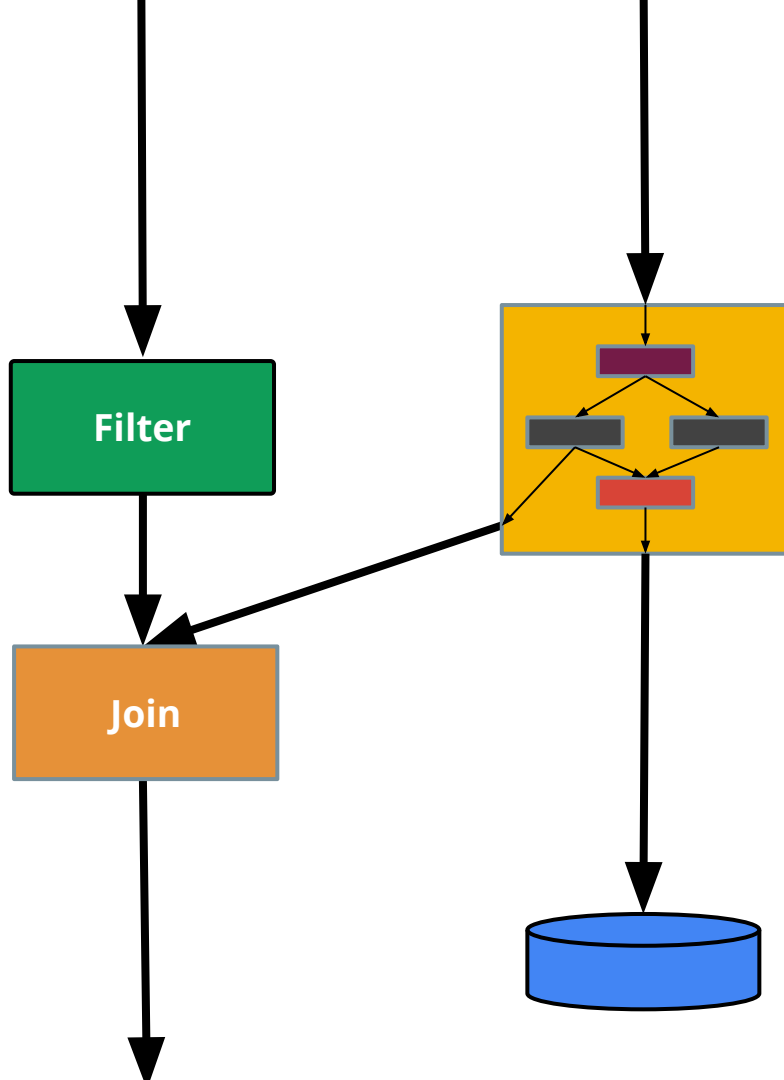




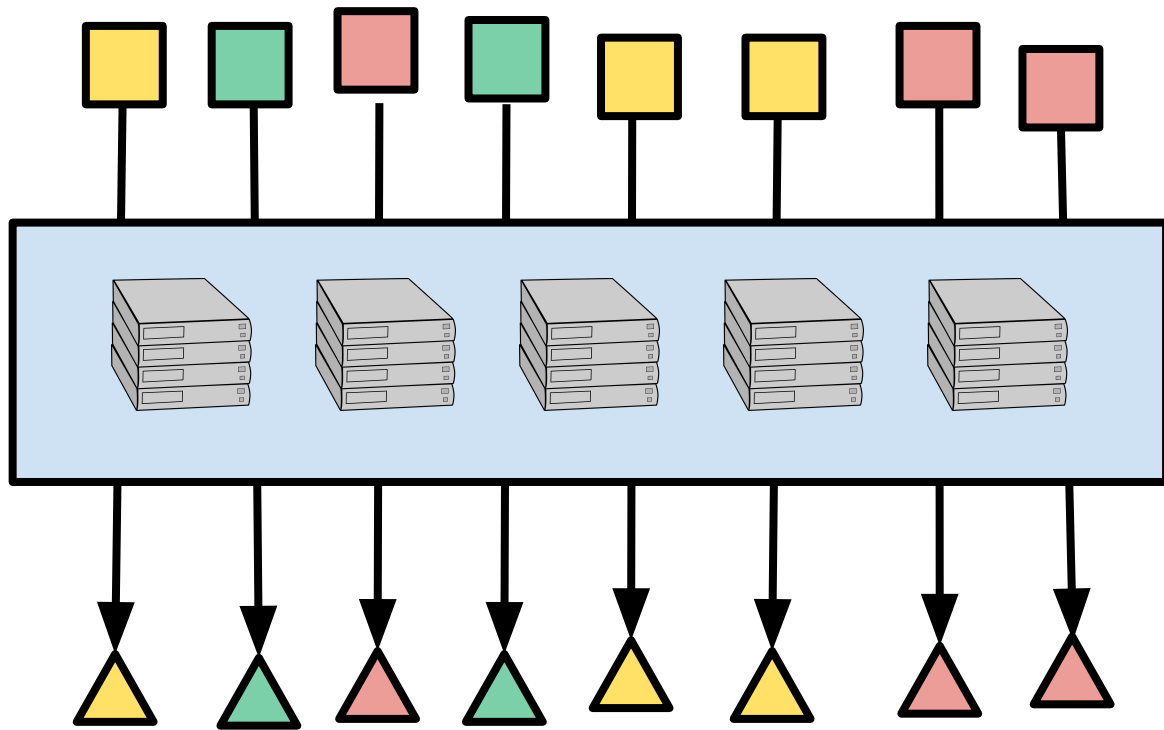
State & Timers Design Patterns

Reza Rokni
Kenn Knowles





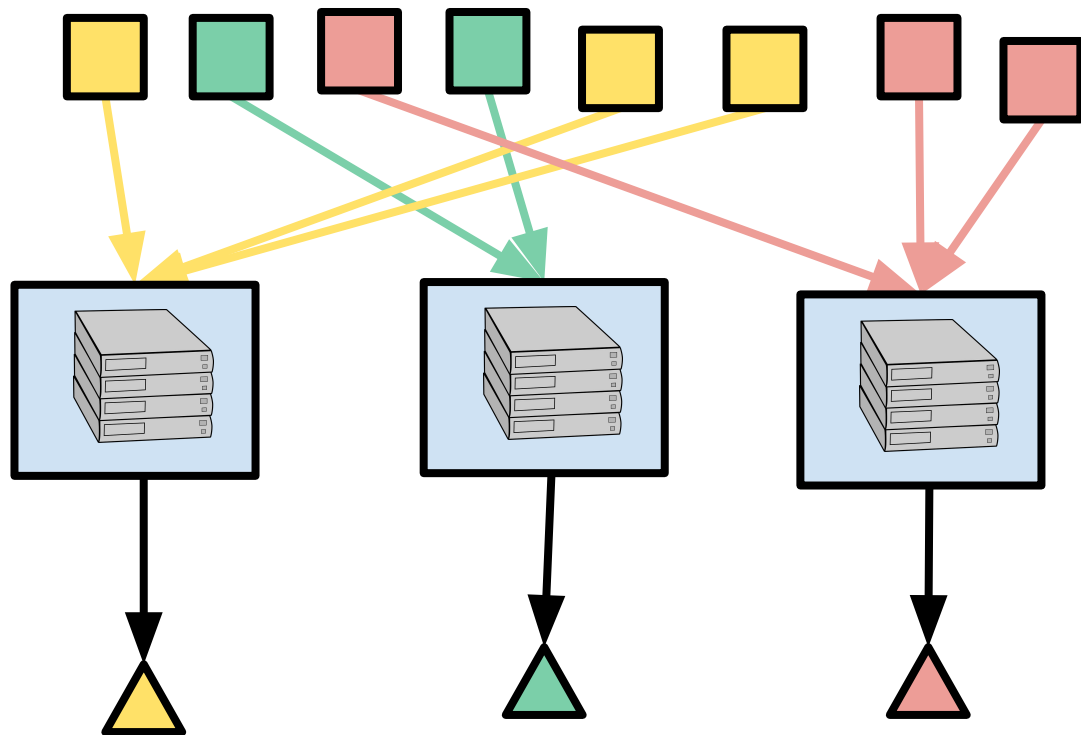
Per element: ParDo (Map, etc)



Every item
processed
independently

Stateless
implementation

Per key: Combine (Reduce, etc)

