



# Markov chain models of coupled calcium channels: K

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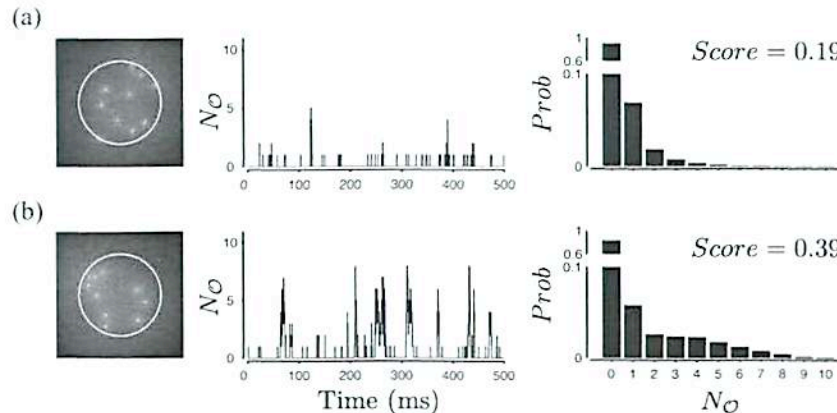
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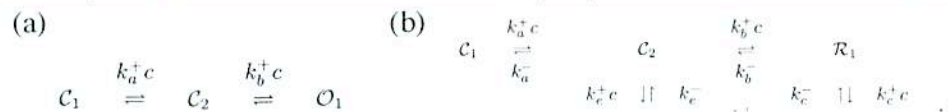
## — Abstract —

Mathematical models of calcium release sites derived from Markov chain models of intracellular calcium channels exhibit collective gating reminiscent of the experimentally observed phenomenon of stochastic calcium excitability (i.e., calcium puffs and sparks). We present a Kronecker structured representation using stochastic automata networks for calcium release site models and perform benchmark stationary distribution calculations using both exact and approximate iterative numerical solution techniques that leverage this structure. We find multi-level methods provide excellent convergence with modest additional memory requirements. When an exact solution is not feasible, iterative approximate methods based on the power method may be used, with performance similar to Monte Carlo estimates.

## — Calcium Release Sites: Puffs and Sparks —



**Figure 1:** (a) Left:  $[Ca^{2+}]$  near  $3 \times 3 \mu m$  endoplasmic reticulum membrane with 12  $Ca^{2+}$ -regulated  $Ca^{2+}$  channels modeled as 3-state Markov chains (see Fig. 2(a)) with positions randomly chosen from a uniform distribution on a disc of radius  $2 \mu m$  (source amplitude 0.05 pA). Buffered  $Ca^{2+}$  diffusion is modeled as in [2]. Middle: Stochastic dynamics of the number of open channels at the release site ( $N_O$ ) that does not include robust puffs/sparks. Right: Probability distribution of the number of open channels leading to a low puff/spark Score (see Eq. 2) of 0.19. (b) Different random channel positions result in a release site that exhibits robust  $Ca^{2+}$  puff/sparks (middle) and an elevated Score of 0.39 (right).

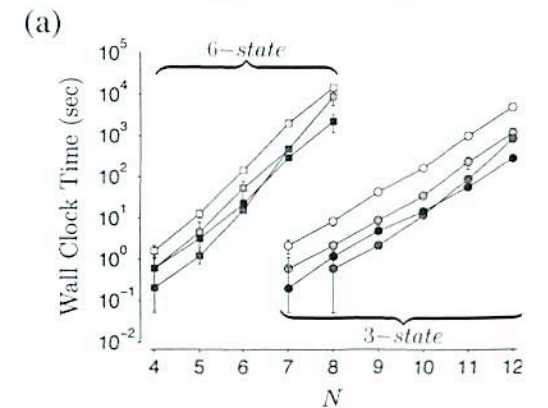


## — Exact Method

Solver	Max Res
JOR	9.49E-13
JOR_AD	9.44E-13
ARNOLDI	2.42E-13
BICGSTAB	8.66E-13
PRE_ARNOLDI	8.62E-15
BSOR_BICGSTAB	8.22E-15
ML_JOR_F_DYN	5.87E-13

