## Discrete-Event Simulation: Multi-Stream Lehmer RNGs

Lawrence M. Leemis and Stephen K. Park, Discrete-Event Simulation - A First Course, Prentice Hall, 2006

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#### Introduction

- Typical DES models have many stochastic components.
  - e.g., arrivals and services
- Have a unique randomness for each stochastic component
- ▶ One option: multiple RNGs
  - Often considered a poor option
- ▶ One RNG with multiple "streams" of random numbers
  - One stream per stochastic component
  - Considered a better option
  - Method: partition output from a RNG into multiple streams
    - We have been using the Lehmer RNG from the authors of the textbook

#### Lehmer RNG

▶ Lehmer RNG used in *ssq2* and *sis2* 

```
▶ in C/C++: rng.h and rng.c
 double Random(void);
  void PutSeed(long x);
 void GetSeed(long *x);
  void TestRandom(void);
▶ in Java: Rng.java
  class Rng {
      [...]
      public double random() {[...]}
      public void putSeed(long x) {[...]}
      public long getSeed() {[...]}
      public void testRandom() {[...]}
```

## Revisit Simulation Program ssq2

 ssq2 has two stochastic components: arrival process and service process, e.g.,

```
double GetArrival(void) {
    static double arrival = START;

    arrival += Exponential(2.0);
    return (arrival);
}

double GetService(void) {
    return (Uniform(1.0, 2.0));
}
```

# Lehmer RNG: Partition Output into Multiple Streams: 1st Approach

- First attempt: partition output from our Lehmer RNG into multiple streams
  - Method: allocate and retain internal state of the RNG for each stochastic process

# Lehmer RNG: Partition Output into Multiple Streams: 1st Approach

- Allocate a different generator state variable to each process and retain it before switching to the other process
  - ▶ In the Lehmer RNG, the generator state is fully represented by the seed
  - Allocate to the service process its own static variable and initialized with a value.

#### GetService with Unique Seed

```
double GetService(void) {
    double s;
    static long x = 12345;
    PutSeed(x);
    s = Uniform(1.0, 2.0);
    GetSeed(&x);
    return (s);
}
```

# Lehmer RNG: Partition Output into Multiple Streams: 1st Approach

 Allocate to arrival process its own static variable and initialized with a different value from the service process

```
double GetArrival(void) {
    static double arrival = START;
    static long x = 54321;
    PutSeed(x);
    arrival += Exponential(2.0);
```

GetArrival with Unique Seed

GetSeed(&x);

return (arrival);

where x represents the current state of the service process

### Modified Arrival and Service Processes

- As modified, arrival and service times are drawn from different streams of random numbers
- Provided the streams do not overlap, the processes are uncoupled
- ▶ Although the choice of seed for each stream is deceivingly simple, the choices may in fact be *poor* ones.
- ► Execution time cost is negligible (see Example 3.2.3 in next slide)

#### Exericse L6-1

Complete Example 3.2.3 as instructed as follows,

- ▶ Make a copy of *ssq2*. To ease the discussion, call the new copy *ssq2b*
- Replace GetArrival and GetService by those introduced in the 1st approach
- ▶ Set *LAST* to 1,000,000 jobs
- Compile and run the program. Measure the execution time of new and original programs
- ► Compare the results with from those using *ssq2* (without any modification). Are they the same?
- ► Compare execution time of the new and original programs, how much slower is the new program?

Submit your work (new and old programs, results from both programs, execution times, answer to the questions).

#### Exericse L6-1: Hints

To measure program execution time in Windows, use *PowerShell*. Below is an example that measures the execution time of a *dir* command,

#### Measure Program Execution Time

```
C:> powershell
PS C:> Measure-Command { dir | Out-Default }
  Directory: C:\
 Mode
                     LastWriteTime
                                      Length Name
  [...]
  Days
  Hours
  Minutes
  Seconds
  Milliseconds
                   : 129
  Ticks
                   : 1299779
  TotalDays
                   : 1.50437384259259E-06
  TotalHours : 3.6104972222222E-05
  TotalMinutes : 0.00216629833333333
  TotalSeconds : 0.1299779
  Total Milliseconds: 129.9779
```

See

https://technet.microsoft.com/en-us/library/hh849910.aspx

#### Exercise L6-1: Hints

To measure program execution time in Linux, use the *time* command. Below is an example that measures the execution time of a *ls* command,

#### Measure Program Execution Time

```
$ time Is ssq2.c
ssq2.c
real 0m0.012s
user 0m0.004s
sys 0m0.008s
```

For more detail on the command, see its manual page (i.e., man time).

