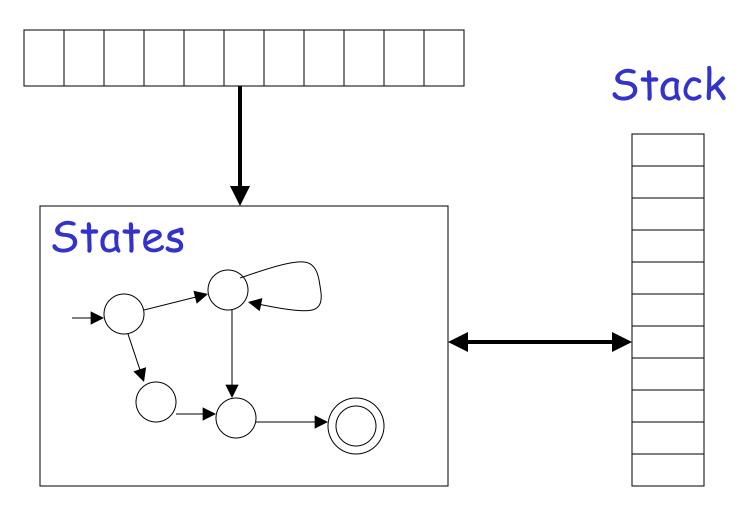
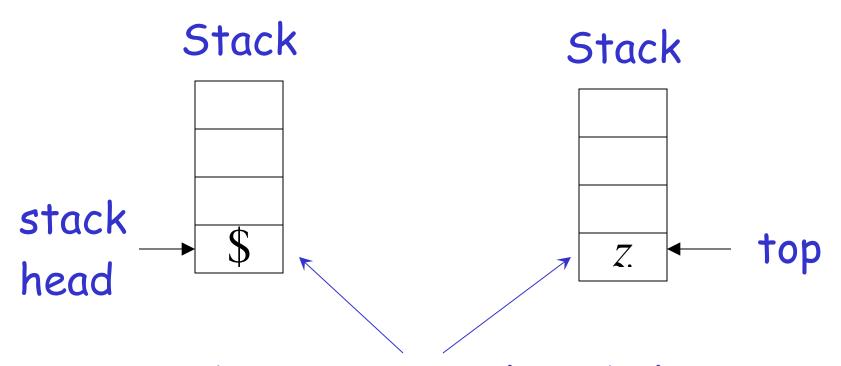
Pushdown Automaton -- PDA

