Type-Targeted Testing

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```
data Tree
    = Leaf
    | Node Int Tree Tree

insert :: Int -> Tree -> Tree
delete :: Int -> Tree -> Tree

Did | get it "right"?
```

An Oracle for Testing

QuickCheck

QuickCheck

How is this possible? Valid trees are a **sparse subset** of all trees!

QuickCheck With Statistics

QuickCheck With Statistics

73% of trees were **empty** and 21% only had one element

QuickCheck: Non-Trivial Trees

QuickCheck: Non-Trivial Trees

Less than 10% of generated trees were valid

Custom Generators

```
newtype BST = BST Tree
instance Arbitrary BST where
arbitrary = ...
prop_insert_bst x (BST t) = ...
```

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```

Must repeat for any further restrictions on input domain!

SmallCheck

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Only 133 input trees were valid!

SmallCheck: How Small?

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Exponential blowup in input space confines search to **very small** inputs

How can we **systematically** generate only **valid** inputs?

Target

Generates tests from refinement types via query-decode-check loop

- 1. Translate input types into SMT query
- 2. **Decode** SMT model into concrete values
- 3. Run function and **check** that result inhabits output type

Refinement Types

```
{v:t | p}
```

The set of values v of type t satisfying a predicate p

Refinement Types

```
type Nat = {v:Int | 0 <= v}
type Pos = {v:Int | 0 < v}
type Rng N = {v:Int | 0 <= v && v < N}</pre>
```

The natural numbers, positive integers, and integers in a range

Refinement Types

$$x:Nat -> \{v:Nat | v = x + 1\}$$

Functions that take a natural number and increment it by one

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Step 1: Query

```
type Nat = {v:Int | 0 <= v}
type Rng N = {v:Int | 0 <= v && v < N}
rescale :: r1:Nat -> r2:Nat -> s:Rng r1 -> Rng r2
```

Step 1: Query

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Represent preconditions directly in logic

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rescale :: r1:Nat -> r2:Nat -> s:Rng r1 -> Rng r2
```

Represent preconditions directly in logic

$$C_0 \doteq 0 \leq r_1 \wedge 0 \leq r_2 \wedge 0 \leq s < r_1$$

Step 2: Decode

```
type Nat = {v:Int | 0 <= v}
type Rng N = {v:Int | 0 <= v && v < N}
rescale :: r1:Nat -> r2:Nat -> s:Rng r1 -> Rng r2
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Represent preconditions directly in logic

$$C_0 \doteq 0 \leq r_1 \wedge 0 \leq r_2 \wedge 0 \leq s < r_1$$

A model $[r_1 \mapsto 1, r_2 \mapsto 1, s \mapsto 0]$ maps to a concrete test case

```
>>> rescale 1 1 0
```

```
type Nat = \{v:Int \mid 0 \le v\}
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Force new test by adding refutation constraint $\neg(r_1 = 1 \land r_2 = 1 \land s = 0)$

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rescale :: r1:Nat -> r2:Nat -> s:Rng r1 -> Rng r2
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Represent preconditions directly in logic, excluding 1st test

$$C_1 \doteq 0 \le r_1 \land 0 \le r_2 \land 0 \le s < r_1 \land \neg (r_1 = 1 \land r_2 = 1 \land s = 0)$$

```
type Nat = {v:Int | 0 <= v} type Rng N = {v:Int | 0 <= v && v < N} rescale :: r1:Nat -> r2:Nat -> s:Rng r1 -> Rng r2 Represent preconditions directly in logic, excluding 1st test C_1 \doteq 0 \leq r_1 \wedge 0 \leq r_2 \wedge 0 \leq s < r_1 \wedge \neg (r_1 = 1 \wedge r_2 = 1 \wedge s = 0) A model [r_1 \mapsto 1, r_2 \mapsto 0, s \mapsto 0] maps to a concrete test case >>> rescale 1 0 0
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After substituting \mathbf{v} and \mathbf{r2}: 0 \le 0 \land 0 < 0 INVALID
rescale 1 0 0 fails the postcondition check!
```

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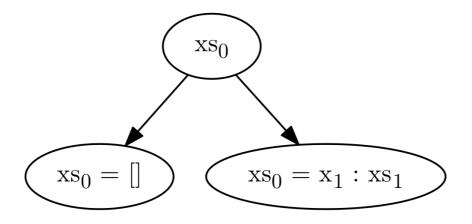
How should we handle **structured data**?

