

# Subnetting and Classless Addressing



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# Acknowledgements

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- ❑ Some pictures used in this presentation were obtained from the Internet
- ❑ The instructor used the following references
  - Larry L. Peterson and Bruce S. Davie, Computer Networks: A Systems Approach, 5th Edition, Elsevier, 2011
  - Andrew S. Tanenbaum, Computer Networks, 5th Edition, Prentice-Hall, 2010
  - James F. Kurose and Keith W. Ross, Computer Networking: A Top-Down Approach, 5th Ed., Addison Wesley, 2009
  - Larry L. Peterson's (<http://www.cs.princeton.edu/~llp/>) Computer Networks class web site

# Outline

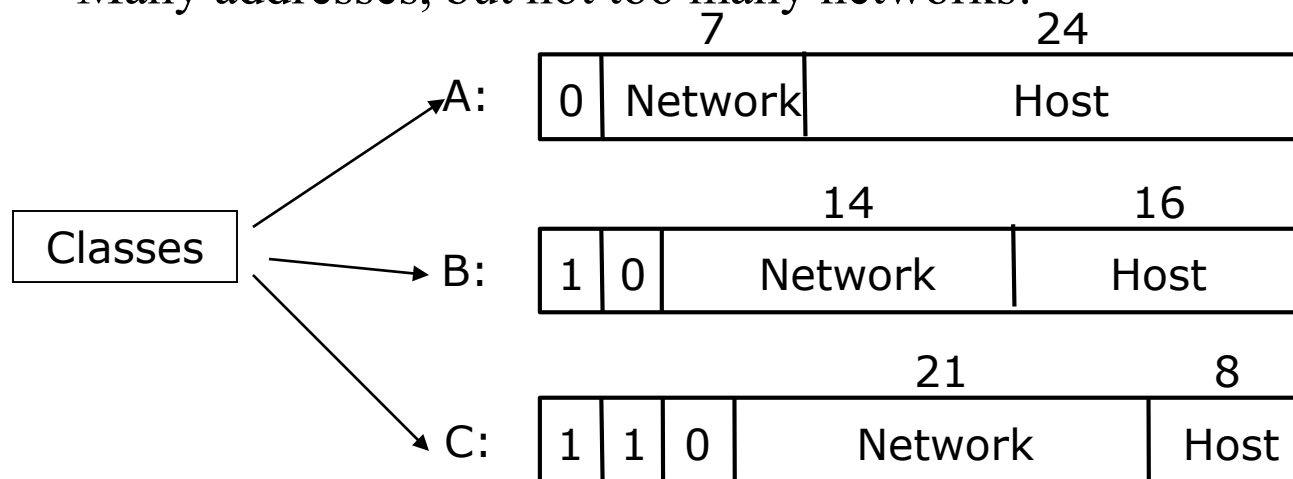
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- ❑ Problem to scale to global network
  - Many networks organized in hierarchical manner
  - Scarcity of IP address
- ❑ Solution
  - Subnetting
  - Supernetting (classless routing)

