

Little Piplup

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1 Settings

1.1 C++

2 Data Structures

2.1 Segment Tree

To deal with queries on intervals, we use segment tree.

```
int arr[SIZE];
int tree[TREE_SIZE];
int makeTree(int left,int right,int node)
{
    if (left == right)
        return tree[node] = arr[left];
    int mid = (left + right) / 2;
    tree[node] += makeTree(left, mid, node * 2);
    tree[node] += makeTree(mid + 1,right, node * 2 +1);
    return tree[node];
}
```

Updating segment Tree

```
void update(int left,int right,int node, int change_node ,int diff)
{
    if (!(left <= change_node &&change_node <= right))
        return; //No effect on such nodes.
    tree[node] += diff; // This part must be changed with tree function.
    if (left != right)
    {
        int mid = (left + right) / 2;
        update(left, mid, node * 2, change_node, diff);
        update(mid +1,right, node * 2 +1, change_node, diff);
    }
}
```

Answering queries via segment tree

```
/*
Our Search range : start to end
Node has range left to right
We may answer query in O(log n) time.
*/
int Query(int node, int left, int right, int start, int end)
{
    if (right < start || end < left)
```

```
        return 0; //Node is out of range
    if (start <= left && right <= end)
        return tree[node]; //If node is completely in range
    int mid = (left + right) / 2;
    return Query(node * 2, left, mid, start, end)
        +Query(node*2+1,mid+1,right,start,end);
}
```

2.2 Fenwick Tree

2.3 Disjoint Set Union

3 Mathematics

3.1 Useful Mathematical Formula

- Catalan Number : Number of valid parantheses strings with n pairs

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

3.2 Number of Integer Partition

```
def partitions(n):
    parts = [1]+[0]*n
    for t in range(1, n+1):
        for i, x in enumerate(range(t, n+1)):
            parts[x] += parts[i]
    return parts[n]
```

3.3 Extended Euclidean Algorithm

```
int Extended_Euclid(int a, int b, int *x, int *y)
{
    if (a == 0)
    {
        *x = 0;
        *y = 1;
        return b;
    }
    int x1, y1;
    int EEd = Extended_Euclid(b%a, a, &x1, &y1);
    *x = y1 - (b/a) * x1;
    *y = x1;
    return EEd;
}
```

3.4 Fast Modulo Exponentiation

Calculating $x^y \bmod p$ in $\mathcal{O}(\log y)$ time.

```
/*
Fast Modulo Exponentiation algorithm
Runs on  $\mathcal{O}(\log y)$  time,
calculate  $x^y \bmod p$ 
*/

ll modpow(ll x, ll y, ll p)
```

```
{
    ll res = 1;
    x = x % p;
    while (y > 0)
    {
        if (y & 1)
            res = (res*x) % p;
        y = y>>1;
        x = (x*x) % p;
    }
    return res;
}
```

3.5 Miller-Rabin Primality Testing

Base values of a chosen so that results are tested to be correct up to 10^{14} .

```
bool MRwitness(ll n, ll s, ll d, ll a)
{
    ll x = modpow(a, d, n);
    ll y = -1;

    while (s)
    {
        y = (x * x) % n;
        if (y == 1 && x != 1 && x != n-1)
            return false;
        x = y;
        s--;
    }
    return (y==1);
}

bool Miller_Rabin(ll n)
{
    if (n<2)
        return false;
    if (n == 2 || n == 3 || n == 5 || n == 7 || n == 11 || n == 13 || n == 17)
        return true;
    if (n%2 == 0 || n%3 == 0 || n%5 == 0)
        return false;
    ll d = (n-1) / 2;
    ll s = 1;
    while (d%2==0)
```

```

{
    d /= 2;
    s++;
}
int candidate[7] = {2,3,5,7,11,13,17};
bool result = true;
for (auto i : candidate)
{
    result = result & MRwitness(n,s,d,i);
    if (!result)
        break;
}
return result;
}