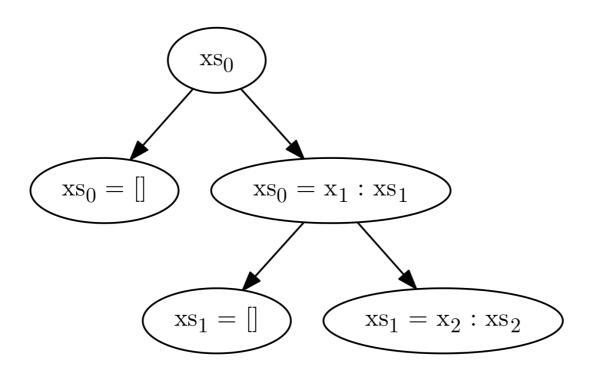
Containers

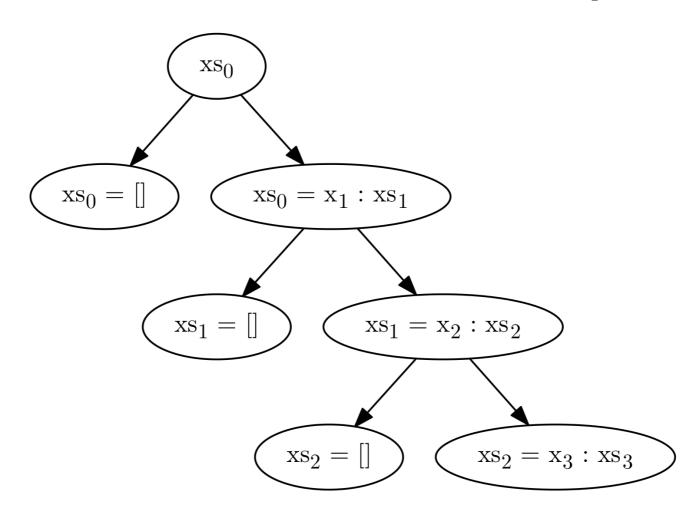
```
type Weight = Pos
type Score = Rng 100
average :: [(Weight, Score)] -> Score
```

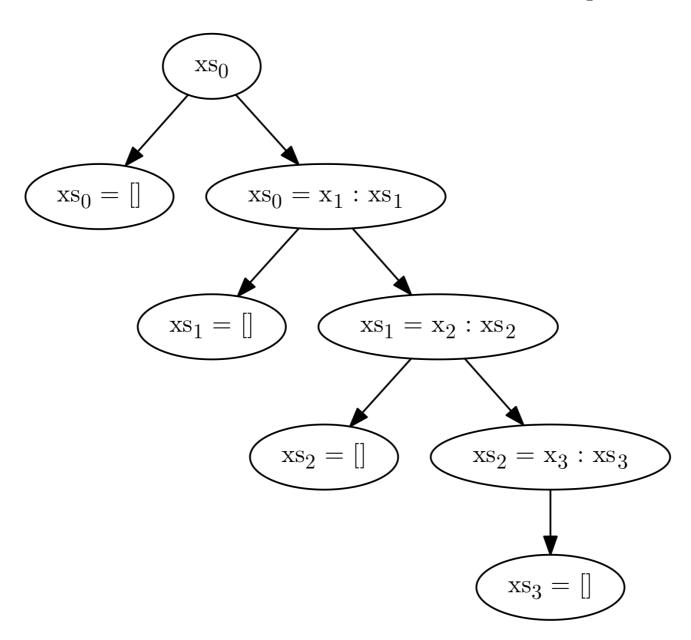
How to generate lists via SMT solver?











