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Project: Homework 1

1. A) $7n^2+20n-91 \in O(n^2)$, using C=8

Solve for k:

$$7n^{2} + 20n - 91 \le 8n^{2}$$

$$20n - 91 \le n^{2}$$

$$n^{2} - 20n + 91 \ge 0$$
Find roots:
$$n = \underbrace{(20 \pm (400 - 364)^{1/2})}_{2}$$

$$= \underbrace{20 \pm 6}_{2}$$

$$n = 13 \text{ or } n = 7$$

For all $n \ge 13$, the inequality holds.

Thus, the smallest value of K is 13.

1. B) $7n^5 + 18n^4 - 3n^3 + 6n^2 - 2n + 5 = O(n^5)$

For polynomials:

Remove negative terms

$$7n^5 + 18n^4 - 3n^3 + 6n^2 - 2n + 5 \le 7n^5 + 18n^4 + 6n^2 + 5$$
 (true for n≥0)

Raise to the highest power

$$7n^5+18n^4$$
 - $3n^3+6n^2-2n+5 \le 7n^5+18n^5+6n^5+5n^5$ (true for n≥1) $7n^5+18n^4$ - $3n^3+6n^2-2n+5 \le 36~n^5$

Thus, C = 36 and K = 1

2. Solution

2	· Proof by induction (Limit of Logarithm	ic ctions)
	Base Case: n=0:	
	$\lim_{x\to\infty} \frac{(\ln x)^2}{\ln x} = \lim_{x\to\infty} \frac{1}{\ln x} = 0$ which is true.	
	Inductive Step: Assume for n=k	
	$\lim_{x\to\infty} (\ln x)^{k} = 0$	
	Prove for n=K+1:	
	$\lim_{x \to \infty} \frac{(\ln x)^{k+1}}{x} = \lim_{x \to \infty} \ln x \cdot (\ln x)^{k}$	
	Apply L'Hopfital's Rule:	
	Desivative of (In x)k+1 95 (k+1)(In.x)k.	1
	Desivative of x 161.	
	$\lim_{\alpha \to \infty} \frac{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}}{(k+1)} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^k \cdot \frac{1}{\alpha}} = \frac{(k+1) \lim_{\alpha \to \infty} (1 + 1)}{(k+1)(\ln \alpha)^$	$nx)^{k} = 0$
	Thus, by induction, the statement	holds.

3. Pseudocode:

```
\begin{split} BINARYSEARCH(\,A[1,\,\ldots,\,N],\,target): \\ low &:= 0 \\ high := n-1 \\ while low &\leq high do: \\ mid := Floor (\,(low+high)\,/\,2\,) \\ if \,A[mid] &= target \,then: \\ return \,mid \\ else \,if \,A[mid] &< target \,then: \\ low &:= mid+1 \\ else: \\ high := mid+1 \\ return-1 \end{split}
```

4. For Algorithm 1

```
For loop a – iterates for n 

For loop b – iterates for n 

For loop c – iterates for n 

For loop d – iterates for n 

For loop e – iterates for n^4 - n^3 

While loop f – iterates for (log n) 

i) a: O(n), b: O(n), c: O(n), d: O(n), e: O(n^4-n^3) = O(n^4), f: O(logn)
```

ii) Total runtime = a [
$$(b*c*d) + (e*f)$$
]
= O $(n [(n*n*n*) + (n^4(logn)])$
= O $(n^4 + n^5 logn)$

iii) Simplified expression = $O(n^5(logn))$

For Algorithm 2

For loop a – iterates for n

For loop b – iterates for log m

While loop c – iterates for m

While loop d – iterates for log m

Call subfunction – m⁵ log m

For loop e – iterates for 3m

For loop f – iterates for $7n^2$

For loop g – iterates for constant time so 1

- i) a: O(n), b: O(logm), c: O(m), d: O(logm), Function call: m^5 logm, e: O(3m)= O(m), f: O(7 n^2)= O(n^2), g: 1
- ii) Total runtime = a [(b*c*d) + Function call] + (e*f*g) = O (n [(logm * m * logm) + m⁵logm] + [m * n² * 1]) = O (mn(logm)² + m⁵n(logm) + mn²)
- iii) Simplified expression = $O(m^5n(log m))$