**Table 1:** Benchmark calculations for 10 3-state channels computed using 8GB RAM solving Eq. 1. Description of solvers [1]: JOR, Jacobi iteration/disaggregation; ARNOLDI, the method of Arnoldi; BICGSTAB, Arnoldi with Neumann pre-conditioning; BSOR\_BICGSTAB, the bic preconditioning; ML\_JOR\_F\_DYN, multi-level method with JOR smoother

## — Exact Methods: 5



**Figure 3:** (a): Circles and error bars show the mean  $\pm$  SD of wall clock time for JOR (black), PRE\_ARNOLDI (blue), BSOR\_BICGSTAB (red), and ML\_JOR\_F\_DYN (green).

# Cronecker representations and iterative solution methods

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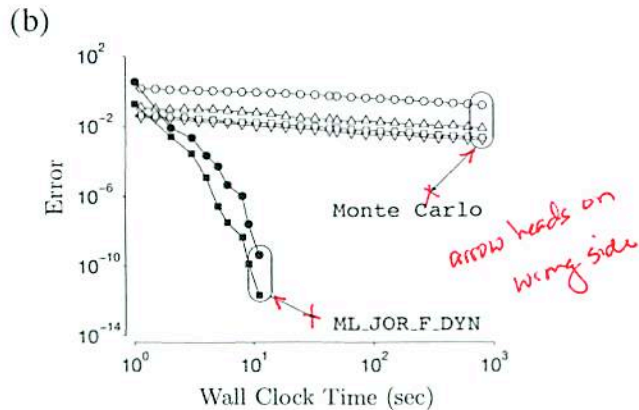


## as: Benchmarking —

Sum Res	CPU (s)	Wall (s)	Iters
5.16E-12	279	279	1840
5.13E-12	415	415	1550
4.04E-11	214	215	1440
4.89E-11	146	148	602
1.82E-12	26	27	160
5.29E-13	19	19	52
1.68E-10	15	15	46

tested using Linux PCs with dual core 3.8GHz EM64T Xeon processors and the following methods: JOR\_AD, the method of Jacobi with aggregation; the biconjugate gradient stabilized method; PRE\_ARNOLDI, the method of conjugate gradient stability method with block successive over-relaxation preconditioning, F cycle, and dynamic ordering.

## Scalability and Error —



Wall clock time for five release site configurations of the 3-state model (Fig. 2(a)) (black circles) and ML\_JOR\_F\_DYN (black circles). Sources and error bars give results for

## — Approximate Methods: Partitioning —

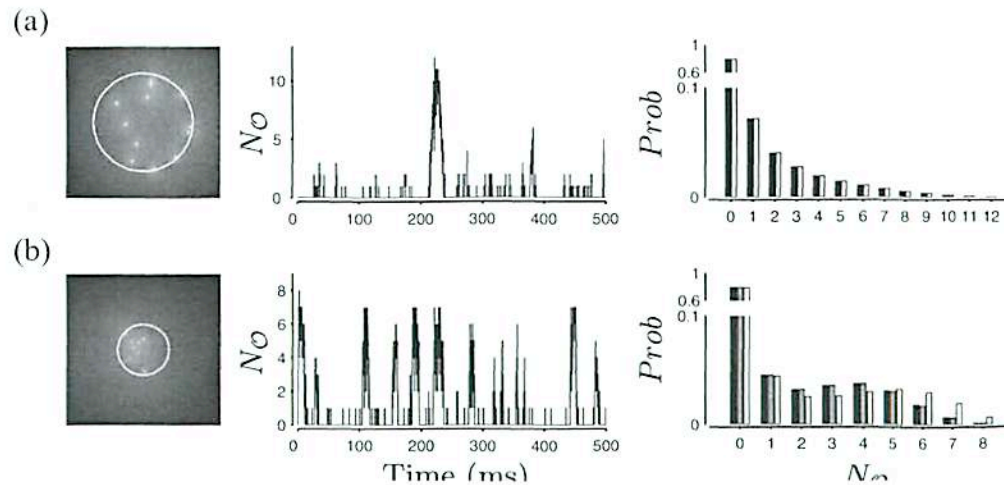
$ P $	Optimal $\mathcal{P}$ for 3-state model	$\epsilon_z$	$\epsilon_{score}$	$\nu(3)$	$\nu(12)$
1	$\{C_1 C_2 O_1\}$	1.66E+0	9.38E-1	3.33E-1	6.77E-5
2	$\{C_1 C_2\} \{O_1\}$	5.40E-3	3.50E-3	1.33E+0	1.39E-1
3	$\{C_1\} \{C_2\} \{O_1\}$	6.00E-8	4.38E-9	3.00E+0	1.20E+1

