

L9: Cipher Techniques: Problems



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Acknowledgement

- ❑ Many slides are from or are revised from the slides of the author of the textbook
 - Matt Bishop, Introduction to Computer Security, Addison-Wesley Professional, October, 2004, ISBN-13: 978-0-321-24774-5. [Introduction to Computer Security @ VSU's Safari Book Online subscription](#)
 - <http://nob.cs.ucdavis.edu/book/book-intro/slides/>

Outline

- ❑ Precomputing possible messages
- ❑ Misordered blocks
- ❑ Statistical regularities

Use Ciphers: A Challenge

- ❑ Cryptographic systems are sensitive to the environment they are being used
- ❑ Using cryptographic systems over a network introduces problems
- ❑ Using a good cipher is not enough, how to use the cipher matters greatly

- ❑ What can go wrong if we naively use ciphers?

Threats in Network Environment

- ❑ Knowledge of the environment and threats in the environment
 - Is the set of possible messages small?
 - Do the messages exhibit regularities that remain after encipherment
 - Can an active wiretapper rearrange or change parts of the message?
- ❑ Three common problems
 - Precomputation, misordered blocks, and statistical regularities

Attack 1. Precomputation

- ❑ Precomputing possible messages or *forward searches*
- ❑ Set of possible messages M small
- ❑ Public key cipher f used
- ❑ Idea: precompute set of possible ciphertexts $f(M)$ and build table $(m, f(m))$ where $m \in M$
- ❑ When ciphertext $f(m)$ appears, use table to find m

Forward Search Attack: Example

- ❑ Cathy knows Alice will send Bob one of two messages using a Public Key Cryptosystem
 - Enciphered BUY or enciphered SELL
- ❑ Using public key e_{Bob} , Cathy precomputes a table
 - $c_1 = f(m_1) = \{\text{BUY}\}_{e_{\text{Bob}}}$
 - $c_2 = f(m_2) = \{\text{SELL}\}_{e_{\text{Bob}}}$
- ❑ Looking up intercepted enciphered message, Cathy sees Alice send Bob m_2 .
- ❑ Cathy knows Alice send SELL

Obscure Threats

- ❑ Example: digitized sound (Simmons, 1982)
 - Initial calculations suggest 2^{32} such plaintexts
 - Seems like far too many possible plaintexts
 - Analysis of redundancy in human speech reduced this to about 100,000 ($\approx 2^{17}$)
 - This is small enough to worry about precomputation attacks

Notes on Precomputation

- ❑ Chosen plaintext attack against symmetric cryptosystems
 - Derive key
 - e.g., Hellman, 1980
- ❑ Precomputation attack against public key cryptosystems
 - Drive plaintext messages
 - Does not reveal private key

Misordered Blocks

- ❑ Parts of a ciphertext message can be deleted, replayed or reordered (Denning, 1982)

Misordered Blocks: Example

- ❑ Alice sends Bob message
 - $n_{Bob} = 77, e_{Bob} = 17, d_{Bob} = 53$
 - Message is LIVE (11 08 21 04)
 - Enciphered message is 44 57 21 16
- ❑ Eve intercepts it, rearranges blocks
 - Now enciphered message is 16 21 57 44
- ❑ Bob gets enciphered message, deciphers it
 - He sees EVIL

Notes on Misordered Blocks

- ❑ Digitally signing each block will not stop this attack
 - The parts are not bound to one another
- ❑ Two approaches to counter the attack
 1. Generate a cryptographic checksum of the *entire* message and sign it
 2. Place sequence numbers in each block of message, so recipient can tell intended order. Then you sign each block

Statistical Regularities

- If plaintext repeats, ciphertext may too

Statistical Regularities: Example

□ Example using DES:

■ input (in hex):

3231 3433 3635 3837 3231 3433 3635 3837

■ corresponding output (in hex):

ef7c 4bb2 b4ce 6f3b ef7c 4bb2 b4ce 6f3b

Notes on Statistical Regularities

❑ Code book mode (CBM)

- Each part is enciphered separately, so the same plaintext always produces the same ciphertext
- Each part is effectively looked up in a list of plaintext-ciphertext pairs
- It is the cause of the statistical regularity

❑ Approach to counter the attack

- Cascade blocks together (chaining, more details later)

What These Mean

- ❑ Use of *strong* cryptosystems, *well-chosen* (or random) keys *not enough* to be secure
- ❑ Other factors:
 - Protocols directing use of cryptosystems
 - Ancillary information added by protocols
 - Implementation (not discussed here)
 - Maintenance and operation (not discussed here)

Summary

- ❑ Discussed three attacks
 - Precomputation (forward search)
 - Misordered blocks
 - Statistical regularities
- ❑ Strong cryptosystems and random keys not enough
- ❑ Careful engineering matters