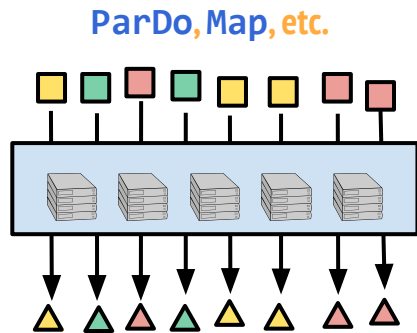


Items grouped by some key and combined

Stateful streaming implementation (buffering until trigger)

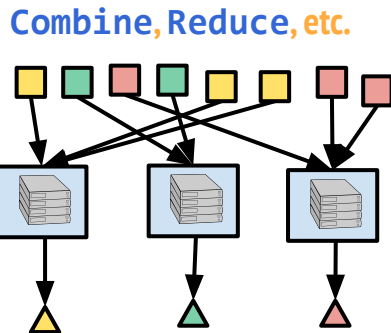
But your code doesn't work with state, just associative & commutative function

It "just works" with massive out-of-order streams



"Parse incoming events
and filter out bad data"

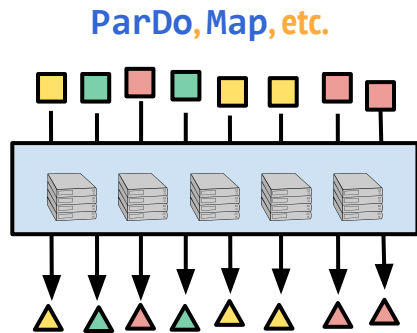
"Put events in 10 minute windows
sliding every 2 minutes"



"Sum per hour and output when
you have the whole hour"

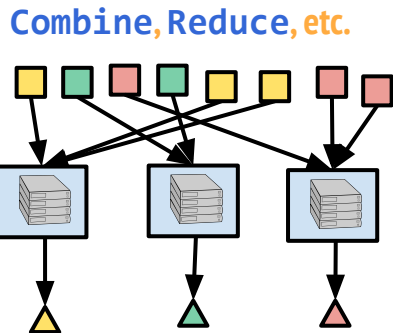
"Group into sessions and
emit as fast as possible"

But what if you need more control?



"I'm getting tied in knots trying to do this with windows & triggers"

"Triggers aren't specific enough for my use case"

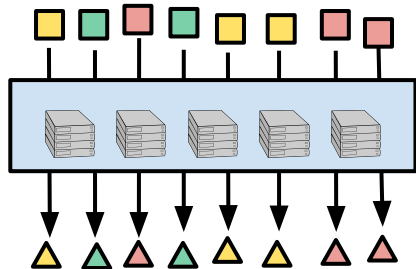


"My aggregation is not an associative & commutative operator"

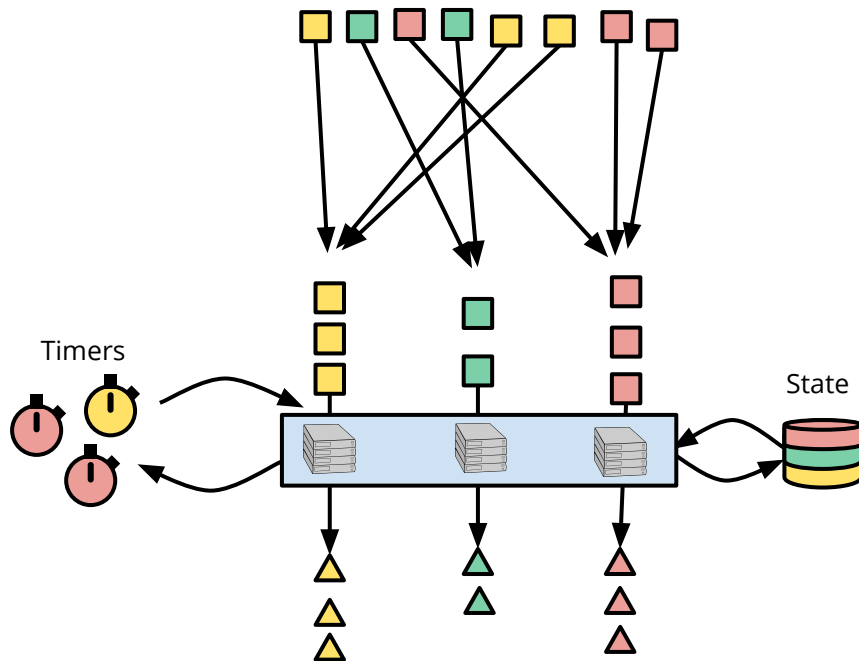
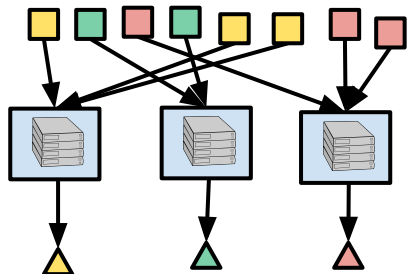
"I need to output even when data isn't coming in"

Beam primitive: state & timers

ParDo, Map, etc.



Combine, Reduce, etc.



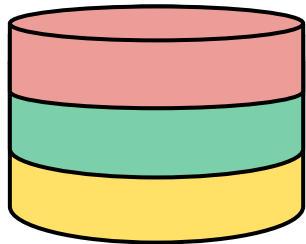
What it looks like (in Java)

```
new DoFn<KV<String, Foo>, Baz>() {  
  
    @StateId("counter")  
    private final StateSpec<Object, ValueState<Integer>> indexSpec = StateSpecs.value();  
  
    @ProcessElement  
    public void processElement(@StateId("counter") ValueState<Integer> counter, ...) {  
        int current = firstNonNull(counter.read(), 0);  
        counter.write(current+1);  
        ...  
    }  
}
```


What it looks like (in Java)

```
new DoFn<KV<String, Foo>, Baz> {  
  
    @TimerId("flush-timer")  
    private TimerSpec flushSpec = TimerSpecs.timer(TimeDomain.PROCESSING_TIME);  
  
    @ProcessElement  
    public void process(@TimerId("flush-timer") Timer flushTimer, ...) {  
        flushTimer.setForNowPlus(...); ...  
    }  
  
    @OnTimer("flushTimer")  
    public void onTimer(...) {  
        ... // read state, output values  
    }  
}
```

