


# Numerical Linear Algebra – Matrix interpretation

CS370 Lecture 31 – March 29, 2017

# Cost of Factorization

```
For  $k = 1, \dots, n$   
  For  $i = k + 1, \dots, n$   
     $mult := a_{ik} / a_{kk}$   
     $a_{ik} := mult$   
    For  $j = k + 1, \dots, n$   
       $a_{ij} := a_{ij} - mult * a_{kj}$   
    EndFor  
  EndFor  
EndFor
```

2 FLOPs (1 subtraction, 1 multiply)  
in the innermost loop.



Summing over all the loops we get:

$$\sum_{k=1}^n \sum_{i=k+1}^n \sum_{j=k+1}^n 2 = \frac{2n^3}{3} + O(n^2)$$

The above requires using the following  
sum identities...

$$\sum_{i=1}^n i = \frac{n(n+1)}{2}$$

and

$$\sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}$$

# Cost of Triangular solve

Show that the *total* FLOP count of backward substitution is:  $n^2 + O(n)$  FLOPs.

For  $i = n, \dots, 1$

$x_i := z_i$

For  $j = i + 1, \dots, n$

$x_i := x_i - u_{ij} * x_j$

EndFor

$x_i := x_i / u_{ii}$

EndFor

We previously used these identities:

$$\sum_{i=1}^n i = \frac{n(n+1)}{2}$$

and

$$\sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}$$

# Costs for Solving Linear Systems

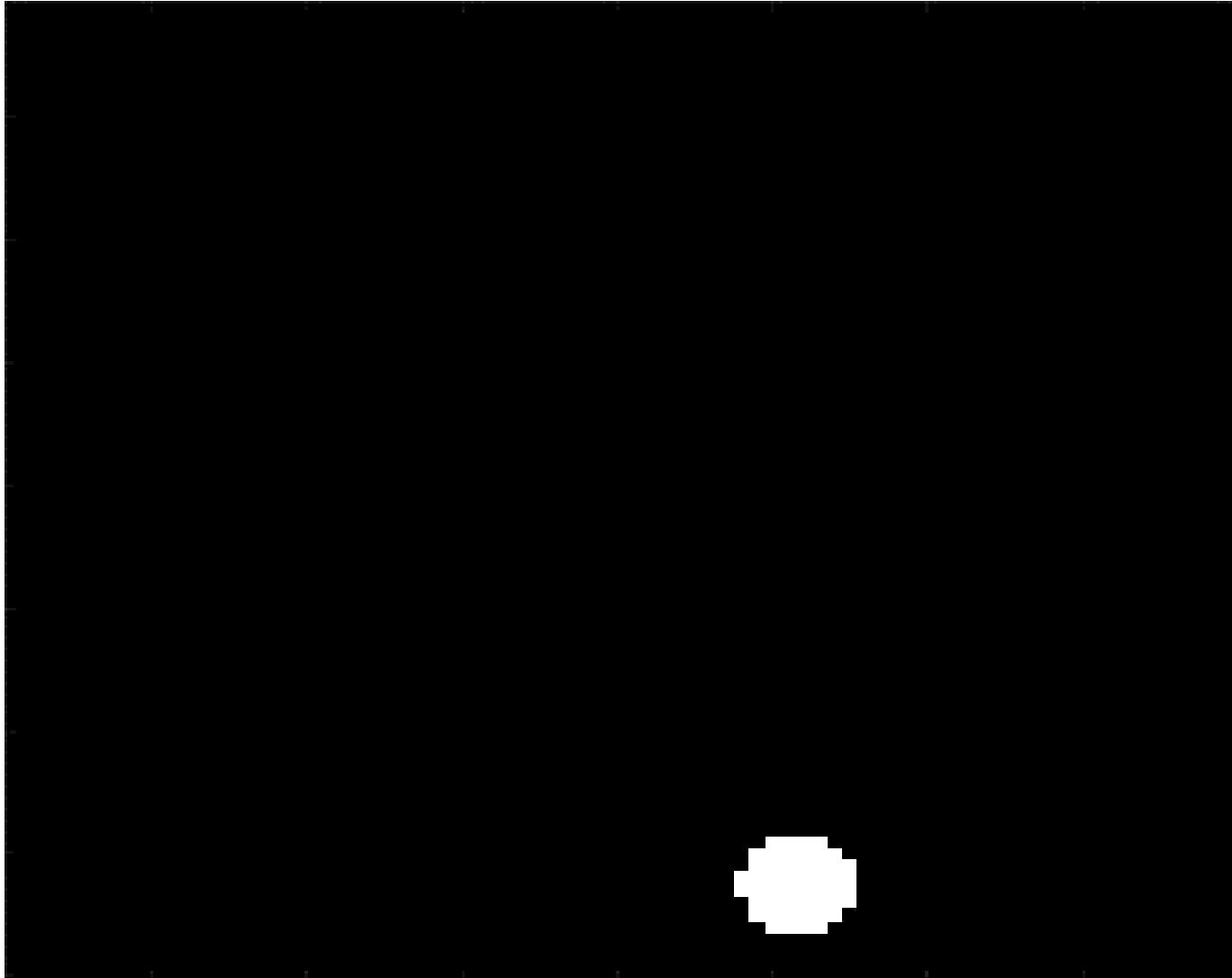
LU factorization costs approximately  $\frac{2n^3}{3} + O(n^2)$  FLOPs.