Input String

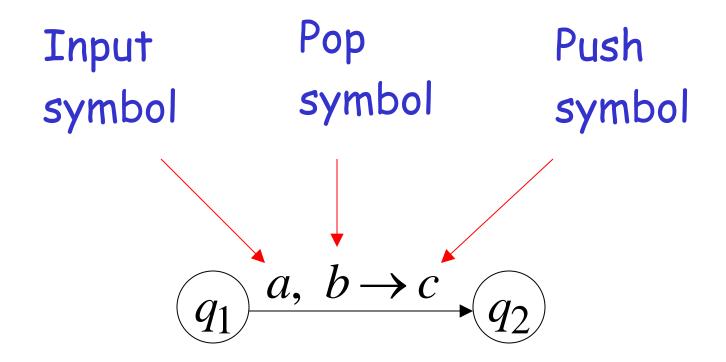


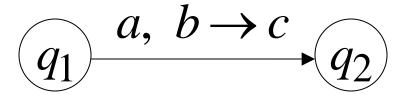
Initial Stack Symbol

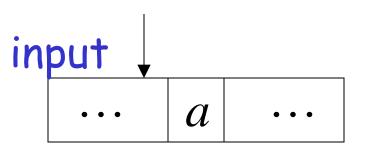


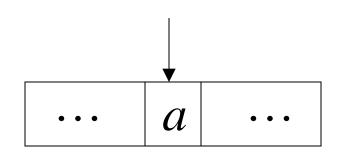
bottom special symbol Appears at time 0

The States

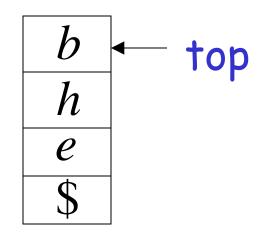




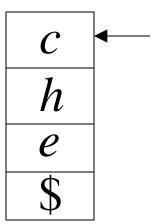


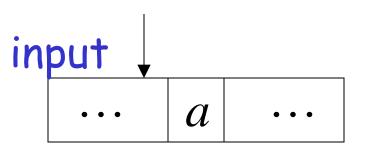


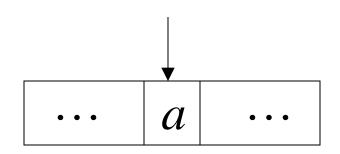
stack



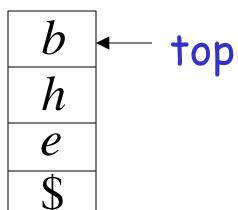


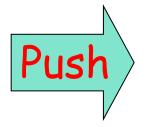


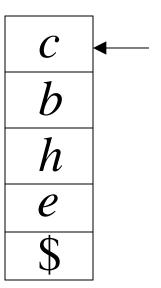


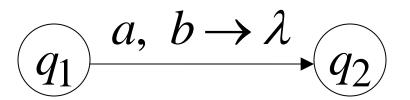


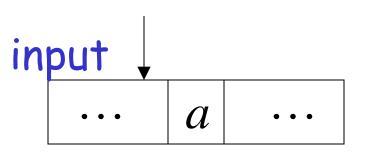


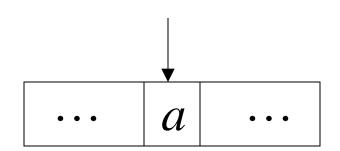




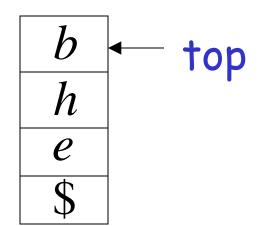




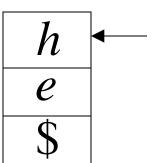


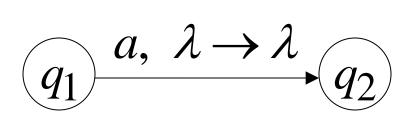


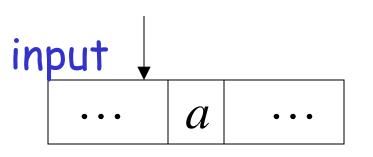
stack

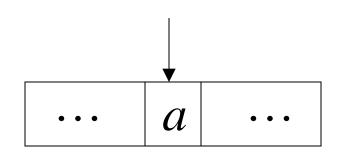








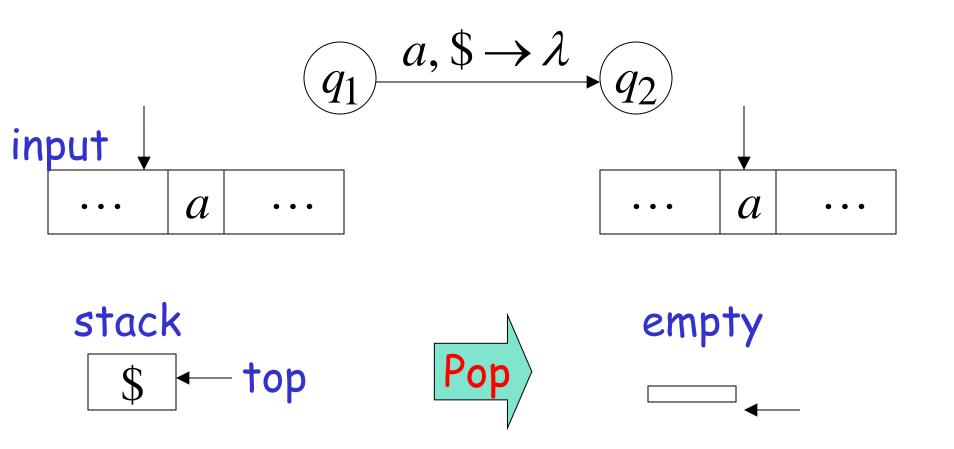




stack

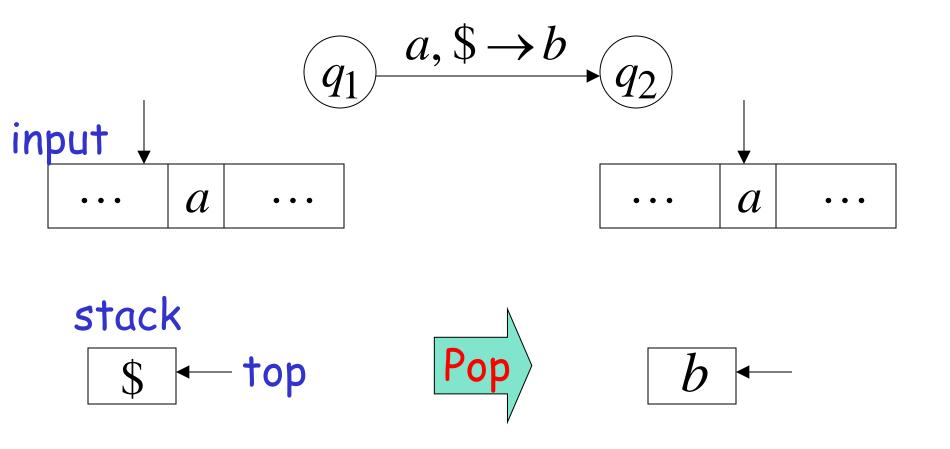


Empty Stack



The automaton HALTS No possible transition after q_2

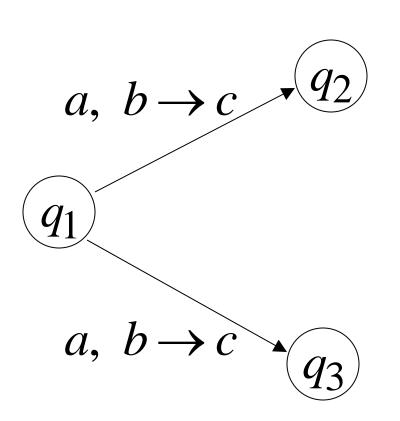
A Possible Transition



Non-Determinism

PDAs are non-deterministic

Allow non-deterministic transitions

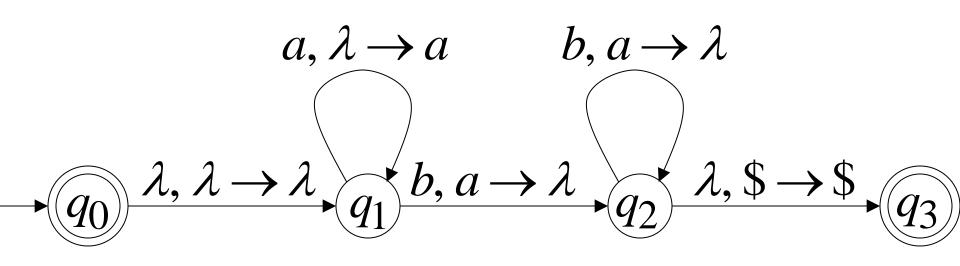


$$\lambda$$
 – transition

Example PDA: Basic Idea?

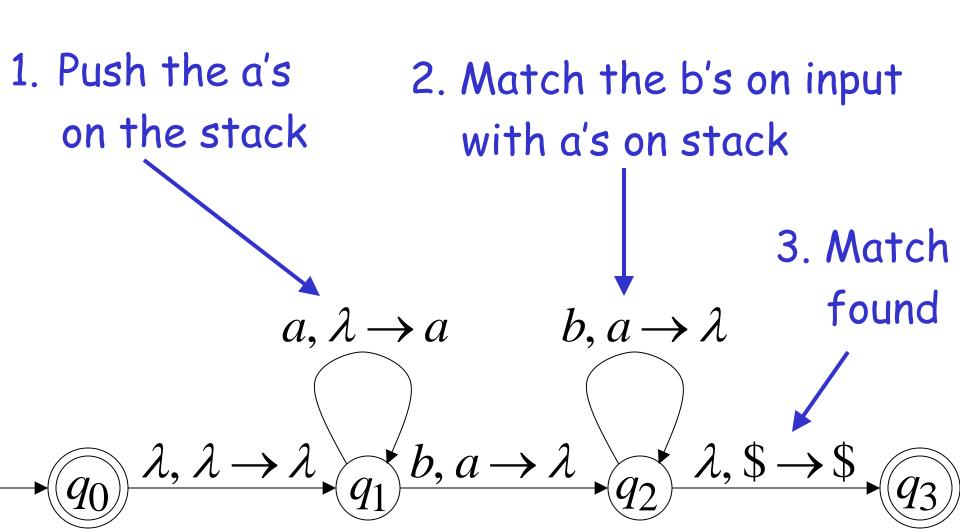
PDA M

$$L(M) = \{a^n b^n : n \ge 0\}$$



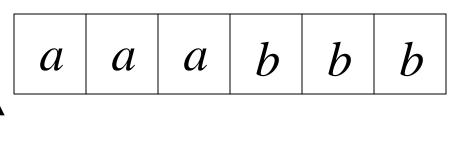
$$L(M) = \{a^n b^n : n \ge 0\}$$

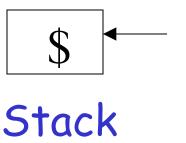
Basic Idea:

