

Discrete-Event Simulation: Multi-Stream Lehmer RNGs

Lawrence M. Leemis and Stephen K. Park, Discrete-Event Simulation - A First Course, Prentice Hall, 2006

Hui Chen

Computer Science
Virginia State University
Petersburg, Virginia

February 22, 2016

Table of Contents

- 1 Introduction
- 2 Streams
- 3 Jump Multipliers
 - Numerical Examples
- 4 SSQ with Multiple Job Types

Introduction

- ▶ Typical DES models have many stochastic components.
 - ▶ e.g., arrivals and services
- ▶ Have a unique randomness for each stochastic component
- ▶ One option: multiple RNGs
 - ▶ Often considered a *poor* option
- ▶ One RNG with multiple “streams” of random numbers
 - ▶ One stream per stochastic component
 - ▶ Considered a *better* option
 - ▶ Method: partition output from a RNG into multiple streams
 - ▶ We have been using the Lehmer RNG from the authors of the textbook

Lehmer RNG

- ▶ Lehmer RNG used in *ssq2* and *sis2*

- ▶ in C/C++: *rng.h* and *rng.c*

```
double Random(void);  
void PutSeed(long x);  
void GetSeed(long *x);  
void TestRandom(void);
```

- ▶ in Java: *Rng.java*

```
class Rng {  
    [...]   
    public double random() {[...]}  
    public void putSeed(long x) {[...]}  
    public long getSeed() {[...]}  
    public void testRandom() {[...]}  
}
```

Revisit Simulation Program *ssq2*

- ▶ *ssq2* has two stochastic components: arrival process and service process, e.g.,

```
double GetArrival(void) {  
    static double arrival = START;  
  
    arrival += Exponential(2.0);  
    return (arrival);  
}
```

```
double GetService(void) {  
    return (Uniform(1.0, 2.0));  
}
```

Lehmer RNG: Partition Output into Multiple Streams: 1st Approach

- ▶ First attempt: partition output from our Lehmer RNG into multiple streams
 - ▶ Method: allocate and retain internal state of the RNG for each stochastic process

Lehmer RNG: Partition Output into Multiple Streams: 1st Approach

- ▶ Allocate a different generator state variable to each process and retain it before switching to the other process
 - ▶ In the Lehmer RNG, the generator state is fully represented by the *seed*
 - ▶ Allocate to the service process its own *static* variable and initialized with a value.

GetService with Unique Seed

```
double GetService(void) {  
    double s;  
    static long x = 12345;  
    PutSeed(x);  
    s = Uniform(1.0, 2.0);  
    GetSeed(&x);  
    return (s);  
}
```

Lehmer RNG: Partition Output into Multiple Streams: 1st Approach

- ▶ Allocate to arrival process its own *static* variable and initialized with a *different* value from the service process

GetArrival with Unique Seed

```
double GetArrival(void) {  
    static double arrival = START;  
    static long x = 54321;  
    PutSeed(x);  
    arrival += Exponential(2.0);  
    GetSeed(&x);  
    return (arrival);  
}
```

where x represents the current state of the service process

Modified Arrival and Service Processes

- ▶ As modified, arrival and service times are drawn from different streams of random numbers
- ▶ Provided the streams do not overlap, the processes are uncoupled
- ▶ Although the choice of seed for each stream is deceptively simple, the choices may in fact be *poor* ones.
- ▶ Execution time cost is negligible (see Example 3.2.3 in next slide)

Exercise L6-1

Complete Example 3.2.3 as instructed as follows,

- ▶ Make a copy of *ssq2*. To ease the discussion, call the new copy *ssq2b*
- ▶ Replace *GetArrival* and *GetService* by those introduced in the *1st approach*
- ▶ Set *LAST* to 1,000,000 jobs
- ▶ Compile and run the program. Measure the execution time of new and original programs
- ▶ Compare the results with from those using *ssq2* (without any modification). Are they the same?
- ▶ Compare execution time of the new and original programs, how much slower is the new program?

Submit your work (new and old programs, results from both programs, execution times, answer to the questions).

