

1 New Random Edge Weights at Each Iteration

Define the PageRank iteration $x^{(k+1)} = \alpha(P + Q_k)x^{(k)} + (1 - \alpha)v$, where each Q_k is independently drawn from the probability distribution function f , and the other variables are defined as usual. We require that f be defined such that $P + Q_k$ is always a valid PageRank matrix. For example, P and $P + Q_k$ must be column stochastic, so the elements of any column of Q_k must sum to zero. Furthermore, without loss of generality, assume that $E[f] = 0$. (If $E[f]$ was not zero, we could redefine P to be $P + E[f]$ and f to be $f - E[f]$.)

By the linearity of expectation, $E[x^{(k+1)}] = \alpha PE[x^{(k)}] + (1 - \alpha)v$, and of course $E[x^{(0)}] = x^{(0)}$, so computing $E[x^{(k+1)}]$ is equivalent to a traditional PageRank problem without randomness. Determining the likely accuracy of this iteration is more complicated. (Note: To ease notation, I use $y = x^{(k+1)}$ and $x = x^{(k)}$ for the rest of this section.) My first attempt was to compute $\text{Var}(\|y\|_2)$, as follows:

$$\begin{aligned} \text{Var}(\|y\|) &= E[\|y\|^2] - E[\|y\|]^2 \\ &= E\left[\sum_{i=1}^n y_i^2\right] - E[\|y\|]^2 \\ &= \left(\sum_{i=1}^n E[y_i^2]\right) - E[\|y\|]^2 \end{aligned}$$

I don't know how to compute the last term, but knowing $E[y_i^2]$ would allow us to compute $\text{Var}(y_i)$, since $\text{Var}(y_i) = E[y_i^2] - E[y_i]^2$, and we know how to compute $E[y_i]$. More generally, I derived an algorithm for computing $E[yy^T]$ given $E[xx^T]$. (This allows us to recursively compute $E[yy^T]$.) We can read the expected values of the squares of the elements of y off the diagonal of this matrix. Furthermore, $E[yy^T] - E[y]E[y]^T$ is the covariance matrix of the elements of y , since $(E[yy^T] - E[y]E[y]^T)_{i,j} = E[y_i y_j] - E[y_i]E[y_j] = \text{Cov}(y_i, y_j)$.

Suppose we're given Q , and let $A = P + Q$, so $y = \alpha Ax + (1 - \alpha)v$. Then,

$$\begin{aligned} E[yy^T | Q] &= E[(\alpha Ax + (1 - \alpha)v)(\alpha x^T A^T + (1 - \alpha)v^T) | Q] \\ &= \alpha^2 AE[xx^T]A^T + \alpha(1 - \alpha)(AE[x]v^T + vE[x]^T A^T) + (1 - \alpha)^2 vv^T \end{aligned}$$

Since $E[Q] = 0$ and $E[A] = E[P] + E[Q] = P$, we have the following:

$$E[yy^T] = \alpha^2 PE[xx^T]P^T + \alpha^2 E[QE[xx^T]Q^T] + \alpha(1 - \alpha)(PE[x]v^T + vE[x]^T P^T) + (1 - \alpha)^2 vv^T$$

Since we know $E[x]$ and $E[xx^T]$, the only term we don't know is $E[QE[xx^T]Q^T]$. For all integers q, r, s, t from 1 to n , let $c_{q,r,s,t} = \text{Cov}(f_{q,r}, f_{s,t}) = E[f_{q,r}f_{s,t}] - E[f_{q,r}]E[f_{s,t}] = E[f_{q,r}f_{s,t}]$, the covariance of two elements of the output of f .

(These values can be precomputed before any iterations of this algorithm.) Then $(QE [xx^T] Q^T)_{q,s} = Q_{q,:} E [xx^T] (Q_{s,:})^T = \sum_{r,t=1}^n Q_{q,r} E [xx^T]_{r,t} Q_{s,t}$, and its expected value is $\sum_{r,t=1}^n c_{q,r,s,t} E [xx^T]_{r,t}$. The entire matrix $E [QE [xx^T] Q^T]$ can be computed by iterating over r and t from 1 to n .

Some remaining problems include finding a closed form for this algorithm, and/or an approximation algorithm that makes it easier to understand how the variance of f results in variance in the PageRank vector, especially as k becomes large. I decided to hold off on that for now in case what I've done so far is on the wrong track.

