}_{t,f'}^k$  due to the maximality of  $f$ , so  $\text{cf}(u, k, t) \neq f'$ . Second, if  $f' < f$ , Since it does not satisfy the

maximality in the definition of core frequency,  $\text{cf}(u, k, t) \neq f'$ .  $\square$

**Lemma 11.** *For any skyline  $(k, t, f)$ -core,  $f$  is a core frequency of at least one vertex, and thus is preserved in the CF-Index.*

*Proof.* Assuming there is a  $(k, t, f)$ -core, whose  $f$  is not preserved in the CF-Index, then there must exist a vertex  $v$  that satisfies  $\text{cf}(v, k, t) > f$ , so  $(k, t, \text{cf}(v, k, t))$ -core dominates  $(k, t, f)$ -core.  $\square$

**Lemma 12.** *For any skyline  $(k, t, f)$ -core, let  $(t', f)$  be preserved in the partition of  $k$ , there must be a successive pair  $(t + 1, f')$  following  $(t', f)$  in a record of the partition.*

*Proof.* For two consecutive entries in the partition of  $k$  and  $v$ ,  $(t_1, f_1)$  and  $(t_2, f_2)$ , since for any  $t_1 \leq t < t_2 - 1$ , core frequency satisfies  $\text{cf}(v, k, t) = \text{cf}(v, k, t_2 - 1) = f_1$ , so  $(k, t_2 - 1, f_1)$ -core dominates any other  $(k, t, f_1)$ -core, which means  $t = t_2 - 1, t' = t_1, f = f_1, f' = f_2$  in the Lemma.  $\square$

**Lemma 13.** *Let  $\text{Skyline}_k$  contain all pairs of  $(t, f)$ . If a triple  $(k_1, t_1, f_1)$  is dominated by another  $(k_2, t_2, f_2)$ , it can only be one of the two cases: 1)  $k_1 = k_2, t_1 < t_2, f_1 \leq f_2$ . 2)  $k_1 < k_2, t_1 = t_2, f_1 = f_2$ .*

*Proof.* Consider the following cases, and we will prove in turn the reasons why they should not be taken into account.

1)  $k_1 = k_2, t_1 = t_2, f_1 \leq f_2$  2)  $k_1 < k_2, t_1 = t_2, f_1 < f_2$  3)  $k_1 < k_2, t_1 < t_2, f_1 \leq f_2$  For the first case, since during the traversal of CF-index to obtain the candidate skyline tuples in the algorithm, for fixed  $k$  and  $t$ , only the largest  $f$  is retained, therefore case 1 does not exist. For the second case, assuming this situation exists, then there must exist a point  $u$  that satisfies  $\text{cf}(u, k_2, t_2) = f_2$ . Due to the monotonicity of core frequency, we have  $\text{cf}(u, k_1, t_1) \geq f_2$ , which contradicts the assumptions of case 2. For the third case, we similarly assume that it holds, then there must exist a point  $u$  that satisfies  $\text{cf}(u, k_2, t_2) = f_2$  and  $\text{cf}(u, k_1, t_2) = f_3$  in which  $f_3 \geq f_2$ . So there exists a tuple  $(k_1, t_4, f_4)$  in the candidate skyline tuples that satisfies  $t_4 \geq t_2 > t_1, f_4 \geq f_3 \geq f_1$ , which means  $(k_1, t_1, f_1)$  have been removed in the first situation of the algorithm.  $\square$