$ P $	Optimal $\mathcal{P}$ for 6-state model	$\epsilon_z$	$\epsilon_{score}$	$\nu(3)$	$\nu(8)$
1	$\{C_1 C_2 R_1 R_2 R_3 O_1\}$	4.84E-1	5.40E-1	8.30E-2	2.86E-5
2	$\{C_1 C_2 R_1 R_2 R_3\} \{O_1\}$	3.50E-2	1.36E-2	3.33E-1	3.70E-3
3	$\{C_1 C_2\} \{R_1 R_2 R_3\} \{O_1\}$	7.10E-3	5.50E-3	7.50E-1	6.25E-2
4	$\{C_1\} \{C_2\} \{R_1 R_2 R_3\} \{O_1\}$	1.90E-4	7.60E-6	1.33E+0	4.68E-1
5	$\{C_1\} \{C_2\} \{R_1 R_2\} \{R_3\} \{O_1\}$	5.07E-7	4.12E-8	2.08E+0	2.23E+0
6	$\{C_1\} \{C_2\} \{R_1\} \{R_2\} \{R_3\} \{O_1\}$	4.21E-8	4.16E-8	3.00E+0	8.00E+0

Table 2: Optimal partitioning strategies for APP\_POWER listed with their relative errors and memory requirements for both  $N = 3$ - and 6-state channels with states  $\mathcal{S} = \{C_1, C_2, O_1\}$  and  $\mathcal{S} = \{C_1, C_2, R_1, R_2, R_3, O_1\}$ , respectively (see Fig. 2). The last column also lists relative memory requirements  $\nu(N)$  for larger  $N$ , where  $\nu(N) = Z_A/Z_E$ , with  $Z_A = NM/P^{N-1}$  and  $Z_E = M^N$  the memory requirements for the approximate and exact methods, respectively. For the 3-state model, the optimal  $\mathcal{P}$  is the best of 1, 3, 1 possibilities when  $|P| = 1, 2, 3$ . For the 6-state model, the optimal  $\mathcal{P}$  is the best of 1, 31, 90, 65, 15, 1 possibilities when  $|P| = 1, 2, 3, 4, 5, 6$ .

## — Approximate Methods: Scalability and Error —







**Figure 2: (a)** 3-state single channel model with  $\text{Ca}^{2+}$ -mediated activation that has two closed ( $\mathcal{C}_1, \mathcal{C}_2$ ) and one open ( $\mathcal{O}_1$ ) state. Parameters in  $\mu\text{M}^{-1} \text{ms}^{-1}$ :  $k_a^+ = 1.5$ ,  $k_b^+ = 150$ ; in  $\text{ms}^{-1}$ :  $k_a^- = 50$ ,  $k_b^- = 1.5$ . **(b)** 6-state single channel model with  $\text{Ca}^{2+}$ -mediated activation and inactivation. Parameters in  $\mu\text{M}^{-1} \text{ms}^{-1}$ :  $k_a^+ = 1.5$ ,  $k_b^+ = k_d^+ = 0.015$ ,  $k_c^+ = k_e^+ = 300$ ,  $k_f^+ = 3.0$ ; in  $\text{ms}^{-1}$ :  $k_a^- = 49.5$ ,  $k_b^- = k_d^- = 0.2475$ ,  $k_c^- = k_e^- = 6.0$ ,  $k_f^- = 0.03$ .

