

$$\text{NWD}(m, n) = \text{NWD}\left(n, m \bmod n\right) \quad (*)$$

IN:  $m \in \mathbb{Z}, n \in \mathbb{N}$

OUT:  $\text{NWD}(m, n)$

$0, 1, \dots, n-1$

$d \leq m$   
 $d' \leq n$

while  $d \neq 0$ :

$d, d' \leftarrow d', d \bmod d'$

return  $d$

$d'_{\text{nowe}} \leq d' - 1$

NIEZMIENNIK

$$\text{NWD}(m, n) = \text{NWD}(d, d')$$

1) program służy do

$$d'_{\text{nowe}} < d'$$

2) (\*)  $\Rightarrow$  NIEZMIENNIK

$$3) \text{NWD}(m, n) = \text{NWD}(d, d')$$

$$\rightarrow \text{NWD}(m, n) = \text{NWD}(d, d')$$

$$\wedge \quad d' = 0$$

$$\text{NWD}(m, n) = \text{NWD}(d, 0) = d$$

$$m, n \sim 2^{1000}$$

$$\text{NWD}(m, n) \sim \sqrt{2^{1000}} \approx 2^{500}$$

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$$(d, d') \rightsquigarrow (d_{\text{nowe}}, d'_{\text{nowe}})$$

$$\begin{cases} d_{\text{nowe}} = d \\ d'_{\text{nowe}} = d \bmod d' \end{cases}$$

$$d > d'$$

$$(m, n) = (d_0, d'_0) \rightsquigarrow (d_1, d'_1) \rightsquigarrow (d_2, d'_2) \rightsquigarrow \dots \rightsquigarrow (d_i, d'_i) \rightsquigarrow$$

$$\underline{d > d'} \Rightarrow q \geq 1 \quad d \bmod d' = d'_{\text{nowe}} \rightsquigarrow (d_{i+1}, 0)$$

$$\begin{aligned} d &= qd' + r \geq d' + r \geq \\ &\geq d_{\text{nowe}} + d'_{\text{nowe}} \geq \underline{2d'_{\text{nowe}}} \quad 0 \leq r \leq d' - 1 \\ &\geq 2d'_{\text{nowe}} \end{aligned}$$

$$\boxed{d \cdot d'} \geq 2d'_{\text{nowe}} \quad d = 2 \boxed{d'_{\text{nowe}} \quad d_{\text{nowe}}}$$

$$(d_j, d'_j) \rightsquigarrow (d_{j+1}, d'_{j+1})$$

$$d_j \cdot d'_j \geq 2d_{j+1} \cdot d'_{j+1}$$

$$\left. \begin{array}{c} \overbrace{(d_{i+1}, 0)} \\ \underbrace{\phantom{(d_{i+1}, 0)}} \end{array} \right\}$$

$$(m, n) = (d_0, d'_0) \rightsquigarrow (d_1, d'_1) \rightsquigarrow (d_2, d'_2) \rightsquigarrow \dots \rightsquigarrow \boxed{(d_i, d'_i)}$$

$$mn = d_0 d'_0 \geq 2 d_1 d'_1 \geq 4 d_2 d'_2 \geq 8 d_3 d'_3 \geq \dots \geq$$

$$\geq \underbrace{2^i d_i d'_i}_{> 0} \geq 2^i$$

$$mn \geq 2^i \quad | \log_2()$$

$$\log_2(mn) \geq i$$

$$\boxed{i \leq \log_2 m + \log_2 n}$$

$$m \in \mathbb{Z}, \quad n \in \mathbb{N}$$

$$\text{LICABA KROKSL} \leq \underline{\log_2 |m| + \log_2 n + 2}$$

$$m, n \approx 2^{1000}$$

$$\text{Liebe Ende} \leq 1000 + 1000 + 2 = 2002$$

$$\begin{array}{r}
 660 \\
 525 \\
 135 \\
 120 \\
 \hline
 15 \\
 0
 \end{array}
 \quad
 \begin{aligned}
 &\leftarrow 135 = 660 - 1 \cdot 525 \\
 &\leftarrow 120 = 525 - 3 \cdot 135 \\
 &\leftarrow 15 = 135 - 1 \cdot 120 \\
 &\boxed{\text{NWD}(660, 525)} = 15 = 135 - 1 \cdot 120 = 135 - 1 \cdot (525 - 3 \cdot 135) = \\
 &= -1 \cdot 525 + 4 \cdot 135 = -1 \cdot 525 + 4(660 - 1 \cdot 525) = \\
 &= \boxed{4} \cdot 660 + \boxed{-5} \cdot 525
 \end{aligned}$$


