Q1. What is the time complexity of below code and how?

```
void fun (int n)
  {
    int j=1, i=0;
    while (i<n) {
      i+=j;
      j++;
    }
  }
```

↳

$j=1$	$i=1$	
$j=2$	$i=\cancel{2}\ 1+2$	m-level
$j=3$	$i=1+2+3$	

for (i)

$\because 1+2+3+\ldots+<n$

$\because 1+2+3+m<n$

$\because \frac{m(m+1)}{2}<n$

$m \approx \sqrt{n}$

By summation method

$\Rightarrow \sum_{i=1}^{m} 1 \Rightarrow 1+1+\ldots+\sqrt{n}$ times

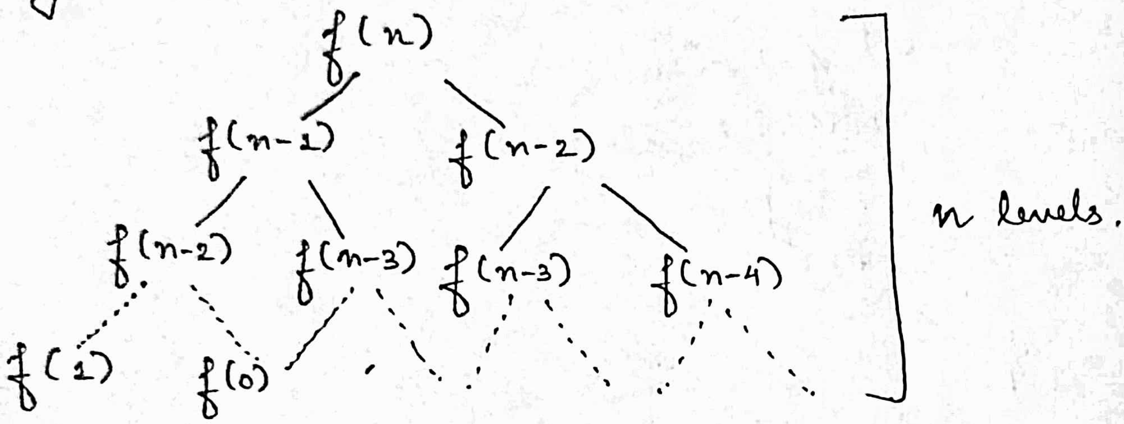
$\boxed{T(n)=\sqrt{n}}$ → Ans

Q2. Write recurrence relation for function that prints Fibanacci series. Solve it to get the time complexity. What will be the space complexity and why?

→ For Fibonacci Series

$f(n) = f(n-1) + f(n-2)$ $\qquad f(0) = 0$, $f(1) = 1$

By forming a tree



∵ At every function call we get 2 function calls

∴ for n levels

We have = $2 \times 2 \ldots n$ times

$\therefore \boxed{T(n) = 2^n}$

MAXIMUM SPACE

Considering Recursive Stack :

No. of calls maximum = n

For each call we have space complexity $O(1)$

$\boxed{\therefore T(n) = O(n)}$

Without considering Recursive Stack :

each call we have time complexity $O(1)$

$\therefore \boxed{T(n) = O(1)}$

Q3.

Write programs which have complexity:
$n(\log n)$, n^3, $\log(\log n)$

1) $n \log n$ → Quick Sort

```
void quicksort (int arr[], int low, int high)
{
    if (low < high)
    {
        int pi = partition (arr, low, high);
        quicksort (arr, low, pi-1);
        quicksort (arr, pi+1, high);
    }
}
int partition (int arr[], int low, int high)
{
    int pivot = arr[high];
    int i = (low -1);
    for ( int j= low ; j<= high -1; j++)
    {
        if (arr[i] < pivot)
        {
            i++ ;
            swap (&arr[i], &arr[j]);
        }
    }
    swap (& arr [i+1], & arr [high]);
    return (i+1);
}
```

2) n^3 → Multiplication of 2 square matrix

```
for (i=0; i<r1; i++) {
    for(j=0; j<c2; j++)
        for (k=0; k<c1; k++)
        {
            res[i][j]+ = a[i][k] * b[k][j];
        }
```

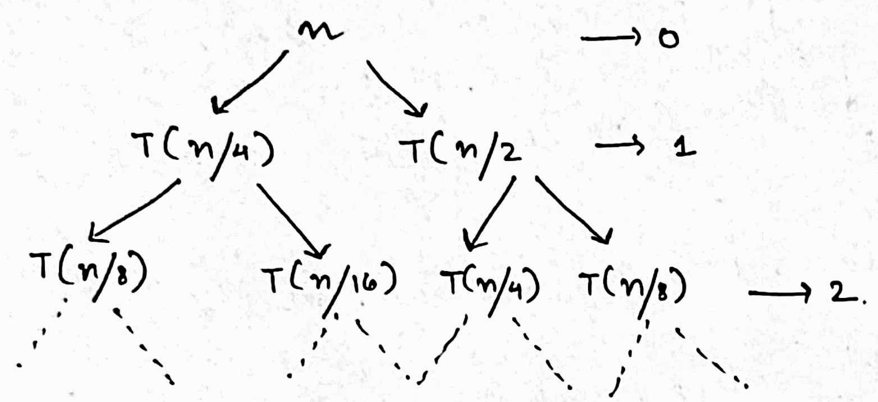
3) **log(log n)**

```
for (i=2; i<n; i=i*i)
  {
     count ++;
  }
```

