GaussDB-Vector:

A Large-Scale Persistent Real-Time Vector Database for LLM Applications

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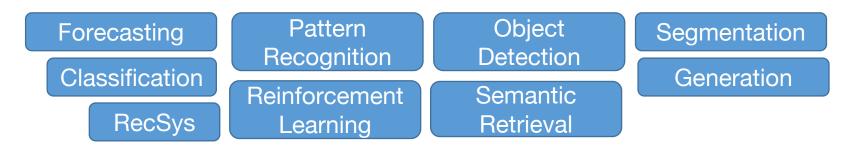
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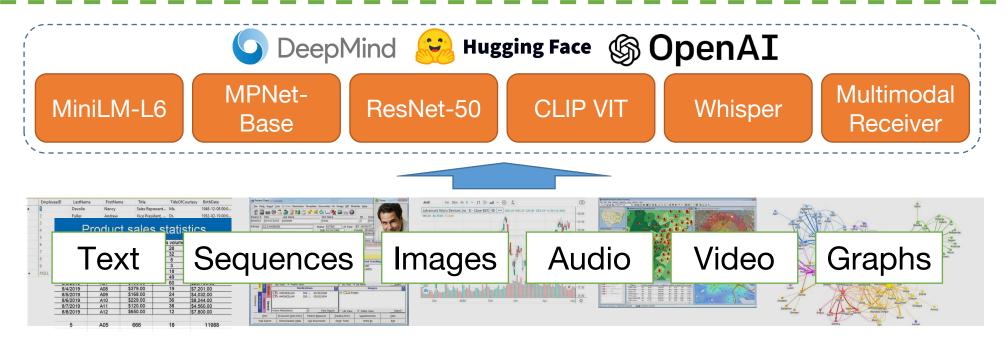




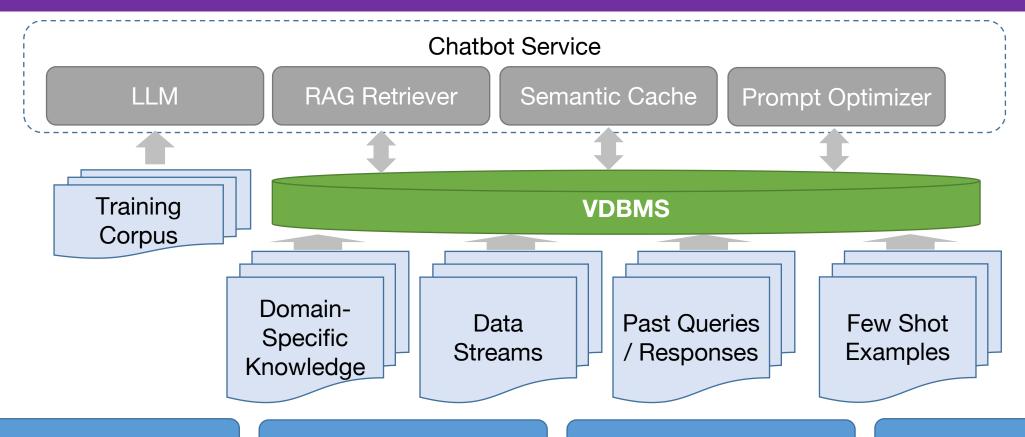
Embeddings: Building Blocks of Future

More and more applications rely on embedding vectors that are retrieved via similarity search





Vector Databases and LLMs



Search Capabilities

 Support various queries, e.g. k-NN, Range NN, Predicated (Filtered) Search

Performance

- Low latency queries
- High throughput
- Memory efficient

Quality

- High recall
- High semantic relevance

System-Level Features

- Scalable
- Available
- Elastic
- Consistent
- Persistent

Vector Databases: Key Challenges

More and more applications rely on embedding vectors that are retrieved via similarity search

Goal: Store embeddings in a DB and retrieve for whatever downstream task



- 1) Size: Existing hardware not for large data units
- 2) Structure: Lack of structure (e.g. attrs) makes embeddings hard to index

Performance

Latency

- Fundamentally high I/O due to more frequent paging
- Naive disk-resident HNSW is not I/O optimal
- High latency of filtered search

