

$$\text{NWD}(m, n) = \text{NWD}(n, m \bmod n) \quad (*)$$

IN: $m \in \mathbb{Z}, n \in \mathbb{N}$

OUT: $\text{NWD}(m, n)$

\uparrow
 $0, 1, \dots, n-1$

$d \leftarrow m$
 $d' \leftarrow n$

while $d' \neq 0$:

$\rightarrow d, d' \leftarrow d', d \bmod d'$

return d

NIEZMIENNIK
 $\text{NWD}(m, n) = \text{NWD}(d, d')$

1) program się kończy, bo

$\Rightarrow d'_{\text{nowe}} < d'$

2) $(*) \Rightarrow$ NIEZMIENNIK

3) $\text{NWD}(m, n) = \text{NWD}(d, d')$

$\Rightarrow \text{NWD}(m, n) = \text{NWD}(d, d')$

$\wedge d' = 0$

$\text{NWD}(m, n) = \text{NWD}(d, 0) = d$

$m, n \sim 2^{1000}$

$\text{NWD}(m, n) \sim \sqrt{2^{1000}} \approx 2^{500}$

ILE TEST OBROTU PETLI

$$(d, d') \rightsquigarrow (d_{\text{move}}, d'_{\text{move}})$$

$$\begin{cases} d_{\text{move}} = d' \\ d'_{\text{move}} = d \bmod d' \end{cases}$$

$$d \geq d'$$

$$(m, n) = (d_0, d'_0) \rightsquigarrow (d_1, d'_1) \rightsquigarrow (d_2, d'_2) \rightsquigarrow \dots \rightsquigarrow (d_i, d'_i) \rightsquigarrow$$

$$\rightsquigarrow (d_{i+1}, 0)$$

$d \geq d' \Rightarrow q \geq 1$ $d \bmod d' = d'_{\text{move}}$

$$\begin{aligned} d &= qd' + r \geq d' + r \geq \\ &\geq d_{\text{move}} + d'_{\text{move}} \geq \underline{2d'_{\text{move}}} \end{aligned}$$

$$0 \leq r \leq d' - 1$$

$$\boxed{d \cdot d'} \geq 2d'_{\text{move}} d' = 2 \boxed{d'_{\text{move}} d_{\text{move}}}$$

$$(d_j, d'_j) \rightsquigarrow (d_{j+1}, d'_{j+1})$$

$$d_j d'_j \geq 2d_{j+1} d'_{j+1}$$

$$\boxed{(d_{i+1}, 0)}$$

$$(m, n) = (d_0, d'_0) \rightsquigarrow (d_1, d'_1) \rightsquigarrow (d_2, d'_2) \rightsquigarrow \dots \rightsquigarrow \boxed{(d_i, d'_i)}$$

$$mn = d_0 d'_0 \geq 2 d_1 d'_1 \geq 4 d_2 d'_2 \geq 8 d_3 d'_3 \geq \dots \geq$$

$$\geq \underbrace{2^i d_i d'_i}_{>0} \geq 2^i$$

$$mn \geq 2^i \quad | \log_2()$$

$$\log_2(mn) \geq i$$

$$\boxed{i \leq \log_2 m + \log_2 n}$$

$$m \in \mathbb{Z}, n \in \mathbb{N}$$

$$\text{LICBA KROKOL} \leq \underline{\log_2 |m| + \log_2 n + 2}$$

$$m, n \sim 2^{1000}$$

$$\text{Licba Krokolo} \leq 1000 + 1000 + 2 = 2002$$

$$\begin{array}{r|l}
 660 & \\
 525 & \\
 135 & \leftarrow 135 = 660 - 1 \cdot 525 \\
 120 & \leftarrow 120 = 525 - 3 \cdot 135 \\
 \underline{15} & \leftarrow 15 = 135 - 1 \cdot 120 \\
 0 &
 \end{array}$$

$$\boxed{\text{NWD}(660, 525)} = 15 = 135 - 1 \cdot 120 = 135 - 1 \cdot (525 - 3 \cdot 135) =$$

$$= -1 \cdot 525 + 4 \cdot 135 = -1 \cdot 525 + 4(660 - 1 \cdot 525) =$$

$$= \boxed{4} \cdot 660 + \boxed{-5} \cdot 525$$

$$\text{NWD}(m, n) = s \cdot m + t \cdot n \quad ?$$

$\begin{array}{c}
 d_0 \\
 d_1 \\
 d_2 \\
 \vdots \\
 d_i
 \end{array}
 \left| \begin{array}{c}
 \uparrow \\
 \uparrow \\
 \uparrow \\
 \uparrow \\
 \uparrow
 \end{array} \right.$

