# **ECONOMICS 8100**

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# Part 1. Budget

# 1. Consumption Set X

**Assumptions:** (Universe of Choice Objects): X

**Bundles:** Elements of X.  $x \in X$ 

# Assumptions about X.

- 1.  $\emptyset \neq X \subseteq \mathbb{R}^n_+$ .
- 2. X is closed.
- 3. X is convex.
- 4.  $0 \in X$ .

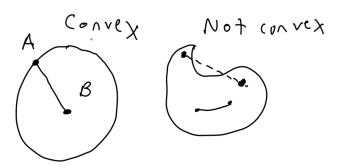


FIGURE 1.1. Examples of a Convex/Non-Convex Set.

# 2. Budget Set B

# Budget Set: $B \subseteq X$

X defines the scope of the model. B is what an  $individual\ consumer\ chooses\ among.$ 

**Example.** Budget Set with Prices and Income

$$B = \{x | x \in X \& x_1 p_1 + x_2 p_2 \le m\}$$

# Example. Ice Cream Bowls

Every ice cream bowl x has some non-negative number of scoops of Vanilla, Chocolate, Strawberry.

$$X = \mathbb{R}^3_+$$

Budget B is the set of bowls with no more than one scoop of ice cream.

$$B = \left\{ x | x \in R_+^3 \& \sum_{i=1}^3 x_i \le 1 \right\}$$

This is the unit-simplex in  $\mathbb{R}_3$ .

 $(1,0,0) \in B$ . (On the boundary.)

 $(0.5, 0.5, 0) \in B$ . (On the boundary.)

 $(0.25, 0.25, 0.25) \in B$ . (In the interior.)

 $(2,0,0) \notin B$ 

#### Part 2. Preference

#### 3. The Preference Relation

Preference Relation is a **Binary Relation**.

Formally, a binary relation on set X is a subset of the Cartesian product X with itself.

$$\succeq \subseteq X \times X$$

Another way to denote an ordered pair is "in" the relation:

If  $(x, y) \in \succeq$  we can also write  $x \succeq y$ .

Informally we say "x" is at least as good as "y", or "x" preferred "y".

Axioms of  $\succeq$ .

**Axiom 0** (reflexive):  $\forall x \in X, x \succeq x$ . This is implied by axiom 1.

**Axiom 1** (complete):  $\forall x, x' \in X$ , either  $x \succeq x'$  or  $x' \succeq x$  (or both).

The consumer has "some" preference over every pair of objects.

**Axiom 2** (transitive):  $\forall x, x', x'' \in X$  if  $x \succ x'$  and  $x' \succ x'' \Rightarrow x \succ x''$ .

≥ is a "weak order" if it is complete and transitive.

# 4. Relations and Sets Related to ≥

### **Subrelations:**

 $\sim$  is the indifference relation.  $x \succeq y$  and  $y \succeq x \Leftrightarrow x \sim y$ .

 $\succ$  is the strict relation.  $x \succeq y$  and not  $y \succeq x \Leftrightarrow x \succ y$ .

# Related Sets:

 $\succeq (x)$  "upper contour set"

#### 5. From Preferences to Choice

# Choice Correspondence.

We will assume that from a budget set B a consumer "chooses" choice set C according to their preference  $\succeq$ .  $C = \{x | x \in B \& \forall x' \in B, x \succeq x'\}$ .

Informally, C is the set of objects that are at least as good as anything else in the set.

# **Example With Transitive Preferences**

 $X = \{a, b, c\}. \ a \succeq b, c \succeq a, c \succeq b.$ 

$$C(\{a\}) = a, C(\{b\}) = b, C(\{c\}) = c$$
  
 $C(\{a,b\}) = a, C(\{a,c\}) = c, C(\{b,c\}) = c$ 

$$C\left(\{a,b,c\}\right)=c$$

- 6. Cycles Lead to Empty Choice Sets
- 6.1. The Problem with Intransitive Preferences.  $X = \{a, b, c\}$ .  $a \succeq b, c \succeq a, b \succeq c$ . This is intransitive!

Choice correspondence:

$$C: P\left(X\right)/\emptyset \to X$$
 
$$C\left(\left\{a\right\}\right) = a, C\left(\left\{b\right\}\right) = b, C\left(\left\{c\right\}\right) = c$$
 
$$C\left(\left\{a,b\right\}\right) = a, C\left(\left\{a,c\right\}\right) = c, C\left(\left\{b,c\right\}\right) = b$$
 
$$C\left(\left\{a,b,c\right\}\right) = \emptyset$$

This consumer cannot make a choice from the set  $\{a, b, c\}$ .

6.2. Cycles and Empty Choices. Notice in the previous example,  $a \succ b, a \succ c, c \succ a$ . We have proved (essentially) that if there is a cycle, there is an empty choice set.

In fact, suppose, there is an empty choice set  $\mathbf{and}\ X$  is finite. There must be a cycle.

$$\forall x \in B, \# (\succsim (x)) < \# (B)$$

By completeness,  $\forall x \exists x' \in X : x' \succ x$ . Choose an  $x_1$ , let  $x_2$  be any element of  $\succ (x_1)$ . We have  $x_2 \succ x_1$ . If there is an  $x_3 \in \succ (x_2)$  such that  $x_1 \succ x_3$  we have identified a cycle. Otherwise, we continue with an inductive step. Suppose we have  $x_n \succ \dots \succ x_1 . \succ (x_n)$  is non-empty. Either it contains an element  $x_{n+1}$  such that there is an  $x_i \succ x_{n+1}$  in which case we have identified a cycle or it does not and we continue with another inductive step. Either we find a cycle or reach the  $N_{th}$  step

with  $x_N \succ x_{n-1} \succ ... \succ x_1$ .  $\succ (x_N)$  is non-empty.

So, the cycle condition is equivalence to a non-empty choice set. Transitivity of  $\succeq$  implies transitivity of  $\succ$  which implies no cycles (try this last step at home). But do no-cycles imply transitivity of  $\succeq$ ? No. Here is a counter-example:

$$x \succ y, y \sim z, z \succ x$$

# 7. Intransitivity: Empty Choices, Incoherent Choices: Pick One.

So if no-cycles of the strict preference is equivalent to non-empty choice (in finite sets), and transitvity of  $\succeq$  is not equivelent to no-cycles, why do we assume it?

Finite non-emptyness: For any B with  $\#(B) \in \mathbb{I}$ ,  $C(B) \neq \emptyset$ 

**Coherence**: For every x, y and B, B' such that  $x, y \in B \cap B'$ ,  $x \in C(B) \land y \notin C(B) \Rightarrow y \notin C(B')$ .

Suppose there is an intransitive  $\succeq$ . There exists either a B where  $C(B) = \emptyset$  or there exists a x, y, B, B' where the choice correspondence is incoherent. By intransitivity:

1) 
$$x \succ y, y \succ z, z \succ x$$
  

$$C(\{x, y, z\}) = \emptyset$$
2)  $x \sim y, y \sim z, z \succ x$   
3)  $x \sim y, y \succ z, z \succ x$ 

For both of these we have incoherent choice from the following sets:

$$x \notin C(\{x, y, z\})$$

$$y \in C(\{x, y, z\})$$

$$x \in C(\{x, y\})$$

$$4) x \succ y, y \sim z, z \succ x$$

Can you find the incoherent choice?