Quality

Recall

 Accuracy degradation of IVF due to data drift following updates

System-Level Features

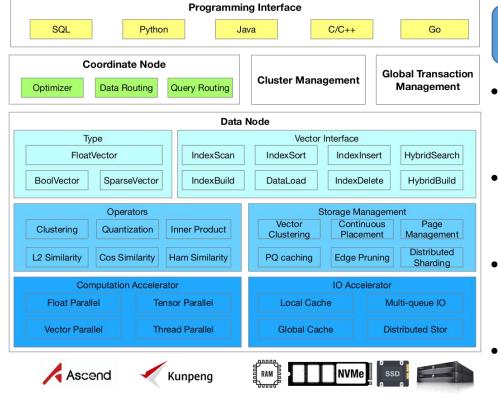
Scalability

 Unclear sharding and query routing strategy for distributed deployments

GaussDB-Vector

Build a distributed persistent real-time database based on IVF & DiskANN

- Architecture: Two-Tiered Distributed Database to provide scalability / availability
- Techniques: IVF & DiskANN but make it easier to build, update, and search



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Index Storage & Construction

- Two-Phase Neighborhood Pruning
- Tailored Storage
 Structures to
 reduce I/O
- Data Sharding with Drift Detection & Auto Rebalancing
- Balanced Tree Hybrid Index

Search

- Hot Nodes Buffering
- Cardinality-Based Query Routing for Distributed Search
- I/O Efficient Search Algorithms

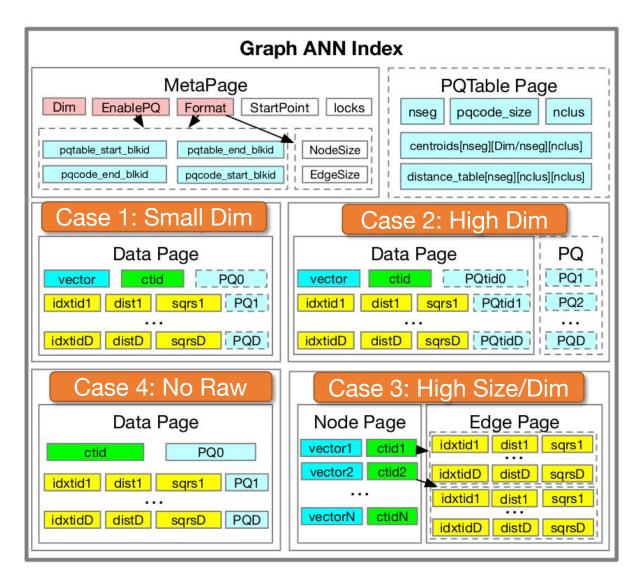
Update

- Asynchronous

 Delete
- Incremental
 Index Updates
 for data
 freshness

Tailored Storage Structures

Use Tailored Storage Structures based on dataset size/dimensions to reduce I/O



Case 1: Small Dimensionality

- Raw Vectors + PQ Codes stored together
- Each node is stored with all its neighbors on a single page to reduce I/O during neighborhood scan

Case 2: High Dimensionality

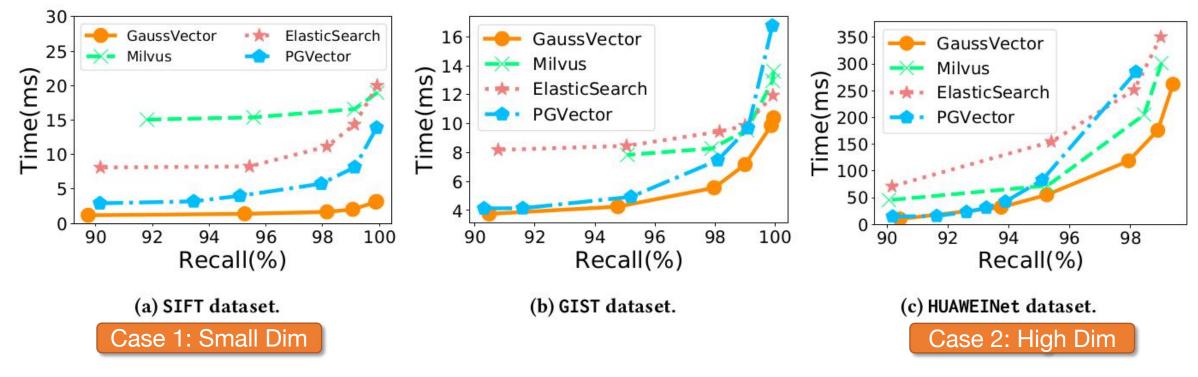
- Raw Vectors are too large and cannot be stored together with PQ Codes
- PQ Codes stored in separately pages and cached in memory