Rekursiver Algorithmus Euklidese

$$d, d' \leftarrow m, n$$

while $d' \neq 0$

$$d, d' \leftarrow d', d \bmod d'$$

$$\text{NWD}(m, n) = d_{i+1} = s \cdot m + t \cdot n$$

? ?

$$\begin{aligned}
 & \left[\begin{array}{l} d_k = s_k m + t_k n \\ d_{k+1} = s_{k+1} m + t_{k+1} n \end{array} \right. \quad \textcircled{2} \\
 & d_k = q_{k+1} d_{k+1} + \boxed{r_k}, \quad q_{k+1} = d_k \text{ div } d_{k+1} \\
 & d_{k+2} = \overbrace{d_k \bmod d_{k+1}} = \\
 & \quad = d_k - \underbrace{(d_k \text{ div } d_{k+1})}_{q_{k+1}} \cdot d_{k+1} = \\
 & \quad = d_k - q_{k+1} d_{k+1} = \\
 & \quad = s_k m + t_k n - q_{k+1} (s_{k+1} m + t_{k+1} n) \\
 & \quad = \underbrace{(s_k - q_{k+1} s_{k+1})}_{s_{k+2}} m + \underbrace{(t_k - q_{k+1} t_{k+1})}_{t_{k+2}} n
 \end{aligned}$$

$$\textcircled{2} \begin{bmatrix} d_k = s_k m + t_k n \\ d_{k+1} = s_{k+1} m + t_{k+1} n \end{bmatrix}$$

$$d_k = q_{k+1} d_{k+1} + r_k$$

$$\underline{d_{k+2}} = s_{k+2} m + t_{k+2} n$$

$$s_{k+2} = \underbrace{s_k - q_{k+1} s_{k+1}}_{\uparrow} \quad \underbrace{t_k - q_{k+1} t_{k+1}}_{\nwarrow} = t_{k+2}$$

$$\underline{d_{k+2} = d_k - q_{k+1} d_{k+1}}$$

$$d, d' \leftarrow m, n$$

$$\left\{ \begin{array}{l} d = m = 1 \cdot m + 0 \cdot n \\ d' = n = 0 \cdot m + 1 \cdot n \end{array} \right.$$