Triangular solves cost  $n^2 + O(n)$  FLOPs *each*, so total is  $2n^2 + O(n)$ .

Factorization cost dominates when  $n$  is large, and scales worse.

Given an existing factorization, solving for new RHS's is cheaper:  $O(n^2)$ .

# Cost Example: Simple 2D smoke simulation



Compare speed between:

(1) Doing full LU factorization and forward/backward solve each step.

(2) Factoring the matrix once at the start, and reusing the L & U factors.

As expected from the FLOP counts, (2) is dramatically faster.

# Justifying the Factorization View of G.E.

We viewed row-swapping as a matrix; do the same for row-subtraction!

Zeroing a (sub-diagonal) entry of a column by row subtraction can be written as applying a specific matrix  $M$  such that

$$MA^{old} = A^{new}$$

where

- $A^{old}$  is the original matrix.
- $A^{new}$  is the matrix *after* subtracting the specific row.

# Row Subtraction via Matrices

e.g. The operation...

$$(2^{\text{nd}} \text{ row}) := (2^{\text{nd}} \text{ row}) - \frac{a_{2,1}}{a_{1,1}} (1^{\text{st}} \text{ row})$$

can be written as a matrix:

$$M = \begin{bmatrix} 1 & 0 & 0 & 0 \\ -\frac{a_{2,1}}{a_{1,1}} & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

$M$  is the identity matrix, but with a zero entry replaced by the (negative of the) necessary multiplicative factor.

# Row Subtraction via Matrices

Example:

$$\begin{matrix} & M & & A^{old} & & A^{new} \\ \begin{bmatrix} 1 & 0 & 0 \\ -3/2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} & \begin{bmatrix} 2 & 3 & 4 \\ 3 & 4 & -1 \\ -2 & 2 & 1 \end{bmatrix} & = & \begin{bmatrix} 2 & 3 & 4 \\ 0 & -1/2 & -7 \\ -2 & 2 & 1 \end{bmatrix} \end{matrix}$$

So, the whole process of factorization can be viewed as a sequence of matrix (left-)multiplications applied to  $A$ .