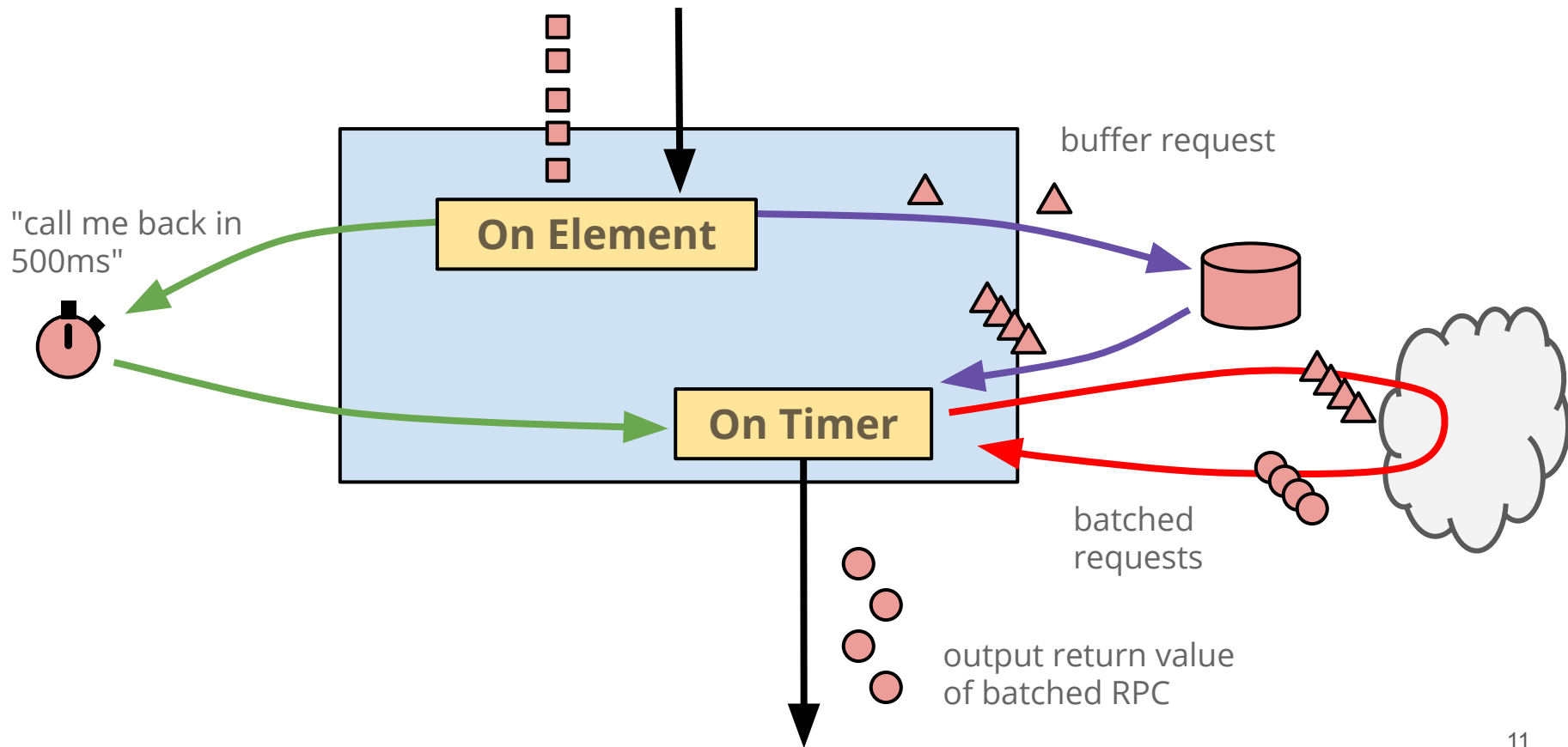
State is per key and window



Key	Window	MEDIAN_IDLE	MAIN_ACTIVITY	...
"kenn"	9am - 10am	10m	"hack"	
	12pm - 1pm	25m	"eat"	
	11pm - 12am	60m	"sleep"	
"reza"	8am - 9am	20m	"flight"	
	11am - 12pm	3m	"hack"	
...	...			

Bonus: automatically garbage collected when a window expires
(vs manual clearing of per-key state)

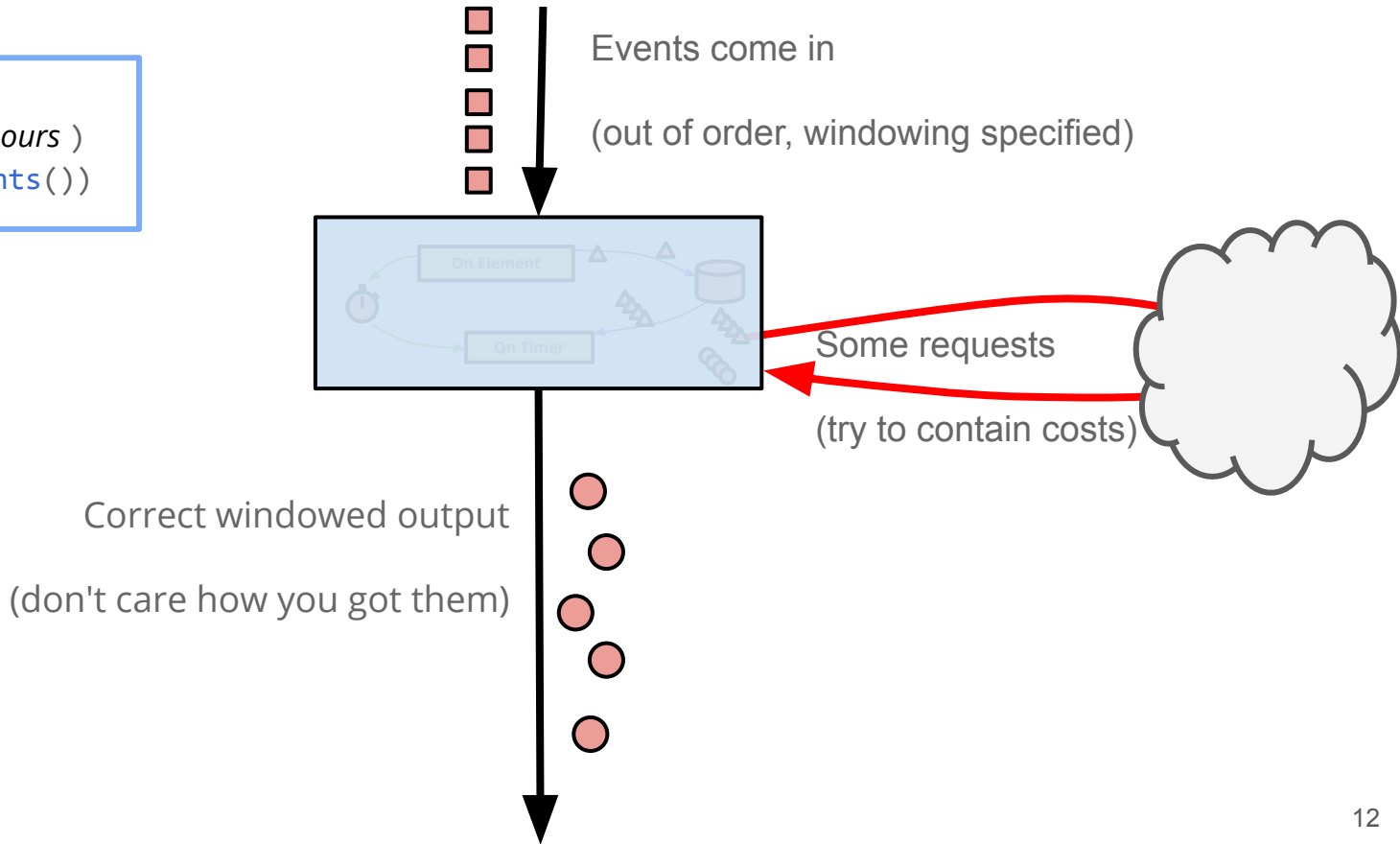
Example: time-batched requests



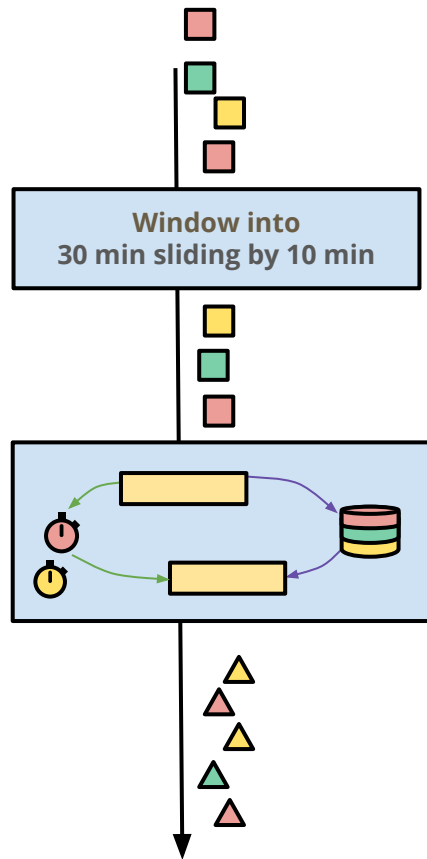
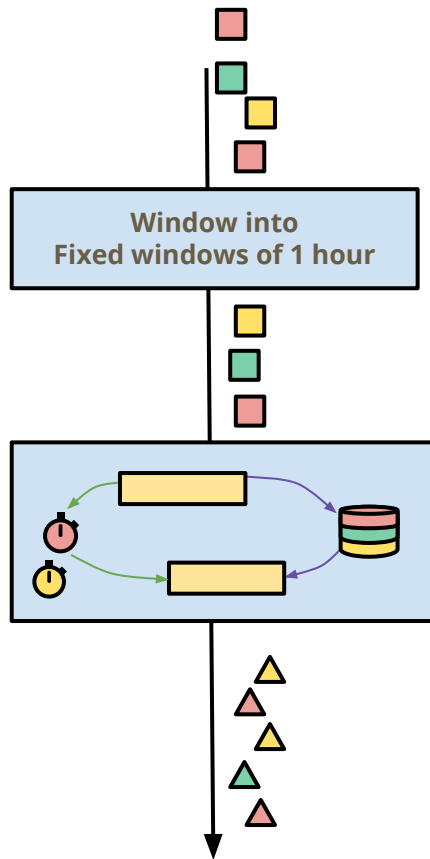
User's view of your transform