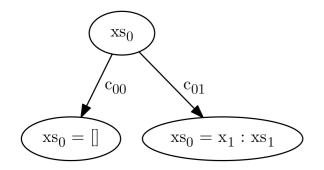
A single set of constraints describes all possible inputs



Choice variables *c* **guard** other constraints

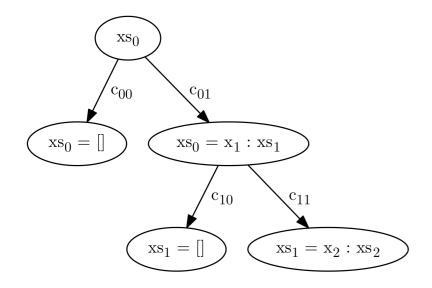
 $C_{\mathsf{list}} \doteq$

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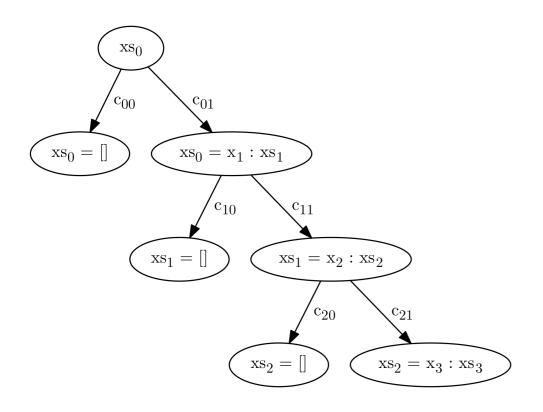
$$\mathsf{C}_{\mathsf{list}} \doteq (c_{00} \Rightarrow xs_0 = []) \land (c_{01} \Rightarrow xs_0 = x_1 : xs_1) \land (c_{00} \oplus c_{01})$$

A single set of constraints describes all possible inputs



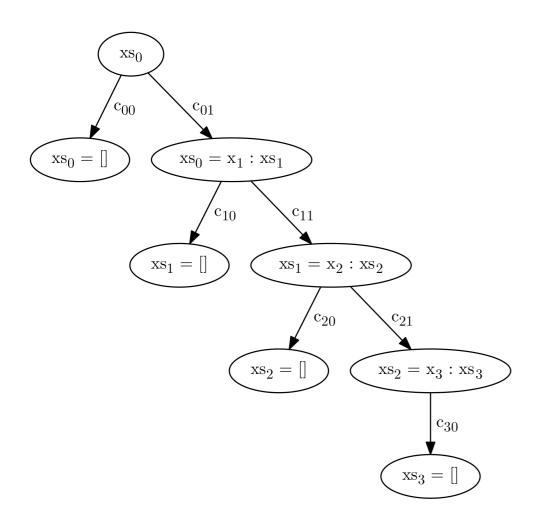
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$$\land (c_{10} \Rightarrow xs_1 = []) \land (c_{11} \Rightarrow xs_1 = x_2 : xs_2) \land (c_{01} \Rightarrow c_{10} \oplus c_{11})$$

A single set of constraints describes all possible inputs



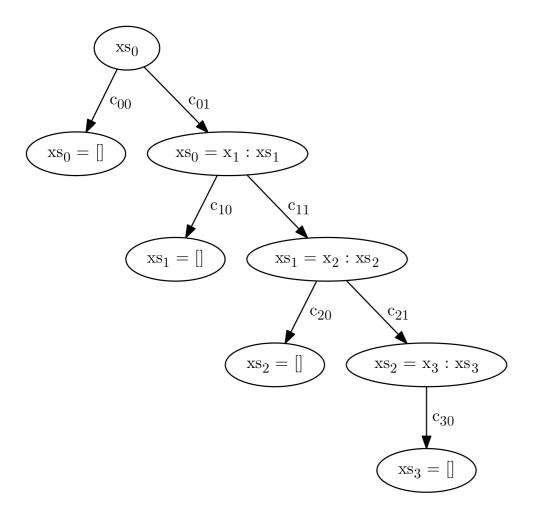
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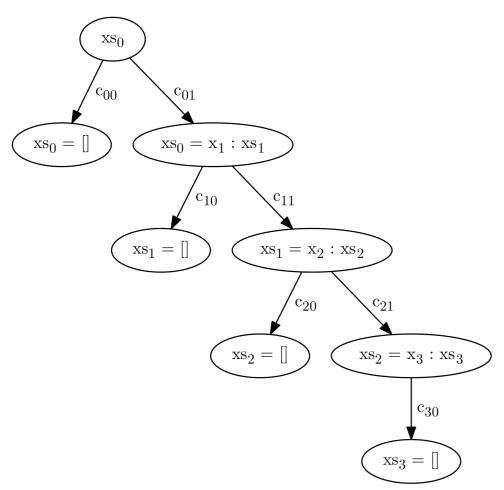
Choice variables c guard other constraints