Exercise L6-1: Hints

To measure program execution time in Windows, use *PowerShell*. Below is an example that measures the execution time of a *dir* command,

Measure Program Execution Time

```
C:> powershell
PS C:> Measure-Command {dir | Out-Default}
Directory: C:\
Mode                LastWriteTime         Length Name
-----
[...]
Days                : 0
Hours               : 0
Minutes             : 0
Seconds             : 0
Milliseconds        : 129
Ticks               : 1299779
TotalDays           : 1.50437384259259E-06
TotalHours          : 3.61049722222222E-05
TotalMinutes        : 0.00216629833333333
TotalSeconds        : 0.1299779
TotalMilliseconds   : 129.9779
```

See

<https://technet.microsoft.com/en-us/library/hh849910.aspx>

Exercise L6-1: Hints

To measure program execution time in Linux, use the *time* command. Below is an example that measures the execution time of a *ls* command,

Measure Program Execution Time

```
$ time ls ssq2.c
ssq2.c

real    0m0.012s
user    0m0.004s
sys     0m0.008s
$
```

For more detail on the command, see its manual page (i.e., *man time*).

Streams using Multiple Seeds: Discussion

- ▶ Objective: allocate a unique stream of random numbers of each stochastic component
 - ▶ Examples of stochastic components: arrival and service processes
- ▶ Discussed approach: using *multiple seeds* of RNGs to produce *multiple unique streams* of random numbers
- ▶ Potential problem: assignment of initial seeds (or initial state)
 - ▶ Initial states should be chosen to produce *disjoint* streams
 - ▶ If states are picked at whim, no *guarantee* of disjoint streams
 - ▶ Some initial states could be just a few calls to *Random()* away from one another

Jump Multipliers

- ▶ Objective: produce multiple *disjoint* streams of random numbers
- ▶ Theorem 3.2.1 is the key to creating streams

Theorem 3.2.1

Given $g(x) = ax \bmod m$ and integer j with $j = 1, 2, \dots, m-1$, the associated *jump function* is

$$g^j(x) = (a^j \bmod m)x \bmod m \quad (1)$$

and has the jump multiplier $a^j \bmod m$

if $(g(\cdot))$ generates x_0, x_1, x_2, \dots then $g^j(\cdot)$ generates x_0, x_j, x_{2j}, \dots

Numerical Examples

- ▶ If $m = 31$ and $a = 3$ and $j = 6$, the jump multiplier is $a^j \bmod m = 3^6 \bmod 31 = 16$
- ▶ If $x_0 = 1$ then $g(x) = 3x \bmod 31$ generates

1, 3, 9, 27, 19, 26, 16, 17, 20, 29, 25, 13,
8, 24, 10, 30, 28, 22, 4, 12, 5, 15, 14, 11,
2, 6...

- ▶ The jump function $g^6(x) = 16x \bmod 31$ generates 1, 16, 8, 4, 2, ...
i.e., the first sequence is x_0, x_1, x_2, \dots ; the second is x_0, x_6, x_{12}, \dots

Program for Numerical Example

```

#include <stdio.h>
#include <math.h>

int main() {
    long m = 31, a = 3, j = 6, jm, i, n = 32, x = 1;

    jm = (long)pow(a, j) % m;

    printf("%2ld", x);
    for (i = 0; i < n; i++) {
        printf(" ", %2ld", x = a * x % m);
    }

    printf("\n\nJump Multiplier = %ld\n\n", jm);

    x = 1;
    printf("%2ld", x);
    for (i = 0; i < n/6; i++) {
        printf(" ", %2ld", x = jm * x % m);
    }
    printf("\n");

    return 0;
}

```

```

import java.io.*;
import java.lang.Math;
import java.text.*;

public class Ex3_2_4{

    public static void main(String[] args) {
        String format = "%2d";

        long m = 31, a = 3, j = 6, jm, i, n = 32, x = 1;

        jm = (long)Math.pow(a, j) % m;

        System.out.format(format, x);
        for (i = 0; i < n; i++) {
            System.out.format(" ", " +
                format, x = a * x % m);
        }
        System.out.format("\n\nJump Multiplier = " + format +
            "\n\n", jm);

        x = 1; System.out.format(format, x);
        for (i = 0; i < n/6; i++) {
            System.out.format(" ", " +
                format, x = jm * x % m);
        }
        System.out.print("\n");
    }
}