## Streams using Multiple Seeds: Discussion

- ▶ Objective: allocate a unique stream of random numbers of each stochastic component
  - Examples of stochastic components: arrival and service processes
- Discussed approach: using multiple seeds of RNGs to produce multiple unique streams of random numbers
- ▶ Potential problem: assignment of initial seeds (or initial state)
  - ▶ Initial states should be chosen to produce *disjoint* streams
  - ▶ If states are picked at whim, no guarantee of disjoint streams
  - Some initial states could be just a few calls to Random() away from one another

# Jump Multipliers

- ▶ Objective: produce multiple *disjoint* streams of random numbers
- ▶ Theorem 3.2.1 is the key to creating streams

#### Theorem 3.2.1

Given  $g(x) = ax \mod m$  and integer j with j = 1, 2, ..., m - 1, the associated *jump function* is

$$g^{j}(x) = (a^{j} \mod m)x \mod m \tag{1}$$

and has the jump multiplier  $a^j \mod m$  if  $(g(\cdot))$  generates  $x_0, x_1, x_2, \ldots$  then  $g^j(\cdot)$  generates  $x_0, x_j, x_{2j}, \ldots$ 

## Numerical Examples

- If m = 31 and a = 3 and j = 6, the jump multiplier is  $a^j \mod m = 36 \mod 31 = 16$
- ▶ If  $x_0 = 1$  then  $g(x) = 3x \mod 31$  generates

► The jump function  $g^6(x) = 16x \mod 31$  generates  $1, 16, 8, 4, 2, \ldots$  i.e., the first sequence is  $x_0, x_1, x_2, \ldots$ ; the second is  $x_0, x_6, x_{12}, \ldots$ 

## Program for Numerical Example

```
#include < stdio.h>
#include <math.h>
int main() {
    long m = 31, a = 3, j = 6, jm, i, n =
          32. \times = 1:
    jm = (long)pow(a, j) \% m;
    printf("%2Id", x);
    for (i = 0: i < n: i ++) {
         printf(", %2Id", x = a * x % m);
    printf("\n\nJump Multiplier = \%ld\n\n".
         im);
    x = 1:
    printf("%2Id", x);
    for (i = 0; i < n/6; i ++)
         printf(", \%2Id", x = im * x \% m):
    printf("\n");
    return 0:
```

```
import iava.io.*:
import iava.lang.Math:
import java.text.*;
public class Ex3_2_4{
    public static void main(String[] args) {
        String format = "%2d":
        long m = 31, a = 3, j = 6, jm, i, n
             = 32. x = 1:
       im = (long) Math.pow(a, j) % m;
        System.out.format(format, x);
        for (i = 0; i < n; i ++) {
            System.out.format(", " +
                 format x = a * x % m:
        System.out.format("\n\nJump
             Multiplier = " + format +
             "\n\n", jm);
        x = 1; System.out.format(format, x);
        for (i = 0; i < n/6; i ++) {
            System.out.format(", " +
                 format, x = im * x % m);
        System.out.print("\n");
```

# **Using Jumper Function**

- 1. Compute the jump multiplier  $g^j(\cdot) = a^j \mod m$ , which is a one time cost.
- 2.  $g^{j}(\cdot)$  permits jumping from  $x_0$  to  $x_j$  to  $x_{2j}$  to ...
- 3. User supplies one initial seed
- 4. If j is chosen well,  $g^j(\cdot)$  can "plant" additional initial seeds
- 5. Each planted seed corresponds to a different stream
- 6. Each planted seed is separated by j calls to Random()

# Maximal Modulus-Compatible Jump Multiplier

#### Definition 3.2.1

Given a Lehmer random-number generator with prime modulus m, full-period modulus-compatible multiplier a, and a requirement for s disjoint streams as widely separated as possible, the maximal jump multiplier is  $a^j \mod m$ , where j is the largest integer less than  $\lfloor m/s \rfloor$  such that  $a^j \mod m$  is modulus-compatible with m.