input

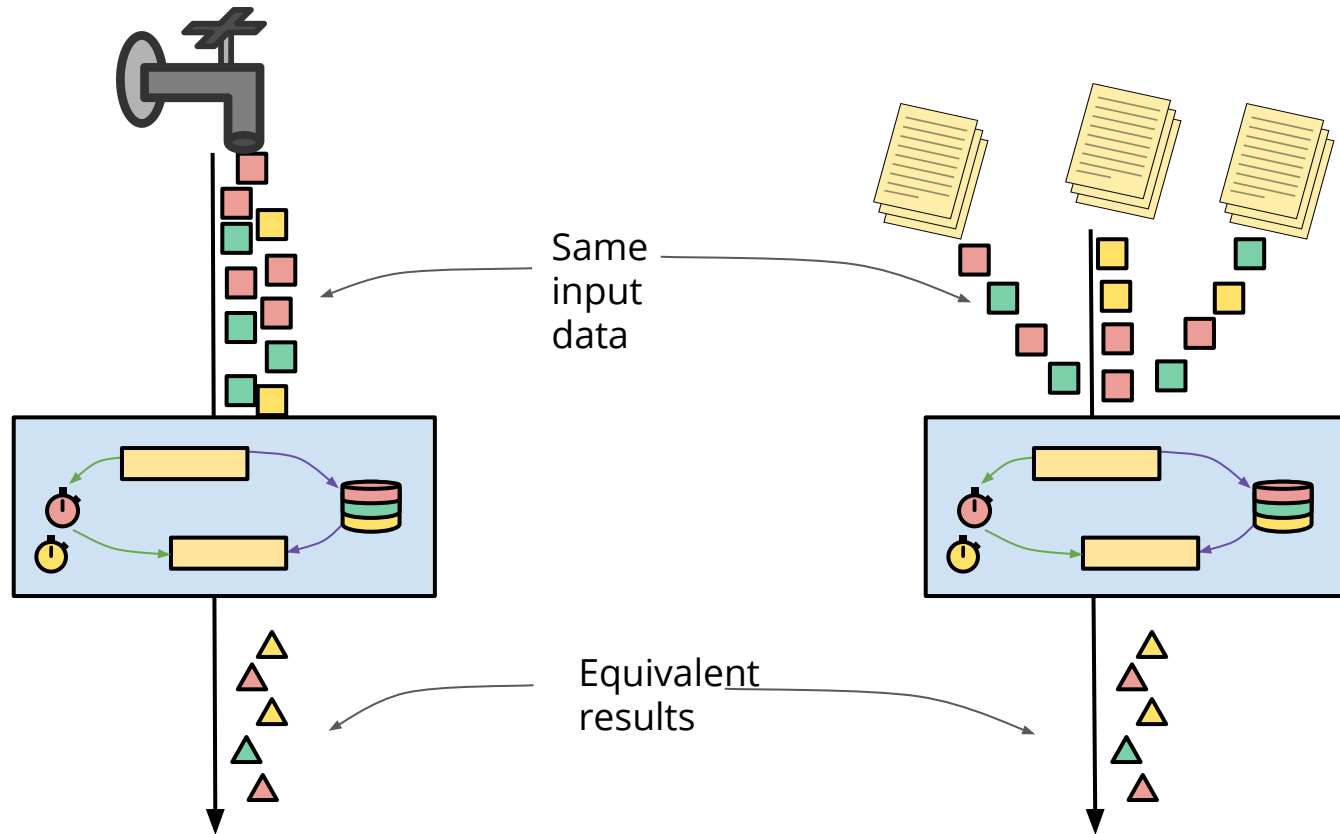
```
.apply(Window.into( hours )  
.apply(new EnrichEvents()))
```



Event time windowing still "just works"



Unified present & historical processing



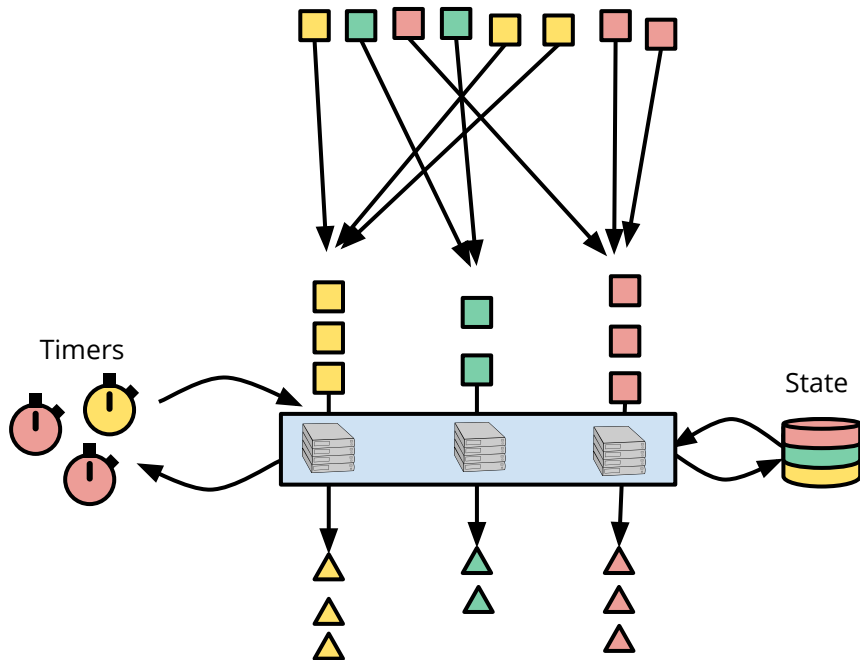
Types of state

Type	Strength	Dataflow Runner
Value	Read/write any value (but always the whole value)	Yes
Bag	Cheap append No ordering on read	Yes
Combining	Associative/commutative compaction	Yes
Not yet supported in Dataflow runner	Membership checking	No
	Map	No

Types of timers

Type	Uses
Processing time	<ul style="list-style-type: none">TimeoutsRelative times ("in 5 minutes")Periodically output based on state
Event time	<ul style="list-style-type: none">Output based on completeness of input dataAbsolute times ("when the data is complete up to 5:00am")Final/authoritative outputsDon't leave data behind in state!

Isn't this just "the only" primitive? (no)



You: express more complex logic

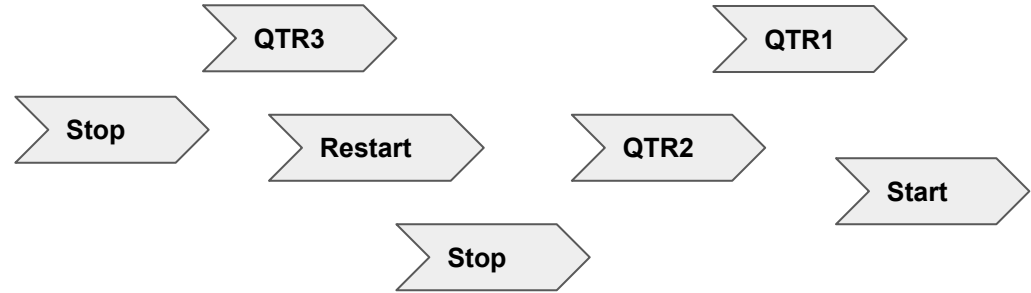
Runner: fewer optimization/execution choices

Most of your pipeline should still be the simpler stuff

What else can you do with state & timers

- **Domain-specific triggering** ("output when five people who live in Seattle have checked in")
- **Slowly changing dimensions** ("update FX rates for currency ABC")
- **Stream joins** ("join-matrix" / "join-biclique")
- **Fine-grained aggregation** ("add odd elements to accumulator A and event elements to accumulator B")
- **Per-key workflows** (like user sign up flow w/ reminders & expiration)

Walk through of a non-trivial example!



Events coming from player devices

Warning boilerplate ahead ... look away now...

State and Timers is a raw set of API's

Dependent on use case, lots of boilerplate may be needed;

- Order is not guaranteed..
- The Timer API does not support read() operation
- OnTimer() does not have key context
- Order is not guaranteed..
- DoFn.StartBundle and DoFn.FinishBundle can not access State
- Did we mention Order is not guaranteed?

Laying out the rules...

Implement some rules around events coming from our video application:

The application has the following Events:

```
@DefaultCoder(SchemaCoder.class)
public enum UserEventType {
    START(0),
    QTR_1(1),
    QTR_2(2),
    QTR_3(3),
    END(4),
    PAUSE(5),
    RESTART(6),
    STOP(7);
}
```

Laying out the rules...