```


Using Jumper Function

1. Compute the jump multiplier $g^j(\cdot) = a^j \bmod m$, which is a one time cost.
2. $g^j(\cdot)$ permits jumping from x_0 to x_j to x_{2j} to \dots
3. User supplies one initial seed
4. If j is chosen well, $g^j(\cdot)$ can “plant” additional initial seeds
5. Each planted seed corresponds to a different stream
6. Each planted seed is separated by j calls to *Random()*

Maximal Modulus-Compatible Jump Multiplier

Definition 3.2.1

Given a Lehmer random-number generator with prime modulus m , full-period modulus-compatible multiplier a , and a requirement for s disjoint streams as widely separated as possible, the maximal jump multiplier is $a^j \bmod m$, where j is the largest integer less than $\lfloor m/s \rfloor$ such that $a^j \bmod m$ is modulus-compatible with m .

Example 3.2.6

Jump multipliers for $(a, m) = (48271, 2^{31} - 1) = (48271, 2147483647)$ RNG

# of streams s	$\lfloor m/s \rfloor$	jump size j	jump multiplier $a^j \bmod m$
1024	$\lfloor \frac{2^{31}-1}{1024} \rfloor = 2097151$	2082675	$48271^{2082675} \bmod 2147483647 = 97070$
512	$\lfloor \frac{2^{31}-1}{512} \rfloor = 4194303$	4170283	$48271^{4170283} \bmod 2147483647 = 44857$
256	$\lfloor \frac{2^{31}-1}{256} \rfloor = 8388607$	8367782	$48271^{8367782} \bmod 2147483647 = 22925$
128	$\lfloor \frac{2^{31}-1}{128} \rfloor = 16777215$	16775552	$48271^{16775552} \bmod 2147483647 = 40509$

Jump Multiplier: A Simple Search Program

Jump Multiplier Search Program in C

```
#include <stdio.h>

long find_m_compatible(long upper, long a, long m);
long modular_pow(long long base, long long exponent, long long modulus);

int main()
{
    long a = 482711, m = 21474836471,
        nstreams[] = {10241, 5121, 2561, 1281}, i, n;

    for (i = 0; i < sizeof(nstreams)/sizeof(long); i++) {
        n = find_m_compatible(m/nstreams[i], a, m);
        printf("%ld_%ld_%ld_%ld\n", nstreams[i], m/nstreams[i], n,
            modular_pow((long long)a, (long long)n, (long long)m));
    }

    return 0;
}
```

Jump Multiplier: A Simple Search Program

Definition 2.2.1

The multiplier a is modulus-compatible with the prime modulus m if and only if $r < q$ where $r = m \bmod a$ and $q = \lfloor m/a \rfloor$

Functions in Jump Multiplier Search Program

```
long modular_pow(long long base, long long exponent, long long modulus);

long find_m_compatible(long upper, long a, long m)
{
    long i, n, r, q;

    for (i = upper; i >= 1; i --) {
        n = modular_pow((long long)a, (long long)i, (long long)m);
        r = m % n;
        q = m / n;
        if (r < q) { /* if n is modulus-compatible with m */
            return i;
        }
    }

    return 0;
}
```

Jump Multiplier: A Simple Search Program

This function is to compute $base^{exponent} \bmod modulus$

Functions in Jump Multiplier Search Program

```
/*
 * Reference
 * http://www.sanfoundry.com/cpp-program-implement-modular-exponentiation-algorithm/
 */

long modular_pow(long long base, long long exponent, long long modulus)
{
    long long result = 1ll;

    while (exponent > 0ll) {
        if (exponent % 2ll == 1ll) {
            result = (result * base) % modulus;
        }

        exponent = exponent >> 1ll;

        base = (base * base) % modulus;
    }

    return (long) result;
}
```

Jump Multiplier: A Simple Search Program

Jump Multiplier Search Program in Java

```
public class JumpMultiplier {  
    public static void main(String[] args) {  
        JumpMultiplier jm = new JumpMultiplier();  
  
        long a = 482711, m = 21474836471,  
            nstreams[] = {10241, 5121, 2561, 1281}, n;  
  
        for (int i = 0; i < nstreams.length; i++) {  
            n = jm.find_m_compatible(m/nstreams[i], a, m);  
            System.out.format("%d %d %d %d\n", nstreams[i], m/nstreams[i], n,  
                jm.modular_pow(a, n, m));  
        }  
    }  
}
```

.....