Q4. Solve the following recurrence relation

$$T(n) = T(n/4) + T(n/2) + cn^2$$

→



At level

$0 \rightarrow Cn^2$

$1 \rightarrow \frac{n^2}{4^2} + \frac{n^2}{2^2} = \frac{C5n^2}{16}$

$2 \rightarrow \frac{n^2}{8^2} + \frac{n^2}{16^2} + \frac{n^2}{4^2} + \frac{n^2}{8^2} = \left(\frac{5}{16}\right)^2 n^2 c$

⋮

max level $= \frac{n}{2^k} = 1$

$= k = \log_2 n$

$T(n) = c\{n^2 + (5/16)n^2 + (5/16)^2 n^2 + \dots + (5/16)^{\log n} n^2)$

$T(n) = Cn^2\left[1 + \left(\frac{5}{16}\right) + \left(\frac{5}{16}\right)^2 + \dots + \left(\frac{5}{16}\right)^{\log n}\right]$

$T(n) = Cn^2 \times 1 \times \left(\frac{1 - (5/16)^{\log n}}{1 - (5/16)}\right)$

$T(n) = Cn^2 \times \frac{11}{5} \times \left(1 - \left(\frac{5}{16}\right)^{\log n}\right)$

$T(n) = O(n^2 c)$

$\underline{O(Cn^2)}$ Ans.

Q5. What is the time complexity of following fun() ?

```
int fun (int n) {
  for (int i=1; i<=n; i++) {
    for (int j=1; j<n; j+=i) {
      // Some O(1) task
} }}
```

→ for

i	j
1	1
2	$1+3+5$
3	$1+4+7$
⋮	$1+5+9$
n	

$j = (n-1)/i$ times

$$\sum_{i=1}^{n} \frac{(n-1)}{i}$$

$$\therefore T(n) = \frac{(n-1)}{1} + \frac{(n-1)}{2} + \frac{(n-1)}{3} + \dots + \frac{(n-1)}{n}$$

$$T(n) = n\left[1 + 1/2 + 1/3 + \dots + 1/n\right] - 1 \times \left[1 + 1/2 + 1/3 + \dots + 1/n\right]$$

$$= n \log n - \log n$$

$$\underline{T(n) = O(n \log n)} \rightarrow \text{Ans.}$$

Q6. What should be time complexity of

```
for (int i=2; i<=n; i=pow(i,k))
{
    // Some O(1)
}
```

where k is a constant

→ for

i
2^1
2^k
2^{k^2}
2^{k^3}
⋮
2^{k^m}

where

$$2^{k^m} <= n$$

$$k^m = \log_2 n$$

$$m = \log_k \log_2 n$$

$$\therefore \sum_{i=1}^{m} 1$$

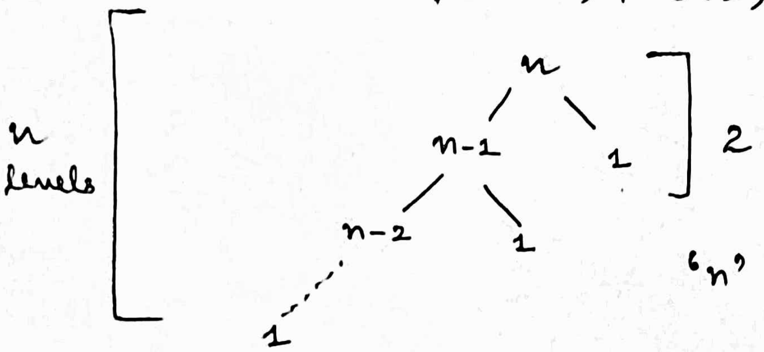
$1 + 1 + 1 \dots m$ times

$$\underline{T(n) = O(\log_k \log n)} \rightarrow \text{Ans.}$$

Q7. Write a recurrence relation when quick sort repeatedly divides array into 2 parts of 99% and 1%. Derive time complexity in this case. Show the recurrence tree while deriving time complexity & find difference in heights of both extreme parts. What do you understand by this analysis?

→ Given algorithm divides array in 99% and 1% part

$\therefore T(n) = T(n-1) + O(1)$



'n' work is done at each level

$$T(n) = \left(T(n-1) + T(n-2) + \dots + T(1) + O(1)\right) \times n$$

$$= n \times n$$

$$\therefore \boxed{T(n) = O(n^2)}$$

lowest height $= 2$

highest height $= n$

$$\boxed{\therefore \text{difference} = n-2} \quad n > 1$$

The given algorithm produces linear result

Q8) Arrange following in increasing order of rate of growth : (7)

a) $n, n!, \log n, \log\log n$, root (n), $\log(n!)$, $n\log n$, $\log^2(n)$, 2^n, 2^{2^n}, 4^n, n^2, 100

→ $100 < \log\log n < \log n < (\log n)^2 < \sqrt{n} < n < n\log n < \log(n!) < n^2 < 2^n < 4^n < 2^{2^n}$

b) $2(2^n)$, $4n$, $2n$, 1, $\log(n)$, $\log(\log(n))$, $\sqrt{\log(n)}$, $\log 2n$, $2\log(n)$, n, $\log(n!)$, $n!$, $n2$, $n\log(n)$

→ $1 < \log\log n < \sqrt{\log n} < \log n < \log 2n < 2\log n < n < n\log n$
$< 2n < 4n < \log(n!) < n2 < n! < 2^{2^n}$

c) 8^{2n}, $\log_2(n)$, $n\log_6(n)$, $n\log_2(n)$, $\log(n!)$, $n!$, $\log_8(n)$, 96, $8n2$, $7n3$, $5n$

→ $96 < \log_8 n < \log_2 n < 5n < n\log_6(n) < n\log_2 n < \log(n!) < 8n^2 <$
$7n^3 < n! < 8^{2n}$