## Stochastic Linear Bandits An Empirical Study

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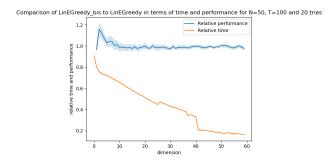
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## 1 Problem 1: Linear epsilon greedy

1. q1

2. q2

3. According to the documentation of numpy, the complexity of the pinv function is  $O(min(nm^2, n^2m))$ . In our problem, the matrix is squared, of size d so the complexity is  $O(d^3)$ . This can create problems when facing high-dimensional problems. We have therefore decided to implement a class LinearEpsilonGreedybis, in which we have changed the estimation of theta. Instead of estimating  $\theta$  through the least square estimator, we decided to estimate it through this estimator:  $\hat{\theta} = \sum_{t=1}^{T} \langle \theta, A_t \rangle A_t$ . We didn't manage to find theoretical guarantees about the expected value of this estimator, as  $\mathbb{E}(\hat{\theta}) = \sum_{t=1}^{T} \mathbb{E}(\langle \theta, A_t \rangle A_t)$ , which can't be precised without assumptions on the distribution of  $A_t$ . However, we have tested it on different problems, and it seems to obtain the same results as the one obtained with the least square estimator. Computing  $\hat{\theta}$  has a complexity in O(d), as we only have to compute scalars products of dvectors. The figure 1 underlines the gain in computational time, while the performances are the same.



**Figure 1:** Comparison of the performances and rime of execution of LinearEpsilonGreedy and the LinearEpsilonGreedy bis, with N=50, T=200 and 20 tries.

## 2 Problem 2: LinUCB and LinTS

- 1. q1
- 2. q2
- 3. q3