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$$\begin{aligned} \mathsf{C}_{\mathsf{data}} &\doteq (c_{01} \Rightarrow x_1 = (w_1, s_1) \ \land \ 0 < w_1 \ \land \ 0 \leq s_1 < 100) \\ &\land (c_{11} \Rightarrow x_2 = (w_2, s_2) \ \land \ 0 < w_2 \ \land \ 0 \leq s_2 < 100) \\ &\land (c_{21} \Rightarrow x_3 = (w_3, s_3) \ \land \ 0 < w_3 \ \land \ 0 \leq s_3 < 100) \end{aligned}$$

Full constraint $C \doteq C_{list} \wedge C_{data}$

A single set of constraints describes all possible inputs

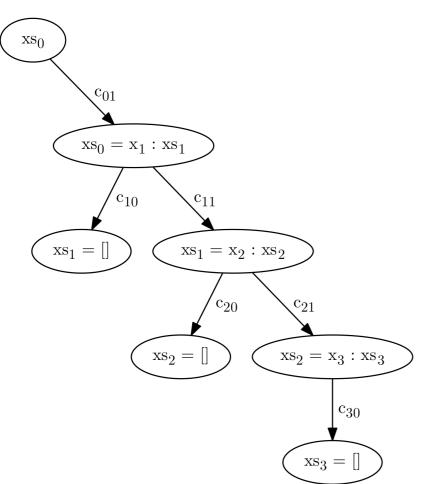


Follow choice variables to reconstruct the list

$$[c_{00} \mapsto \mathsf{false}, \ c_{01} \mapsto \mathsf{true}, \ x_1 \mapsto (w_1, s_1), \ w_1 \mapsto 1, \ s_1 \mapsto 2,$$

$$c_{10} \mapsto \mathsf{true}, \ c_{11} \mapsto \mathsf{false}, \ x_2 \mapsto (w_2, s_2), \ w_2 \mapsto 3, \ s_2 \mapsto 4, \dots]$$

A single set of constraints describes all possible inputs



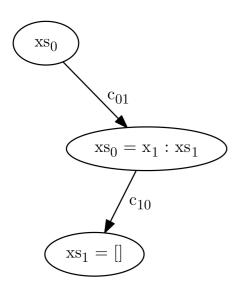
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$$c_{01} \mapsto \mathsf{true} \Rightarrow xs_0 = x_1 : xs_1$$

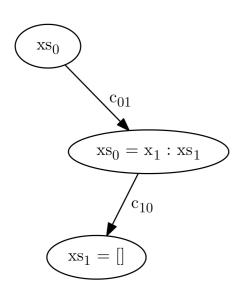
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A single set of constraints describes all possible inputs

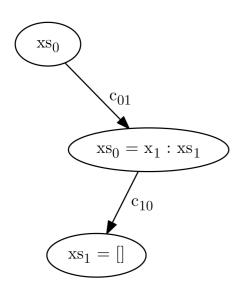


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Realized value: [(1,2)]

A single set of constraints describes all possible inputs



Follow choice variables to reconstruct the list

$$\begin{array}{l} [\ c_{00} \mapsto \ \mathsf{false},\ c_{01} \mapsto \ \mathsf{true},\ x_1 \mapsto (w_1,s_1),\ w_1 \mapsto 1,\ s_1 \mapsto 2, \\ c_{10} \mapsto \ \mathsf{true},\ c_{11} \mapsto \ \mathsf{false},\ x_2 \mapsto (w_2,s_2),\ w_2 \mapsto 3,\ s_2 \mapsto 4,\dots \,] \\ \\ c_{01} \mapsto \ \mathsf{true} \Rightarrow xs_0 = x_1 : xs_1 \\ c_{10} \mapsto \ \mathsf{true} \Rightarrow xs_1 = []$$

Realized value: [(1,2)]