Case 3: High Size/Dimensionality

- Raw Vectors are too large and cannot be stored with PQ Codes nor Neighbors
- Store Nodes and Neighborhoods separately

Tailored Storage Structures

Use Tailored Storage Structures based on dataset size/dimensions to reduce I/O



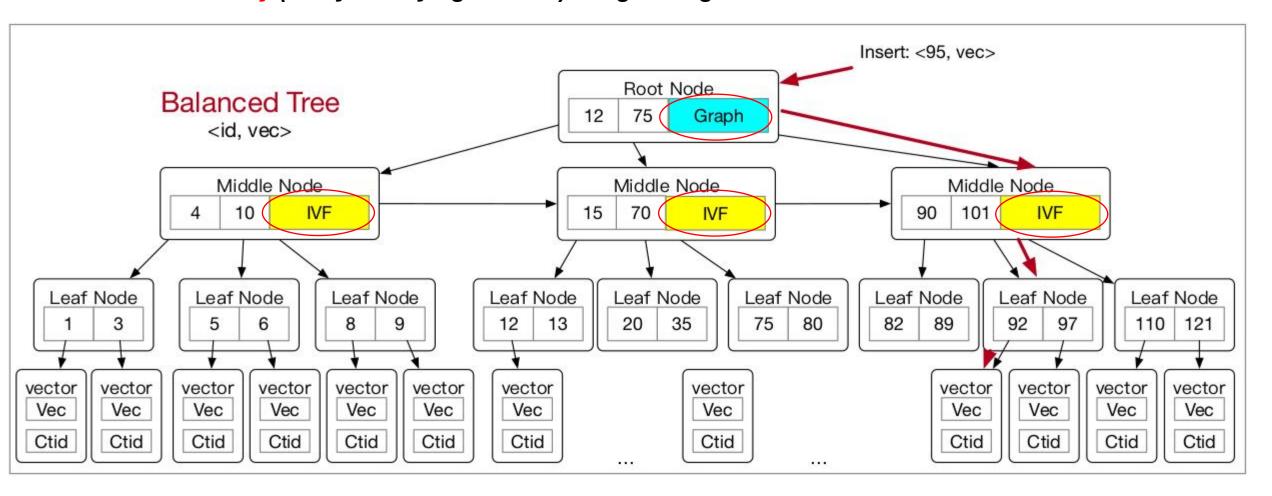
ElasticSearch & Milvus require loading multiple disk segments into memory

GaussDB-Vector adapts to high-dim scenarios to maintain high search performance

Balanced Tree Hybrid Index

Use Balanced Tree with intermediate sub-indexes to support fast pre- / single-stage filtering

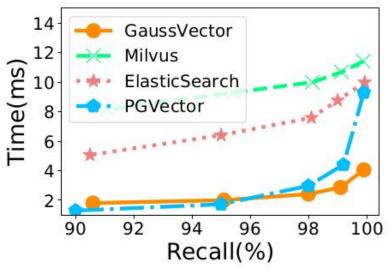
- High Selectivity (few satisfying records): Sequential Scan over leaf nodes
- Low Selectivity (many satisfying records): Single-Stage Scan over sub-indexes

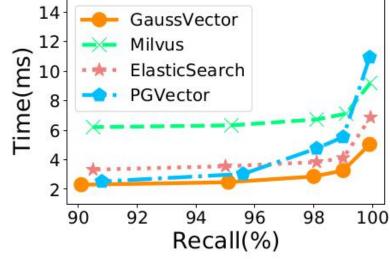


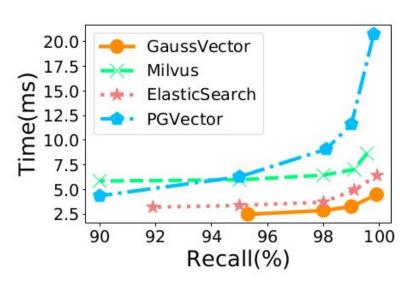
Balanced Tree Hybrid Index

Use Balanced Tree with intermediate sub-indexes to support fast pre- / single-stage filtering

- High Selectivity (few satisfying records): Sequential Scan over leaf nodes
- Low Selectivity (many satisfying records): Single-Stage Scan over sub-indexes







(a) SIFT dataset, 90% scalar selectivity.