## — Model Setup: Exact Methods —

The generator matrix for a single 3- or 6-state channel model takes the form

$$Q = K_- + (c_\infty I + c_d I_{\mathcal{O}}) K_+$$

where  $K_-$  and  $K_+$  are  $M \times M$  matrices that collect the unimolecular ( $k_i^-$ ) and bimolecular ( $k_i^+$ ) transition rates,  $M$  is the number of states in the single channel model,  $I$  is the  $M \times M$  identity matrix,  $I_{\mathcal{O}} = \text{diag}\{e_{\mathcal{O}}\}$ , and  $e_{\mathcal{O}}$  is an  $M \times 1$  vector indicating open states of the single channel model [2]. For example, for the 3-state model of Fig. 2(a) we have

$$K_- = \begin{pmatrix} 0 & 0 & 0 \\ k_a^- & -k_a^- & 0 \\ 0 & k_b^- & -k_b^- \end{pmatrix}, \quad K_+ = \begin{pmatrix} -k_a^+ & k_a^+ & 0 \\ 0 & -k_b^+ & k_b^+ \\ 0 & 0 & 0 \end{pmatrix}$$

In the case of  $N$  channels coupled at the  $\text{Ca}^{2+}$  release site, the expanded generator matrix—i.e., the SAN descriptor—is given by

$$Q^{(N)} = \bigoplus_{n=1}^N X_\infty + \sum_{i,j=1}^N \bigotimes_{n=1}^N X_{ij}^n \quad (1)$$

where  $X_\infty = K_- + c_\infty K_+$  and

$$X_{ij}^n = \begin{cases} I_{\mathcal{O}} & \text{for } i \neq j, i = n \\ c_{ij} K_+ & \text{for } i \neq j, j = n \\ c_d I_{\mathcal{O}} K_+ & \text{for } i = j = n \\ I & \text{otherwise.} \end{cases}$$

We are interested in calculating both the stationary distribution given by

$$\pi^{(N)} Q^{(N)} = \mathbf{0} \quad \text{subject to} \quad \pi^{(N)} \mathbf{e}^{(N)} = 1$$

and a coarser measure called the puff/spark *Score* defined as

$$\text{Score} = \frac{\text{Var}[f_{\mathcal{O}}]}{\text{E}[f_{\mathcal{O}}]} = \frac{1}{N} \frac{\text{Var}[N_{\mathcal{O}}]}{\text{E}[N_{\mathcal{O}}]}. \quad (2)$$



The authors thank Buchholz and Dayar for sharing their implementation of Nsolve. This material is

using: *SSR* (white), *FRS* (gray), *BSR* (black), *CS* (the 6-state model (Fig. 2(b))). Single-channel parameters as in Fig. 2 C GB RAM. **(b)**: Convergence of response measures for a release site  $\text{Ca}^{2+}$  and open symbols, respectively). Circles and squares give 1- and  $\infty$ -no puff/spark *Score* for Monte Carlo (mean of 50 simulations shown) or the *lower pointing triangles* give the relative error in the probability th

## — Model Setup: Approximate Methods —

For one 3-state channel, the state-space  $\mathcal{S}$  with  $\{\mathcal{C}_1, \mathcal{C}_2, \mathcal{O}_1\}$  and  $\mathcal{P} = \{\mathcal{C}, \mathcal{O}\}$  with  $\mathcal{C} = \{\mathcal{C}_1, \mathcal{C}_2\}$  the induced partitioning on the expanded state-space  $\mathcal{S} = s^1 s^2 \dots s^N \in \mathcal{S}^{(N)} \mapsto p = p^1 p^2 \dots p^N \in \mathcal{P}^{(N)}$  where the transition matrix  $Q$  is dec