#### Example 3.2.6

Jump multipliers for $(a, m) = (48271, 2^{31} - 1) = (48271, 2147483647)$ RNG			
# of streams $s$	$\lfloor m/s \rfloor$	jump size <i>j</i>	
1024	$\left\lfloor \frac{2^{31}-1}{1024} \right\rfloor =$	2082675	48271 <sup>2082675</sup>
	2097151		$mod\ 2147483647 = 97070$
512	$\lfloor \frac{2^{31}-1}{512} \rfloor =$	4170283	48271 <sup>4170283</sup>
	4194303		$mod\ 2147483647 = 44857$
256	$\lfloor \frac{2^{31}-1}{256} \rfloor =$	8367782	48271 <sup>8367782</sup>
	8388607		$mod\ 2147483647 = 22925$
128	$ \frac{2^{31}-1}{129} =$	16775552	48271 <sup>16775552</sup>
	16777215		$mod\ 2147483647 = 40509$

### Jump Multiplier Search Program in C

#### Definition 2.2.1

The multiplier a is modulus-compatible with the prime modulus m if and only if r < qwhere  $r = m \mod a$  and  $q = \lfloor m/a \rfloor$ 

#### Functions in Jump Multiplier Search Program

```
long modular-pow(long long base, long long exponent, long long modulus);
long find_m_compatible(long upper, long a, long m)
    long i, n, r, q;
    for (i = upper; i >= 1; i --) {
        n = modular_pow((long long)a, (long long)i, (long long)m);
        r = m \% n:
        q = m / n;
        if (r < q) { /* if n is modulus-compatibl with m */
            return i:
    return 0;
```

This function is to compute base exponent mod modulus

#### Functions in Jump Multiplier Search Program

```
* Reference
* http://www.sanfoundry.com/cpp-program-implement-modular-exponentiation-algorithm/
* */
long modular_pow(long long base, long long exponent, long long modulus)
    long long result = 111;
    while (exponent > 0 II) {
        if (exponent % 2|| == 1||) {
            result = (result * base) % modulus;
        exponent = exponent >> 1 | 1 |;
        base = (base * base) % modulus:
    return (long) result;
```

## Jump Multiplier Search Program in Java

## Jump Multiplier Search Program in Java

```
. . . . . . . . .
 private long modular_pow(long base, long exponent, long modulus) {
     long result = 11;
     while (exponent > 01) {
         if (exponent % 2| == 1|) {
             result = (result * base) % modulus;
         exponent = exponent >> 11;
         base = (base * base) % modulus;
     return result;
```

## Jump Multiplier Search Program in Java

```
. . . . . . . . .
 private long find_m_compatible(long upper, long a, long m) {
     long i, n, r, q;
     for (i = upper; i >= 1; i --) {
         n = modular_pow(a, i, m):
         r = m \% n:
         q = m / n;
         if (r < q) { /* if n is modulus-compatibl with m */
             return i:
     return 0;
```

#### Exericse L6-2

- ▶ Use the program discussed above to compute the jump multiplier table similar to slide 18.
  - You will enter, compile, and run the programs.
- Submit the program(s) and the table in Blackboard under L6-2

# Library rngs

- lacktriangledown rngs is an upward-compatible multi-stream replacement for  $\emph{rng}$
- By default, provides 256 streams, indexed 0 to 255 (0 is the default)
- Only one stream is active at any time
- Six available functions:
  - Random(void)
  - PutSeed(long x): superseded by PlantSeeds
  - ► GetSeed(long \*x)
  - TestRandom(void)
  - ▶ SelectStream(int s): used to define the active stream
  - ▶ PlantSeeds(long x): "plants" one seed per stream
- ▶ Henceforth, *rngs* is the library of choice

## Example 3.2.7: ssq2 Revisited

- ▶ Use *rngs* functions for *GetArrival*, *GetService*
- ▶ Include *rngs.h* and use *PlantSeeds*(12345)

#### GetArrival Method

```
double GetArrival(void) {
    static double arrival = START;
    SelectStream(0);
    arrival += Exponential(2.0);
    return (arrival);
}
```

#### GetService Method

```
double GetService(void) {
    SelectStream(2);
    return (Uniform(1.0, 2.0));
}
```

#### Exericse L6-3

Complete Example 3.2.7 as instructed as follows,

- ▶ Make a copy of *ssq2*. To ease the discussion, call the new copy *ssq2c*
- Replace GetArrival and GetService by those using library rngs. Find library rngs in Blackboard.
- ▶ Set *LAST* to 1,000,000 jobs
- Compile and run the program. Measure the execution time of new and original programs
- ► Compare the results with from those using *ssq2* (without any modification) and those using *ssq2b*(Exercise L6-1). Are they the same?
- ► Compare execution time of the new and originial programs, how much slower is the new program?

Submit your work (new and old programs, results from both programs, execution times, answer to the questions).

