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## **AOA Experiment 9**

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Aim:

To implement & analyze Naive String matching algorithm:

## Implementation:

```
#include<stdio.h>
#include<string.h>
void Naive_string_matcher(char *T, char *P)
{
    int m = strlen(T);
    int n = strlen(P);
    for(int s = 0; s ≤ m - n; s++)
    {
        int x;
        for(x = 0; x < n; x++)
            if(T[s + x] ≠ P[x])
            {
                break;
            }
        if(x == n)
        {
            printf("Pattern Occurs with shift:\t%d\n", s);
        }
    }
}
int main()
{
    char str[30], pattern[30];
    printf("\nEnter a Text:\t");
    scanf("%s", str);
    printf("\nEnter a Pattern:\t");
    scanf("%s", pattern);
    Naive_string_matcher(str, pattern);
    return 0;
}
```

Output:

```
Enter a Text:  moon

Enter a Pattern:  on
Pattern Occurs with shift:  2

-----
Process exited after 5.402 seconds with return value 0
Press any key to continue . . .
```

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## Exp 09:

### Theory

- This algorithm follows the brute force approach.
- Let an array  $T[1:n]$  is the given base string of length  $n$  and an array  $P[1:m]$  is the pattern of length  $m$  that is to be searched for in  $T$ .
- The algorithm slides the pattern  $P$  across the base string  $T$  looking for a match.
- Let  $s$  is the number of times  $P$  has been shifted when the algorithm find  $P$  in  $T$ , it outputs  $s$  by indicating that "Pattern occurs with shift  $s$ ".
- This algorithm does not give an optimal solution since it ignores the information gained about the base string  $T$  for one value of  $s$  while considering the other values of  $s$ .
- We can find substring by checking once for the string. It does not occupy extra space to perform operations.

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### Algorithm :

- The naive algorithm finds all valid shifts using a loop that checks the condition  $P[1 \dots m] = T[s+1 \dots s+m]$  for each of the  $n-m+1$  possible values of  $s$

Naive-String-Matcher ( $T, P$ )

$n = T.length$

$m = P.length$

for  $s = 0$  to  $n-m$

if  $P[1 \dots m] = T[s+1 \dots s+m]$

print "Pattern occurs with shift  $s$ "

### Analysis :

- The time complexity of the naive string matching algorithm is  $O((n-m+1)m)$ , where  $n = |T|$  = the length of a string being searched &  $m = |P|$  = the length of a pattern substring being compared.

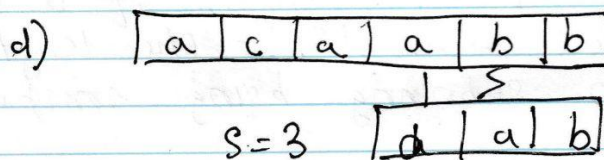
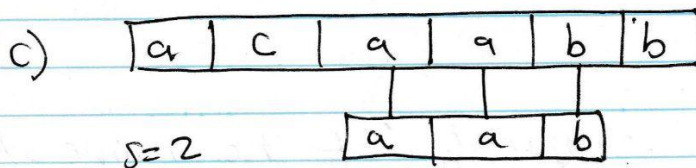
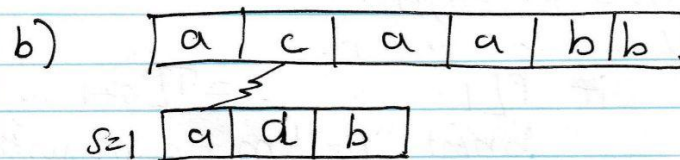
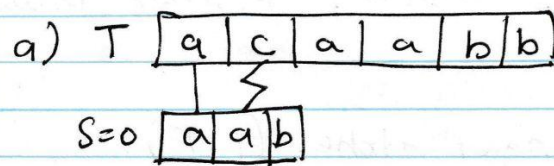
### Example

- Consider a string matching problem instance:  $T = \text{acabb}$  &  $P = \text{aab}$ . The working of the naive string matching algorithm is as shown
- The vertical straight line connects the



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matching character & a jagged line connects the first mismatching character if any. In this example, the algorithm finds the matching pattern with shift  $s=2$



Conclusion:

- The implementation of Naïve string matching algorithm with the analysis was successful.