Given those events we would like to implement these rules

Rule 0 Output and GC When we have END or no activity for 10 Mins no allowed lateness

Rule 1 Output every time we see a QTR_3

Rule 2 Output when we see STOP but no other activity for 2 Min's

Rule 3 Output when we see multiple RESTART's without any QTR updates during GC

Rule 0 - Setup State Objects

Rule 0 Output and GC When we have END or no activity for 10 Mins (no allowed lateness)

```
@TimerId("ttl")
private final TimerSpec ttl = TimerSpecs.timer(TimeDomain.EVENT_TIME);

@StateId("key")
private final StateSpec<ValueState<String>> key = StateSpecs.value();

@StateId("maximumTimestampObserved")
private final StateSpec<ValueState<Instant>> maximumTimestampObserved = StateSpecs.value();
```

Rule 0 - Setup State Objects

Rule 0 Output and GC When we have END or no activity for 10 Mins no allowed lateness

```
@ProcessElement
public void processElement(
    @Element KV<String, UserEventType> event,
    @TimerId("ttl") Timer ttl,
    @StateId("key") ValueState<String> key,
    @StateId("maximumTimestampObserved") ValueState<Instant> maximumTimestampObserved,
    ...
)
```


Rule 0 - Process Events

Rule 0 Output and GC When we have END or no activity for 10 Mins no allowed lateness

```
@ProcessEvent
...
// Current limitation of OnTimer() is it does not store the key
key.write(event.getKey());

// Set last value processed Do NOT EXPECT ORDERED INPUTS!
if (maximumTimestampObserved.read() == null) {
    lastProcessedEvent.write(timestamp);
} else {
    if (maximumTimestampObserved.read().isBefore(timestamp)) {
        maximumTimestampObserved.write(timestamp);
    }
}

// Rule 0 GC after 10 min's of no activity.. no allowed lateness
ttl.set(maximumTimestampObserved.read().plus(Duration.standardMinutes(10)));
```

Rule 0 - Garbage Collect

Rule 0 Output and GC When we have END or no activity for 10 Mins no allowed lateness

```
@OnTimer("ttl")
public void ttl(
    @StateId("key") ValueState<String> key,
    @StateId("maximumTimestampObserved") ValueState<Instant> maximumTimestampObserved,
    ....
    OutputReceiver<KV<String, SystemEventType>> o) {
    ....

    o.output(KV.of(key.read(), SystemEventType.GC));

    maximumTimestampObserved.clear();
    key.clear();
    ...
}
```

Rule 1 - Simple filter

Rule 1 Output every time we see a QTR_3

```
@ProcessElement
```

```
...
```

```
// Check for Rule 1 : Output when we are at QTR_3
```

```
if (event.getValue().equals(UserEventType.QTR_3)) {  
    o.output(KV.of(event.getKey(), SystemEventType.RULE_1));  
}
```

Rule 2 - Setup State Objects

Rule 2 Output when we see STOP but no other activity for 2 Min's

```
@StateId("key")
private final StateSpec<ValueState<String>> key = StateSpecs.value();

@StateId("maximumTimestampObserved")
private final StateSpec<ValueState<Instant>> maximumTimestampObserved = StateSpecs.value();

@TimerId("rule2")
private final TimerSpec rule2 = TimerSpecs.timer(TimeDomain.EVENT_TIME);
```

Rule 2 - Catch Stop & Set Timer

Rule 2 Output when we see STOP but no other activity for 2 Min's

Note: Timer object does not have reset / delete option

```
// Check for Rule 2 Output when we see STOP but no RESTART for 2 Min's
if (event.getValue().equals(UserEventType.STOP)) {
    rule2.set(timestamp.plus(Duration.standardMinutes(2)));
}
```

Rule 2 - Process Stop

Rule 2 Output when we see STOP but no other activity for 2 Min's

```
@OnTimer("rule2")
public void rule2(
    @StateId("key") ValueState<String> key,
    @StateId("maximumTimestampObserved") ValueState<Instant> maximumTimestampObserved,
    OnTimerContext otc) {
    // If no new events since STOP timer was set output event
    // Assume if an event came in with time == stop time, STOP is later
    if (otc.timestamp().isAfter(maximumTimestampObserved.read())
        && !otc.timestamp().isEqual(maximumTimestampObserved.read()))
        otc.output(KV.of(key.read(), SystemEventType.RULE_2));
    }
}
```

Rule 3 - Setup State Objects

Rule 3 Output when we see multiple RESTART's without any QTR updates, at Rule 0

```
@StateId("key")
private final StateSpec<ValueState<String>> key = StateSpecs.value();

@StateId("events")
private final StateSpec<BagState<TimestampedValue<Integer>>> events = StateSpecs.bag();

@StateId("maximumTimestampObserved")
private final StateSpec<ValueState<Instant>> maximumTimestampObserved = StateSpecs.value();
```

Rule 3 - Process Restart

Rule 3 Output when we see multiple RESTART's without any QTR updates, at Rule 0

We can use Rule 0 GC function rather than setup a new OnTimer()

```
// Check for Rule 3 : Output when we see RESTART more than once in session without a QTR
// update

if (event.getValue().equals(UserEventType.RESTART)) {
    eventsBag.add(TimestampedValue.of(1, timestamp));
}
```


Rule 3 - Test in GC

Rule 3 Output when we see multiple RESTART's without any QTR updates at Rule 0

```
@OnTimer("ttl")
public void ttl(
    @StateId("key") ValueState<String> key,
    @StateId("maximumTimestampObserved") ValueState<Instant> maximumTimestampObserved,
    @StateId("events") BagState<TimestampedValue<Integer>> eventsBag,
    OutputReceiver<KV<String, SystemEventType>> o) {

    // Implement Rule 3
    Instant maxQTRTime = Instant.EPOCH;
    int restartCount = 0;
    for (TimestampedValue<Integer> events : eventsBag.read()) {
        if (events.getValue() == 0) {
            maxQTRTime =
                (maxQTRTime.isAfter(events.getTimestamp())) ? events.getTimestamp() : maxQTRTime;
        }
    }
}
```

Rule 3 - Test in GC

Rule 3 Output when we see multiple RESTART's without any QTR updates at Rule 0

We can use Rule 0 GC function rather than setup a new OnTimer()

```
for (TimestampedValue<Integer> events : eventsBag.read()) {
    if (events.getValue() == 0) {
        maxQTRTime =
            (maxQTRTime.isAfter(events.getTimestamp())) ? events.getTimestamp() : maxQTRTime;}}
for (TimestampedValue<Integer> events : eventsBag.read()) {
    if (events.getValue() == 1) {
        // Is this restart after max QTR time?
        restartCount += (maxQTRTime.isBefore(events.getTimestamp())) ? 1 : 0;}}

if (restartCount > 0) {
    o.output(KV.of(key.read(), SystemEventType.RULE_3));
}
// If there was two or more RESTARTS after last QTR then send out a event
o.output(KV.of(key.read(), SystemEventType.GC));
```

Summary

- Use State and Timers when you need more control
- State is per key per window
- State and Timer Types

Thank you!

Questions?



Solving for Time Series

Time Series

A **time series** is a series of data points indexed (or listed or graphed) in time order. Most commonly, a time series is a sequence taken at successive equally spaced points in time.

https://en.wikipedia.org/wiki/Time_series

Timeseries

A **time series** is a series of data points indexed (or listed or graphed) in time order. Most commonly, a time series is a sequence taken at successive equally spaced points in time.

https://en.wikipedia.org/wiki/Time_series

Words Order / Sequence...
Not easy to deal with ...

Timeseries

A **time series** is a series of data points indexed (or listed or graphed) in time order. Most commonly, a time series is a sequence taken at successive equally spaced points in time.

https://en.wikipedia.org/wiki/Time_series

What happens when we need to count when there is no data.