(b) SIFT dataset, 50% scalar selectivity.

(c) SIFT dataset, 10% scalar selectivity.

Many Satisfying Records

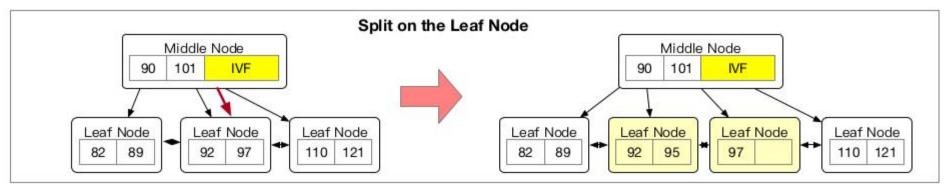
ElasticSearch & Milvus require merging across many segments

Few Satisfying Records

Pre-Filtering via index followed by Sequential Scan beats singlestage scan over segments

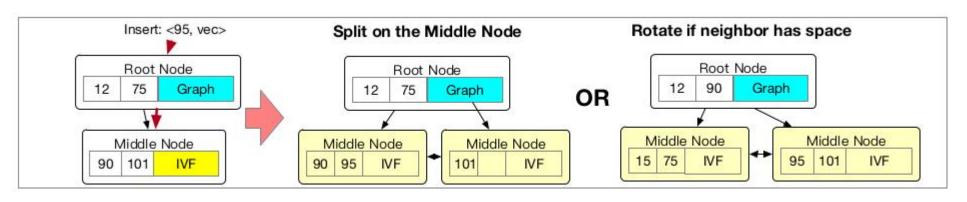
Incremental Index Updates

Balanced Tree with intermediate sub-indexes enables incremental updates for data freshness



Leaf node split requires no index update

(a) Splitting a leaf node following an insert

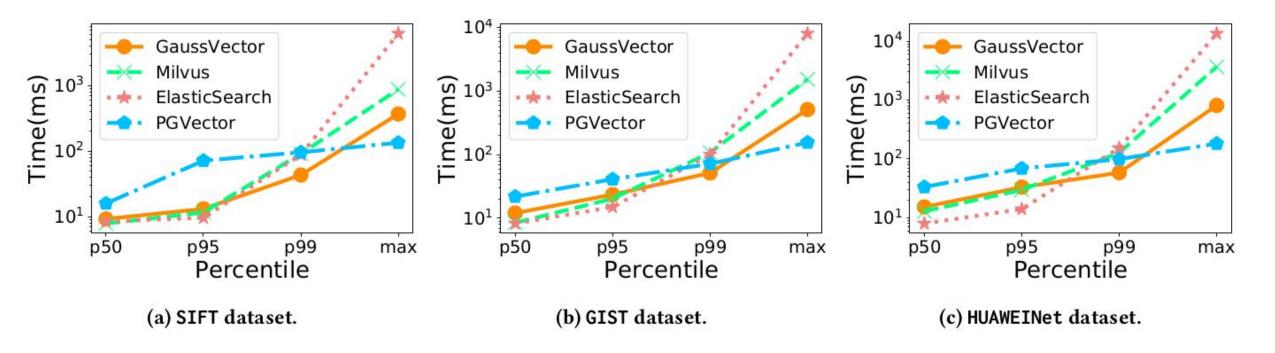


Inner node with IVF index is easy to update

(b) Splitting an inner node following an insert

Incremental Index Updates

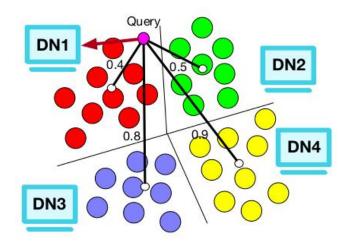
Balanced Tree with intermediate sub-indexes enables incremental updates for data freshness