The matrix left at the end is  $U$ , so e.g., in 3x3 case, we have shown:

$$M^{(3)}M^{(2)}M^{(1)}A = U$$



# Row Subtraction via Matrices

The matrix left at the end is  $U$ , so e.g., in 3x3 case, we have shown:

$$M^{(3)}M^{(2)}M^{(1)}A = U$$

$$\text{Therefore } A = \left(M^{(3)}M^{(2)}M^{(1)}\right)^{-1}U = \underbrace{\left(M^{(1)}\right)^{-1}\left(M^{(2)}\right)^{-1}\left(M^{(3)}\right)^{-1}}_L U.$$

Define  $L = \left(M^{(1)}\right)^{-1}\left(M^{(2)}\right)^{-1}\left(M^{(3)}\right)^{-1}$  and we have our factorization!

But what is  $\left(M^{(k)}\right)^{-1}$ ?

# Inverse of $M^{(i)}$

The inverse of this simple matrix form is the same matrix, but with the off-diagonal entry ***negated***.

e.g.

$$\begin{bmatrix} 1 & 0 & 0 \\ -3/2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}^{-1} = \begin{bmatrix} 1 & 0 & 0 \\ 3/2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

You can easily verify that:

$$\begin{bmatrix} 1 & 0 & 0 \\ -3/2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 3/2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

This is why we stored  $a_{i,k}/a_{k,k}$  (not  $-a_{i,k}/a_{k,k}$ ) during row subtraction.

# Summary: Row Subtraction via Matrices

The effect of row reduction is to left-multiply  $A$  by a series of matrices  $M^{(k)}$  that comprise  $L^{-1}$  to get  $U$ , recording the entries of  $L$  as we go.

$$M^{(3)}M^{(2)}M^{(1)}A = L^{-1}A = U$$

Interleaving permutation matrices  $P^{(k)}$  before each  $M^{(k)}$  similarly leads to the  $PA = LU$  factorization (but it's a bit trickier to see this.)

\*Note: The course notes combine multiple row subtraction operations for a given column into a single matrix operation  $M$ ; the net result is the same.

# Costs: Solving $Ax = b$ by Matrix Inversion.

An obvious alternative for solving  $Ax = b$  is:

1. Invert  $A$  to get  $A^{-1}$ .
2. Multiply  $A^{-1}b$  to get  $x$ .

One can show that the above is actually **more** expensive (in FLOPs) than using our “factor and triangular solve” strategy.

It also generally incurs more F.P. error.

*Most numerical algorithms avoid ever computing  $A^{-1}$ .*

# Summary: Gaussian elimination

Linear systems like  $Ax = b$  can be efficiently solved by:

1. LU factorization of  $A$ , followed by...
2. Forward/backward substitution.

Adding row pivoting avoids div-by-zero and reduces floating point error.

Total cost is  $\frac{2n^3}{3} + O(n^2) = O(n^3)$  FLOPs, dominated by the factorization.

If RHS changes, the factorization can be reused at lower cost,  $O(n^2)$ .

# Final topic: Norms and Conditioning



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# Norms and Conditioning

**Norms** are measurements of “size”/magnitude for vectors or matrices.

We will start by defining some norms.

Then, we will use these norms to explore **conditioning** of matrices.

Conditioning describes how the output of a function/operation/matrix changes due to changes in input.

# Vector Norms

There are many reasonable norms for a vector  $\mathbf{x} = [x_1, x_2, \dots, x_n]^T$ .  
Common choices are:

1-norm (or taxicab/manhattan norm):  $||\mathbf{x}||_1 = \sum_{i=1}^n |x_i|$

2-norm (or Euclidean norm):  $||\mathbf{x}||_2 = \sqrt{\sum_{i=1}^n x_i^2}$

$\infty$ -norm (or max norm):  $||\mathbf{x}||_\infty = \max_i |x_i|$



# $p$ -norms

These are collectively called  $p$ -norms, for  $p = 1, 2, \infty$  and can be written:

$$||\mathbf{x}||_p = \left( \sum_{i=1}^n |x_i|^p \right)^{\frac{1}{p}}$$

(Only holds *in the limit* for  $\infty$  case).

Matlab has a **norm()** command that implements these vector norms.