```

3.6 Pollard-Rho Factorization

```

ll PollardRho(ll n)
{
    srand (time(NULL));
    if (n==1)
        return n;
    if (n % 2 == 0)
        return 2;
    ll x = (rand()%(n-2))+2;
    ll y = x;
    ll c = (rand()%(n-1))+1;
    ll d = 1;
    while (d==1)
    {
        x = (modpow(x, 2, n) + c + n)%n;
        y = (modpow(y, 2, n) + c + n)%n;
        y = (modpow(y, 2, n) + c + n)%n;
        d = gcd(abs(x-y), n);
        if (d==n)
            return PollardRho(n);
    }
    return d;
}

```

3.7 Euler Totient

Calculating number of integers below n which is coprime with n .

```

#define ll long long
ll euler_phi(ll n)
{
    ll p=2;
    ll ephi = n;
    while(p*p<=n)
    {
        if (n%p == 0)
            ephi = ephi/p * (p-1);
        while(n%p==0)
            n/=p;
        p++;
    }
    if (n!=1)
    {
        ephi /= n;
        ephi *= (n-1);
    }
    return ephi;
}

```

3.8 Modular Multiplicative Inverse

```

ll modinv(ll x, ll p)
{
    return modpow(x,p-2,p);
}

```

4 Geometry

4.1 Closest Pair Problem

4.2 Smallest Enclosing Circle

4.3 Convex Hull

4.4 Intersection of Line Segment

5 Graphs

5.1 Topological Sorting

Topological sorting with dfs

```
vector<int> graph[V];
bool visited[V];
vector<int> sorted;

void dfs(int root)
{
    visited[root] = 1;
    for (auto it:graph[root])
    {
        if (!visited[it])
            dfs(it);
    }
    sorted.push_back(root);
}

int main()
{
    int n, m;
    scanf("%d%d",&n,&m);
    for (int i = 0; i<m; i++)
    {
        int small, big;
        scanf("%d%d",&small,&big);
        graph[small].push_back(big);
    }
    for (int i = 1; i<=n; i++)
        if (!visited[i])
            dfs(i);
    reverse(sorted.begin(),sorted.end()); // must reverse!
}
```

5.2 Dijkstra

$\mathcal{O}(E \log V)$ Single-Start-Shortest-Path.

Not working for graph with minus weight.

```
const int INF = 987654321;
const int MX = V+something;
struct Edgeout
{
```

```

    int dest, w;
    bool operator<(const Edgeout &p) const
    {
        return w > p.w;
    }
};

vector <Edgeout> edgelist[MX];
int V, E, start;
int dist[MX];

bool relax(Edgeout edge, int u)
{
    bool flag = 0;
    int v = edge.dest, w = edge.w;
    if (dist[v] > dist[u] + w && (dist[u] != INF))
    {
        flag = true;
        dist[v] = dist[u] + w;
    }
    return flag;
}

int dijkstra()
{
    fill(dist, dist + MX, INF);
    dist[start] = 0;
    priority_queue<Edgeout> pq;
    pq.push({start, 0});
    while(!pq.empty())
    {
        Edgeout x = pq.top();
        int v = x.dest, w = x.w;
        pq.pop();
        if (w > dist[v])
            continue;
        for (auto ed : edgelist[v])
            if (relax(ed, v))
                pq.push({ed.dest, dist[ed.dest]});
    }
}

```

5.3 Bellman Ford

$\mathcal{O}(EV)$ Single-Start-Shortest-Path.

Not working for graph with minus cycle → must detect.

```

struct Edge
{
    int u, v, w;
};

vector <Edge> edgelist;
int V, E;
int dist[V+1];

bool relax_all_edge()
{
    bool flag = false;
    for (auto it:edgelist)
    {
        int u = it.u, v = it.v, w = it.w;
        if (dist[v] > dist[u] + w && (dist[u] != INF))
        {
            flag = true;
            dist[v] = dist[u] + w;
        }
    }
    return flag;
}

int bellman_ford()
{
    fill(dist, dist + V + 2, INF);
    dist[1] = 0;
    for (int i = 0; i < V - 1; i++)
    {
        relax_all_edge();
    }
    if (relax_all_edge())
        return -1;
    else
        return 0;
}

```

5.4 Floyd-Warshall

Works on adjacency matrix, in $\mathcal{O}(V^3)$.

```
int d[120][120];
int n;
void Floyd_Warshall()
{
    for (int i = 1; i<=n; i++)
        for (int j = 1; j<=n; j++)
            d[j][k] = MIN(d[j][k], d[j][i]+d[i][k]);
}
```

