

CHINA 中国

THINK OPEN 开放性思维



Disclaimer

- This talk vastly over-simplifies things.
- See notes for full details and resources.

https://github.com/gregkh/presentation-spectre

Spectre

- Hardware bugs
- Valid code can be "tricked" into exposing sensitive data to attacking programs.
- Exploits the speculative execution model of modern CPUs.
- Many different variants.
- Is going to be with us for a very long time!

Different Variants

- 1 Bounds Check Bypass (BCB)
- 2 Branch Target Isolation (BTI)
- 3 Rogue Data Cash Load (RDCL)
- 3a Rogue System Register Read (RSRE)
- 4 Speculative Store Bypass (SSB)
- 5 Lazy Floating Point State Restore (LazyFP)

variant 1 – Bounds check bypass

- Uses the kernel to read memory of another process or virtual machine.
- Fixed by core kernel changes.
- Lots of individual drivers still need to be fixed.

variant 1 – vulnerable code

```
int load_array(int *array, unsigned int user_value)
{
    if (user_value >= MAX_SIZE)
        return 0;

    return array[user_value];
}
```

· variant 1 – vulnerable code

```
int load_dependent_array(int *array1, int *array2,
                         int index)
      int value1, value2;
      value1 = load_array(array1, index);
      value2 = load_array(array2, value1);
      return value2;
```

```
int load_array(int *array, unsigned int user_value)
      if (user_value >= MAX_SIZE)
            return 0;
      user_value = array_index_nospec(user_value, MAX_SIZE);
      return array[user_value];
```

variant 1 – fixed code

```
#define array_index_nospec(index, size)
      typeof(index) _i = (index);
      typeof(size) _s = (size);
      unsigned long _mask = array_index_mask_nospec(_i, _s);\
      BUILD_BUG_ON(sizeof(_i) > sizeof(long));
      BUILD_BUG_ON(sizeof(_s) > sizeof(long));
      (typeof(_i)) (_i & _mask);
```

```
static inline unsigned long
      array_index_mask_nospec(unsigned long index,
                               unsigned long size)
      unsigned long mask;
      asm ("cmp %1,%2; sbb %0,%0;"
                       :"=r" (mask)
                       :"g"(size),"r" (index)
                       :"cc");
      return mask;
```

```
int load_array(int *array, unsigned int user_value)
      if (user_value >= MAX_SIZE)
            return 0;
      user_value = array_index_nospec(user_value, MAX_SIZE);
      return array[user_value];
```

variant 1 – Fix dates*

```
x86
   4.14.14
            17 January 2018
            17 January 2018
   4.9.77
            23 January 2018
   4.4.113
ARM
   4.15.4
            17 February 2018
            22 February 2018
   4.14.21
```



- These are the "first fixed" dates.
- Later kernels get more fixes and improvements.
- Keep updating your kernel & microcode!

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variant 1 – Fix dates again

```
x86
   4.16.11 22 May 2018
   4.14.43 22 May 2018
   4.9.102
            22 May 2018
ARM
   4.16
            1 April 2018
            20 April 2018
   4.9.95
```

variant 2 – Branch target injection

- Abuses the CPU branch predictor.
- Read data from kernel or other virtual machine.
- Fixed by compiler, kernel, & microcode updates.
- "retpoline"

variant 2 – Fix dates*

x86

4.15.9 11 March 2018

4.14.26 11 March 2018

4.9.87 11 March 2018

4.4.121 11 March 2018

*More fixes and optimizations happened in later kernels



- Spectre variant "3"
- Read kernel data from userspace.
- Fixed with page table isolation kernel changes (Kaiser for older kernels).
- Fix slows down enter/exit of the kernel.
- Implemented differently for different kernel releases and distros.



Meltdown – fix dates

17 February 2018

08 April 2018

```
x86
4.14.11 02 January 2018
4.9.75 05 January 2018
4.4.110 05 January 2018
ARM
4.15.4 17 February 2018
```

4.14.20

4.9.93



variant 3a – Rouge system register read

- Abuses the reading of system registers.
- Read data from kernel or other virtual machine.
- Kernel fix for Meltdown solves this problem.

variant 4 – Speculative Store Bypass

- Can execute and read beyond what is expected.
- Read data from kernel or other virtual machine.
- Minor kernel changes.
- Microcode update required for full protection.



variant 4 – Speculative Store Bypass

```
x86
```

```
4.16.11 22 May 2018
```

4.14.43 22 May 2018

4.9.102 22 May 2018

variant 5 – Lazy FP state restore

- Uses the old "lazy floating point restore" method to read memory of another process or virtual machine.
- Details to be published June 27.
- Linux kernel fixed in 2016.

variant 5 – Fix dates

```
x86
4.6 15 May 2016
```

4.4.138 16 June 2018

Why this is such a big deal

- CPU bugs require software & microcode fixes.
- All operating systems are affected.
- Performance will decrease.
- Totally new class of vulnerabilities.
- We will be finding, and fixing, these for a very long time.

Linux's response

- Companies were notified, but not developers.
- Developers notified very late, resulting in delay of fixes.
- Majority of the world runs non-corporate kernels.
- Intel is now working with some developers.

Keeping a secure system

- Take ALL stable kernel updates.
 - Do NOT cherry-pick patches.
- Enable hardening features.
- Update to a newer major kernel version where ever possible.
- Update your microcode/BIOS!