2 Approximations Regarding Random Edge Weights in a Direct Solution

Suppose we have a PageRank Problem $Mx = b$, but now we add randomness, so we try to solve $(M - E)x_e = b$, where E is random and has expected value 0. Assume that E is small enough so that $(M - E)^{-1} = \sum_{k=0}^{\infty} M^{-1} (EM^{-1})^k$. If E is guaranteed to be sufficiently small, then $(M - E)^{-1} \approx M^{-1} + M^{-1}EM^{-1}$. ($M^{-1}EM^{-1}EM^{-1}$ is a good estimate of the error of this first order approximation, for sufficiently small E .) This allows us to approximate the effect of randomness on the solution:

$$\begin{aligned} x_e - x &= (M - E)^{-1}b - x \\ &= (M - E)^{-1}Mx - x \\ &= \left((M - E)^{-1}M - I \right) x \\ &\approx \left((M^{-1} + M^{-1}EM^{-1})M - I \right) x \\ &= (I + M^{-1}E - I)x \\ &= M^{-1}Ex \end{aligned}$$

It follows that $E[x_e - x] \approx M^{-1}E[E]x = 0$, and also, $E[(M - E)^{-1}] \approx M^{-1} + M^{-1}E[E]M^{-1} = M^{-1}$. Regardless of the approximation used, knowing $E[(M - E)^{-1}]$ allows us to find the expected solution of the PageRank problem, since $E[(M - E)^{-1}b] = E[(M - E)^{-1}]b$. If instead we want a second order approximation of $(M - E)^{-1}$, we can compute its expected value as follows:

$$\begin{aligned} E[(M - E)^{-1}] &\approx E[M^{-1} + M^{-1}EM^{-1} + M^{-1}EM^{-1}EM^{-1}] \\ &= M^{-1} + M^{-1}E[EM^{-1}E]M^{-1} \end{aligned}$$

$E[EM^{-1}E]$ depends on the covariances of the elements of E , and can be computed in a manner similar to that in the last paragraph of the previous section. Similarly, a third-order approximation also depends on $M^{-1}E[EM^{-1}EM^{-1}E]M^{-1}$,

which depends on the expected values of products of three elements of E . The same pattern holds for higher-order approximations.

The variance of the first-order approximation of $(M - E)^{-1}$ is $\text{Var}(M^{-1} + M^{-1}EM^{-1}) = \text{Var}(M^{-1}EM^{-1}) = \mathbb{E}[(M^{-1}EM^{-1})^2] - \mathbb{E}[M^{-1}EM^{-1}]^2 = \mathbb{E}[M^{-1}EM^{-2}EM^{-1}] - 0^2 = M^{-1}\mathbb{E}[EM^{-2}E]M^{-1}$, which depends on the covariances of the elements of E . The variance of the second-order approximation can be computed as follows:

$$\begin{aligned} \text{Var}(\cdot) &= \text{Var}(M^{-1} + M^{-1}EM^{-1} + M^{-1}EM^{-1}EM^{-1}) \\ &= \text{Var}(M^{-1}EM^{-1} + M^{-1}EM^{-1}EM^{-1}) \\ &= M^{-2}\text{Var}(E + EM^{-1}E)M^{-2} \\ &= M^{-2}(\mathbb{E}[(E + EM^{-1}E)^2] - \mathbb{E}[E + EM^{-1}E]^2)M^{-2} \\ &= M^{-2}(\mathbb{E}[E^2] + \mathbb{E}[E^2M^{-1}E] + \mathbb{E}[EM^{-1}E^2] + \mathbb{E}[EM^{-1}E^2M^{-1}E] - \mathbb{E}[EM^{-1}E]^2)M^{-2} \end{aligned}$$

If the k th order approximation of $(M - E)^{-1}$ is $\sum_{i=0}^k (M^{-1}E)^i M^{-1}$, then its error is $\sum_{i=k+1}^{\infty} (M^{-1}E)^i M^{-1} = (M^{-1}E)^{k+1} \sum_{i=0}^{\infty} (M^{-1}E)^i M^{-1} = (M^{-1}E)^{k+1} (M - E)^{-1}$. We can bound the relative error of this approximation as follows:

$$\begin{aligned} \frac{\|(M^{-1}E)^{k+1}(M - E)^{-1}\|}{\|(M - E)^{-1}\|} &\leq \frac{(\|M^{-1}\|\|E\|)^{k+1}\|(M - E)^{-1}\|}{\|(M - E)^{-1}\|} \\ &= (\|M^{-1}\|\|E\|)^{k+1} \end{aligned}$$