Jump Multiplier: A Simple Search Program

Jump Multiplier Search Program in Java

.....

```
private long modular_pow(long base, long exponent, long modulus) {  
    long result = 1l;  
  
    while (exponent > 0l) {  
        if (exponent % 2l == 1l) {  
            result = (result * base) % modulus;  
        }  
  
        exponent = exponent >> 1l;  
  
        base = (base * base) % modulus;  
    }  
  
    return result;  
}
```

.....

Jump Multiplier: A Simple Search Program

Jump Multiplier Search Program in Java

```
.....  
  
private long find_m_compatible(long upper, long a, long m) {  
    long i, n, r, q;  
  
    for (i = upper; i >= 1; i --) {  
        n = modular_pow(a, i, m);  
        r = m % n;  
        q = m / n;  
        if (r < q) { /* if n is modulus-compatible with m */  
            return i;  
        }  
    }  
  
    return 0;  
}  
  
}
```


Exercise L6-2

- ▶ Use the program discussed above to compute the jump multiplier table similar to slide 18.
 - ▶ You will enter, compile, and run the programs.
- ▶ Submit the program(s) and the table in Blackboard under L6-2

Library *rngs*

- ▶ *rngs* is an upward-compatible multi-stream replacement for *rng*
- ▶ By default, provides 256 streams, indexed 0 to 255 (0 is the default)
- ▶ Only one stream is active at any time
- ▶ Six available functions:
 - ▶ `Random(void)`
 - ▶ `PutSeed(long x)`: superseded by `PlantSeeds`
 - ▶ `GetSeed(long *x)`
 - ▶ `TestRandom(void)`
 - ▶ `SelectStream(int s)`: used to define the active stream
 - ▶ `PlantSeeds(long x)`: “plants” one seed per stream
- ▶ Henceforth, *rngs* is the library of choice

Example 3.2.7: *ssq2* Revisited

- ▶ Use *rngs* functions for *GetArrival*, *GetService*
- ▶ Include *rngs.h* and use *PlantSeeds(12345)*

GetArrival Method

```
double GetArrival(void) {  
    static double arrival = START;  
    SelectStream(0);  
    arrival += Exponential(2.0);  
    return (arrival);  
}
```

GetService Method

```
double GetService(void) {  
    SelectStream(2);  
    return (Uniform(1.0, 2.0));  
}
```

Exercise L6-3

Complete Example 3.2.7 as instructed as follows,

- ▶ Make a copy of *ssq2*. To ease the discussion, call the new copy *ssq2c*
- ▶ Replace *GetArrival* and *GetService* by those using library rngs. Find library rngs in Blackboard.
- ▶ Set *LAST* to 1,000,000 jobs
- ▶ Compile and run the program. Measure the execution time of new and original programs
- ▶ Compare the results with from those using *ssq2* (without any modification) and those using *ssq2b* (Class Exercise L6-1). Are they the same?
- ▶ Compare execution time of the new and original programs, how much slower is the new program?

Submit your work (new and old programs, results from both programs, execution times, answer to the questions).