# Some Properties of Norms

If the norm is zero, then the vector must be the zero-vector.

$$||\mathbf{x}|| = 0 \rightarrow x_i = 0 \quad \forall i.$$

The norm of a scaled vector must satisfy

$$||\alpha \mathbf{x}|| = |\alpha| \cdot ||\mathbf{x}|| \quad \text{for scalar } \alpha.$$

The triangle inequality holds:

$$||\mathbf{x} + \mathbf{y}|| \leq ||\mathbf{x}|| + ||\mathbf{y}||$$

# Defining Norms for Matrices

Matrix norms are often defined/“induced” as follows, *using* p-norms of vectors:

$$||A|| = \max_{||x|| \neq 0} \frac{||Ax||}{||x||}$$

A bit tricky... clearly we can't try out all possible  $x$  to determine this!

But, there are simpler equivalent definitions in some cases:

$$||A||_1 = \max_j \sum_{i=1}^n |A_{ij}|$$

(max absolute column sum)

$$||A||_\infty = \max_i \sum_{j=1}^n |A_{ij}|$$

(max absolute row sum)

# Matrix 2-norm (or *spectral* norm)

Using the vector 2-norm, we get the matrix 2-norm, or spectral norm.

$$||A||_2 = \max_{||x|| \neq 0} \frac{||Ax||_2}{||x||_2},$$

The matrix's 2-norm relates to the eigenvalues.

Specifically, if  $\lambda_i$  are the eigenvalues of  $A^T A$ , then

$$||A||_2 = \max_i \sqrt{|\lambda_i|}$$

# Some Matrix Norm Properties

$$||A|| = 0 \leftrightarrow A_{ij} = 0 \forall i, j.$$

$$||\alpha A|| = |\alpha| \cdot ||A|| \text{ for scalar } \alpha$$

$$||A + B|| \leq ||A|| + ||B||$$

$$||A\mathbf{x}|| \leq ||A|| \cdot ||\mathbf{x}||$$

$$||AB|| \leq ||A|| \cdot ||B||$$

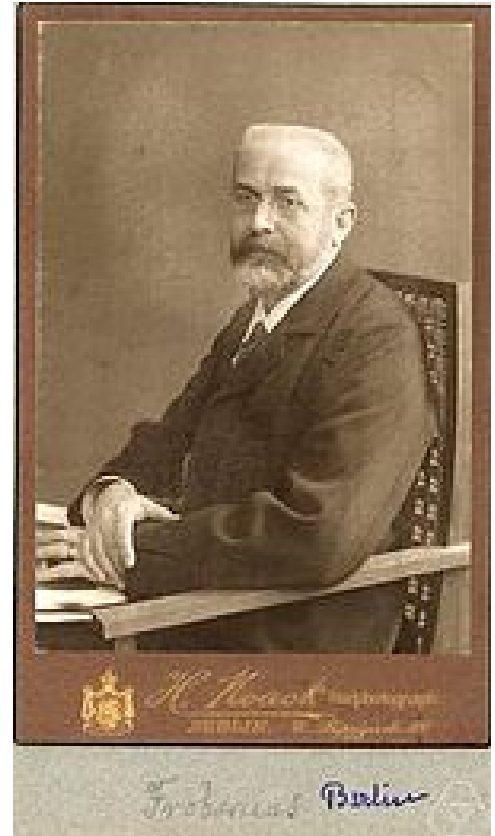
$$||I|| = 1$$

# Bonus: Frobenius Norm

Another fun matrix norm that is not “induced” by a standard vector norm is the *Frobenius* norm.

$$\|A\|_F = \sqrt{\sum_{i=1}^n \sum_{j=1}^n A_{ij}^2}$$

(Like a “2D” version of the 2-norm for vectors.)



Dr. Frobenius

# Next Time: Conditioning of Linear Systems

Conditioning describes how the output of a function/operation/matrix changes due to changes in input.

Conditioning is indicative of how difficult a problem is to solve, *independent* of the algorithm / numerical method used.

Norms are a useful tool to help characterize the conditioning of linear systems.