Cool: Language Overview

Classroom Object Oriented Language (Cool)

- Designed by Alex Aiken from Stanford for teaching compiler construction
- Supports modern language features
 - abstraction, reuse (inheritance), static typing, memory management
- Many features are left out
 - feasible to implement in a semester

Good News

- No arrays
- No floating point operations
- No "static" modifier
- No interfaces
- No method overloading (but still allow overriding)
- No exceptions
- No packages

Better News

Cool language is still rich enough for doing interesting things

Classes and Objects

- Cool programs are sets of class definitions
- A special class Main with a special method main
- A class is a collection of attributes and methods
- Instance of a class are objects
- The expression "new Point" creates a new object of class Point

```
class Point {
    x : Int <- 0;
    y : Int; (* use default value *)
};
```

Methods

- Manipulating attributes
- Refer to the current object using self

```
class Point {
  x : Int <-0;
  y : Int <- 0;
  movePoint (newx : Int, newy : Int) : Point {
    {       x <-newx;
            y <- newy;
            self;
      } -- close block expression
    }; -- close method
}; -- close class</pre>
```

```
class A {
    p : Point;
    q : Point;
    foo() : Int {{
        p <- new Point;
        q <-p.movePoint(1,2);
        1;
     }};
};</pre>
```

Information Hiding in Cool

- Methods are global (public)
- Attributes are local to a class (private)
- Attributes can only be accessed by the class's methods

```
class Point {
  x : Int <-0;
  y : Int <- 0;
  getx () : Int { x };
  setx (newx : Int) : Int { x <- newx };
};</pre>
```

Inheritance

 We can extend points to colored points using subclass (class hierarchy)

Cool Types

- Every class is a type
- Base classes
 - Int
 - Bool (true, false)
 - String
 - Object (root of the class hierarchy)
- All variables must be declared
- Compiler infers types for expressions

Cool Type Checking

 Type safety: a well-typed program cannot result in type errors at runtime

```
x : A;x <- new B;</li>(* well typed if A is an ancestor of B in the class hierarchy *)
```

Method Invocation and Inheritance

- Methods are invoked by dispatch
- Understanding dispatch in the presence of inheritance is a subtle aspect of OO languages

```
p : Point; --p has static type Point
p <- new ColorPoint; --p has dynamic type ColorPoint
p.movePoint(1,2); -- invoke ColorPoint version of movePoint
```

(* static dispatch: invoke the version of movePoint defined in Point *) p@Point.movePoint(1,2);

Cool memory management

- Memory is allocated every time new is invoked
- Memory is deallocated automatically when an object is not reachable anymore
 - done by the garbage collector (GC)
 - part of Cool runtime system

Other Expressions

- Every expression has a type and a value
- Expression language
 - Assignment x <- E
 - Loops: while E loop E pool
 - Conditionals: if E then E else E fi
 - Case: case E of x: Type E; esac
 - Let let x : Type <- E in E
 - Arithmetic and logical operations
 - Primitive I/O out_string(s), in_string(), ...