Gap filling

	TS-1	TS-2	TS-3	TS-4	TS-5
$t_0 \rightarrow t_1$	●	●	●	□	●
$t_1 \rightarrow t_2$	●	●	□	□	●
$t_2 \rightarrow t_3$	●	□	●	□	□
$t_3 \rightarrow t_4$	●	●	□	□	□
$t_4 \rightarrow t_5$	●	●	●	●	□

● Data Points □ No Data

Gap filling

	TS-1	TS-2	TS-3	TS-4	TS-5
$t_0 \rightarrow t_1$	●	●	●	●	●
$t_1 \rightarrow t_2$	●	●	●	●	●
$t_2 \rightarrow t_3$	●	●	●	●	●
$t_3 \rightarrow t_4$	●	●	●	●	●
$t_4 \rightarrow t_5$	●	●	●	●	●

● Data Point ● Generated Data

Options

- Option 1 : Create a heartbeat message external to Pipeline
 - Bounded / UnBounded source creating an impulse at every time interval.
- Option 2 : Create a heartbeat message internal to Pipeline
 - Use of impulses from utility classes, in java Generate Sequence.
- Option 3 : Looping timers
 - Make use of self setting Timers using the Apache Beam Timer API

Option 1 External HeartBeat

$t_0 \rightarrow t_1$	●
$t_1 \rightarrow t_2$	●
$t_2 \rightarrow t_3$	●
$t_3 \rightarrow t_4$	●
$t_4 \rightarrow t_5$	●

External heartbeats

Timeseries Source

$t_0 \rightarrow t_1$	●	●	●	□	●
$t_1 \rightarrow t_2$	●	●	□	□	●
$t_2 \rightarrow t_3$	●	□	●	□	□
$t_3 \rightarrow t_4$	●	●	□	□	□
$t_4 \rightarrow t_5$	●	●	●	●	□

Flatten

$t_0 \rightarrow t_1$	●	●	●	●	●
$t_1 \rightarrow t_2$	●	●	●	●	●
$t_2 \rightarrow t_3$	●	●	●	●	●
$t_3 \rightarrow t_4$	●	●	●	●	●
$t_4 \rightarrow t_5$	●	●	●	●	●

Option 1 External HeartBeat - Needs Fanout

$t_0 \rightarrow t_1$	●
$t_1 \rightarrow t_2$	●
$t_2 \rightarrow t_3$	●
$t_3 \rightarrow t_4$	●
$t_4 \rightarrow t_5$	●

External heartbeats

Fanout 1 -> k1,k2,k3

Timeseries Source

$t_0 \rightarrow t_1$	●	●	●	□	●
$t_1 \rightarrow t_2$	●	●	□	□	●
$t_2 \rightarrow t_3$	●	□	●	□	□
$t_3 \rightarrow t_4$	●	●	□	□	□
$t_4 \rightarrow t_5$	●	●	●	●	□

Flatten

$t_0 \rightarrow t_1$	●	●	●	●	●
$t_1 \rightarrow t_2$	●	●	●	●	●
$t_2 \rightarrow t_3$	●	●	●	●	●
$t_3 \rightarrow t_4$	●	●	●	●	●
$t_4 \rightarrow t_5$	●	●	●	●	●

Option 1 External HeartBeat - Changing keys

$t_0 \rightarrow t_1$	●
$t_1 \rightarrow t_2$	●
$t_2 \rightarrow t_3$	●
$t_3 \rightarrow t_4$	●
$t_4 \rightarrow t_5$	●

External heartbeats

Dynamic Key List

Fanout 1 -> k1,k2,k3

Timeseries Source

$t_0 \rightarrow t_1$	●	●	●	□	●
$t_1 \rightarrow t_2$	●	●	□	□	●
$t_2 \rightarrow t_3$	●	□	●	□	□
$t_3 \rightarrow t_4$	●	●	□	□	□
$t_4 \rightarrow t_5$	●	●	●	●	□

Flatten

$t_0 \rightarrow t_1$	●	●	●	●	●
$t_1 \rightarrow t_2$	●	●	●	●	●
$t_2 \rightarrow t_3$	●	●	●	●	●
$t_3 \rightarrow t_4$	●	●	●	●	●
$t_4 \rightarrow t_5$	●	●	●	●	●

Option 1 External HeartBeat - Changing keys

$t_0 \rightarrow t_1$	●
$t_1 \rightarrow t_2$	●
$t_2 \rightarrow t_3$	●
$t_3 \rightarrow t_4$	●
$t_4 \rightarrow t_5$	●

External heartbeats

Dynamic Key List

Fanout 1 -> k1,k2,k3

Flatten

$t_0 \rightarrow t_1$	●	●	●	●	●
$t_1 \rightarrow t_2$	●	●	●	●	●
$t_2 \rightarrow t_3$	●	●	●	●	●
$t_3 \rightarrow t_4$	●	●	●	●	●
$t_4 \rightarrow t_5$	●	●	●	●	●

Timeseries Source

Distinct<K>

$t_0 \rightarrow t_1$	●	●	●	□	●
$t_1 \rightarrow t_2$	●	●	□	□	●
$t_2 \rightarrow t_3$	●	□	●	□	□
$t_3 \rightarrow t_4$	●	●	□	□	□
$t_4 \rightarrow t_5$	●	●	●	●	□

Option 2 Internal HeartBeat

$t_0 \rightarrow t_1$	●
$t_1 \rightarrow t_2$	●
$t_2 \rightarrow t_3$	●
$t_3 \rightarrow t_4$	●
$t_4 \rightarrow t_5$	●

GenerateSequence
heartbeats

Dynamic Key List

Fanout 1 -> k1,k2,k3

Flatten

$t_0 \rightarrow t_1$	●	●	●	●	●
$t_1 \rightarrow t_2$	●	●	●	●	●
$t_2 \rightarrow t_3$	●	●	●	●	●
$t_3 \rightarrow t_4$	●	●	●	●	●
$t_4 \rightarrow t_5$	●	●	●	●	●

Timeseries Source

Distinct<K>

$t_0 \rightarrow t_1$	●	●	●	□	●
$t_1 \rightarrow t_2$	●	●	□	□	●
$t_2 \rightarrow t_3$	●	□	●	□	□
$t_3 \rightarrow t_4$	●	●	□	□	□
$t_4 \rightarrow t_5$	●	●	●	●	□

Option 3 - Looping Timer

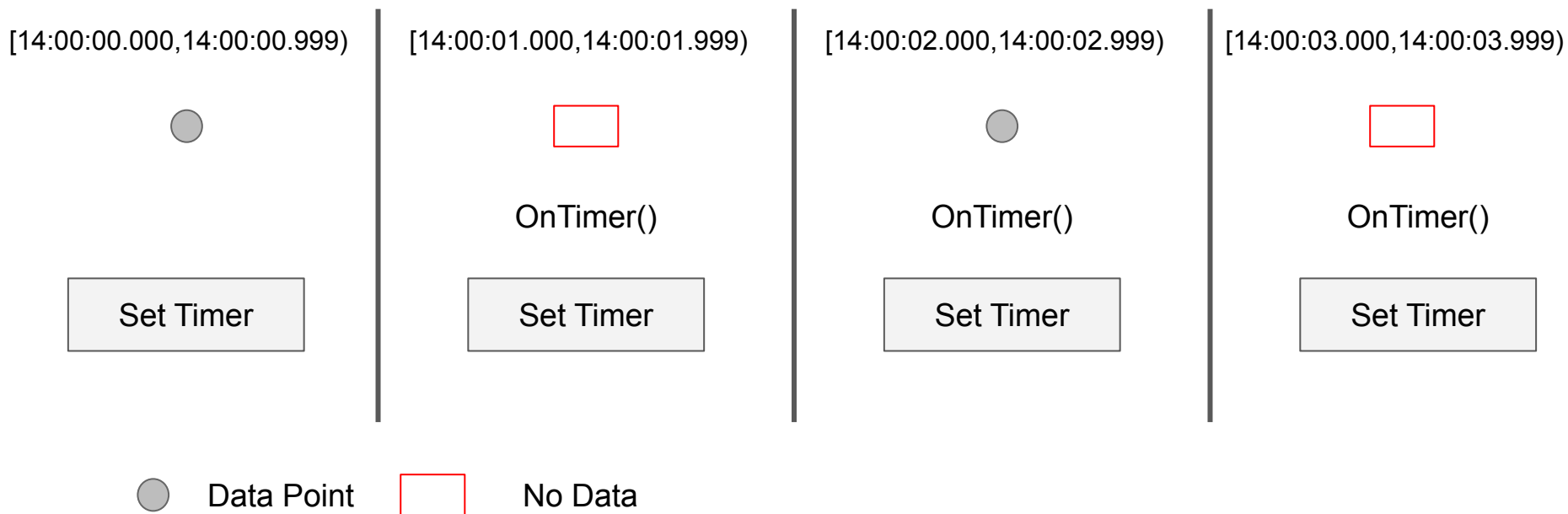
What is a Timer:

public interface Timer

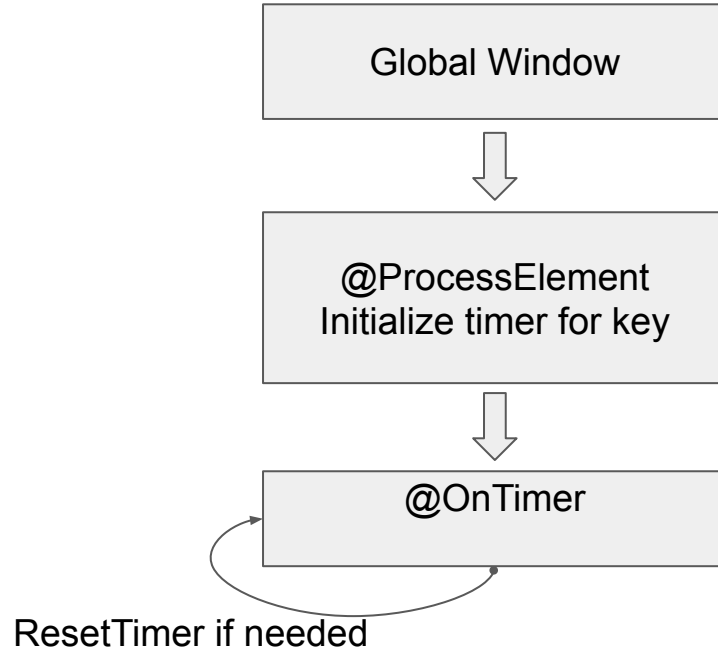
A timer for a specified time domain that can be set to register the desire for further processing at particular time in its specified time domain.