If $M = I - \alpha P$, then $\|M^{-1}\| = \|\sum_{k=0}^{\infty} (\alpha P)^k\| \leq \sum_{k=0}^{\infty} (\alpha\|P\|)^k = \frac{1}{1 - \alpha\|P\|}$, so $(\|M^{-1}\|\|E\|)^{k+1} \leq \left(\frac{\|E\|}{1 - \alpha\|P\|}\right)^{k+1}$. Since the P is column stochastic in the PageRank problem, its 1-norm is 1, so the relative error in the 1-norm is bounded by $\left(\frac{\|E\|_1}{1 - \alpha}\right)^{k+1}$.

3 Miscellaneous Observations

Most of these observations don't seem useful at the moment, but it seemed like a good idea to keep them in case they prove useful later.

The Sherman Morrison Woodbury formula gives $(M + E)^{-1} = M^{-1} - M^{-1}(E^{-1} + M^{-1})^{-1}M^{-1}$, which looks similar to the first order approximation for $(M + E)^{-1}$.

Suppose we have a PageRank iteration $x_{k+1} = \alpha P x_k + (1 - \alpha)v$. Let $e_k = x_k - x_*$, where x_* is the fixed point of the iteration. Subtracting the equation $x_* = \alpha P x_* + (1 - \alpha)v$ from the original iteration equation yields $e_{k+1} = \alpha P e_k$, so $e_k = (\alpha P)^k e_0$.

Suppose we have a PageRank iteration $x_{k+1} = \alpha P x_k + (1 - \alpha)v$, with initial guess x_0 . I prove by induction that $x_k = (\alpha P)^k x_0 + (I - (\alpha P)^{k+1})(I - \alpha P)^{-1}(1 - \alpha)v$. A simple computation shows this is true for $k = 0$. If it's true for k , then

$$\begin{aligned}
x_{k+1} &= \alpha P x_k + (1 - \alpha)v \\
&= \alpha P ((\alpha P)^k x_0 + (I - (\alpha P)^{k+1})(I - \alpha P)^{-1}(1 - \alpha)v) + (1 - \alpha)v \\
&= (\alpha P)^{k+1} x_0 + (\alpha P(I - \alpha P)^{-1} - (\alpha P)^{k+1}(I - \alpha P)^{-1} + I)(1 - \alpha)v \\
&= (\alpha P)^{k+1} x_0 + (\alpha P - (\alpha P)^{k+1} + I - \alpha P)(I - \alpha P)^{-1}(1 - \alpha)v \\
&= (\alpha P)^{k+1} x_0 + (I - (\alpha P)^{k+2})(I - \alpha P)^{-1}(1 - \alpha)v
\end{aligned}$$

Suppose we have $x = \alpha P x + (1 - \alpha)v$. Note that $(I - \alpha P)^{-1} = \sum_{k=0}^{\infty} (\alpha P)^k = I + \sum_{k=1}^{\infty} \alpha^k P^k$. Therefore $x = (I - \alpha P)^{-1}(1 - \alpha)v = (I + \sum_{k=1}^{\infty} \alpha^k P^k)(1 - \alpha)v = v + \sum_{k=1}^{\infty} \alpha^k (P^k v - P^{k-1} v) = v + \sum_{k=1}^{\infty} \alpha^k P^{k-1} (P - I)v = (I + \sum_{k=1}^{\infty} (\alpha^k P^{k-1})(P - I))v = (I + \alpha(I - \alpha P)^{-1})(P - I)v$.

$(M - E)^{-1} = M^{-1} + M^{-1}E(M - E)^{-1}$, so $E[(M - E)^{-1}] = M^{-1} + M^{-1}E[E(M - E)^{-1}]$. This would be useful if E is drawn from a probability distribution for which $E[E(M - E)^{-1}]$ is easier to compute than $E[(M - E)^{-1}]$.