5.5 Bipartite Checking

```
vector <int> graph[20200];
vector <pair <int, int>> edge;
bool visited[202020];
int color[202020];

void dfs(int root)
{
    for (auto it:graph[root])
    {
        if (!visited[it])
        {
            visited[it] = true;
            color[it] = color[root]%2+1;
            dfs(it);
        }
    }
}

bool is_bipartite(vector <int> &mygraph, int v, int e)
{
    for (int i = 1; i<=v; i++)
    {
        if (!visited[i])
        {
            visited[i] = 1;
            color[i] = 1;
            dfs(i);
        }
    }
    for (int i = 0; i<e; i++)
```

```
        if (color[edge[i].first]==color[edge[i].second])
            return false;
    return true;
}
```

5.6 MST Kruskal Algorithm

5.7 MST Prim Algorithm

5.8 Strongly Connected Component

5.9 Ford-Fulkerson Algorithm

6 Dynamic

6.1 Largest Sum Subarray

Computes sum of largest sum subarray in $\mathcal{O}(N)$

```
void consecsum(int n)
{
    dp[0] = number[0];
    for (int i = 1; i < n; i++)
        dp[i] = MAX(dp[i-1]+number[i], number[i]);
}

int maxsum(int n)
{
    consecsum(n);
    int max_sum = -INF;
    for (int i = 0; i < n; i++)
        dp[i] > max_sum ? max_sum = dp[i] : 0;
    return max_sum;
}
```

6.2 Knapsack

6.3 Longest Common Subsequence

```
//input : two const char*
//output : their LCS, in c++ std::string type
string lcsf(const char *X, const char *Y)
{
    int m = (int)strlen(X);
    int n = (int)strlen(Y);
    int L[m+1][n+1];
    for (int i=0; i<=m; i++)
    {
        for (int j=0; j<=n; j++)
        {
            if (i == 0 || j == 0)
                L[i][j] = 0;
            else if (X[i-1] == Y[j-1])
                L[i][j] = L[i-1][j-1] + 1;
            else
                L[i][j] = max(L[i-1][j], L[i][j-1]);
        }
    }
}
```

```
int index = L[m][n];

char lcsstring[index+1];
lcsstring[index] = 0;

int i = m, j = n;
while (i > 0 && j > 0)
{
    if (X[i-1] == Y[j-1])
    {
        lcsstring[index-1] = X[i-1];
        i--; j--; index--;
    }
    else if (L[i-1][j] > L[i][j-1])
        i--;
    else
        j--;
}
string lcsstr = lcsstring;
return lcsstr;
}
```

6.4 Edit Distance

7 String

7.1 KMP Algorithm

7.2 Manacher's Algorithm

7.3 Trie

7.4 Rabin-Karp Hashing

7.5 Aho-Corasick Algorithm

8 Miscellaneous

8.1 Useful Bitwise Functions in C++

```
int __builtin_clz(int x); // number of leading zero
int __builtin_ctz(int x); // number of trailing zero
int __builtin_clzll(ll x); // number of leading zero
int __builtin_ctzll(ll x); // number of trailing zero
int __builtin_popcount(int x); // number of 1-bits in x
int __builtin_popcountll(ll x); // number of 1-bits in x

lsb(n): (n & -n); // last bit (smallest)
floor(log2(n)): 31 - __builtin_clz(n | 1);
floor(log2(n)): 63 - __builtin_clzll(n | 1);

// compute next perm. ex) 00111, 01011, 01101, 01110, 10011, 10101..
ll next_perm(ll v)
{
    ll t = v | (v-1);
    return (t + 1) | (((~t & -~t) - 1) >> (__builtin_ctz(v) + 1));
}
```

8.2 List of Useful Numbers

< 10 ^k	prime	# of prime	< 10 ^k	prime
1	7	4	10	9999999967
2	97	25	11	9999999977
3	997	168	12	99999999989
4	9973	1229	13	99999999971
5	99991	9592	14	999999999973
6	999983	78498	15	999999999989
7	9999991	664579	16	9999999999937
8	99999989	5761455	17	9999999999997
9	999999937	50847534	18	99999999999989

9 Debugging Checkpoints