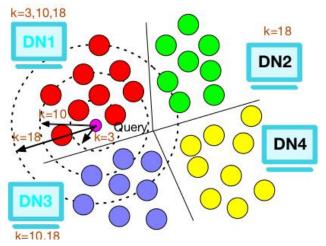
- Large max latency due to inner node / root node splits but otherwise competitive with baselines
- Insert / search latency balanced by adjusting indexing granularity

Data Sharding & Distributed Search

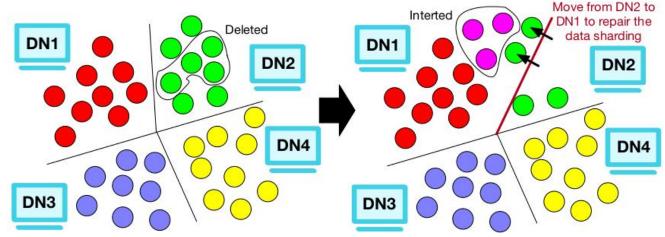
Balanced Tree with intermediate sub-indexes enables incremental updates for data freshness



(a) Insert into Nearest Shard



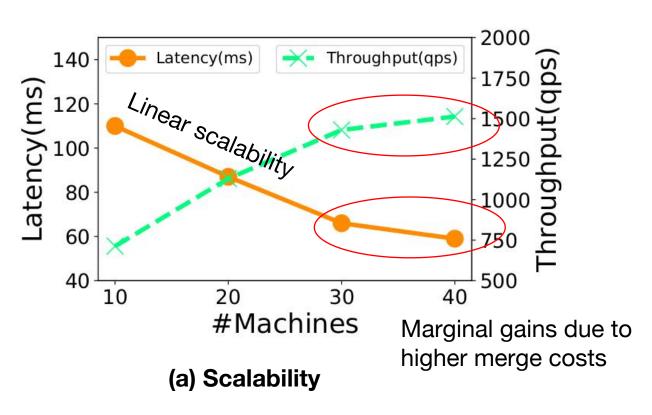
(b) Search Shards based on Estimated Cardinality



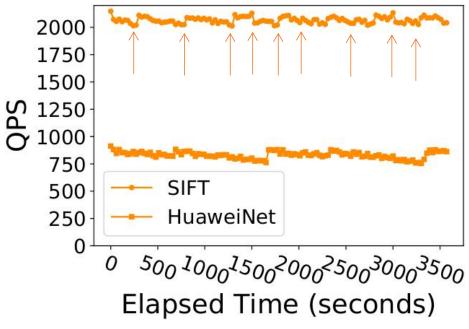
(c) Background Monitoring & Remigration

Data Sharding & Distributed Search

Balanced Tree with intermediate sub-indexes enables incremental updates for data freshness



Rebalancing Spikes



(b) Update Throughput

Summary & Conclusion

GaussDB-Vector integrates vector search capabilities within existing data management system

	Query Processor	Storage & Indexing	Query Optimizer	Query Executor
Milvus milvus	 k-NN Range NN Predicated k-NN Multi-vector search Grouping Search Full Text Search Reranking 	 Replicas Data Partitions Multiple index types Blob storage (backup & persistence) 	Cost-based predicated query planning	 Scatter-gather R/W Disagg. CPU/Store Disagg. Log-backed durable writes Cache & SIMD acceleration GPU acceleration Tunable consistency
PostgreSQL (pgvector)	 k-NN Range NN Predicated k-NN Full Text Search Relational queries 	Page-Based StorageShards & ReplicasHNSW, IVFQuantization	Cost-based predicated query planning	 MVCC Scatter-gather Log-backed durable writes etc.
GaussDB- Vector	 k-NN Range NN Predicated k-NN Relational queries 	 Page-Based Storage Shards & Replicas HNSW, IVF, Hybrid Index Quantization 	Cost-based predicated query planning	 GaussDB Transaction Manager Scatter-gather Log-backed durable writes Cache & SIMD acceleration NPU/GPU acceleration