Only refute constraints that contribute to realized value

$$\neg(c_{00} = \mathsf{false} \land c_{01} = \mathsf{true} \land x_1 = (w_1, s_1) \land w_1 = 1 \land s_1 = 2 \land c_{10} = \mathsf{true})$$

```
best :: k:Nat -> {xs:[Score] | k <= len xs}
    -> {v:[Score] | k = len v}
```

best takes a list of **at least k** scores, and returns a list with **exactly k** scores

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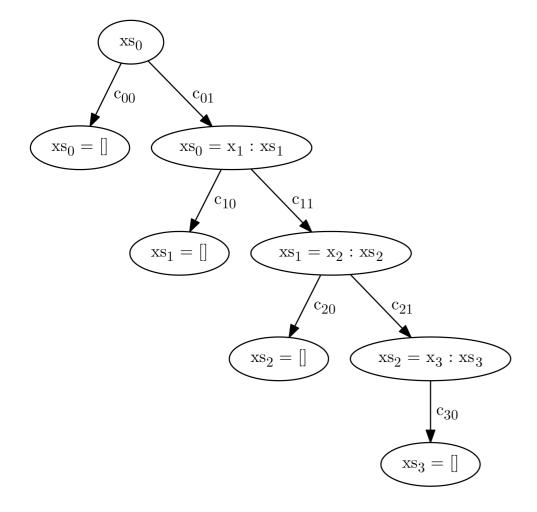
best takes a list of at least k scores, and returns a list with exactly k scores

```
measure len :: [a] -> Nat
len [] = 0
len (x:xs) = 1 + len xs
```

len is a **logical function** that describes the length of a list.

