# CS 354 - Machine Organization & Programming Thursday, February 22, 2018

Project p2 (6%): DUE at 10 pm THIS SUNDAY, February 25th

Note: It is better to submit a working program using indexing than a non-working program attempting to use address arithmetic.

Project p3 (6%): Assigned Tomorrow

Homework hw3 (1.5%): Assigned Tomorrow

## **Last Time**

C's Heap Allocator (stdlib.h)
Posix brk (unistd.h)
Allocator Design
Simple View of Heap
Free Block Organization
Implicit Free List

# **Today**

Implicit Free List (from last time)
Placement Policies
Free Block - Too Much/Too Little
Coalescing Free Blocks
Footers

#### **Next Time**

Heap Caveats

**Read:** B&O 9.9.11, 9.9.13

**Skim:** B&O 9.9.12

# **Placement Policies**

### **Placement Policies**

 First Fit (FF): start search from stop at

fail if

mem util:

thruput:

Next Fit (NF): start search from

stop at fail if

mem util:

thruput:

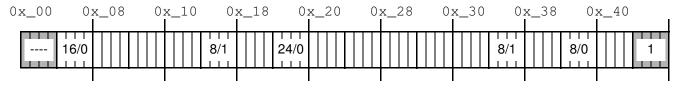
Best Fit (BF): start search from

stop at fail if

mem util:

thruput:

# **Heap Allocation Run 3** using a Placement Policy:



→ Given the original heap above and the placement policy, what address is ptr assigned?

→ Given the original heap above and the address of block most recently allocated, what address is ptr assigned using NF?

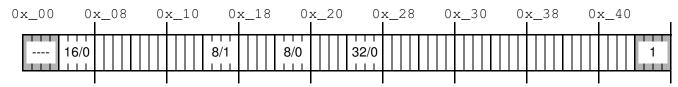
# Free Block - Too Much/Too Little

What happer	ns if free b	lock chosen	is bigger tha	n the request?

•

•

## Heap Allocation Run 4 using Splitting and using FF



→ Given the heap to be modified by the 4 mallocs below, what address is assigned to each pointer?

If there is a new free block, what is its address and size in bytes?

```
1) ptr1 = malloc(sizeof(char));
2) ptr2 = malloc(11 * sizeof(char));
3) ptr3 = malloc(2 * sizeof(int));
4) ptr4 = malloc(5 * sizeof(int));
```

What happens if no free block is large enough to satisfy the request?

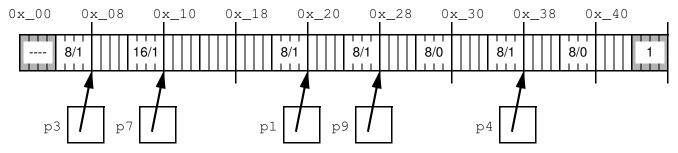
1st.

2nd.

3rd.

# **Coalescing Free Blocks**

# **Heap Allocation Run 5**



→ What's the problem resulting from the following heap operations using FF?

```
1) free(p9); p9 = NULL;
2) free(p1); p1 = NULL;
3) p1 = malloc(4 * sizeof(int));
```

## **Problem:**

## Solution:

## immediate

## <u>delayed</u>

→ Given the original heap above, what is the size in bytes of the freed heap block?

```
1) free(p7); p7 = NULL;
```

- → Given a pointer to a payload, how do you find its block header?
- → Given a pointer to a payload, how do you find the block header of the next block?
- → Given the modified heap above, what is the size in bytes of the freed heap block when immediate coalescing is used?

```
2) free(p3); p3 = NULL;
3) free(p1); p1 = NULL;
```

> Given a pointer to a payload, how do you find the block header of the previous block?

# **Footers**

# Heap Free Block Layout with Header and Footer

31 3|210 bits Footer (AKA Boundary Tag)

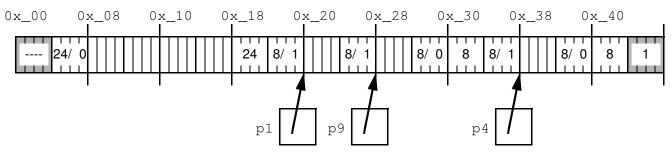
Header = block size |0pa

IPossibly More Free WordsI

Footer = only block size

- → Why don't allocated blocks need footers?
- → Given a pointer to a payload, how do you get to the header of a previous block that's free?

Heap Allocation Run 6 with given Free List using Immediate Coalescing and Free Block Footers



- → Given the heap above, what is the size in bytes of the freed heap block?
  - 1) free(p1);
- → Given the modified heap above, what is the size in bytes of the freed heap block?
  - 2) free(p4);
- ➤ Is coalescing done in a fixed number of steps (constant time) or is it dependent on the number of blocks (linear time)?