# Is $2^{32}$ too small a number?

## □ IPv4 address

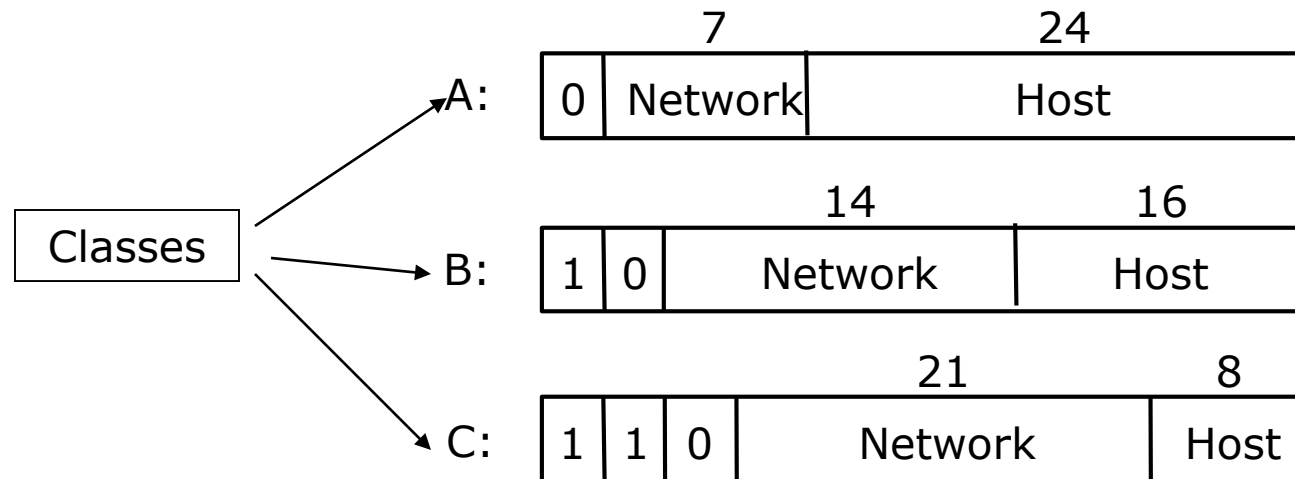
- 32 bit integers  $\sim 2^{32} = 4,294,967,296$
- Many addresses, but not too many networks!



- Testimony: <http://www.iana.org/assignments/ipv4-address-space/>
- Examples
  - A network of two nodes needs a class C network
  - A network of 256 nodes needs a class B network

# Can the number of networks be too many?

- ❑ How many class B networks are there?
- ❑ Potentially how big a routing table can be?



# Subnetting

- ❑ Add another level to address/routing hierarchy: *subnet*
- ❑ *Subnet masks* define variable partition of host part of class A and B addresses
- ❑ Subnets visible only within site

Network number	Host number
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Class B address

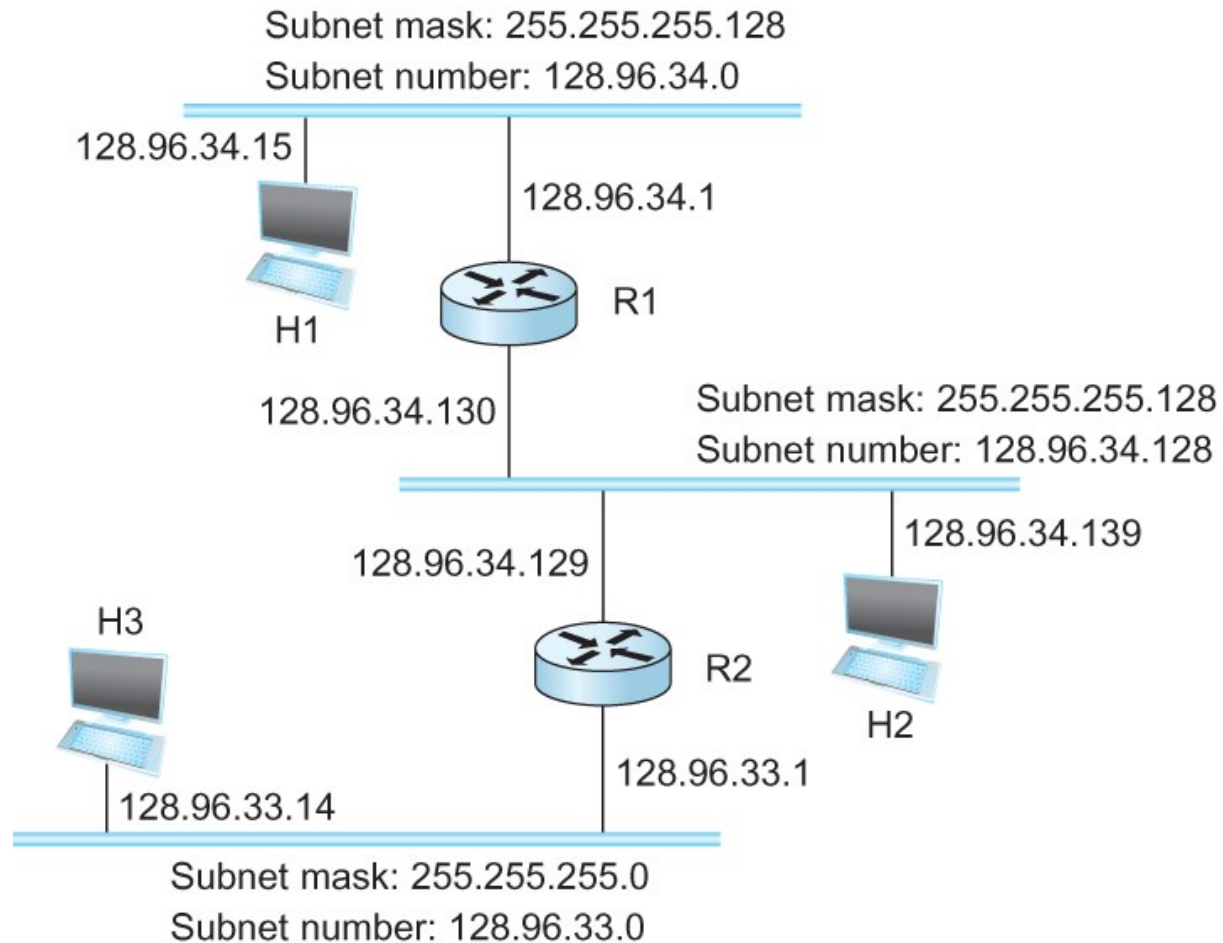
111111111111111111111111	00000000
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Subnet mask (255.255.255.0)