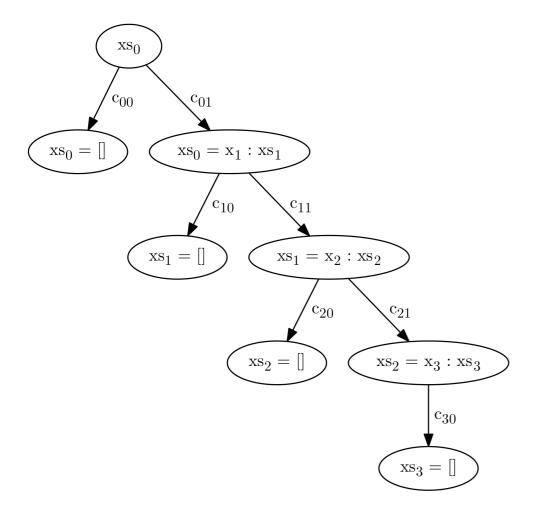
We instantiate measure definition each time we unfold [] or (:)

```
best :: k:Nat -> {xs:[Score] | k <= len xs}
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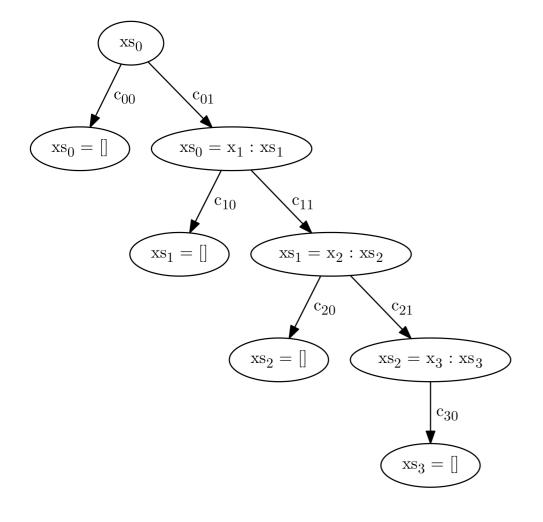
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$$\begin{aligned} \mathsf{C}_{\mathsf{size}} &\doteq (c_{00} \Rightarrow \mathsf{len} \ xs_0 = 0) \land (c_{01} \Rightarrow \mathsf{len} \ xs_0 = 1 + \mathsf{len} \ xs_1) \\ &\land (c_{10} \Rightarrow \mathsf{len} \ xs_1 = 0) \land (c_{11} \Rightarrow \mathsf{len} \ xs_1 = 1 + \mathsf{len} \ xs_2) \\ &\land (c_{20} \Rightarrow \mathsf{len} \ xs_2 = 0) \land (c_{21} \Rightarrow \mathsf{len} \ xs_2 = 1 + \mathsf{len} \ xs_3) \\ &\land (c_{30} \Rightarrow \mathsf{len} \ xs_3 = 0) \end{aligned}$$

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Enforce relation between **k** and **xs** by adding $k \leq \text{len } xs_0$

$$\mathsf{C} \doteq \ \mathsf{C}_{\mathsf{list}} \land \mathsf{C}_{\mathsf{data}} \land \mathsf{C}_{\mathsf{size}} \land 0 \leq k \leq \mathsf{len} \ xs_0$$

Demo: Targeting **BST**s

Evaluation

Our Claims

- 1. Target handles highly structured inputs automatically
- 2. Target generates tests that provide high code coverage

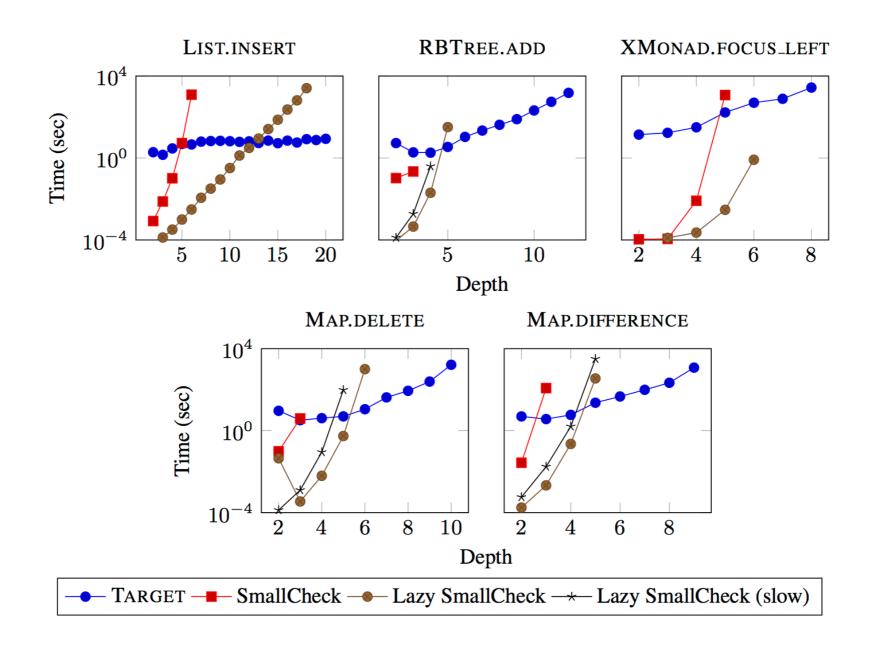
Benchmarks

- 1. Data.Map: checked balancing and ordering invariants
- 2. **RBTree**: checked red-black and ordering invariants
- 3. XMonad.StackSet: checked uniqueness of windows

Compared Target against QuickCheck and SmallCheck

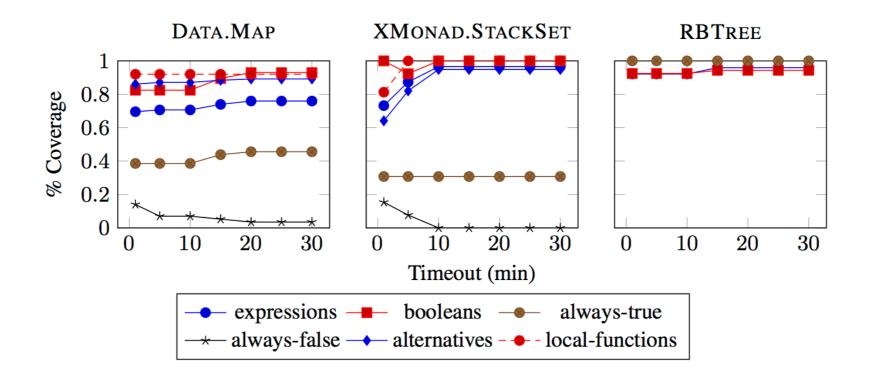
Evaluation: Results

Target checks larger inputs than brute-force



Evaluation: Results

Target provides high coverage with low investment



Takeaway

Target - a new approach for automatically testing functions with preconditions.

- guarantees inputs satisfy preconditions
- vs QuickCheck: does not require custom generators
- vs SmallCheck: defers the onset of input explosion