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$$\text{NWD}(m, n) = s \cdot m + t \cdot n \quad ?$$

$d_0, d_1, d_2, \dots, d_i$       ↑  
 ↓

Rücksichtiger Algorithmus Euklides

$d, d' \leftarrow m, n$   
 $\text{while } d' \neq 0$   
 $d, d' \leftarrow d', d \bmod d'$

$d_k = s_k m + t_k n$   
 $d_{k+1} = s_{k+1} m + t_{k+1} n$   
 $d_k = q_{k+1} d_{k+1} + \boxed{r_k}, q_{k+1} = d_k \text{ div } d_{k+1}$   
 $d_{k+2} = \boxed{d_k \bmod d_{k+1}} =$   
 $= d_k - \underbrace{(d_k \text{ div } d_{k+1})}_{q_{k+1}} \cdot d_{k+1} =$   
 $= d_k - q_{k+1} d_{k+1} =$   
 $= s_k m + t_k n - q_{k+1} (s_{k+1} m + t_{k+1} n)$   
 $= \underbrace{(s_k - q_{k+1} s_{k+1})}_s m + \underbrace{(t_k - q_{k+1} t_{k+1})}_t n$

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$$\text{NWD}(m, n) = d_{i+1} = s \cdot m + t \cdot n \quad ? \quad ?$$

$$\left. \begin{array}{l} d_k = s_k u + t_k u \\ d_{k+1} = s_{k+1} u + t_{k+1} u \end{array} \right\}$$

$$d_k = q_{k+1} d_{k+1} + r_k$$

$$\cancel{d_{k+2}} = s_{k+2}^m + t_{k+2}^n$$

$$s_{k+2} = \underbrace{s_k - q_{k+1}s_{k+1}}_{t_k - q_{k+1}t_{k+1}} \quad t_k - q_{k+1}t_{k+1} = t_{k+2}$$

$$d_{k+2} = d_k - q_{k+1} d_{k+1}$$

$$d, d' \in \underline{m, n} \quad \left\{ \begin{array}{l} d = m = 1 \cdot m + 0 \cdot n \\ d' = n = 0 \cdot m + 1 \cdot n \end{array} \right.$$

$$\begin{array}{c} \cancel{s}, \cancel{s} \\ | \quad | \\ t, \quad t \end{array} \leftarrow \begin{array}{c} 1, 0 \\ 0, 1 \end{array}$$

