

Bloom filters

$h1(x) = r$

$h1(y) = r-1$

$f1\ x = 1$

if 0 it definitely doesn't exist, if 1 it might exist.

$h2(x) = s$

$f2\ s = 1$

74000 words = 40kilobytes*10 = 400 KB

unsigned char T[max]; 8 bits in an unsigned char $\log_2(8) = 3$

8*max bits in T

$Bk\ T[k \gg 3]$ is same as $k/8$ kinda you shift over 3 bites on a number which is the same as dividing by 8. Do this because it's faster.

$01 \ll (k \& 07)$

$k \& 07$ gives you the first 3 digits, $k \gg 3$ gives you the numbers after the first 3 digits (aka number divided by 8)

00000111 I want these bits

$(T[k \gg 3] \& 01 \ll (k \& 07))$

ASK ARJUN ABOUT THIS!!!! HE SEEMS TO UNDERSTAND

$(T[k \gg 3] \& (k \& 07)) \& 01$ This allows you to find the bit GETBIT

$T[k \gg 3] |= (01 \ll (k \& 07))$

$T[k \gg 3] \&= \sim(01 \ll (k \& 07))$

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$1001 | 01$

$k \gg 2 | k \& 03$

length of char vector is num of things/8 +1

We're gonna use c++ apparently it looks like Java.

Make a class called bitmap:

Make sure to set everything to 0.

Only ever read the file once.

make dictionary
bad : ungood

Binary trees:

Parent: left child: right child

$l < p < r$

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1 3

is a binary tree, left is less, right is more

```
find(T,X)
    if(T==NULL)
        return not found
    else if(T->key == x)
        return T
    else if(T->key > x)
        return find(T->left,x)
    else
        return find(T->right,x)
```

insert(T,num) //you can always insert as a leaf, always.

As HW write the code for find and insert for Binary trees.

```
inorder(T)
    if(T !=null)
        inorder(T->left)
        print T->key
        inorder(T->right)
preorder(T)
    if(T !=null)
        print T->key
        preorder(T->left)
        preorder(T->right)
postorder(T)
    if(T !=null)
        postorder(T->left)
        postorder(T->right)
        print T->key
```

Maybe look up treaded binary trees

Huffman coding:

this is our alphabet

a b c d e f g

7 42

Huffman coding will be on the final probably

DFS: depth first search

BFS: breadth first search

DFS(preorder traversal) recursion, recursion like stacks

BFS is more like a queue, look at the thing, inque it's children.

While not empty Q:

 x = dequeue

 find? return

 enqueue children

this works for any number of children as does depth first search.

We're doing graphs!!!!

Nodes aka vertices

Lines are edges