Network number	Subnet ID	Host ID
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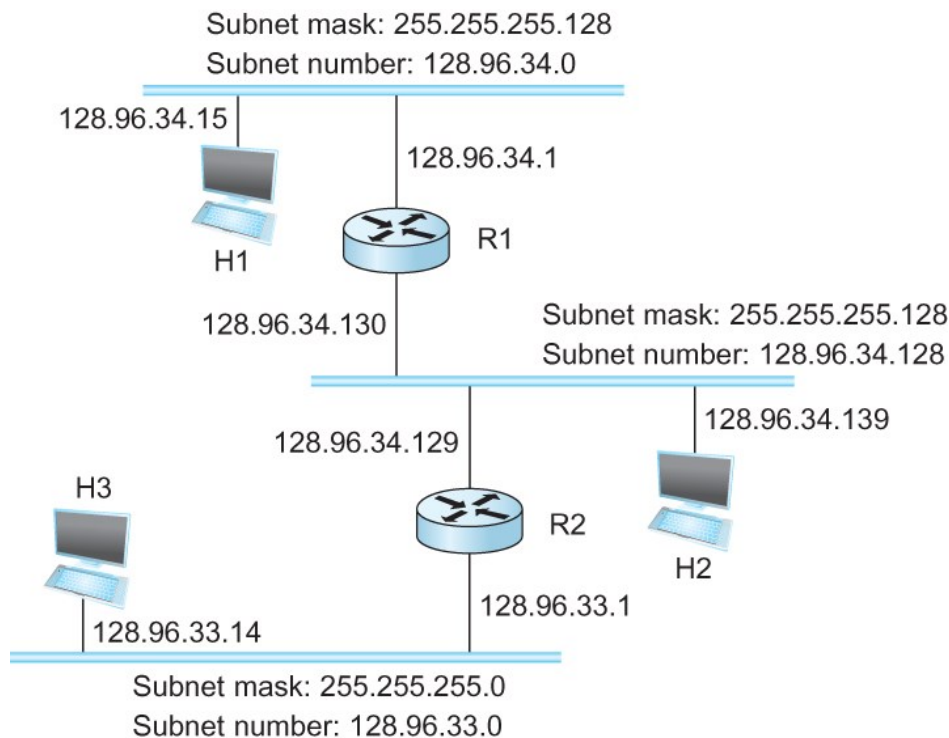
Subnetted address