$$s, s' \leftarrow 1, 0$$

$$t, t' \leftarrow 0, 1$$

while $d \neq 0$

$$\underline{\cancel{d, d' \leftarrow d, d \bmod d'}}$$

$$q \leftarrow d \text{ div } d'$$

$$d, d' \leftarrow d', d - q \cdot d'$$

$$s, s' \leftarrow s', s - q \cdot s'$$

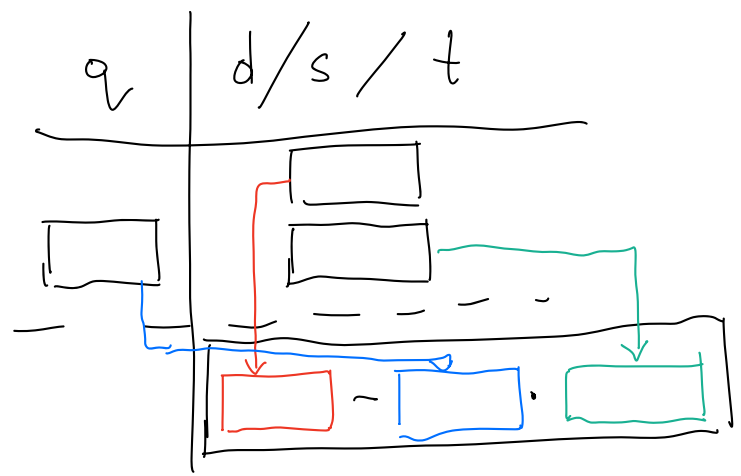
$$t, t' \leftarrow t', t - q \cdot t'$$

$$d = s m + t n \quad d' = s' m + t' n$$

$$\begin{array}{c} \boxed{d}, \boxed{d'} \leftarrow \boxed{d'}, \boxed{d} - q \cdot \boxed{d'} \\ \uparrow \quad \uparrow \quad \quad \uparrow \quad \uparrow \quad \quad \uparrow \\ d, s, t \quad d', s', t' \quad d', s', t' \quad d, s, t \quad d', s', t' \end{array}$$

$$d = s m + t n$$

d	q	s	t
660		1	0
525	1	0	1
135	3	1	-1
120	1	-3	4
15	8	4	-5
0			



$$15 = 4 \cdot 660 - 5 \cdot 525$$

$$\text{NWD}(137, 92) = s \cdot 137 + t \cdot 92$$

d	q	s	t
137		1	0
92	1	0	1
45	2	1	-1
2	22	-2	3
1	2	45	-67
0			

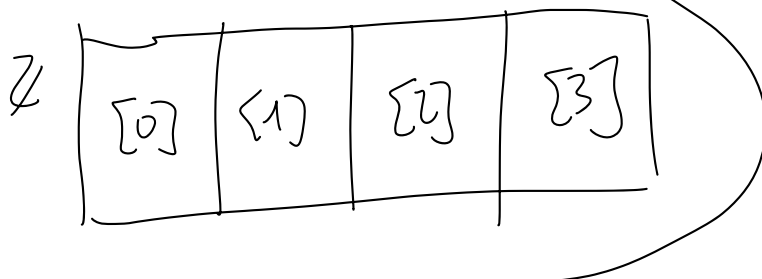
$$\text{NWD}(137, 92) = 1 = 45 \cdot 137 - 67 \cdot 92$$

Relacja przystawania (kongruencji)

$$n \in \mathbb{N}, \\ a, b \in \mathbb{Z}$$

$$X = \mathbb{Z}$$

$$(a \sim b) \Leftrightarrow n \mid a - b$$



$$a \sim b \leadsto a \equiv b \pmod{n}$$

$$a \equiv_n b$$

$$10 \equiv 3 \pmod{7}$$

$$24 \equiv 3 \pmod{7}$$

$$-4 \equiv 10 \equiv 3 \pmod{7}$$

$$(a + b) \bmod n = ((a \bmod n) + (b \bmod n)) \bmod n$$

$$(ab) \bmod n = ((a \bmod n)(b \bmod n)) \bmod n$$

Fakt. $\begin{cases} a \equiv b \pmod{n} \\ c \equiv d \pmod{n} \end{cases}$

$$\Rightarrow a + c \equiv b + d \pmod{n}$$

$$\Rightarrow a \cdot c \equiv b \cdot d \pmod{n}$$

$$\Rightarrow a^k \equiv b^k \pmod{n}$$

$$a \equiv b \pmod{n}$$

Zad. Wyznaczyć resztę z dzielenia
 7^{2025}

przez 15.

$$7^1 \equiv 7 \pmod{15} \quad | \cdot 7$$

$$7^2 \equiv 49 \equiv 4 \pmod{15} \quad | \cdot 7$$

$$7^3 \equiv 4 \cdot 7 \equiv 28 \equiv 13 \pmod{15} \quad | \cdot 7$$

$$7^4 \equiv 13 \cdot 7 \equiv 91 \equiv 1 \pmod{15} \quad | ()^k$$

$$7^{4k} \equiv 1 \pmod{15}, \quad k \in \mathbb{N}$$

$$7^{2025} = 7^{2024+1} = 7^{2024} \cdot 7 = \underbrace{7^{4 \cdot k}}_{1} \cdot 7 \overset{56k}{=} 1 \cdot 7 = 7 \pmod{15}$$