## **Uncoupling Stochastic Processes**

- Per modifications, arrival and service processes are uncoupled
- Consider changing the service process to Uniform(0.0, 1.5) + Uniform(0.0, 1.5)
- Without uncoupling, arrival process sequence would change!
- With uncoupling, the service process sees exactly the same arrival sequence
- Important variance reduction technique

## Single-Server Service Node with Multiple Job Types

- Extend the single-server service node model from Chapter 1
- Consider multiple job types, each with its own arrival and service process
- ► Examples 3.2.8 and 3.2.9: Suppose there are two job types
  - 1. Exponential(4.0) interarrivals, Uniform(1.0, 3.0) service
  - 2. Exponential(6.0) interarrivals, Uniform(0.0, 4.0) service

Use rngs to allocate a different stream to each stochastic process

## Arrival Process for Multiple Job Types

The arrival process generator in program ssq2 can be modified as follows,

#### Example 3.2.8: Arrival Process

```
double GetArrival(int *j) { /* returns job type in j */ const double mean[2] = {4.0, 6.0}; /* two job types */
    static double arrival[2] = {START, START};
    static int init = 1:
    double temp;
                                            /* initialize the arrival arrav */
    if (init) {
        SelectStream (0):
         arrival[0] += Exponential(mean[0]);
        SelectStream (1);
         arrival[1] += Exponential(mean[1]):
         init = 0:
    if (arrival[0] <= arrival[1])
        *i = 0:
                                            /* next arrival is job type 0 */
    else
                                            /* next arrival is job type 1 */
        *i = 1:
                                           /* next arrival time to be returned */
    temp = arrival[*j];
    SelectStream(*i);
    arrival[*i] += Exponential(mean[*i]); /* arrival after next arrival */
    return (temp);
```

## Service Process for Multiple Job Types

The service process generator in program ssq2 can be modified as follows,

#### Example 3.2.9: Service Progress

```
double GetService(int j)
{
    const double min[2] = {1.0, 0.0};
    const double max[2] = {3.0, 4.0};

/*
    * Two RNG streams, i.e., streams 0 and 1 are used in the arrival
    * process generator. We now use streams 2 and 3 for the service
    * process generator. In the following, j should be either 0 or 1.
    * */
    SelectStream(j + 2);
    return (Uniform(min[j], max[j]));
}
```

## Service Process for Multiple Job Types

- ▶ Index *j* matches service time to appropriate job type
- ► All four simulated stochastic processes are uncoupled
- Any process could be changed without altering the random sequence of others!

# Consistency Check

- Additional modification to ssq2
  - job-type-specific statistics-gathering needs to be added in the main method/function
- ▶ How do we know if our modifications are correct? *Use consistency check to increase confidence.* 
  - $\overline{w} = \overline{d} + \overline{s}$
  - $ightharpoonup \overline{I} = \overline{q} + \overline{x}$
  - ► How about average service time of both job types? Since (1.0 + 3.0)/2 = (0.0 + 4.0)/2 = 2.0, we expect  $\overline{s} = 2.0$
  - ▶ How about the net arrival rate of both job types? Since the arrival rates of job types 0 and 1 are 1/4 and 1/6, respectively, we expect the net arrival rate should be 1/4 + 1/6 = 5/12, i.e.,  $\overline{r} = 1/(5/12) = 12/5 = 2.4$ .
  - ► The steady-state utilization should be the ratio of the arrival rate to the service rate, i.e.,  $(5/12)/(1/20 = 5/6 \approx 0.83$

### Exercise L6-4

Modify program *ssq2* to support two job types as discussed.

- ▶ Make a copy of *ssq2*. To ease the discussion, call the new copy *ssq2d*
- ▶ Modify program *ssq2d* as suggested in Examples 3.2.8 and 3.2.9.
- Modify the main method/function to include job-type-specific statisics
- Answer the following questions,
  - What portion of processed jobs are of type 0?
  - ▶ What are  $\overline{w}$ ,  $\overline{d}$ ,  $\overline{s}$   $\overline{l}$ ,  $\overline{q}$ , and  $\overline{x}$  for each job type?
  - ► What did you do to convince yourself that your results are valid (hint: consistency check)?
  - ▶ Why are  $\overline{w}$ ,  $\overline{d}$ , and  $\overline{s}$  the same for both job types, while  $\overline{l}$ ,  $\overline{q}$ , and  $\overline{s}$  are different?

Submit your programs, results, and answers to the questions in Blackboard.

## Summary

- Multiple stream RNGs
  - Generate disjoint random number streams
  - Want the streams are far apart
- More simulation examples to be discussed