# Subnetting: Example



# Subnetting: Example

## □ Forwarding Table at Router R1



SubnetNumber	SubnetMask	NextHop
128.96.34.0	255.255.255.128	Interface 0
128.96.34.128	255.255.255.128	Interface 1
128.96.33.0	255.255.255.0	R2



# Forwarding Algorithm

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```
D = destination IP address
for each entry < SubnetNum, SubnetMask, NextHop>
  D1 = SubnetMask & D
  if D1 = SubnetNum
    if NextHop is an interface
      deliver datagram directly to destination
    else
      deliver datagram to NextHop (a router)
```

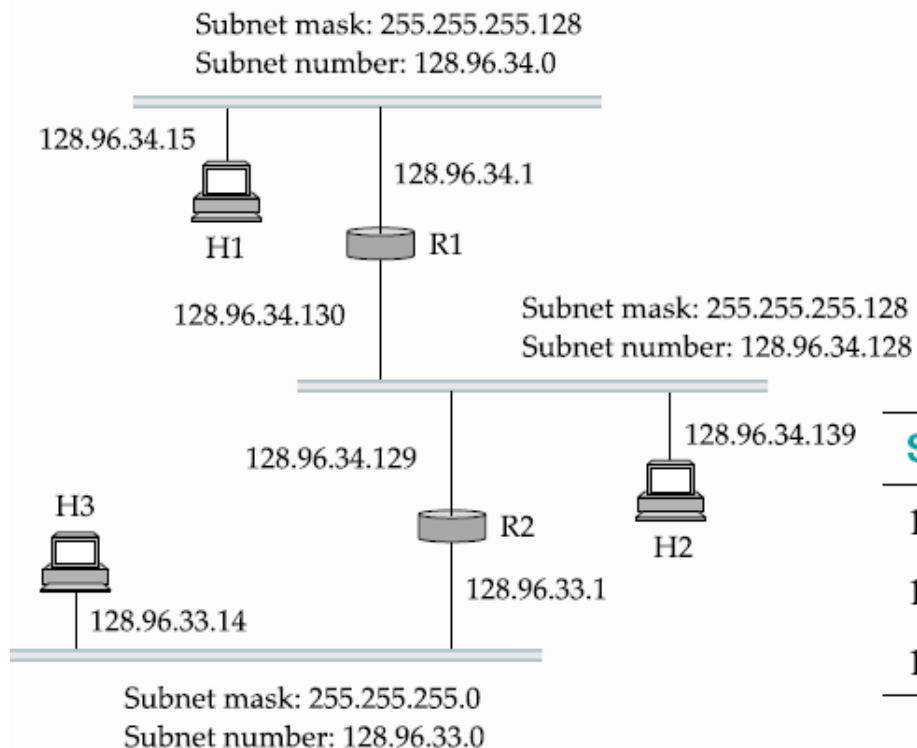
# Subnetting: Discussion

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- ❑ Would use a default router if nothing matches
- ❑ Subnet masks do not have to align with a byte boundary
- ❑ Subnet masks need **not** to be contiguous 1's
  - 255.255.1.0 is OK
    - ❑ 11111111 11111111 00000001 00000000
    - ❑ What is subnet number of IP address 128.96.34.1?  
10000000 01100000 00100010 00000000 &  
11111111 11111111 00000001 00000000 →  
10000000 01100000 00000000 00000000 →  
128.96.0.0 → can not directly tell from the IP address
  - In practice, use contiguous 1's
- ❑ Multiple subnets can be on a single physical network
- ❑ Subnets not visible from the rest of the Internet

# Subnetting: Discussion

- How do you tell whether an IP address is on a given subnet?



SubnetNumber	SubnetMask	NextHop
128.96.34.0	255.255.255.128	Interface 0
128.96.34.128	255.255.255.128	Interface 1
128.96.33.0	255.255.255.0	R2

# Exercise L11-1

- State to what next hop the IP packets a packet addressed to each of the following destinations will be delivered (assuming longest prefix match)

(a) 128.96.171.92

(b) 128.96.167.151

(c) 128.96.163.151

(d) 128.96.169.192

(e) 128.96.165.121

**Table 3.19 Routing Table for Exercise 56**

SubnetNumber	SubnetMask	NextHop
128.96.170.0	255.255.254.0	Interface 0
128.96.168.0	255.255.254.0	Interface 1
128.96.166.0	255.255.254.0	R2
128.96.164.0	255.255.252.0	R3
<default>		R4

# Scaling Problem

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- ❑ Need to address two scaling concerns in the Internet
  - The growth of backbone routing table as more and more network numbers need to be stored in them
  - Potential exhaustion of the 32-bit address space
- ❑ Address assignment efficiency
  - Arises because of the IP address structure with class A, B, and C addresses
  - Forces us to hand out network address space in fixed-size chunks of three very different sizes
    - ❑ A network with two hosts needs a class C address:
      - Address assignment efficiency =  $2/255 = 0.78$
    - ❑ A network with 256 hosts needs a class B address
      - Address assignment efficiency =  $256/65535 = 0.39$

# First Attempt

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- ❑ Exhaustion of IP address space centers on exhaustion of the class B network numbers
- ❑ Solution
  - Say “NO” to any Autonomous System (AS) that requests a class B address unless they can show a need for something close to 64K addresses
  - Instead give them an appropriate number of class C addresses
  - For any AS with at least 256 hosts, we can guarantee an address space utilization of at least 50%
- ❑ What is the problem with this solution?

# Classless Addressing

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- ❑ Problem with this solution
  - Excessive storage requirement at the routers.
- ❑ If a single AS has, say 16 class C network numbers assigned to it,
  - Every Internet backbone router needs 16 entries in its routing tables for that AS
  - This is true, even if the path to every one of these networks is the same
- ❑ If we had assigned a class B address to the AS
  - The same routing information can be stored in one entry
  - Efficiency =  $16 \times 255 / 65,536 = 6.2\%$

# Addressing Scaling Problem

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## ❑ Classless Inter-Domain Routing (CIDR)

- Addresses two scaling concerns in the Internet
  - ❑ The growth of backbone routing table as more and more network numbers need to be stored in them
  - ❑ Potential exhaustion of the 32-bit address space
- CIDR tries to balance the desire to minimize the number of routes that a router needs to know against the need to hand out addresses efficiently.
- CIDR uses aggregate routes
  - ❑ Uses a single entry in the forwarding table to tell the router how to reach a lot of different networks
  - ❑ Breaks the rigid boundaries between address classes



# Classless Addressing: Example

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- ❑ Consider an AS with 16 class C network numbers.
- ❑ Instead of handing out 16 addresses at random, hand out a block of contiguous class C addresses
- ❑ Suppose we assign the class C network numbers from 192.4.16 through 192.4.31
- ❑ Observe that top 20 bits of all the addresses in this range are the same (11000000 00000100 0001)
  - We have created a 20-bit network number (which is in between class B network number and class C number)
- ❑ Requires to hand out blocks of class C addresses that share a common prefix (sometimes, called supnetting)

