Ordinary Differential Equations – Adaptive Time-Stepping CS370 Lecture 16 – Feb 8, 2017

Time Step Control

Time stepping methods advance by a discrete time step size h.

So far, we have held h as a prescribed constant for a whole calculation.

Since our error was $O(h^p)$ for some p, depending on the scheme, smaller $h \to 0$ implied smaller error (both local and global).

Time Step Tradeoffs

If small step h yields small error, is smaller always "better" for practical calculations?

No! Smaller h requires more steps to reach a given time t_{final} ,

$$\#Steps = \frac{t_{final} - t_0}{h} = O(h^{-1}),$$

implying a higher computational cost.

But too large h introduces excessive error! (Not to mention stability issues for some schemes...)

Time Step Tradeoffs

 $h \to 0$: many steps, computationally expensive, more accurate.

 $h \rightarrow \text{large}$: fewer steps, so cheaper, but also less accurate.

 \therefore should minimize effort/cost by choosing largest h such that error is less than a desired error tolerance:

error < tolerance.

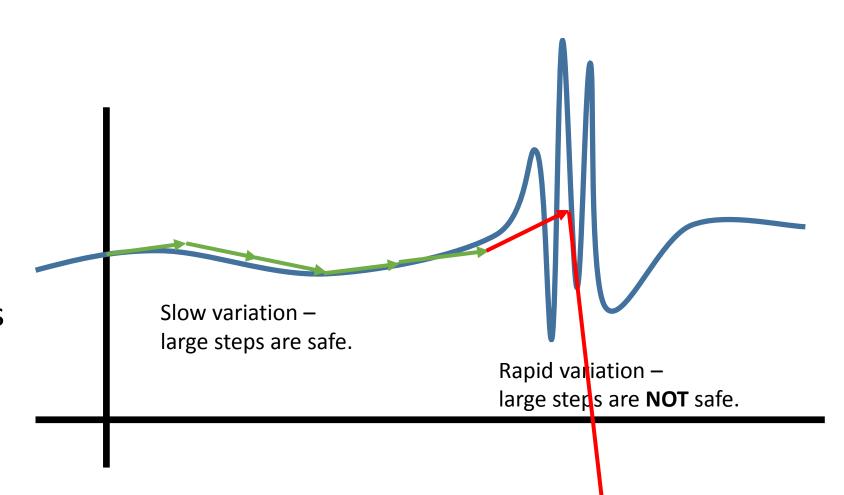
Can we work out an ideal constant h for a given problem? Not necessarily a good idea...

Adaptive Time-Stepping - Motivation

Rate of change of the solution often varies over time!

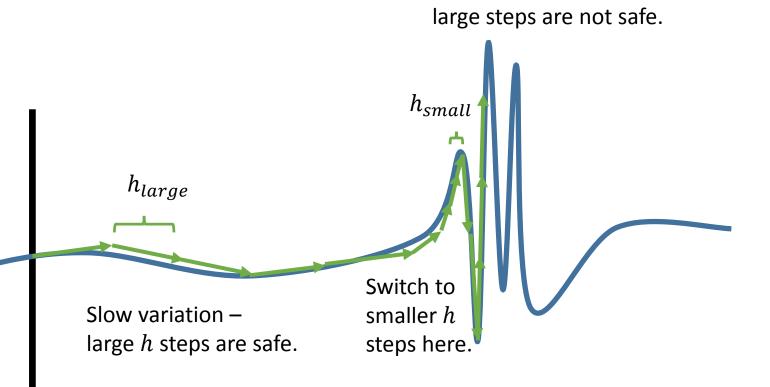
e.g.

- Seasonal variations.
- Sudden strong forces in a physical system.
- Shocks to the stock market.



Adaptive Time-Stepping - Motivation

Idea: adapt the time step during the computation to minimize error while avoiding wasted effort.



Rapid variation –

Adaptive Time-Stepping – Concept

Another "chicken-and-egg" problem?

If we knew the error for a given h, we could choose a good h to always satisfy error < tolerance. But...

...if we already knew the error, we'd know the solution too!

Approach: Use **two** different time-stepping schemes **together**.

Compare their results to *estimate* the error, and adjust h accordingly.

Let's derive a justification for this!

Adaptive Time-Stepping - Basic Algorithm

- 1. Compute approximate solutions with two schemes of different orders.
- 2. Estimate the error by taking their difference.
- 3. While(error > tolerance)
 - Set h := h/2, and recompute the solutions (1.) & error (2.).
- 4. Estimate error coefficient, and predict a good next stepsize h_{new} .
- 5. Repeat until end time is reached.

Example: Matlab

Matlab's **ODE45** routine uses "Runge-Kutta-Fehlberg" scheme. It uses 4th and 5th order Runge-Kutta schemes together:

Matlab's **ODE23** is similar but uses a lower-order pair.

Doesn't this double the cost?

No: evaluation points of f are carefully chosen/weighted to avoid duplicate calls (calling f is often the main cost).

$$k_1 = hf(t_n, y_n)$$

$$k_2 = hf(t_n + \frac{h}{4}, y_n + \frac{k_1}{4})$$

$$k_3 = hf(t_n + \frac{3h}{8}, y_n + \frac{3k_1}{32} + \frac{9k_2}{32})$$

$$k_4 = hf(t_n + \frac{12h}{13}, y_n + \frac{1932k_1}{2197} - \frac{7200k_2}{2197} + \frac{7296k_3}{2197})$$

$$k_5 = hf(t_n + h, y_n + \frac{439k_1}{216} - 8k_2 + \frac{3680k_3}{513} - \frac{845k_4}{4104})$$

$$k_6 = hf(t_n + \frac{h}{2}, y_n - \frac{8k_1}{27} + 2k_2 - \frac{3544k_3}{2565} + \frac{1859k_4}{4104} - \frac{11k_5}{40})$$

$$y_{n+1}^* = y_n + \frac{25k_1}{216} + \frac{1408k_3}{2565} + \frac{2197k_4}{4104} - \frac{k_5}{5} \text{ with error } O(h^4)$$

$$y_{n+1} = y_n + \frac{16k_1}{135} + \frac{6656k_3}{12825} + \frac{28561k_4}{56430} - \frac{9k_5}{50} + \frac{2k_6}{55} \text{ with error } O(h^5).$$

Time Step Control

Summary:

Adapting the timestep h based lets us avoid excess computational cost, while approximately satisfying a given error tolerance.

i.e. we expend **no more effort than we have to** for our answer.