while  $d \neq 0$

~~while  $d \neq 0$   
 $d = d_1, d \bmod d_1$~~

$$d = s_m + t_n \quad d' = \underline{s_m} + \underline{t_n}$$

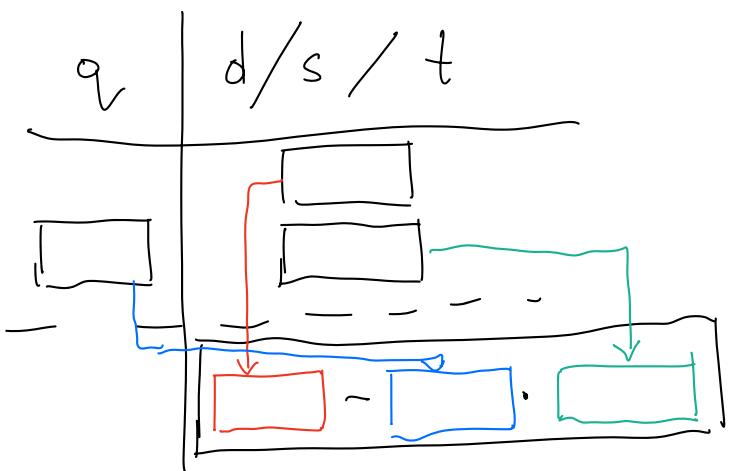
$\vec{v} \leftarrow d \text{ div } d'$

$$d' \leftarrow d' - q \cdot d'$$

$$\begin{array}{c} \text{s} \quad \text{s}' \\ | \quad | \\ \text{t} \quad \text{t}' \end{array} \leftarrow \begin{array}{c} \text{s} \quad \text{s}' \\ | \quad | \\ \text{t} \quad \text{t}' \end{array}, \begin{array}{c} \text{s} - \text{q} \cdot \text{s}' \\ | \\ \text{t} - \text{q} \cdot \text{t}' \end{array}$$

$$d = sm + th$$

d	q	s	t
660		1	0
525	1	0	1
135	3	1	-1
120	1	-3	4
15	8	4	-5
0			



$$15 = 4 \cdot 660 - 5 \cdot 525$$

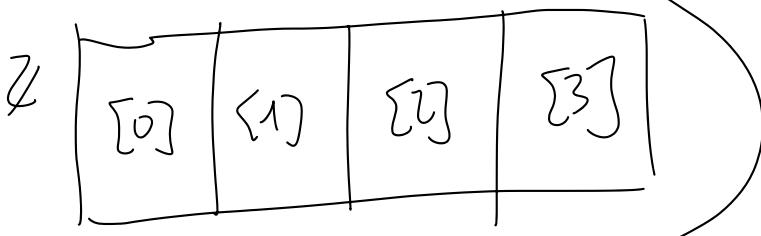
$$\text{NWD}(137, 92) = s \cdot 137 + t \cdot 92$$

d	q	s	t
137		1	0
92	1	0	1
45	2	1	-1
2	22	-2	3
1	2	45	-67
0			

$$\text{NWD}(137, 92) = 1 = 45 \cdot 137 - 67 \cdot 92$$

# Relacje przystawienia (kongruencji)

$n \in \mathbb{N}$        $X = \mathbb{Z}$   
 $a, b \in \mathbb{Z}$        $a \sim b \Leftrightarrow n | a - b$



$a \sim b \rightsquigarrow a \equiv b \pmod{n}$        $a \equiv_n b$

$$10 \equiv 3 \pmod{7}$$

$$24 \equiv 3 \pmod{7}$$

$$-4 \equiv 10 \equiv 3 \pmod{7}$$

$$(a+b) \bmod n = ((a \bmod n) + (b \bmod n)) \bmod n$$

$$(ab) \bmod n = ((a \bmod n)(b \bmod n)) \bmod n$$

Fakt.  $\begin{cases} a \equiv b \pmod{n}, \\ c \equiv d \pmod{n} \end{cases}$

$$\Rightarrow a+c \equiv b+d \pmod{n}$$

$$\Rightarrow a \cdot c \equiv b \cdot d \pmod{n}$$

$$\Rightarrow a^k \equiv b^k \pmod{n}$$

$$a \equiv b \pmod{n}$$

Zed. Hypothen verste & drehen

$\gamma^{2025}$

pfer 15.

$$\gamma^1 \equiv \gamma \pmod{15} \quad | \cdot 4$$

$$\gamma^2 \equiv 4\gamma \equiv 4 \pmod{15} \quad | \cdot 7$$

$$\gamma^3 \equiv 4 \cdot 7 \equiv 28 \equiv 13 \pmod{15} \quad | \cdot 7$$

$$\gamma^4 \equiv 13 \cdot 7 \equiv 91 \equiv 1 \pmod{15} \quad | 0 \downarrow$$

$$\gamma^{4k} \equiv 1 \pmod{15}, \quad k \in \mathbb{N} \quad \text{sof}$$

$$\gamma^{2025} = \gamma^{2024+1} = \gamma^{2024} \cdot \gamma = (\underbrace{\gamma^{4k}}_{1}) \cdot \gamma \equiv 1 \cdot 7 = 7 \pmod{15}$$