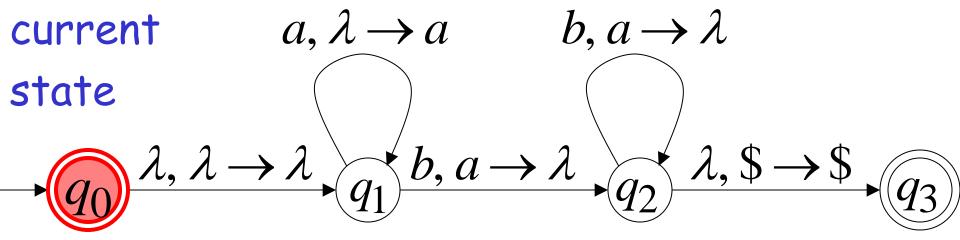


Execution Example: Time 0

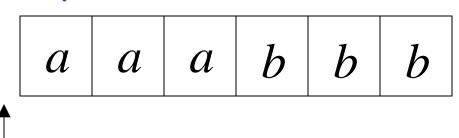
Input

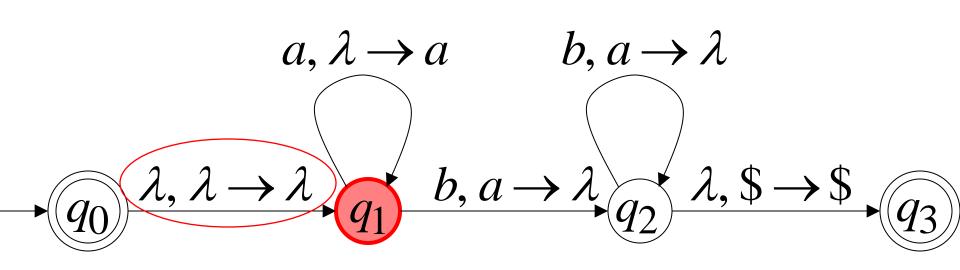




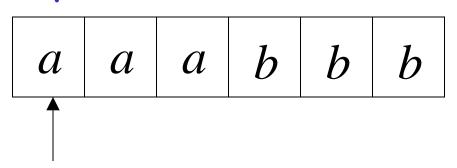


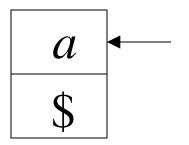
Input

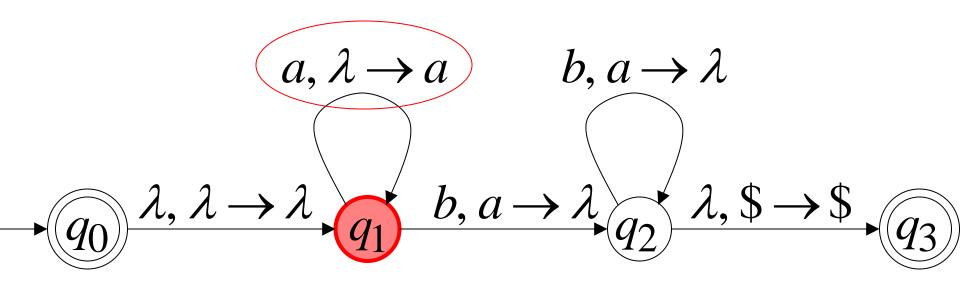




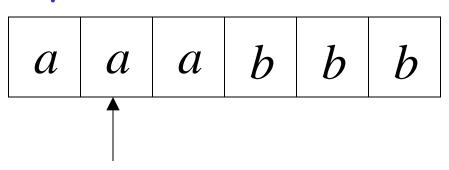
Input

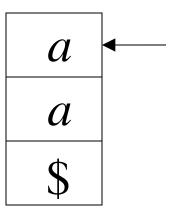


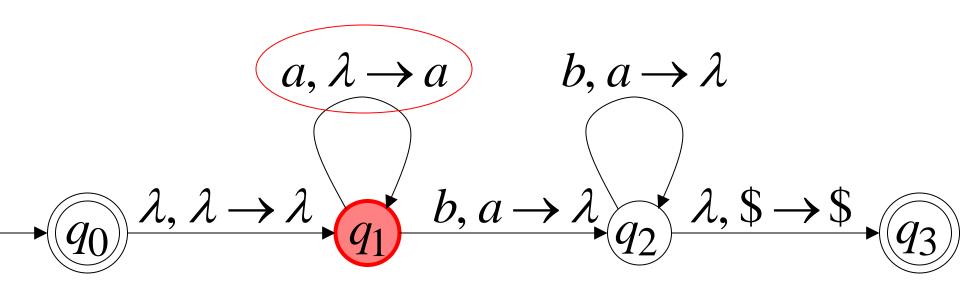




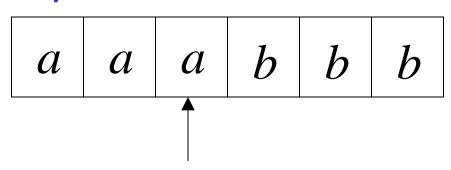
Input

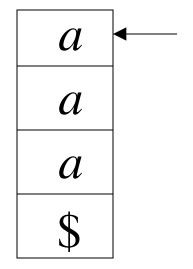


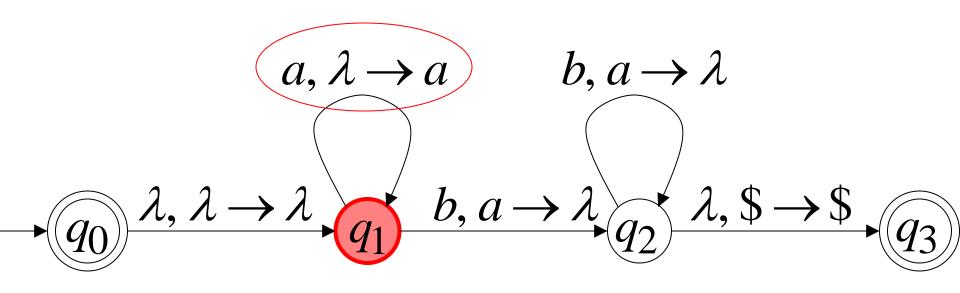




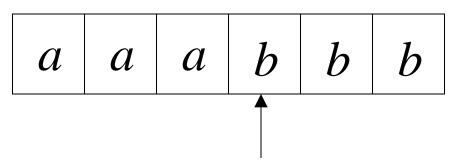
Input

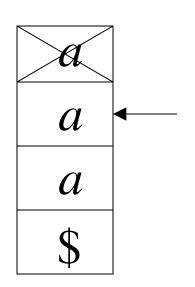


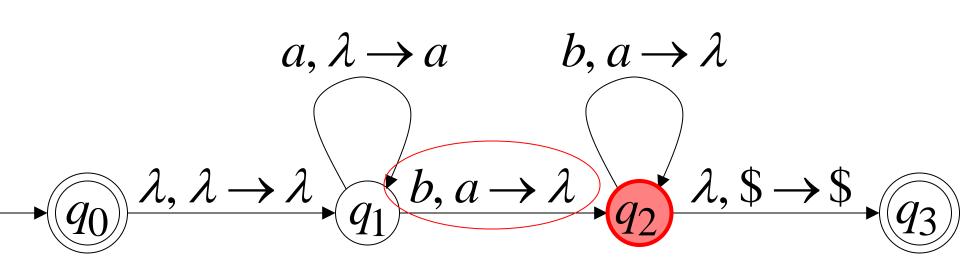




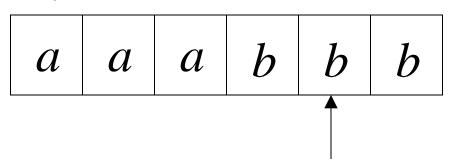
Input

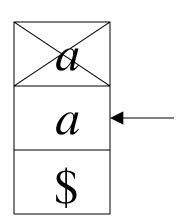


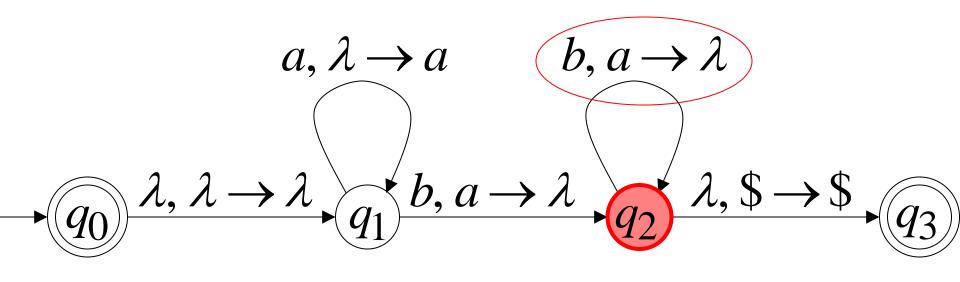




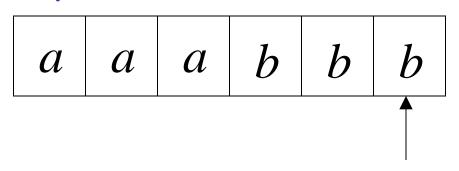
Input

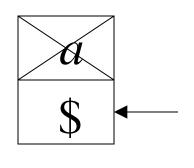


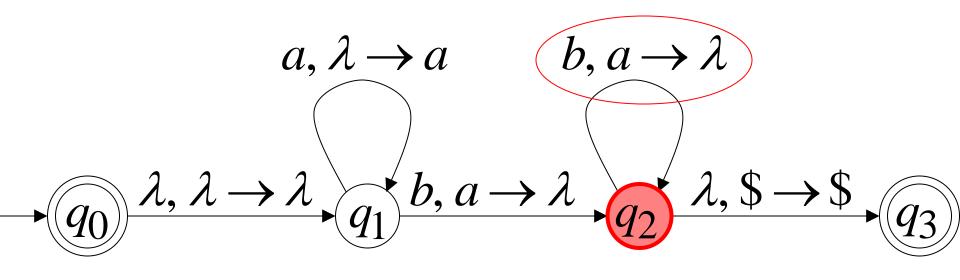




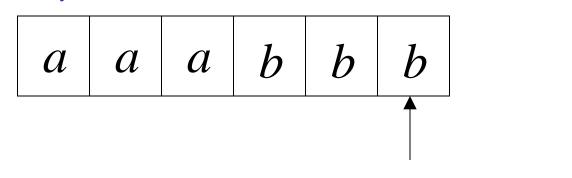
Input

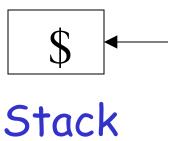


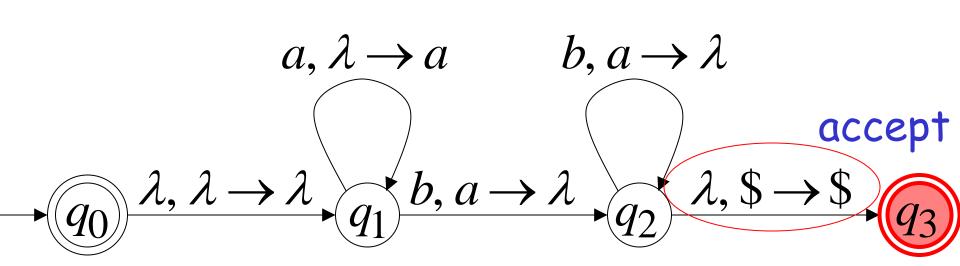




Input







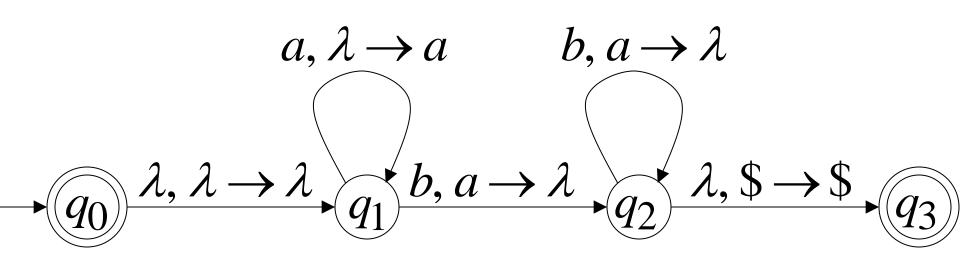
A string is accepted if there is a computation such that:

All the input is consumed AND

The last state is an accepting state

At the end of the computation, we do not care about the stack contents (the stack can be empty at the last state)

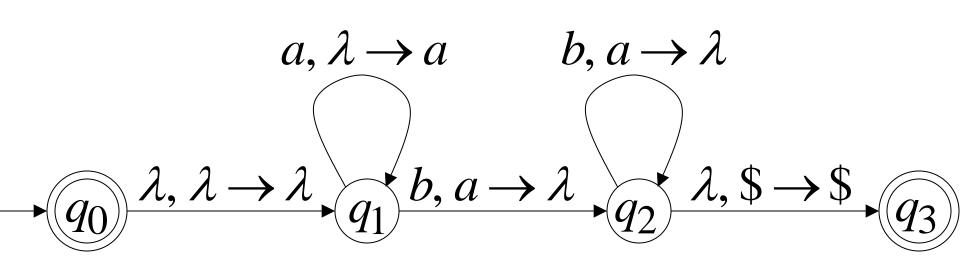
The input string aaabbb is accepted by the PDA:



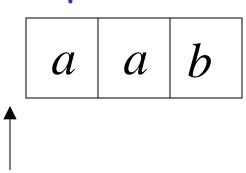
In general,

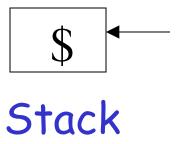
$$L = \{a^n b^n : n \ge 0\}$$

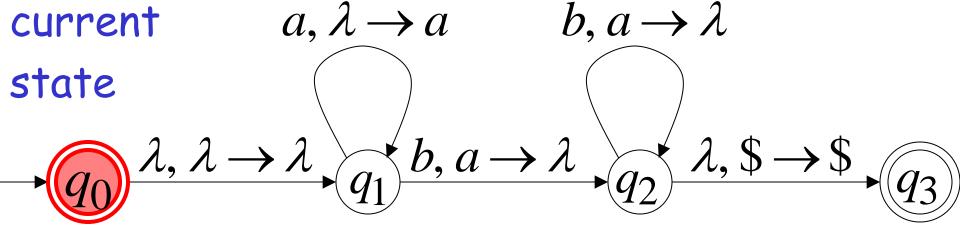
is the language accepted by the PDA:



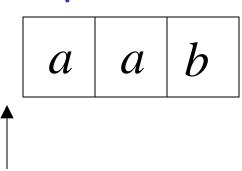
Input

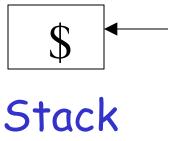


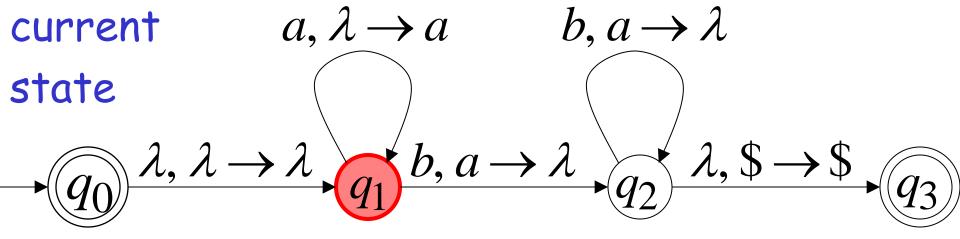




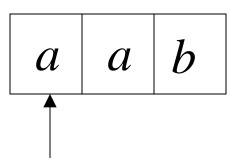
Input

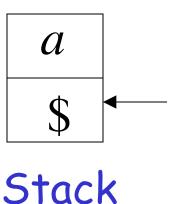




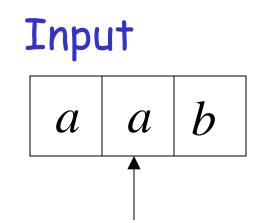


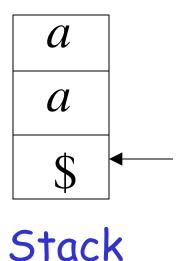
Input

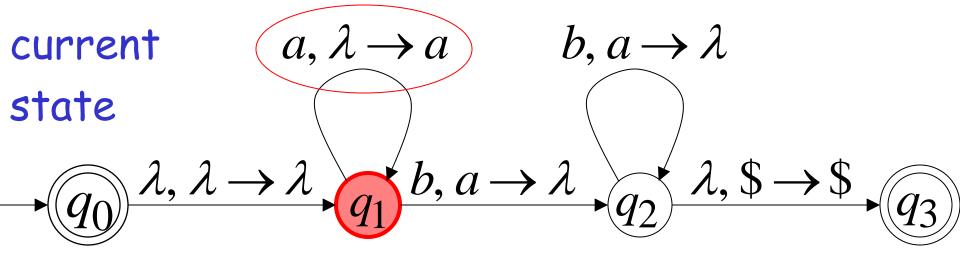




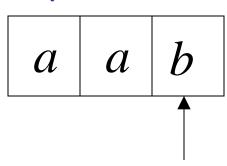
current $a, \lambda \to a$ $b, a \to \lambda$ state $b, a \to \lambda$ $a \to \lambda$ $b, a \to \lambda$

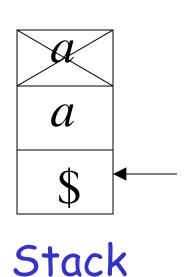


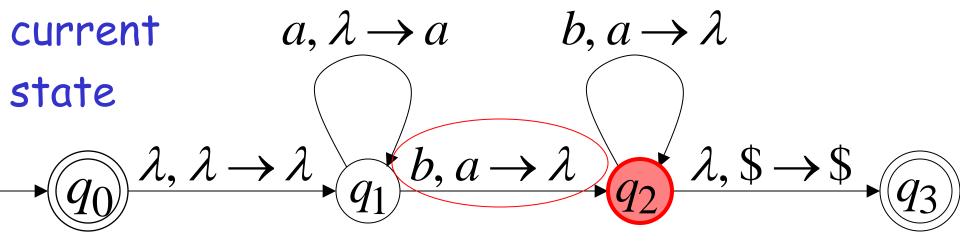




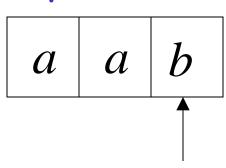
Input

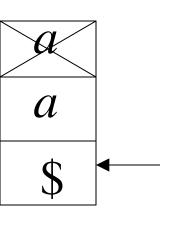






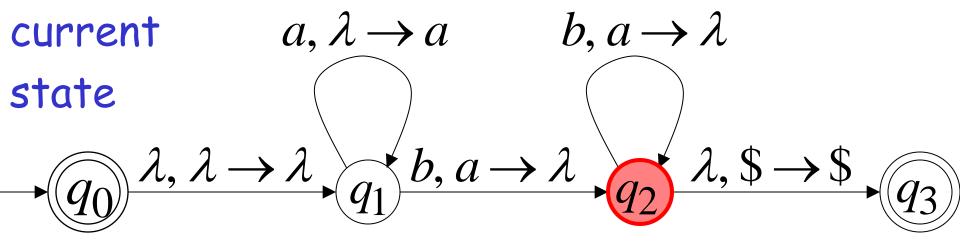




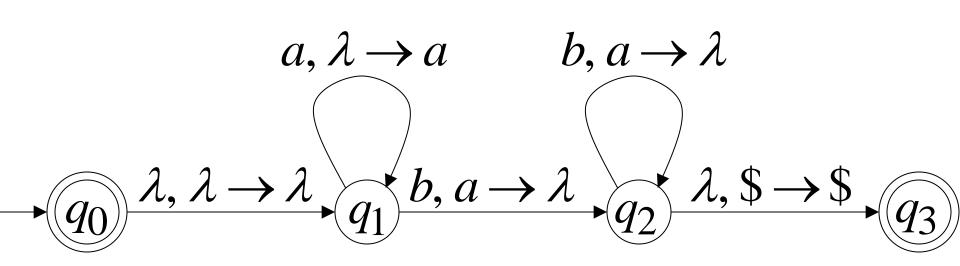


Stack

reject



The input string aab is rejected by the PDA:



A string is rejected if there is **no** computation such that:

All the input is consumed AND

The last state is an accept state

At the end of the computation, we do not care about the stack contents

Another PDA example: Language?

PDA M

$$a, \lambda \rightarrow a$$
 $a, a \rightarrow \lambda$
 $b, \lambda \rightarrow b$ $b, b \rightarrow \lambda$
 $\downarrow q_0$ $\lambda, \lambda \rightarrow \lambda$ $\downarrow q_1$ $\lambda, \$ \rightarrow \$$ $\downarrow q_2$

Another PDA example

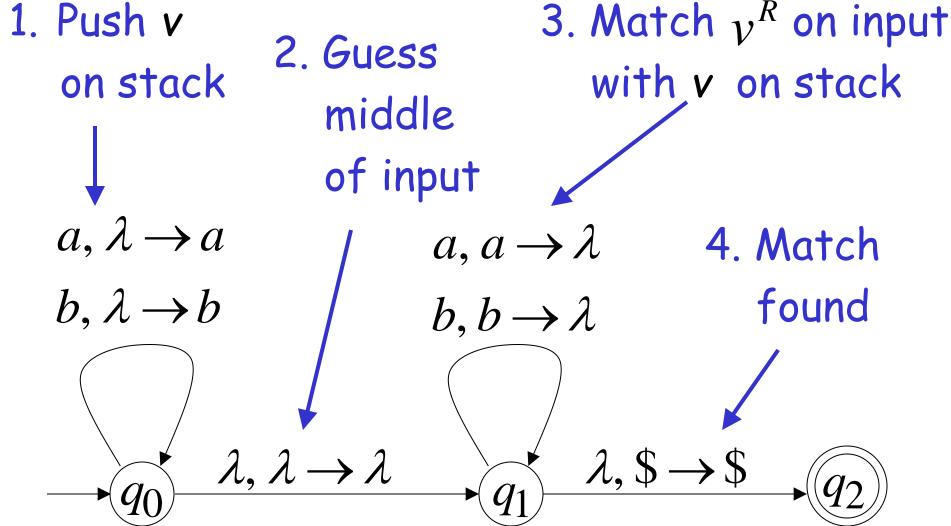
$$L(M) = \{vv^R : v \in \{a,b\}^*\}$$

PDA M

$$a, \lambda \rightarrow a$$
 $a, a \rightarrow \lambda$
 $b, \lambda \rightarrow b$ $b, b \rightarrow \lambda$
 q_0 $\lambda, \lambda \rightarrow \lambda$ q_1 $\lambda, \$ \rightarrow \$$ q_2

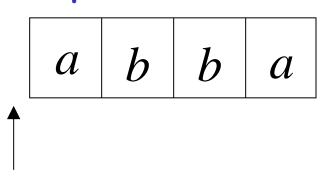
Basic Idea:

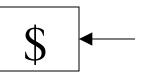
$$L(M) = \{vv^R : v \in \{a,b\}^*\}$$



Execution Example: Time 0

Input



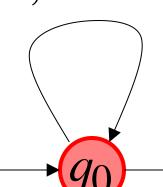


$$a, \lambda \rightarrow a$$

$$b, \lambda \rightarrow b$$

$$a, a \rightarrow \lambda$$

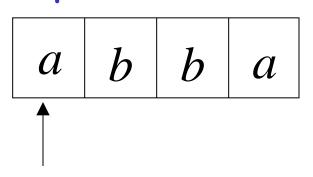
$$b, b \rightarrow \lambda$$

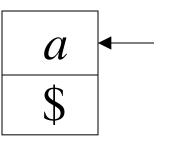


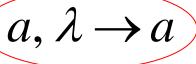
$$\lambda, \lambda \rightarrow \lambda$$

$$\lambda, \$ \rightarrow$$

Input



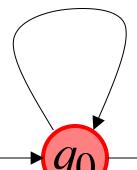




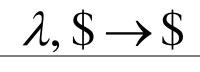
$$(a, a \rightarrow \lambda)$$

$$b, \lambda \rightarrow b$$

$$b, b \rightarrow \lambda$$

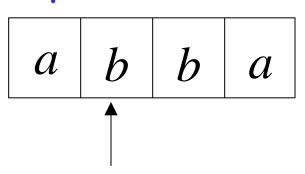


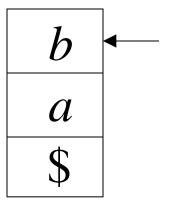
$$\lambda, \lambda \rightarrow \lambda$$





Input





Stack

$$\begin{array}{c}
a, \lambda \to a \\
b, \lambda \to b
\end{array}$$

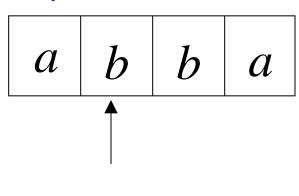
$$\begin{array}{c}
\lambda, \lambda \to \lambda
\end{array}$$

$$a, a \to \lambda$$

$$b, b \to \lambda$$

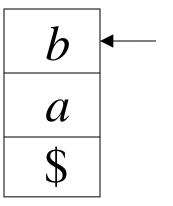
 λ , \$ \rightarrow \$

Input



 $\lambda, \lambda \rightarrow \lambda$

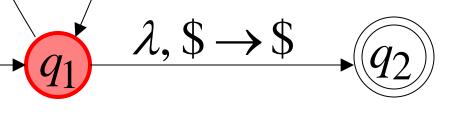
Guess the middle of string



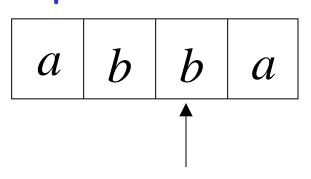
 $a, \lambda \rightarrow a$ Stack $a, a \rightarrow \lambda$

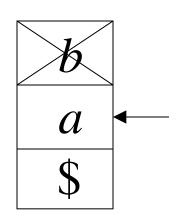
$$b, \lambda \rightarrow b$$

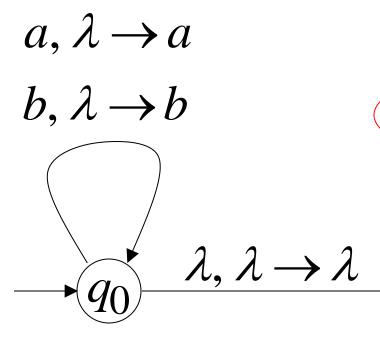
 $b, b \rightarrow \lambda$

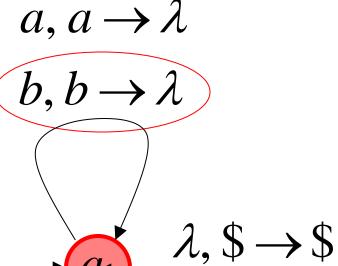


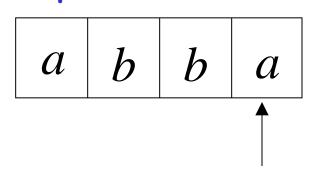
Input

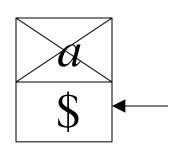








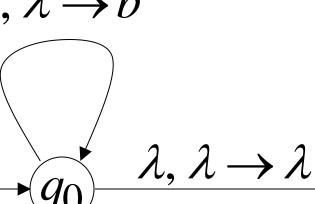


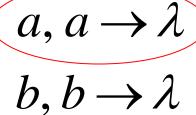


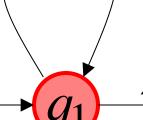


$$a, \lambda \rightarrow a$$

$$b, \lambda \rightarrow b$$



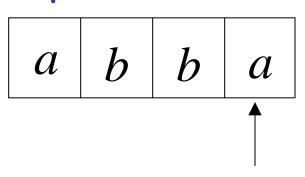


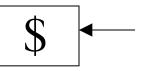


$$\lambda$$
, \$ \rightarrow \$



Input



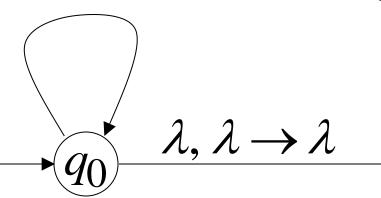


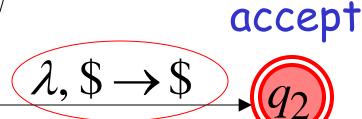
$$a, \lambda \rightarrow a$$

$$b, \lambda \rightarrow b$$

$$a, a \rightarrow \lambda$$

$$b, b \rightarrow \lambda$$

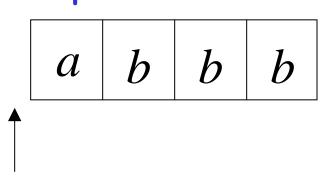


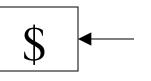


Rejection Example:

Time 0

Input



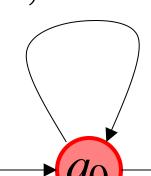


$$a, \lambda \rightarrow a$$

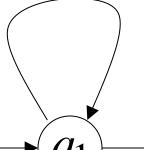
$$b, \lambda \rightarrow b$$

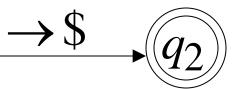
$$a, a \rightarrow \lambda$$

$$b, b \rightarrow \lambda$$

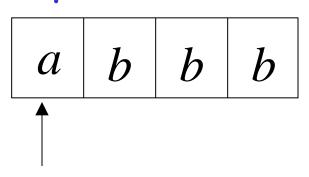


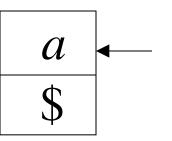
$$\lambda, \lambda \rightarrow \lambda$$

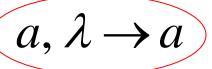




Input



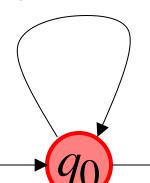




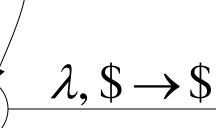
$$a, a \rightarrow \lambda$$

$$b, \lambda \rightarrow b$$

$$b, b \rightarrow \lambda$$

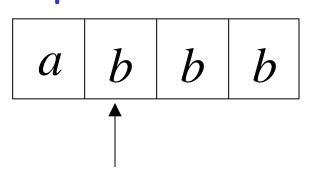


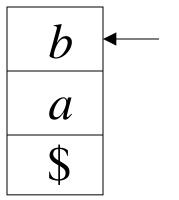
$$\lambda, \lambda \rightarrow \lambda$$



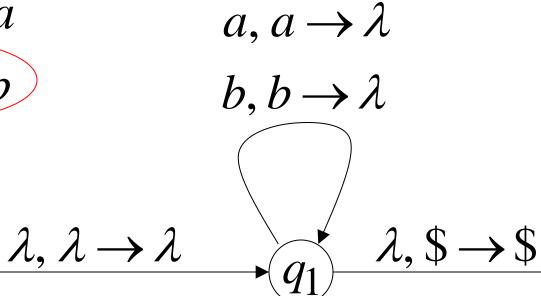


Input

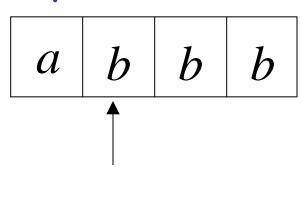




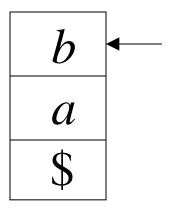
$$\begin{array}{c}
a, \lambda \to a \\
b, \lambda \to b
\end{array}$$



Input



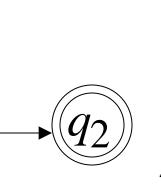
Guess the middle of string



 $a, \lambda \rightarrow a$ Stack $a, a \rightarrow \lambda$

$$b, \lambda \rightarrow b$$

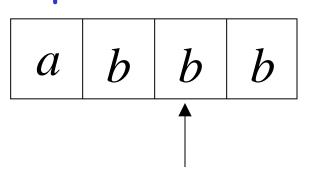
 $b, b \rightarrow \lambda$

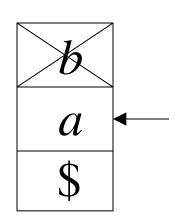


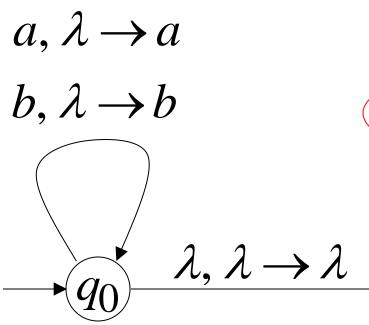
 $\lambda, \lambda \rightarrow \lambda$

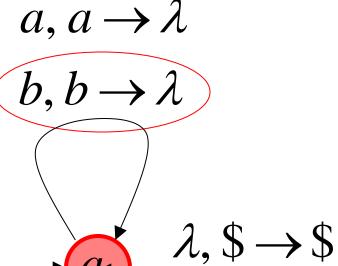
 $\lambda, \$ \rightarrow \$$

Input



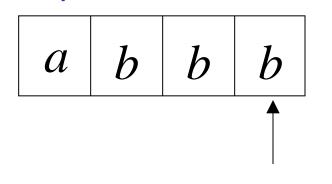






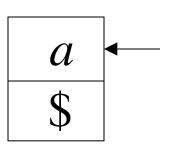
Input

There is no possible transition.