Uncoupling Stochastic Processes

- ▶ Per modifications, arrival and service processes are uncoupled
- ▶ Consider changing the service process to $Uniform(0.0, 1.5) + Uniform(0.0, 1.5)$
- ▶ Without uncoupling, arrival process sequence would change!
- ▶ With uncoupling, the service process sees exactly the same arrival sequence
- ▶ Important variance reduction technique

Single-Server Service Node with Multiple Job Types

- ▶ Extend the single-server service node model from Chapter 1
- ▶ Consider multiple job types, each with its own arrival and service process
- ▶ Examples 3.2.8 and 3.2.9: Suppose there are two job types
 1. Exponential(4.0) interarrivals, Uniform(1.0, 3.0) service
 2. Exponential(6.0) interarrivals, Uniform(0.0, 4.0) service

Use rngs to allocate a different stream to each stochastic process

Arrival Process for Multiple Job Types

The arrival process generator in program *ssq2* can be modified as follows,

Example 3.2.8: Arrival Process

```
double GetArrival(int *j) {           /* returns job type in j */
    const double mean[2] = {4.0, 6.0}; /* two job types */
    static double arrival[2] = {START, START};
    static int init = 1;
    double temp;
    if (init) {                       /* initialize the arrival array */
        SelectStream(0);
        arrival[0] += Exponential(mean[0]);
        SelectStream(1);
        arrival[1] += Exponential(mean[1]);
        init = 0;
    }
    if (arrival[0] <= arrival[1])
        *j = 0;                       /* next arrival is job type 0 */
    else
        *j = 1;                       /* next arrival is job type 1 */
    temp = arrival[*j];               /* next arrival time to be returned */
    SelectStream(*j);
    arrival[*j] += Exponential(mean[*j]); /* arrival after next arrival */
    return (temp);
}
```

Service Process for Multiple Job Types

The service process generator in program *ssq2* can be modified as follows,

Example 3.2.9: Service Progress

```
double GetService(int j)
{
    const double min[2] = {1.0, 0.0};
    const double max[2] = {3.0, 4.0};
    /*
     * Two RNG streams, i.e., streams 0 and 1 are used in the arrival
     * process generator. We now use streams 2 and 3 for the service
     * process generator. In the following, j should be either 0 or 1.
     */
    SelectStream(j + 2);
    return (Uniform(min[j], max[j]));
}
```


Service Process for Multiple Job Types

- ▶ Index j matches service time to appropriate job type
- ▶ All four simulated stochastic processes are uncoupled
- ▶ Any process could be changed without altering the random sequence of others!

Consistency Check

- ▶ Additional modification to *ssq2*
 - ▶ job-type-specific statistics-gathering needs to be added in the *main* method/function
- ▶ How do we know if our modifications are correct? *Use consistency check to increase confidence.*
 - ▶ $\overline{w} = \overline{d} + \overline{s}$
 - ▶ $\overline{l} = \overline{q} + \overline{x}$
 - ▶ How about average service time of both job types?
Since $(1.0 + 3.0)/2 = (0.0 + 4.0)/2 = 2.0$, we expect $\overline{s} = 2.0$
 - ▶ How about the net arrival rate of both job types?
Since the arrival rates of job types 0 and 1 are $1/4$ and $1/6$, respectively, we expect the net arrival rate should be $1/4 + 1/6 = 5/12$, i.e., $\overline{r} = 1/(5/12) = 12/5 = 2.4$.
 - ▶ The steady-state utilization should be the ratio of the arrival rate to the service rate, i.e., $(5/12)/(1/20) = 5/6 \approx 0.83$

Exercise L6-4

Modify program *ssq2* to support two job types as discussed.

- ▶ Make a copy of *ssq2*. To ease the discussion, call the new copy *ssq2d*
- ▶ Modify program *ssq2d* as suggested in Examples 3.2.8 and 3.2.9.
- ▶ Modify the *main* method/function to include job-type-specific statistics
- ▶ Answer the following questions,
 - ▶ What proportion of processed jobs are of type 0?
 - ▶ What are \overline{w} , \overline{d} , \overline{s} , \overline{l} , \overline{q} , and \overline{x} for each job type?
 - ▶ What did you do to convince yourself that your results are valid (hint: consistency check)?
 - ▶ Why are \overline{w} , \overline{d} , and \overline{s} the same for both job types, while \overline{l} , \overline{q} , and \overline{x} are different?

Submit your programs, results, and answers to the questions in Blackboard.

Summary

- ▶ Multiple stream RNGs
 - ▶ Generate disjoint random number streams
 - ▶ Want the streams are far apart
- ▶ More simulation examples to be discussed