

CS370 Lecture 1

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Five Topics in the Course (Jan 4th)

- Floating point numbers and Arithmetic
- Iterpolation, Splines, Parametric Curves
- Initial Value Problems - solve differencial equations
- Discrete Fourier Analysis
- Numerical Linear Algebra - solve equations - google pagerank

Topic 1: Floating Point Arithmetic

Examples where problems come whehn using approximation

eg1. $e^{-5.5} =$

$$e^x = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots + \frac{x^n}{n!} + \dots$$

$$e^{-x} = \frac{1}{e^x}$$

$$e^{-5.5} = \frac{1}{e^{5.5}} = \frac{1}{1 + 5.5 + \frac{5.5^2}{2} + \dots}$$

Now do arithmetic keeping only 5 digits. In both cases infinite sums remain unchanged after 25 terms. There is no sense in going any further - we end up just truncating all of the terms after this as they are smaller than the 5th digit.

Method 1 gives $e^{-5.5} = 0.0026363$

Method 2 gives $e^{-5.5} = 0.0040868$

eg2. $ax^2 + bx + c = 0$

1

$$x_1 = \frac{-b + \sqrt{b^2 - 4ac}}{2a}$$

2

$$x_2 = \frac{-b - \sqrt{b^2 - 4ac}}{2a}$$

$$x_1 = \frac{2a}{-b + \sqrt{b^2 - 4ac}}$$

$$x_2 = \frac{2a}{-b - \sqrt{b^2 - 4ac}}$$

$$x^2 + 621x + 1 = 0$$

use 4 digit arithmetic

$$1: x_1 = -0.200 \quad x_2 = 62.1$$

$$2: x_1 = -0.0161 \quad x_2 = -62.1$$

Floating point number systems (Jan 6th)

$$F(\beta, t, L, u)$$

Floating point number is either

$$0 + / - 0.x_1x_2\dots x_t$$

Where x_n is a digit

- $x_1 \neq 0$
- $0 \leq x_i < \beta$
- $L \leq d \leq u$

eg $F(10, 6, L, u)$ for $\pi = 0.314519x10^1$

Single Precision: $[+/-][8 \text{ bits}][23 \text{ digits}] F(2, 24, -126, 127)$

Double Precision: $[+/-][11 \text{ bits}][52 \text{ digits}] F(2, 53, -1022, 1023)$

Given any real number:

$$+ / - 0.x_1x_2\dots x_tx_{t+1}\dots x\beta^d$$

you can truncate (rounding possible):

$$fl(x) = + / - 0.x_1x_2...x_t$$

Question: How close is $fl(x)$ to x ? Relatively!

$$\delta = \frac{fl(x) - x}{x}$$

Claim: $|\delta| < \beta^{1-t}$ when truncating and $|\delta| < \frac{1}{2}\beta^{1-t}$ when rounding (Let this definition be ϵ)

Truncating:

$$|\delta| = \frac{0.00...0x_{t+1}...x\beta^d}{0.x_1x_2...x\beta^d}$$

$$|\delta| = \frac{0.x_{t+1}x_{t+2}...x\beta^{-t}}{x_1.x_2x_3...x\beta^{-1}}$$

$$|\delta| < \beta^{1-t}$$

In general $|\delta| < \epsilon$

$$fl(x) = x(1 + \delta)$$

$$|\delta| < \epsilon$$

What about arithmetic an errors?

x, y real

$$x \oplus y =$$

eg. $x = 0.1111... F(2,4,L,u)$

$$y = 0.1110...$$

$$fl(x) = 0.1111$$

$$fl(y) = 0.1110$$

when adding them: 1.1101 then truncate to 0.1110×2^1

x, y real

$$x \oplus y = fl(fl(x) + fl(y))$$

$$\begin{aligned}
& \left| \frac{x \oplus y - (x + y)}{x + y} \right| \\
&= \left| \frac{fl(fl(x) + fl(y)) - (x + y)}{x + y} \right|
\end{aligned}$$

since $|\delta_1| < \epsilon$ and $|\delta_2| < \epsilon$ and $|\delta_3| < \epsilon$

$$\begin{aligned}
&= \left| \frac{(x(1 + \delta_1) + y(1 + \delta_2))(1 + \delta_3) - (x + y)}{x + y} \right| \\
&= \left| \frac{x + y + x\delta_1 + y\delta_2 + x\delta_3 + y\delta_3 + x\delta_1\delta_3 + y\delta_2\delta_3 - x - y}{x + y} \right| \\
&= \left| \frac{x\delta_1 + y\delta_2 + x\delta_3 + y\delta_3 + x\delta_1\delta_3 + y\delta_2\delta_3}{x + y} \right| \\
&\leq \frac{|x||\delta_1| + |y||\delta_2| + |x||\delta_3| + |x||\delta_1||\delta_3| + |y||\delta_2||\delta_3|}{|x + y|}
\end{aligned}$$

Since $|a + b| \leq |a| + |b|$ and $|a \times b| = |a||b|$

$$< \frac{(|x| + |y|)}{|x + y|} (2\epsilon + \epsilon^2)$$

$$\begin{aligned}
& \left| \frac{x \oplus y - (x + y)}{x + y} \right| \\
&< \frac{|x| + |y|}{|x + y|} (2\epsilon + \epsilon^2)
\end{aligned}$$

If x and y have the same sign then $|x + y| = |x| + |y|$ and so addition has relative error bounded by $2\epsilon + \epsilon^2$

But if x and y are of opposite sign and perhaps of nearly the same size:

$$\frac{|x| + |y|}{|x + y|}$$

CATASTROPHIC CANCELLATION - Subtraction is deadly.