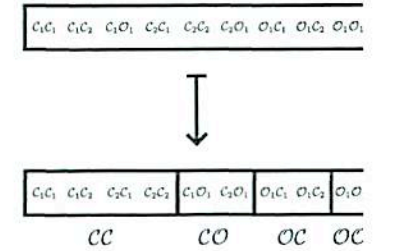
$$Q^{(N)}[p, p] = \bigoplus_{n=1}^N X_\infty[p^n, p^n] + \sum_{i,j=1}^N \zeta$$

$$Q^{(N)}[p, q] = \left( \bigotimes_{n=1}^{k-1} I_{|p^n|} \right) \otimes X_\infty[p^k, q]$$

Compositional representation of  $\pi$  corresponding The approximate numerical method we consider

## — Approximate Method

(a)



**Figure 4: (a)** Permutation of states and partition structure for  $N$  where  $\mathcal{C} = \{\mathcal{C}_1, \mathcal{C}_2\}$  and  $\mathcal{O} = \{\mathcal{O}_1\}$ . The induced partitioning on  $\mathcal{S}^{(2)}$  in both  $\mathcal{S}^{(2)}$  and each partition. **(b)** Block structure of the expanded generator matrix  $Q^{(N)}$  denotes the hierarchical structure of the partitioning.

black gray), and ML\_JOR\_F\_DYN (black). Squares and error bars give results for calculations performed using 2.66 GHz Dual-Core Intel Xeon processors and 2 imposed of 10 3-state channels using ML\_JOR\_F\_DYN and Monte Carlo (filled circles) of the residual errors, upper pointing triangles give the relative error in the compared with the Score given by ML\_JOR\_F\_DYN upon convergence. Similarly, at all  $N$  channels are closed. Parameters as in Fig. 1.

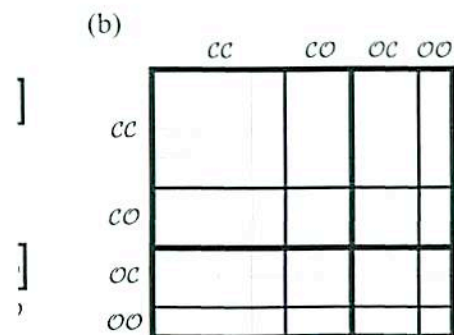
## Approximate Methods —

For a particular partition structure  $\mathcal{P}$  is given by  $\mathcal{S} = \{S_1, S_2\}$ ,  $\mathcal{O} = \{O_1\}$ . In the case of  $N$  3-state channels, space is a mapping from  $3^N$  states to  $2^N$  partitions by  $\mathcal{P}^{(N)}$ . We can use the hierarchical Kronecker representation composed into blocks over partitions  $p \neq q$  such that

$$\bigotimes_{n=1}^N X_{ij}^n[p^n, p^n] + \sum_{i,j=1}^N \bigotimes_{n=1}^N X_{ij}^n[p^n, q^n].$$

to  $p$  is  $\pi[p] = \alpha_p \bigotimes_{n=1}^N \pi_p^n$  with  $\sum_{p \in \mathcal{P}^{(N)}} \alpha_p = 1$ . that leverages this structure is APP\_POWER.