In my head... you get to set an alarm...

Looping Timers



Looping Timer



Looping timers - Advantages

- No need for a fanout operation.
- Timer will auto initiate when a key is seen for the first time.
- You can set a per key time to live.
 - Set a property with the key on the TTL of the looping timer.
 - Store the last seen timestamp value in state
 - Before rest of timers check if $\text{lastSeenTimestamp} + \text{TTL} < \text{OnTimer.timestamp}$
- OnTimer code is also where we can propagate last value seen

Gap filling

	TS-1	TS-2	TS-3	TS-4	TS-5
$t_0 \rightarrow t_1$	●	●	●	●	●
$t_1 \rightarrow t_2$	●	●	●	●	●
$t_2 \rightarrow t_3$	●	●	●	●	●
$t_3 \rightarrow t_4$	●	●	●	●	●
$t_4 \rightarrow t_5$	●	●	●	●	●

● Data Point ● Generated Data

Gap filling - with hold and propagation

	TS-1	TS-2	TS-3	TS-4	TS-5
$t_0 \rightarrow t_1$	●	●	●	●	●
$t_1 \rightarrow t_2$	●	●	●	●	●
$t_2 \rightarrow t_3$	●	●	●	●	●
$t_3 \rightarrow t_4$	●	●	●	●	●
$t_4 \rightarrow t_5$	●	●	●	●	●

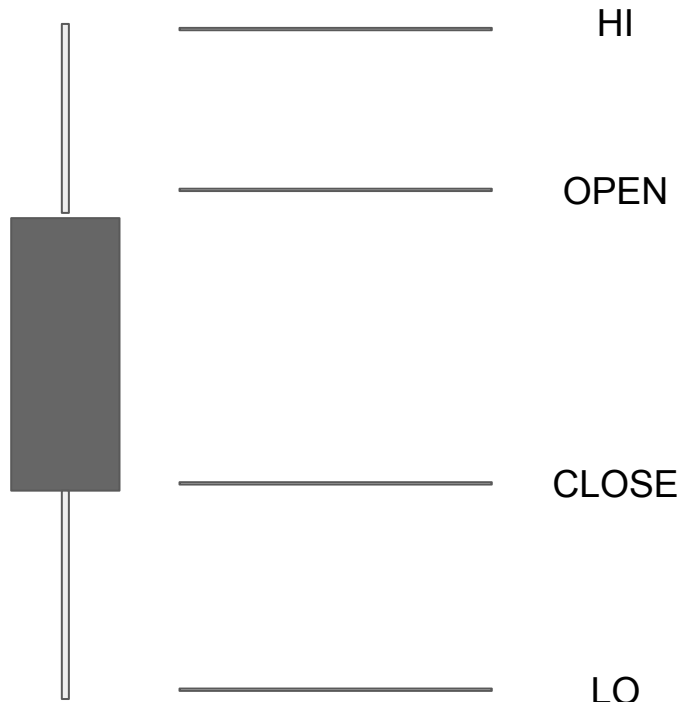


Data Point



Generated Data

Concrete Example - Building candlesticks



*<https://pixabay.com/photos/chart-trading-courses-forex-1905224/>

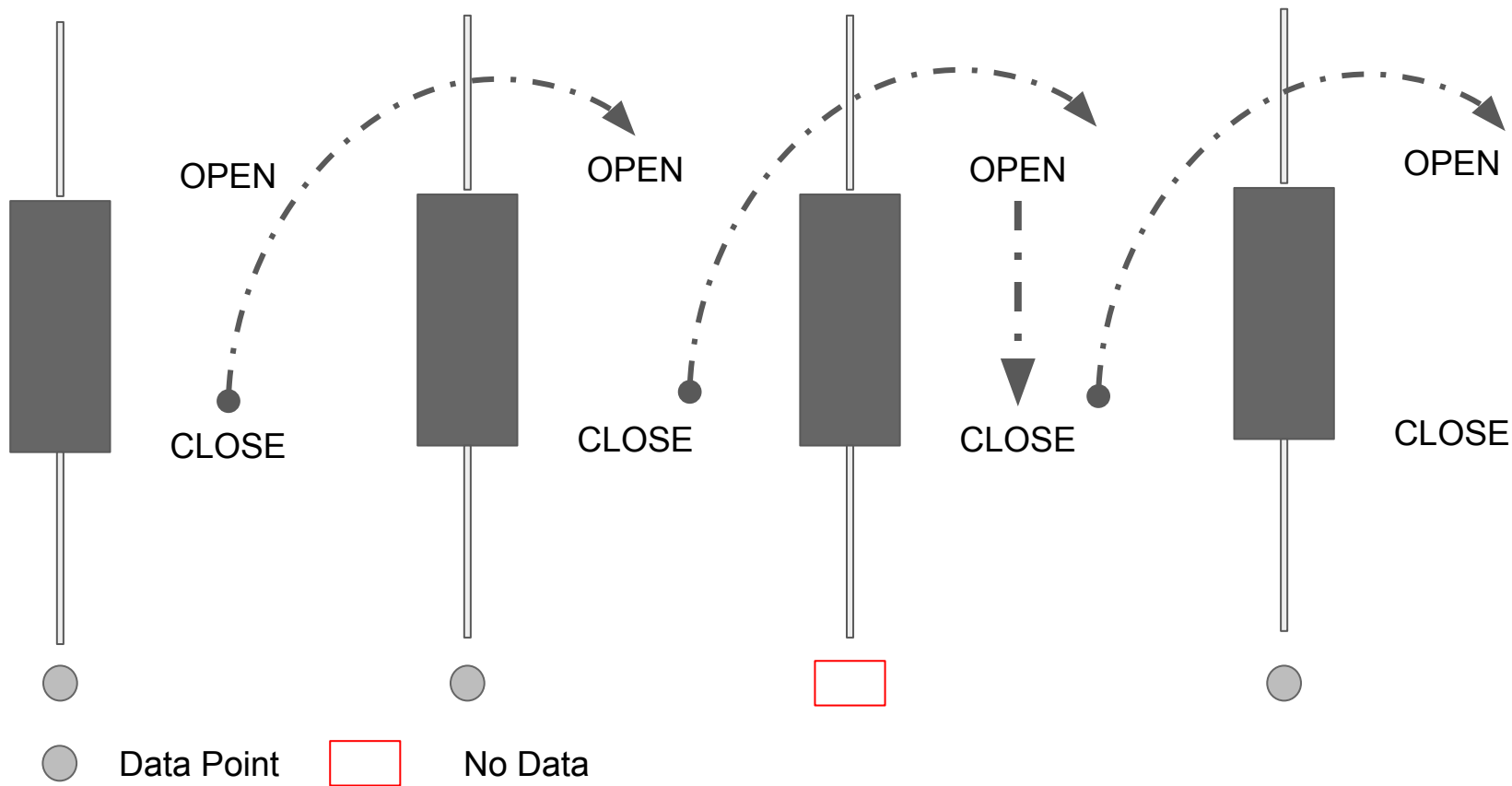
Example data

Time	Price Stock A	Price Stock B
14:00:00.000	1.54	0.98
14:00:00.500	1.53	-
14:00:01.200	-	0.99
14:00:02.200	1.53	-
14:00:03.100	1.52	1.00
14:00:03.300	1.52	-
14:00:03.400	-	0.98