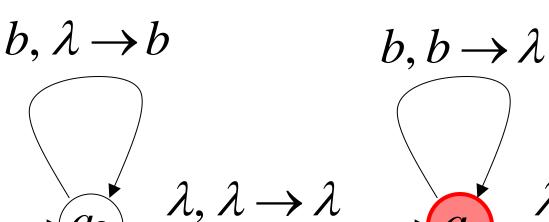
Input is not consumed

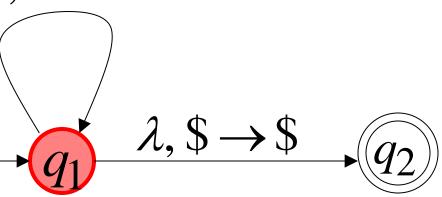
 $a, a \rightarrow \lambda$



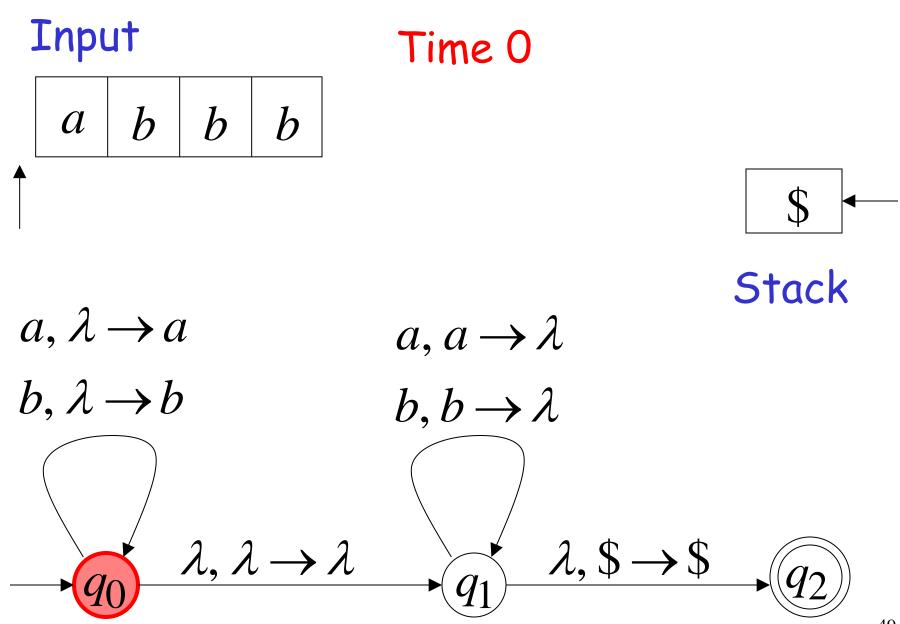
$$a, \lambda \rightarrow a$$

$$b, \lambda \rightarrow b$$

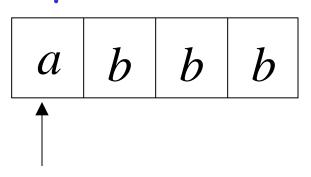


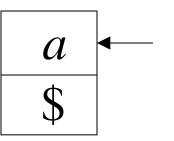


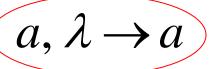
Another computation on same string:



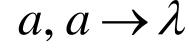
Input



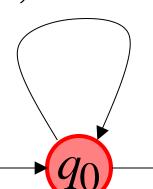




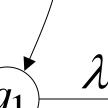
$$b, \lambda \rightarrow b$$



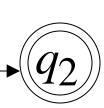
$$b, b \rightarrow \lambda$$



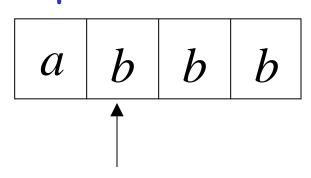
$$\lambda, \lambda \rightarrow \lambda$$

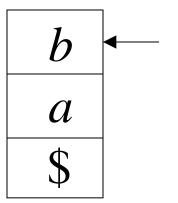


$$\lambda$$
, \$ \rightarrow \$



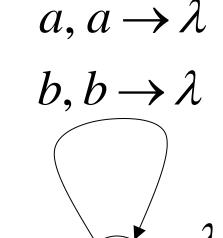
Input

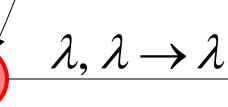




$$a, \lambda \rightarrow a$$

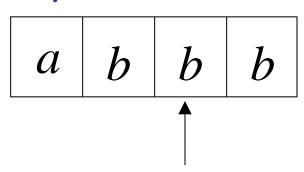
$$b, \lambda \rightarrow b$$

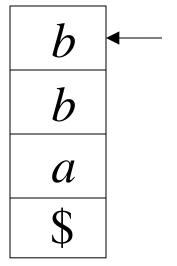






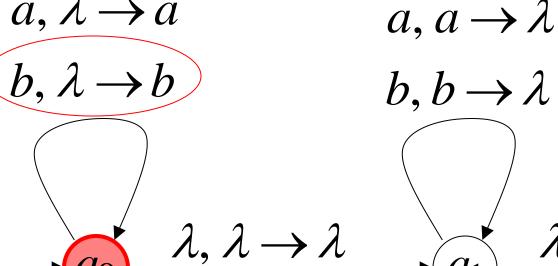
Input



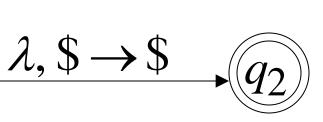


$$a, \lambda \rightarrow a$$

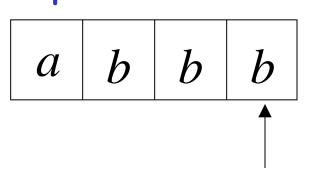
$$b, \lambda \rightarrow b$$

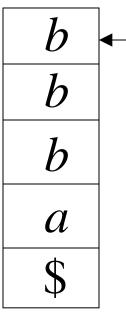


$$b, b \rightarrow \lambda$$



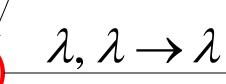
Input





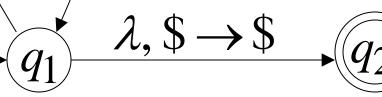
$$a, \lambda \rightarrow a$$

$$b, \lambda \rightarrow b$$

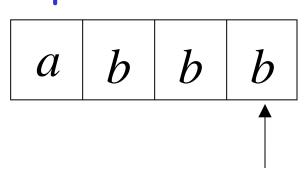


$$a, a \rightarrow \lambda$$

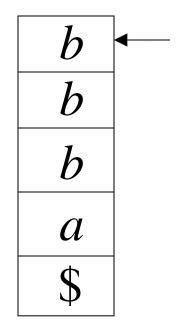
$$b, b \rightarrow \lambda$$



Input

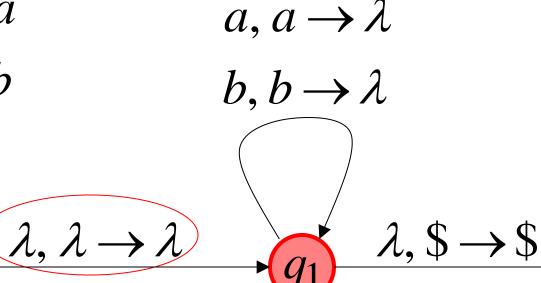


No final state is reached



$a, \lambda \rightarrow a$

$$b, \lambda \rightarrow b$$

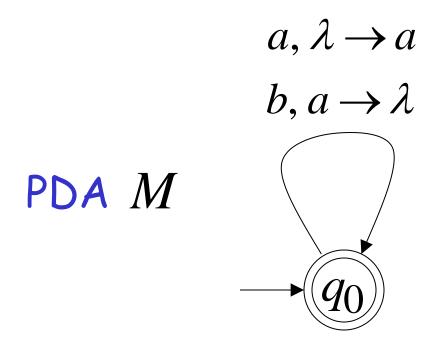


There is no computation that accepts string *abbb*

 $abbb \notin L(M)$

$$a, \lambda \rightarrow a$$
 $a, a \rightarrow \lambda$
 $b, \lambda \rightarrow b$ $b, b \rightarrow \lambda$
 $\downarrow q_0$ $\lambda, \lambda \rightarrow \lambda$ $\downarrow q_1$ $\lambda, \$ \rightarrow \$$ $\downarrow q_2$

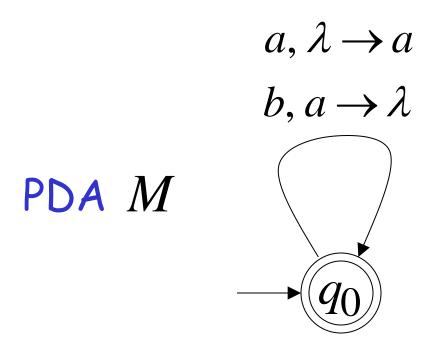
Another PDA example: Language?