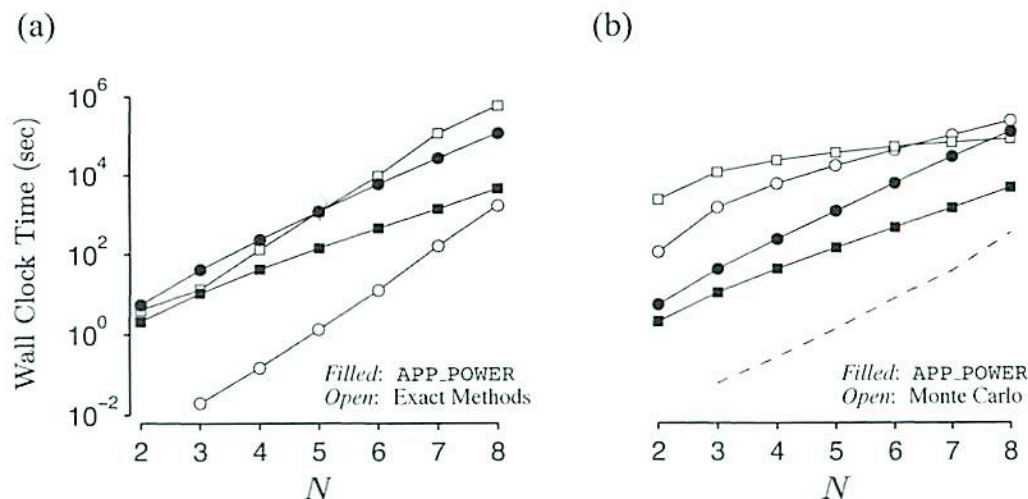
## 5: Hierarchical Structure —



is thus  $\mathcal{P}^{(2)} = \{cc, co, oc, oo\}$ , with the states ordered lexicographically generator matrix  $Q^{(2)}$  when permuted in this manner. The thickness of the lines

based upon work supported by the National Science Foundation under Grants No. 0133132 and 0443843.

**Figure 5:** (a) Statistics for a release site composed of 12 3-state channels. Left: Local  $[Ca^{2+}]$  near  $3 \times 3 \mu m$  ER membrane modeled as in Fig. 1. Middle: Localized  $Ca^{2+}$  elevations reminiscent of  $Ca^{2+}$  puffs/sparks. Right: Probability distribution of the number of open channels calculated exactly using ML\_JOR\_F\_DYN (black bars) and approximately using APP\_POWER with C/O partitioning (white bars). (b) Statistics as in A for 8 6-state channels with black bars denoting ML\_JOR\_F\_DYN and white and grey bars denoting APP\_POWER with C/O and C/R/O partitioning, respectively.



**Figure 6:** (a) Filled and open symbols show the wall clock time for the 6-state model using approximate and exact methods, respectively. Approximate results are shown for two levels of partitioning (C/O, squares and C/R/O, circles) with the APP\_POWER method. Exact solutions are calculated using the POWER method (squares) and ML\_JOR\_F\_DYN method (circles). (b) Results as in A with open symbols corresponding to Monte Carlo estimates of two coarse response measures: the distribution of the number of open channels (circles) and the distribution of probability across the  $M$  states of an arbitrarily selected individual channel (squares). The dashed line shows the projected performance of an approximate multi-level solver that uses ML\_JOR\_F\_DYN rather than POWER as its iterative engine.

## — References —

- [1] W. Stewart, *Introduction to the numerical solution of Markov chains*, Princeton: Princeton University Press (1994).
- [2] V. Nguyen, R. Mathias, and G. D. Smith, *A stochastic automata network descriptor for Markov chain models of instantaneously coupled intracellular  $Ca^{2+}$  channels*, Bulletin of Mathematical Biology, 67 (2005), pp.393–432.
- [3] H. DeRemigio, M. D. LaMar, P. Kemper, and G. D. Smith, *Markov chain models of coupled calcium channels: Kronecker representations and iterative solution methods*, Physical Biology, 5 (2008).



## — Abstract —

Mathematical models of calcium release, since derived from Markov chain models of intracellular calcium channels, exhibit collective gating reminiscent of the experimentally observed phenomenon of stochastic calcium excitability (i.e., calcium puffs and sparks). We present a Kronecker structured representation using stochastic automata networks for calcium release site models and perform benchmark calculations using both exact and approximate iterative numerical solution techniques that leverage this structure. We find multilevel methods provide excellent convergence with modest additional memory requirements. When an exact solution is not feasible, iterative approximate methods based on the power method may be used, with performance similar to Monte Carlo estimates.

## — Calcium Release Sites: Puffs and Sparks —

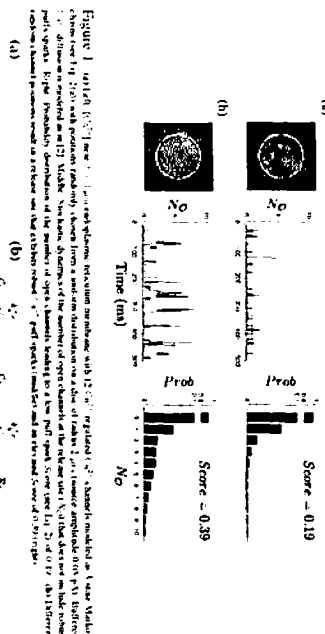


Figure 1: (a) Left: A 3D visualization of a calcium release site model showing a spherical structure with internal compartments. Right: A plot of the probability distribution of the number of open channels (No) versus time (ms). (b) Left: A 3D visualization of a calcium release site model showing a spherical structure with internal compartments. Right: A plot of the probability distribution of the number of open channels (No) versus time (ms).