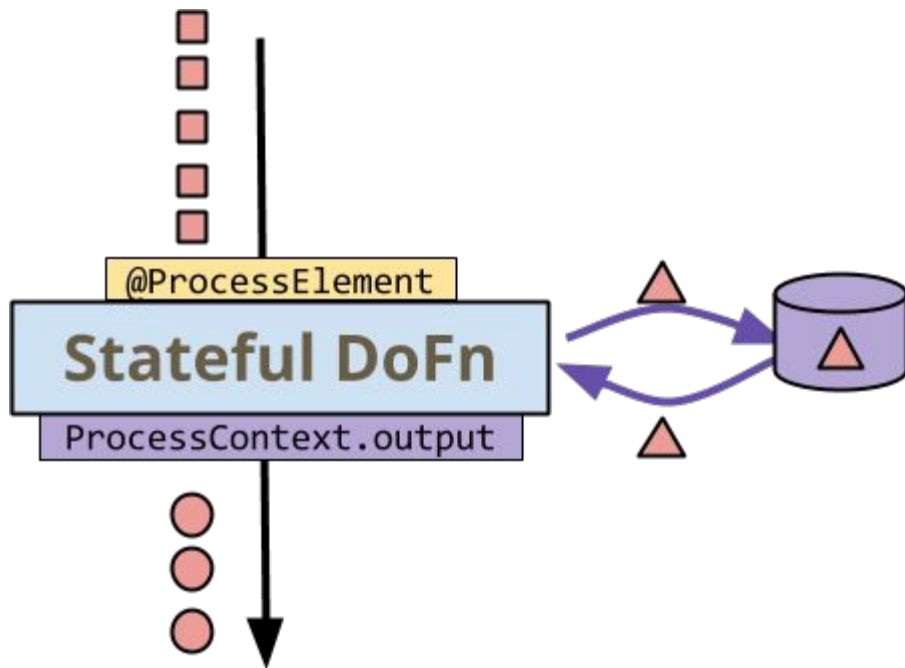
FixedWindow - GBK

Time	Price Stock A	Price Stock B
[14:00:00.000,14:00:00.999)	{1.54,1.53}	{0.98}
[14:00:01.000,14:00:01.999)	No Data	{0.99}
[14:00:02.000,14:00:02.999)	{1.53}	No Data
[14:00:03.000,14:00:03.999)	{1.52,1.52}	{1.00,0.98}

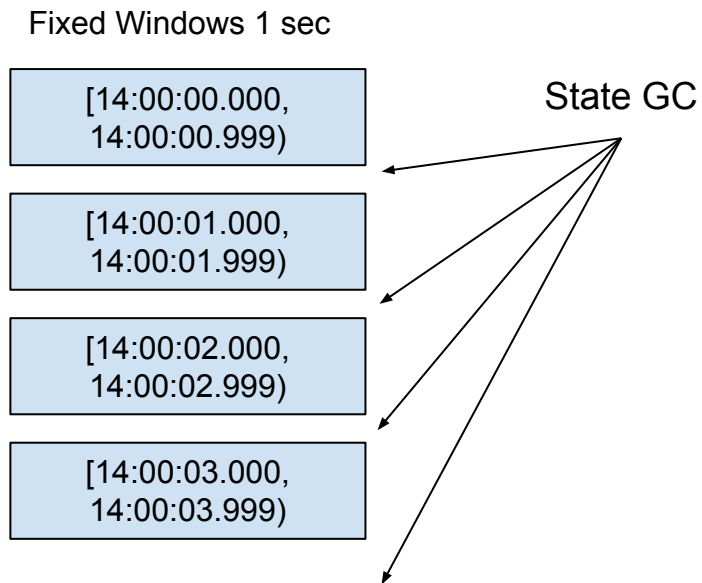
Concrete Example - Building candlesticks



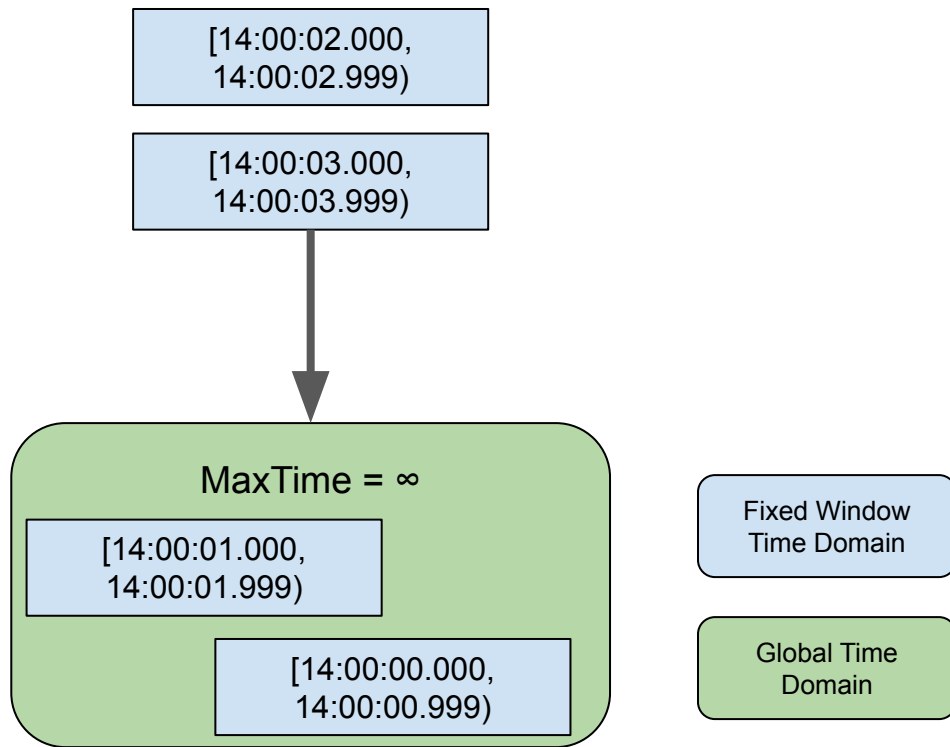
State API - Per Key-Window State



Transfer data across aggregations..



Global Windows ... no order!



Dealing with lack of order

.apply(Fixed Window)

T 0

T 1

T 2

T 3

.apply(Global Window)

BagState<Candle>

T
4

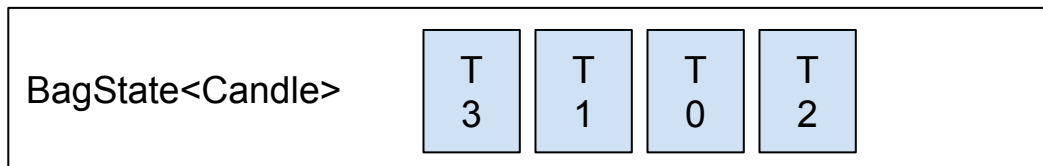
T
2

T
1

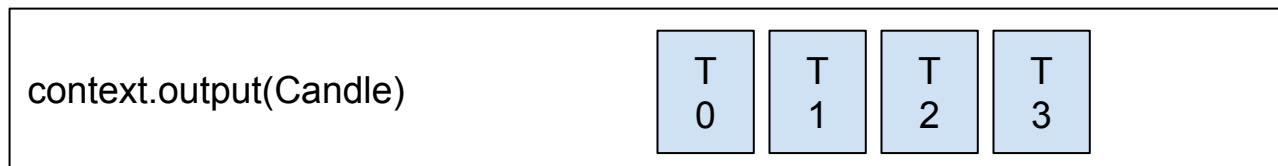
T
3

← No order

Dealing with lack of order



`OnTimer (T0), OnTimer(T1),OnTimer(...)` - Order and process



Interval list - Looping Timer + Global State

Time	Price Stock A	Price Stock B
[14:00:00.000,14:00:00.999)	{1.54,1.53}	{0.98}
[14:00:01.000,14:00:01.999)	Previous Close {1.53} No Data -	Previous Close {0.98} {0.99}
[14:00:02.000,14:00:02.999)	Previous Close {1.53} {1.53} -	Previous Close {0.99} No Data - 0.99
[14:00:03.000,14:00:03.999)	Previous Close {1.53} {1.52,1.52}	Previous Close {0.99} {1.00,0.98}

Final result and limitations

- Limitation 1
 - In memory sort of the key up to the OnTimer timestamp.
 - The accumulation step before is therefore important for data heavy loads. This will act as compression of the data before the global window step.
- Limitation 2
 - In order to GC the lists we need to do a full read-modify-write of whole list.
- Limitation 3
 - Late Data