# Classes Addressing: Notation

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- ❑ Requires to hand out blocks of class C addresses that share a common prefix
- ❑ The convention is to place a /X after the prefix where X is the prefix length in bits
- ❑ For example, the 20-bit prefix for all the networks 192.4.16 through 192.4.31 is represented as 192.4.16/20
- ❑ By contrast, if we wanted to represent a single class C network number, which is 24 bits long, we would write it 192.4.16/24

# Routing and Classes

## Addressing

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- ❑ How do the routing protocols handle this classless addresses
  - It must understand that the network number may be of any length
- ❑ Represent network number with a single pair

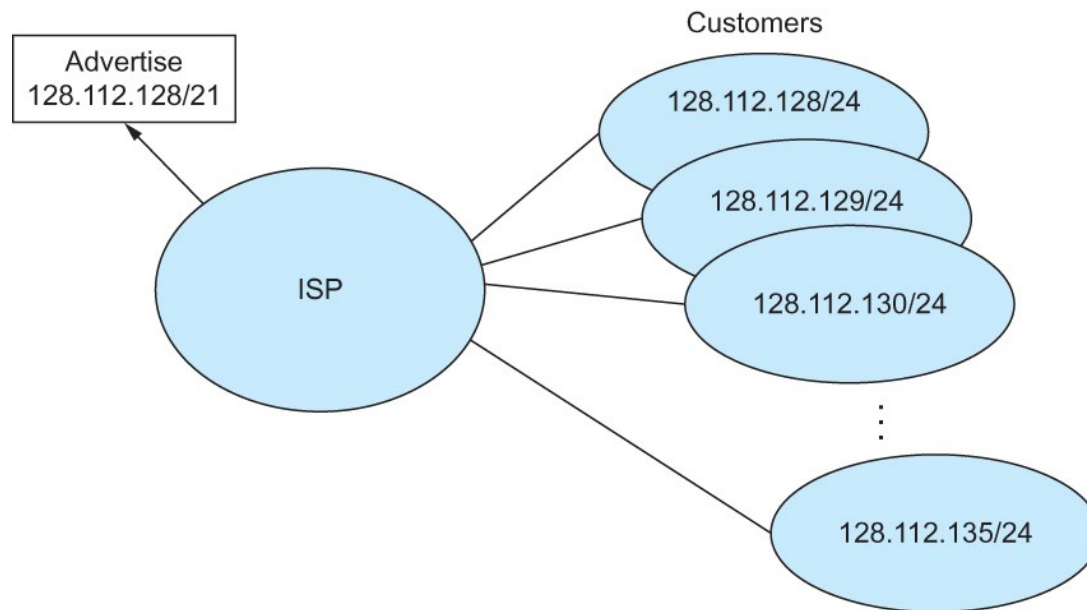
**<length, value>**

- ❑ All routers must understand CIDR addressing

# Routing and Classes Addressing: Example

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## □ Route aggregation with CIDR



# IP Forwarding Revisited (1)

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- ❑ IP forwarding mechanism *assumes* that it can find the network number in a packet and then look up that number in the forwarding table
- ❑ We need to *change* this assumption in case of CIDR

# IP Forwarding Revisited (2)

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- ❑ CIDR means that prefixes may be of any length, from 2 to 32 bits
- ❑ It is also possible to have prefixes in the forwarding tables that overlap
  - Some addresses may match more than one prefix
- ❑ e.g., we might find both 171.69 (a 16 bit prefix) and 171.69.10 (a 24 bit prefix) in the forwarding table of a single router
- ❑ A packet destined to 171.69.10.5 clearly matches both prefixes.
  - The rule is based on the principle of “longest match”
    - ❑ 171.69.10 in this case
- ❑ A packet destined to 171.69.20.5 would match 171.69 and not 171.69.10

# Exercise L11-2

- State to what next hop the IP packets a packet addressed to each of the following destinations will be delivered

(a) C4.4B.31.2E

(b) C4.5E.05.09

(c) C4.4D.31.2E

(d) C4.5E.03.87

(e) C4.5E.7F.12

(f) C4.5E.D1.02

**Table 3.21 Routing Table for Exercise 73**

Net/MaskLength	Nexthop
C4.5E.2.0/23	A
C4.5E.4.0/22	B
C4.5E.C0.0/19	C
C4.5E.40.0/18	D
C4.4C.0.0/14	E
C0.0.0.0/2	F
80.0.0.0/1	G

# Summary

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- ❑ Subnetting
  - Network number and network mask
- ❑ Classless addressing
  - ❑ Network prefix and length of prefix