## — Model Setup: Exact Methods —

$$Q = K_1 + K_2 + \dots + K_N$$

where  $K_1$  and  $K_2$  are  $N \times N$  matrices that collect the unimolecular ( $K_1$ ) and bimolecular ( $K_2$ ) transition rates.  $N$  is the number of states in the single channel model.  $I$  is the  $N \times N$  identity matrix,  $Q = \text{diag}(Q_1)$  and  $Q_2$  is an  $N \times 1$  vector indicating open states of the single channel model [2]. For example, for the 3-state model of Fig. 2(a) we have

$$K_1 = \begin{pmatrix} 0 & 0 & 0 \\ k_{12} & -k_{12} & 0 \\ 0 & k_{21} & -k_{21} \end{pmatrix}, \quad K_2 = \begin{pmatrix} -k_{12} & k_{12} & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

In the case of  $N$  channels coupled at the  $C_{\text{SR}}$  release site, the expanded generator matrix  $\pi \cdot Q$  of the SAN descriptor is given by

$$\pi \cdot Q = \bigoplus_{i=1}^N K_i + \sum_{i,j=1}^N \lambda_{ij} \otimes K_j$$

where  $\lambda_{ij} = K_i - K_j$ , and

$$\lambda_{ij} = \begin{cases} k_{ij} & \text{for } i \neq j, i, j = 1, \dots, N \\ 0 & \text{for } i = j, i, j = 1, \dots, N \end{cases}$$

We are interested in calculating both the stationary distribution given by

$$\pi \cdot Q \cdot \mathbf{1} = 0 \quad \text{subject to} \quad \pi \cdot \mathbf{1} = 1$$

and a common measure called the pull-spark score defined as

$$\text{Score} = \frac{\text{Val}(\pi)}{E(\pi)} = \frac{1}{N} \sum_{i=1}^N \pi_i$$



The authors thank Buchholz and Dayar for sharing their implementation of Nucleo. This material is based upon work supported by the National Science Foundation under Grants (DMS-1133132 and DMS-1448343).

## — Exact Methods: Benchmarking —

Model	Nodes	Non-Zero	CPUs	Wall Time
200	6,400,000	3,100,000	278	18.60
200	6,400,000	3,100,000	415	15.55
200	6,400,000	3,100,000	415	15.55
200	6,400,000	3,100,000	415	15.55
200	6,400,000	3,100,000	415	15.55
200	6,400,000	3,100,000	415	15.55
200	6,400,000	3,100,000	415	15.55
200	6,400,000	3,100,000	415	15.55
200	6,400,000	3,100,000	415	15.55
200	6,400,000	3,100,000	415	15.55

Table 1: Benchmarking results for 10 3-state channels coupled using 1 unit CPU with dual core Xeon E5-2680 v2 (16GB) from previous and with 100 units CPU. The method of choice is 100 units CPU. The benchmark problem is the same as the one used in the previous work [1]. The method of choice is 100 units CPU. The benchmark problem is the same as the one used in the previous work [1]. The method of choice is 100 units CPU. The benchmark problem is the same as the one used in the previous work [1].

## — Exact Methods: Scalability and Error —

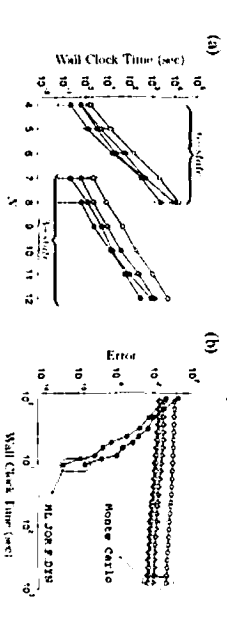


Figure 3: (a) Left: A plot of Wall Clock Time (sec) versus N. Right: A plot of Error versus Wall Clock Time (sec). (b) Left: A plot of Wall Clock Time (sec) versus N. Right: A plot of Error versus Wall Clock Time (sec).