Another PDA example

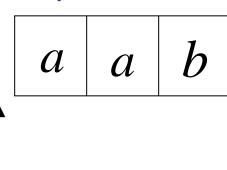
$$L(M) = \{w \in \{a,b\}^*:$$

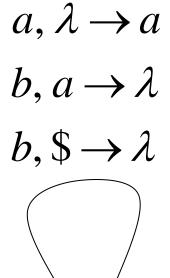
in every prefix v , $n_a(v) \ge n_b(v)$

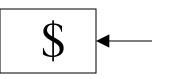


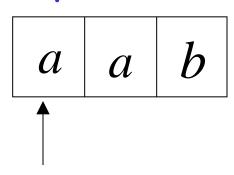
Execution Example: Time 0

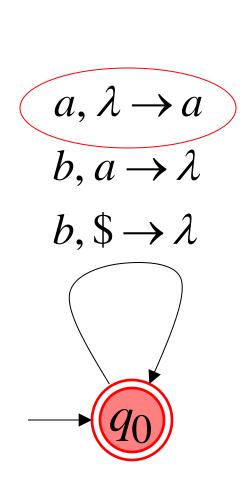
Input

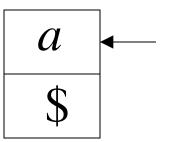






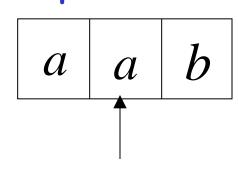


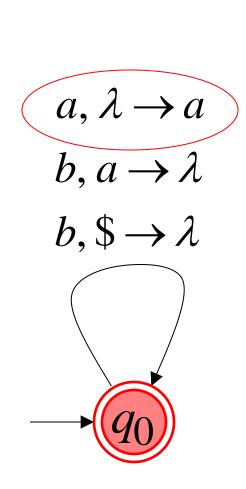


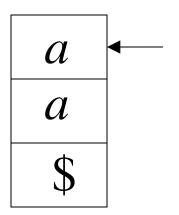


Stack

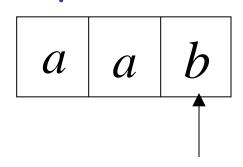
Input

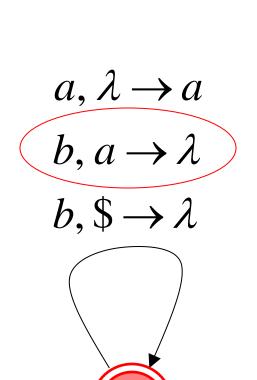


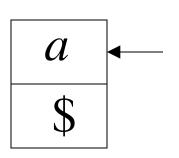




Input





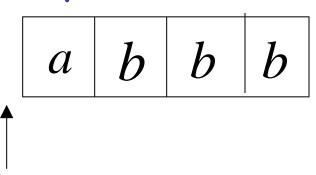


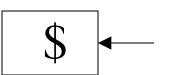
Stack

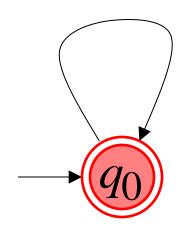
accept

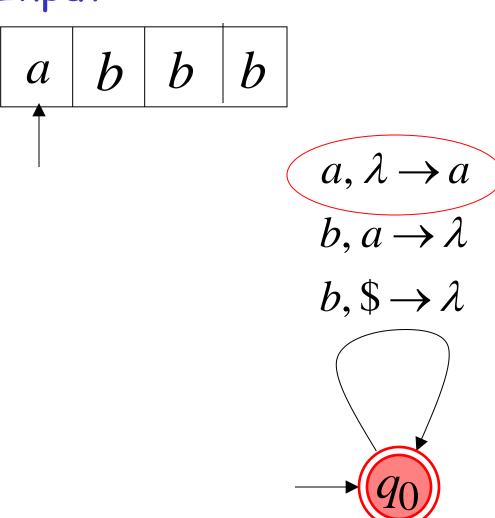
Rejection example: Time 0

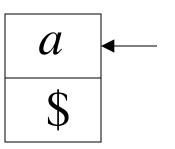
Input





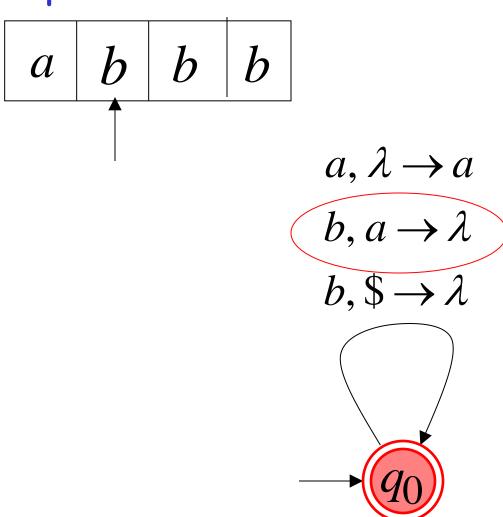


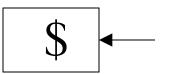


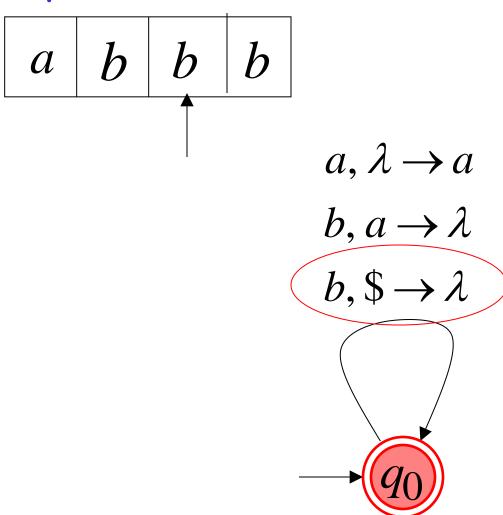


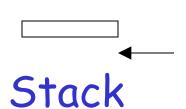
Stack

Input

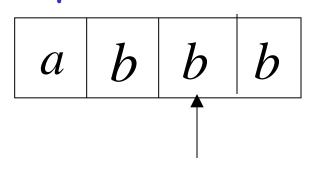


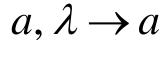


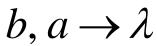




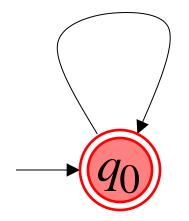
Input

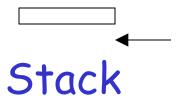






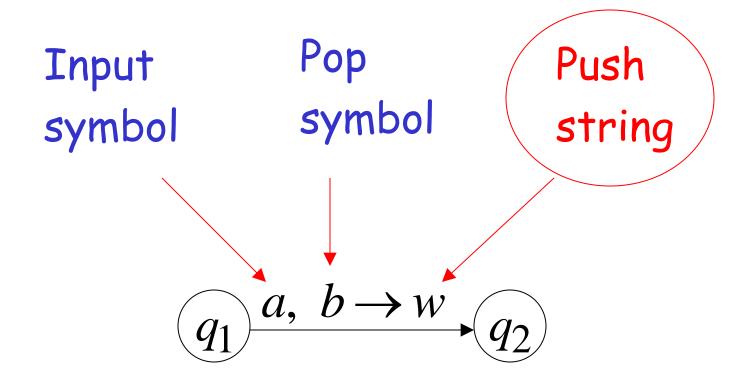
$$b, \$ \rightarrow \lambda$$



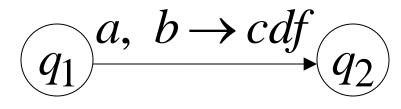


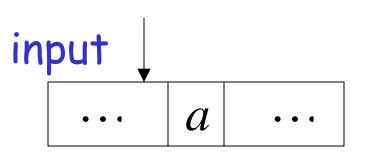
Halt and Reject

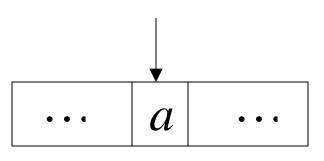
Pushing Strings

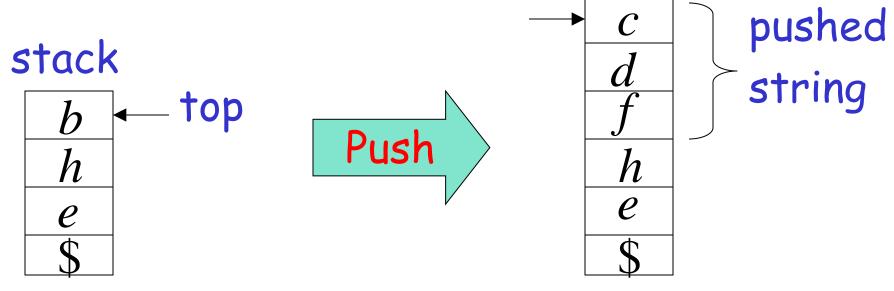


Example:









Another PDA example: Language?