Jan 9th:

Recall: $F(\beta, t, L, u)$

For any non-zero real number

$$x = 0.x_1x_2\dots x\beta^d$$

$$fl(x) = 0.x_1x_2\dots x_t \text{ then the exponent } x\beta^d$$

We should the magnitude of the error:

$$\delta = \frac{fl(x) - x}{x}$$

showed: $|\delta| < \epsilon = \beta^{1-t}$ or $\epsilon = \frac{1}{2}\beta^{1-t}$

ϵ is called machine epsilon.

$$fl(x) = x(1 + \delta) \text{ where } |\delta| < \epsilon$$

Operations

If x and y are two real numbers

$$x \oplus y = fl(fl(x) + fl(y))$$

$$\left| \frac{(x \oplus y) - (x + y)}{x + y} \right| < (2\epsilon + \epsilon^2) \left(\frac{|x| + |y|}{|x + y|} \right)$$

Notice: if x and y are of the same sign then relative error of addition $< 2\epsilon + \epsilon^2$

If x and y are of opposite sign and $x + y \approx 0$ then the upper bound can be significant. This is called catastrophic cancellation

$$x = 0.x_1x_2\dots x_t\dots x\beta^d$$

$$y = -0.x_1x_2\dots \hat{x}_t\dots x\beta^d$$

Subtracted = 0.0.....0? with exponent $x\beta^d$

$$x = 0.1239\dots 1231$$

$$y = -0.1229\dots 1221$$

Subtracted = 0.0010

$$e^{-5.5} = 1 - 5.5 + \frac{5.5^2}{2} - \frac{5.5^3}{6} + \dots = x - \hat{x}$$

$$e^{-5.5} = \frac{1}{e^{5.5}} = \frac{1}{1 + 5.5 + \frac{5.5^2}{2} + \dots}$$

Numerical Algorithm Issues

Problem: given α Solve numerically:

$$I_n = \int_0^1 \frac{x^n}{x + \alpha} dx$$

$n = 0, 1, 2, \dots$

Method

$$\begin{aligned} I_0 &= \int_0^1 \frac{1}{x + \alpha} dx = \ln(x + \alpha) \Big|_0^1 \\ &= \ln(1 + \alpha) - \ln(\alpha) = \ln\left(\frac{1 + \alpha}{\alpha}\right) \end{aligned}$$

$$I_n = \frac{1}{n} - \alpha I_{n-1}$$

For $n \geq 1$:

$$\begin{aligned} I_n &= \int_0^1 \frac{x^{n-1}(x + \alpha - \alpha)}{x + \alpha} dx \\ &= \int_0^1 x^{n-1} dx - \alpha \int_0^1 \frac{x^{n-1}}{x + \alpha} dx \end{aligned}$$

$\alpha = 0.5$

$$I_0 = \ln(3) = 1.09861228866\dots$$

$$I_0^a pp = 1.098612$$

$$I_1^{00^a} pp = 0.00664$$

$\alpha = 2.0$

$$I_0 = \ln(1.5) = 0.405465108108\dots$$

$$I_0^a pp = 0.4054654$$

$$I_1 00^a pp = 2.1 \times 10^{22}$$

He used his own application and went through the solutions to the above.
They are incorrect :S

$$\text{Math: } I_0, I_n = \frac{1}{n} - \alpha I_{n-1}$$

$$\text{CS: } I_0^{app}, I_n^{app} = \frac{1}{n} - \alpha I_{n-1}^{app}$$

$$\begin{aligned} e_0 &= I_0 - I_0^{app} = \left(\frac{1}{n} - \alpha I_{n-1}\right) - \left(\frac{1}{n} - \alpha I_{n-1}^{app}\right) \\ &= -\alpha I_n + \alpha I_{n-1}^{app} = -\alpha e_{n-1} \end{aligned}$$

$$e_n = \alpha e_{n-1}$$

$$= \alpha^2 e_{n-2}$$

$$= -\alpha^3 e_{n-3}$$

...

$$= (-\alpha)^n e_0$$

IF $|\alpha| < 1$ then $|e_n| \rightarrow 0$ as $n \rightarrow \infty$ stable

$|\alpha| > 1$ then $|e_n| \rightarrow \infty$ is $n \rightarrow \infty$ unstable!!!

When $|\alpha| > 1$ then we can work backwards!

$$I_{200}, I_{199}, \dots, I_{101} I_{100}$$

$$I_{n-1} = \frac{1}{\alpha n} - \frac{1}{\alpha} I_n$$

$$e_{n-1} = \frac{1}{2} e_n$$

$$|e_{199}| = \frac{1}{|\alpha|} |e_{200}|$$

$$|e_{100}| = \frac{1}{|\alpha|^{100}} |e_{200}|$$