## — Model Setup: Approximate Methods —

For one 3-state channel, the state-space  $S$  with a particular partition structure  $P$  is given by  $S = \{C_1, C_2, C_3\}$  and  $P = \{C_1, C_2\}$ . In the case of  $N$  3-state channels, the induced partitioning on the expanded state-space is a mapping from  $N$  states to  $2^N$  partitions by  $\pi \cdot Q = \bigoplus_{i=1}^N K_i + \sum_{i,j=1}^N \lambda_{ij} \otimes K_j$ . We can use the *Iterative Method of Kronecker* (IMK) iteration where the transition matrix  $Q$  is decomposed into blocks over partitions  $P$  for  $q$  such that

$$Q \approx \bigoplus_{i=1}^N K_i + \sum_{i,j=1}^N \lambda_{ij} \otimes K_j$$

(approximate representation of  $Q$  corresponding to  $P$ ) is  $\pi \cdot Q \approx \pi \cdot Q \otimes \mathbf{1}$  with  $\sum_{i=1}^N \pi_i = 1$ . The approximate numerical method we consider that leverages this structure is APP-POWER.

## — Approximate Methods: Hierarchical Structure —

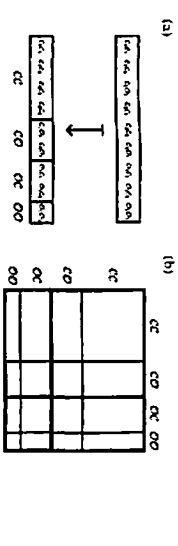


Figure 4: (a) Left: A diagram showing the hierarchical structure of the expanded generator matrix Q. Right: A diagram showing the hierarchical structure of the expanded generator matrix Q.

## — Approximate Methods: Partitioning —

P	Optimal P	Approximate P	Approximate P	Approximate P
1	1	1	1	1
2	2	2	2	2
3	3	3	3	3
4	4	4	4	4
5	5	5	5	5
6	6	6	6	6
7	7	7	7	7
8	8	8	8	8
9	9	9	9	9
10	10	10	10	10

Table 2: Approximate methods compared for APP-POWER and IMK iteration using 1 unit CPU with dual core Xeon E5-2680 v2 (16GB) from previous and with 100 units CPU. The method of choice is 100 units CPU. The benchmark problem is the same as the one used in the previous work [1]. The method of choice is 100 units CPU. The benchmark problem is the same as the one used in the previous work [1]. The method of choice is 100 units CPU. The benchmark problem is the same as the one used in the previous work [1].

## — Approximate Methods: Scalability and Error —

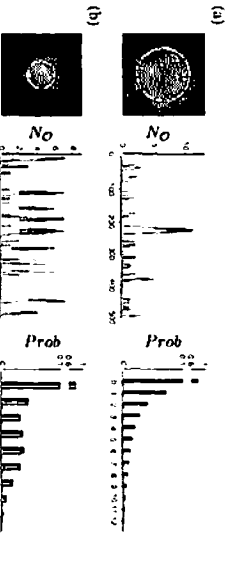


Figure 5: (a) Left: A plot of Wall Clock Time (sec) versus N. Right: A plot of Error versus Wall Clock Time (sec). (b) Left: A plot of Wall Clock Time (sec) versus N. Right: A plot of Error versus Wall Clock Time (sec).

## — References —

- [1] W. Stewart, *Introduction to the numerical solution of Markov chains*, Princeton University Press (1994).
- [2] V. Nguyen, R. Steinhilber, and G. D. Smith, *A stochastic automata network descriptor for Markov chain models of intracellular calcium release*, *Bulletin of Mathematical Biology*, 67 (2005), pp. 393–432.
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