PDA M

Another PDA example

$$L(M) = \{w \in \{a,b\}^*: n_a(w) = n_b(w)\}$$

PDA M

$$a, \$ \rightarrow 0\$$$
 $b, \$ \rightarrow 1\$$
 $a, 0 \rightarrow 00$ $b, 1 \rightarrow 11$
 $a, 1 \rightarrow \lambda$ $b, 0 \rightarrow \lambda$

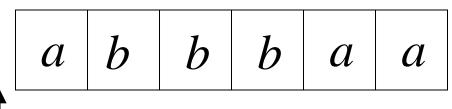
$$\lambda, \$ \rightarrow \$$$

$$q_1$$

$$\lambda, \$ \rightarrow \$$$

Execution Example: Time 0

Input



$$a, \$ \rightarrow 0\$$$
 $b, \$ \rightarrow 1\$$
 $a, 0 \rightarrow 00$ $b, 1 \rightarrow 11$
 $a, 1 \rightarrow \lambda$ $b, 0 \rightarrow \lambda$

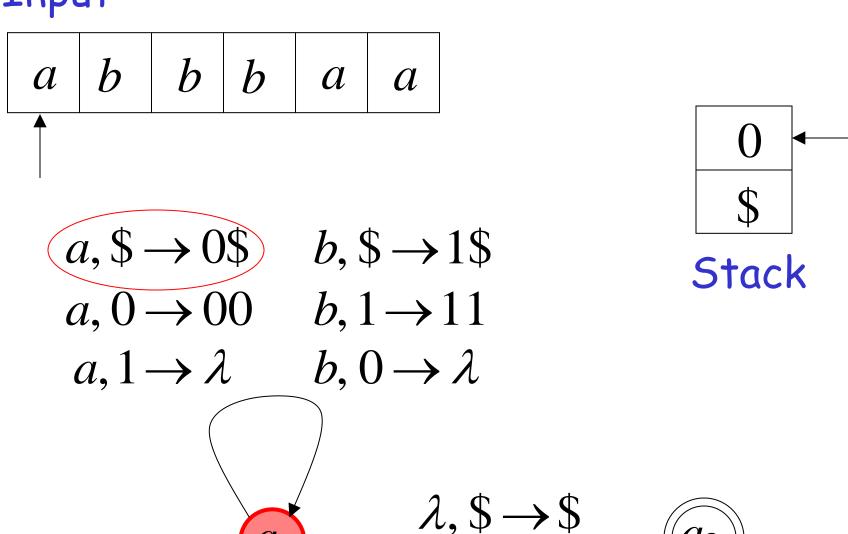
Stack

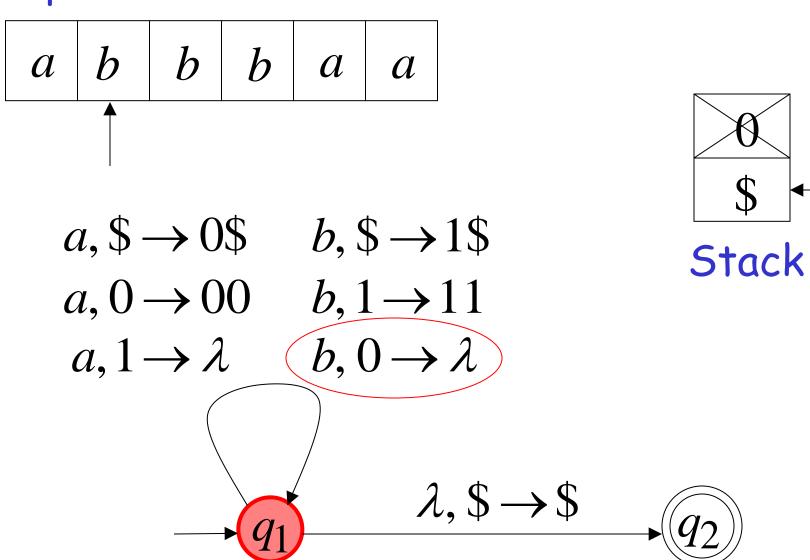
Stack

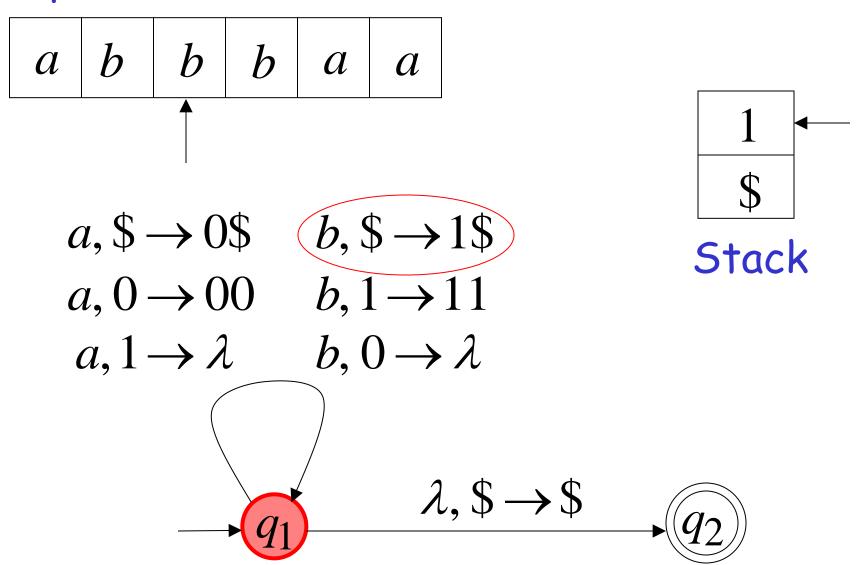
current

state

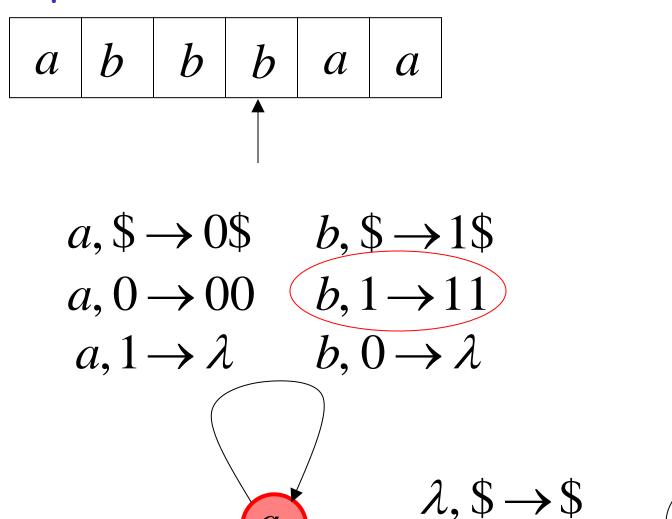
$$\lambda, \$ \rightarrow \$$$

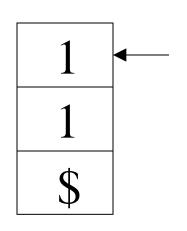




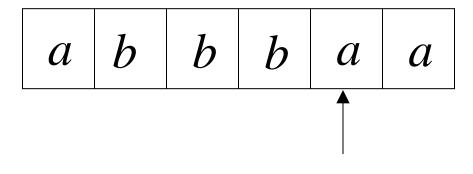


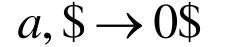
Input











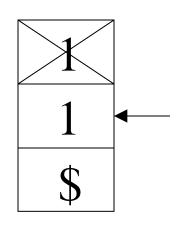
$$b, \$ \rightarrow 1\$$$

$$a, 0 \rightarrow 00$$
 $b, 1 \rightarrow 11$

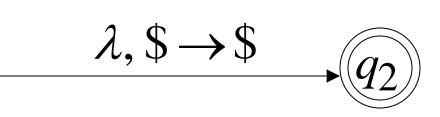
$$b, 1 \rightarrow 11$$

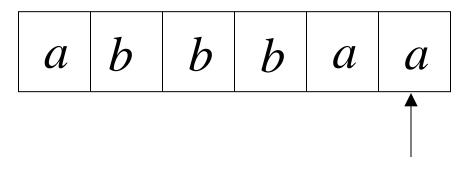
$$(a, 1 \rightarrow \lambda)$$

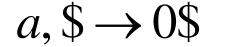
$$b, 0 \rightarrow \lambda$$



Stack







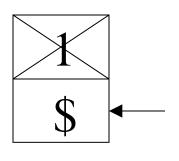
$$b, \$ \rightarrow 1\$$$

$$a, 0 \rightarrow 00$$
 $b, 1 \rightarrow 11$

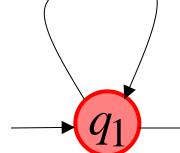
$$b, 1 \rightarrow 11$$

$$(a, 1 \rightarrow \lambda)$$

$$b, 0 \rightarrow \lambda$$

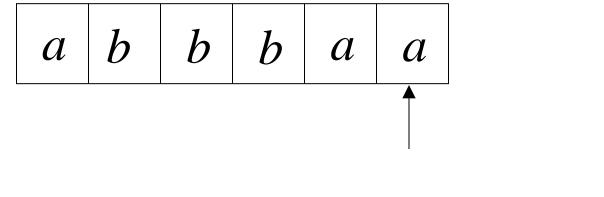


Stack



$$\lambda, \$ \rightarrow \$$$





$$a, \$ \to 0\$$$
 $b, \$ \to 1\$$

$$a, 0 \rightarrow 00$$
 $b, 1 \rightarrow 11$

$$a, 1 \rightarrow \lambda$$
 $b, 0 \rightarrow \lambda$





